

P802.1Qca – D0.2

Editor's Notes

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D0.2 Clauses

- › 1 Overview
 - 1.3 Introduction
 - › 2 Normative references
 - › 3 Definitions
 - › 4 Abbreviations
 - › 27 *Shortest Path Bridging*
 - › 28 *ISIS-SPB Link State Protocol*
 - › 45 Path Control and Reservation
 - use of SPB TLVs
 - 45.1–45.5: see next slide
 - › Bibliography
- } minimal changes

D0.2 Clause 45



> 45.1 Explicit and constrained paths

- 45.1.1 Constrained paths: [SRLG](#)
- 45.1.2 Explicit paths
- 45.1.3 Point-to-point explicit path
- 45.1.4 Explicit tree
- ~~45.1.5 Notification on path status~~

feedback got:

problem not aimed to be solved by Qca

> 45.2 Reservation

- 45.2.1 Basic reservation by IS-IS: [conflict resolution](#)
- [45.2.2 Reservation support for SRP](#)

> 45.3 Redundancy

- 45.3.1 Loop Free Alternates
- [45.3.2 Maximum disjoint paths](#)

> 45.4 Parameters for time synchronization: [updates](#)

> 45.5 Parameters for time scheduling

Use of ISIS-SPB TLVs

SPB Base VLAN Identifiers sub-TLV



- > It will be made clear by the conformance statements which ISIS-SPB TLVs have to be supported for 802.1Qca
- > SPB Base VLAN Identifiers sub-TLV
 - Associates an ECT Algorithm to the VLAN's Base VID
 - Dedicated ECT Algorithm for Qca: Explicit Path ECT Algorithm

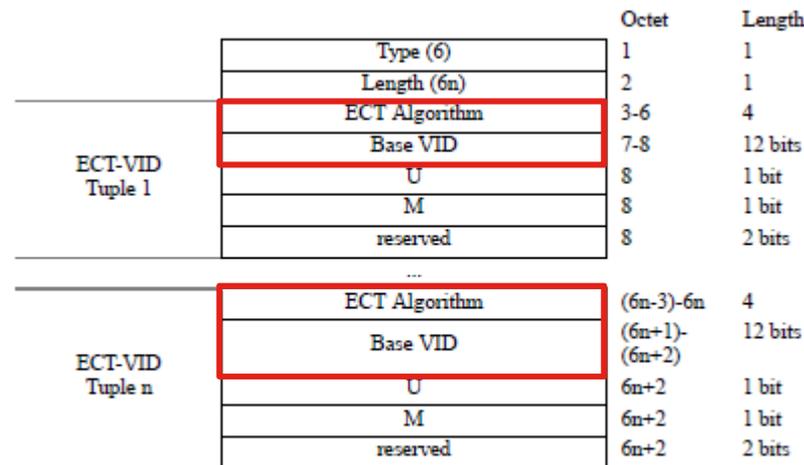


Figure 28-5—SPB Base VLAN-Identifiers sub-TLV

Use of ISIS-SPB TLVs

SPBM Service Identifier and Unicast Address sub-TLV



- › SPBM Service Identifier and Unicast Address sub-TLV
 - Used for the distribution of Individual Addresses within the Domain for VLANs allocated to the SPBM MSTID

		Octet	Length
	Type (3)	1	1
	Length	2	1
	B-MAC Address	3-8	6
	reserved	9	4 bits
	Base VID	9-10	12 bits
I-SID Tuple 1	T	11	1 bit
	R	11	1 bit
	reserved	11	6 bits
	I-SID	12-14	3
	...		
I-SID Tuple n	T	(4n+7)	1 bit
	R	(4n+7)	1 bit
	reserved	(4n+7)	6 bits
	I-SID	(4n+8)- (4n+10)	3

Figure 28-11—SPBM Service Identifier and Unicast Address sub-TLV

Use of ISIS-SPB TLVs

SPBM Group MAC Addresses?



- › Question: Do we want to use locally administered source specific Group MAC addresses on the explicit paths/trees as well in the non-learning VLANs?

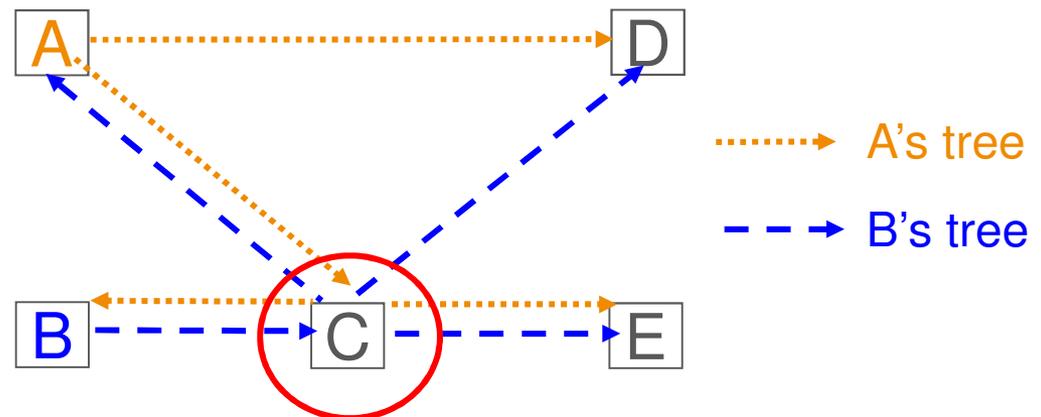
- › Background: Why do we use (S,G) multicast in SPB:

- Different sources use different trees

- A and B are Sources in the same Group

- C should ONLY forward B's frames to D but NOT A's frames

- (*,G) does not work, (S,G) is needed instead



Use of ISIS-SPB TLVs

SPB Instance sub-TLV



- > If Group MAC is source specific



then SPSourceID has to be distributed in the Domain

- > SPSourceID is distributed in the SPB Instance sub-TLV

- > I-SID → Base VID → EP ECT Algorithm

		Octet	Length
	Type (1)	1	1
	Length	2	1
	CIST Root Identifier	3-10	8
	CIST External Root Path Cost	11-14	4
	Bridge Priority	15-16	2
	reserved	17-18	11 bits
	V	18	1 bit
	SPSourceID	18-20	20 bits
	Number of Trees	21	1
VLAN ID Tuple 1	U	22	1 bit
	M	22	1 bit
	A	22	1 bit
	reserved	22	5 bits
	ECT Algorithm	23-26	4
	Base VID	27-28	12 bits
	SPVID	28-29	12 bits
	...		
VLAN ID Tuple n	U	8n+14	1 bit
	M	8n+14	1 bit
	A	8n+14	1 bit
	reserved	8n+14	5 bits
	ECT Algorithm	(8n+15)- (8n+18)	4
	Base VID	(8n+19)- (8n+20)	12 bits
	SPVID	(8n+20)- (8n+21)	12 bits

Figure 28-6—SPB Instance sub-TLV



Use of ISIS-SPB TLVs

SPBM Service Identifier and Unicast Address sub-TLV for multicast

- › If a single shared tree is used for a Group, then (*,G) can be applied
- › Group MAC could be then distributed by the SPBM Service Identifier and Unicast Address sub-TLV

- › Should we allow it?
- › Should we rename the sub-TLV to *SPBM Service Identifier and MAC Address sub-TLV*?

		Octet	Length
	Type (3)	1	1
	Length	2	1
	B-MAC Address	3-8	6
	reserved	9	4 bits
	Base VID	9-10	12 bits
I-SID Tuple 1	T	11	1 bit
	R	11	1 bit
	reserved	11	6 bits
	I-SID	12-14	3
	...		
I-SID Tuple n	T	(4n+7)	1 bit
	R	(4n+7)	1 bit
	reserved	(4n+7)	6 bits
	I-SID	(4n+8)-(4n+10)	3

Figure 28-11—SPBM Service Identifier and Unicast Address sub-TLV

Use of ISIS-SPB TLVs

SPBV MAC Address sub-TLV



- › SPBV MAC Address sub-TLV
 - Can be used for both Individual and Group MAC Address distribution for VLANs allocated to the SPBV MSTID

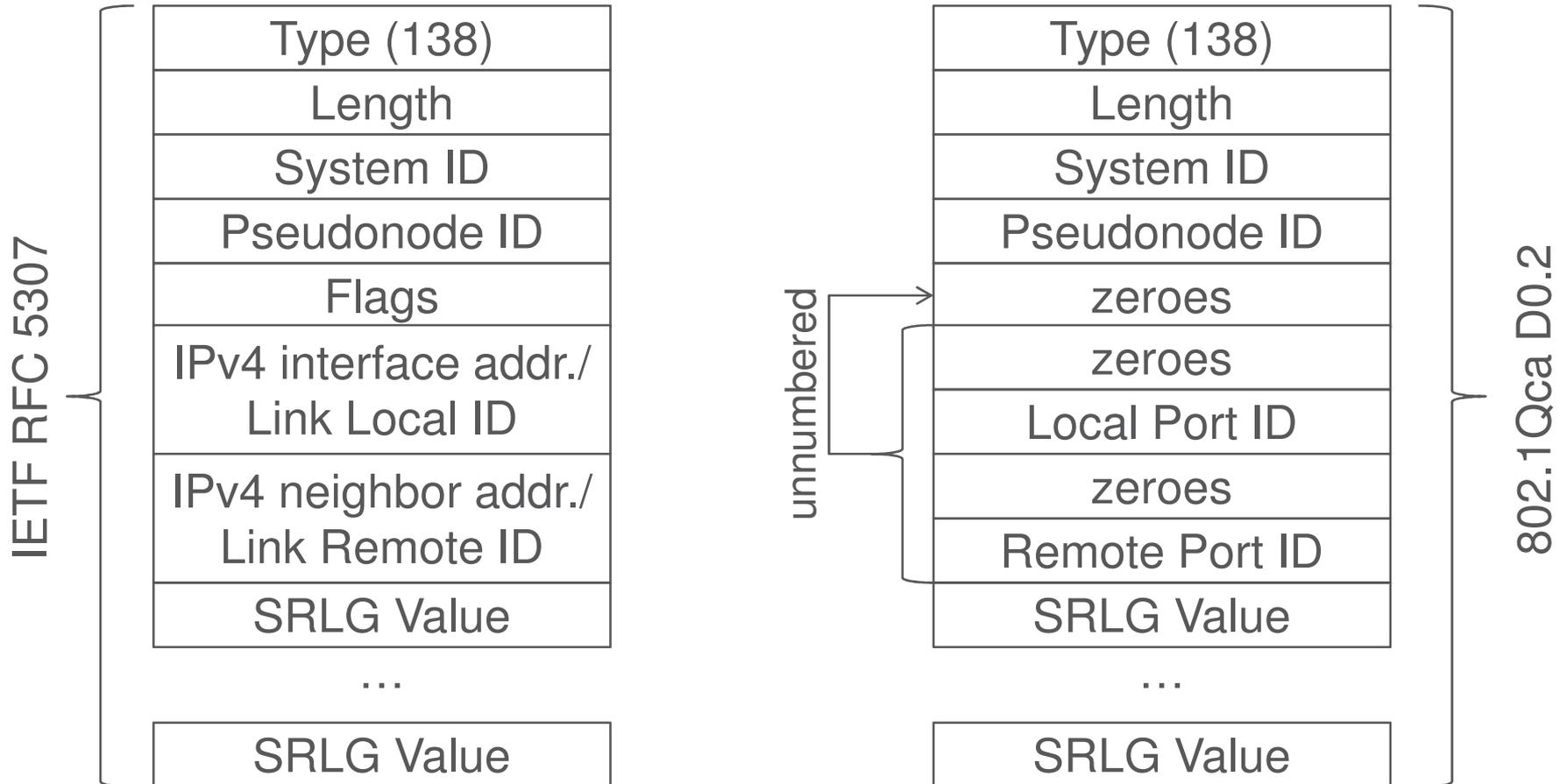
		Octet	Length
MAC Address Tuple 1	Type (4)	1	1
	Length	2	1
	reserved	3	2 bits
	S-R	3	2 bits
	SPVID	3-4	12 bits
	T	5	1 bit
	R	5	1 bit
	reserved	5	6 bits
	MAC Address	6-11	6
	...		
MAC Address Tuple n	T	(7n-2)	1 bit
	R	(7n-2)	1 bit
	reserved	(7n-2)	6 bits
	MAC Address	(7n-1)- (7n+4)	6

Figure 28-10—SPBV MAC Address sub-TLV

45.1.1 Constrained Paths Shared Link Risk Group (SRLG)



› Do we want to use SRLG?



45.2.1 Basic reservation by IS-IS



- › **New:** Conflict resolution
 - Similar to SPVID allocation conflict resolution, see subclause 27.10 and page 6 of <http://www.ieee802.org/1/files/public/docs2009/aq-farkas-allocation-0518-v02.pdf>
- › Existing reservations are not touched
- › Conflict should not occur in case of a single PCE
- › But we can have a safeguard: we can e.g. tie-break on
 - the least number of hops or nodes
 - the lowest Path ID
 - the smallest Reservation Value
 - the numerically least System ID of the originator of the LSP
 - the numerically least LSP sequence number
- › Everybody gets every LSP
 - Winner is known
 - Loser's reservation is rolled back

45.2.1 IS-IS Support for SRP



- › **How should SRP and IS-IS interwork?**

- › See discussion in

- <http://www.ieee802.org/1/files/public/docs2013/ca-farkas-SRP-ISIS-interworking-0713-v01.pdf>

- › Initial sub-TLV for now

	Octet	Length
Type (???)	1	1
Length (16)	2	1
Stream DA	3-8	6
Stream ID	9-14	6
TSpec	14-17	4
Priority	18	4 bits
Ranking	18	4 bits

Figure 45-9—Stream sub-TLV

45.3.2 Disjoint Paths

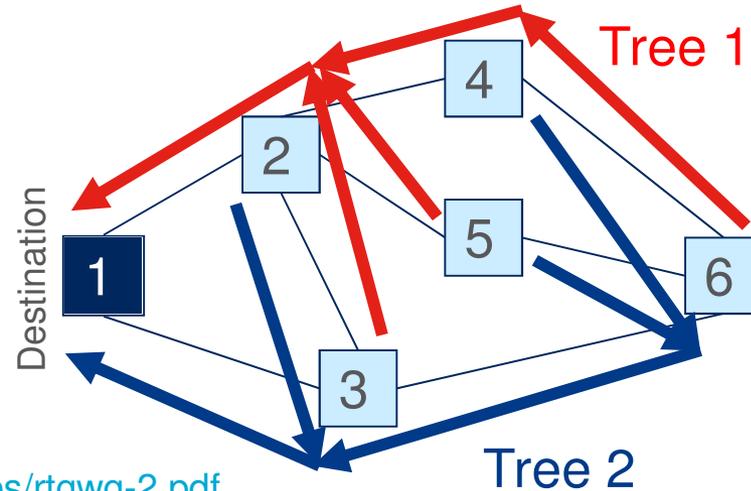


- › Explicit Paths are not updated
 - PCE can compute the wanted paths, e.g. maximally disjoint paths
 - IS-IS installs them and does not touch them any more
 - (subclause 45.1)
- › **Questions:**
- › Do we want to leverage IS-IS automation features?
 - Define standard algorithm for disjoint path?
 - Allow careful update?
- › Maximally redundant trees?
 - Do we want to go for Maximal Redundant Trees (MRT) as well or just point-to-to-point paths?
 - › MRT algorithm: <https://datatracker.ietf.org/doc/draft-ietf-rtgwg-mrt-frr-architecture/>
 - › see next slide for MRT examples

MRT examples

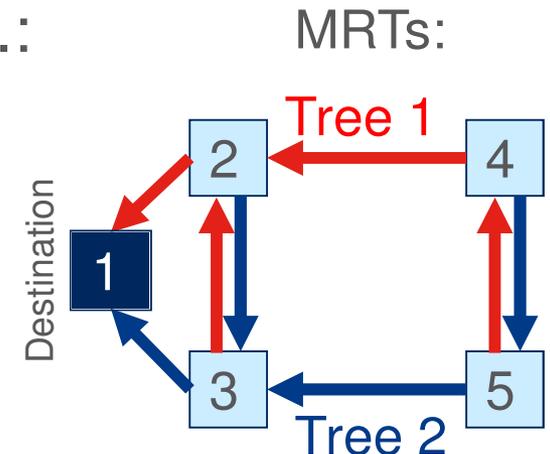
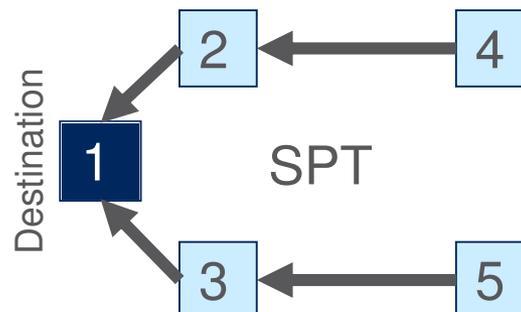
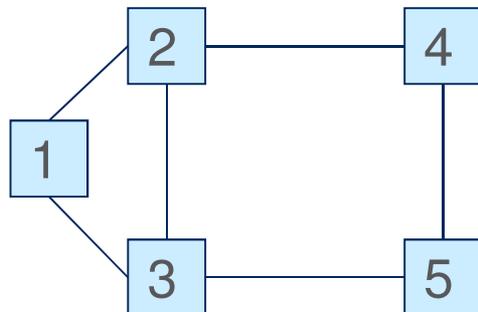


- › An example for MRTs:



<http://www.ietf.org/proceedings/81/slides/rtgw-2.pdf>

- › Note: MRTs may not collide with SPT, e.g.:



45.4 Time Synchronization Parameters



› Updated based on the input got from Geoffrey Garner, Thanks!

› Each clock sends

	Octet
Type (???)	1
Length ()	2
priority1	3
clockClass	
clockAccuracy	
offsetScaledLogVariance	
priority2	
clockIdentity	

Figure 45-10—Clock Attributes sub-TLV

› Grandmasters send

	Octet
Type (???)	1
Length ()	2
grandmasterIdentity	
timeSource	
currentUtcOffset	
reserved	
leap59	
leap61	
currentUtcOffsetValid	
ptpTimescale	
timeTraceable	
frequencyTraceable	

Figure 45-11—Grandmaster sub-TLV