Deadline Application Model & Stream Class Based Scheduling Model

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#### **Interoperability Introduction**

**Interoperability** issues for co-existence and co-operation:

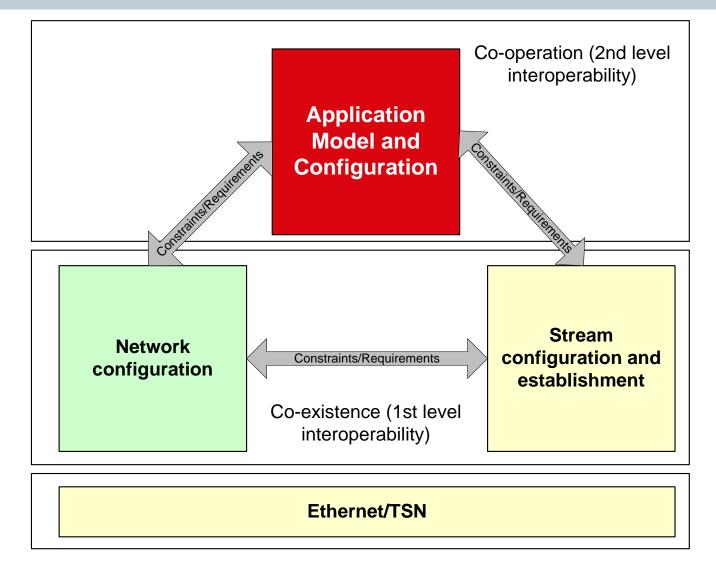
- Application model and configuration
- Network configuration
- Stream configuration and establishment

### **Application Model and Configuration**

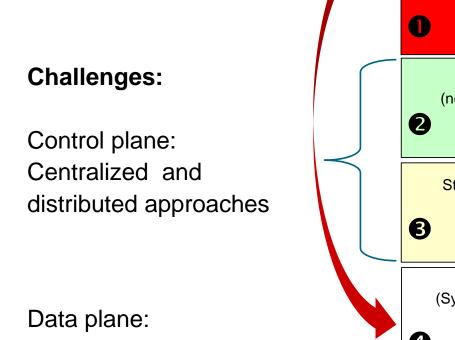
is essential for co-operation and impacts Network configuration and Stream configuration and establishment

and vice versa.

Application Model and Configuration Constraints/Requirements are identical in any location within the production hierarchy (machine, cell, ...).



Scope and Challenges of Interoperability

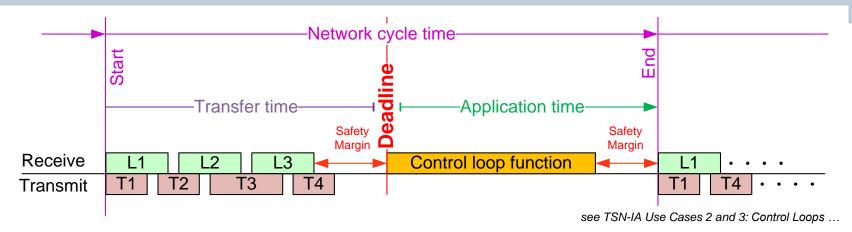


Scope of Vendor, Customer **Application Model and Configuration** (Application Profile, Communication Profile Vendor Model, ...) Middleware (Fieldbus layer, OPC UA, ...) **Network Configuration** (netconf, restconf/YANG, SNMP/MIB, ...) Stream Configuration and Establishment (RAP, restconf, ...) Goals: Ethernet/TSN (Synchronization, Neighborhood detection, bounded latency communication, ...) 4 Application specific requirements

Scope of IEC61158/IEC61784 IEC62541, ... Goal: co-operation as 2<sup>nd</sup> level of interoperability Scope of IEEE 802 and **IEC/IEEE 60802 TSN-IA Profile** 

- co-existence as 1<sup>st</sup> level of interoperability
- Common hardware requirements

#### **Deadline Application Model**



- Applications: include 1...n Talkers and 1...m Listeners
- Deadline: Application defined fixed reference point within network cycle
  Network cycle time is e.g. multiple of 31,25 µs

 Transmit: Application's network access is synchronized to network cycle Start Talker frames are assigned to send cycles, which are multiples of the network cycle time Transmission of frames as a convoy starts at network cycle Start:

1<sup>st</sup> isochronous cyclic, 2<sup>nd</sup> cyclic real-time

Talker frame sequence inside a traffic class convoy follows "longest path first"

 Receive: Listeners expect Talker frames before assigned deadline, in case of isochronous cyclic: aligned to the network cycle / deadline

Stream Class based Scheduling

- → Deadline Application Model is independent of the realizations of areas ② and ③.
- → Deadline Application Model does not imply stream specific timing, which avoids per-stream schedule planning.
- → Stream class based scheduling suffices the Deadline Application Model:
  - Assign streams to traffic classes
  - Transmit streams as convoy at send cycle start
  - 1<sup>st</sup> isochronous cyclic, 2<sup>nd</sup> cyclic real-time
  - Sequence inside a traffic class convoy follows "longest path first"

Stream Class based Scheduling in Network / Bridges

- Stream Destination MAC Address based forwarding on path
- Queue based transmission selection:
  - Strict Priority together with e.g.
    - Preemption, and/or
    - Enhancements for scheduled Traffic (e.g. TAS window per Traffic Class)

#### Stream Class based Scheduling: Scalability, Strengths

- Ad-hoc stream establishment and removal of streams without any impact on running streams
- No fragmentation of transmission interval independent from order of stream establishment
- Perfectly fits to seamless redundancy (different latencies on asymmetric disjoint paths)
- Stream class based scheduling allows:
  - lowest bridge latency
  - best bandwidth utilization
- Bridges do not require synchronization if TAS is not used
- Simplifies link speed transitions if TAS is not used

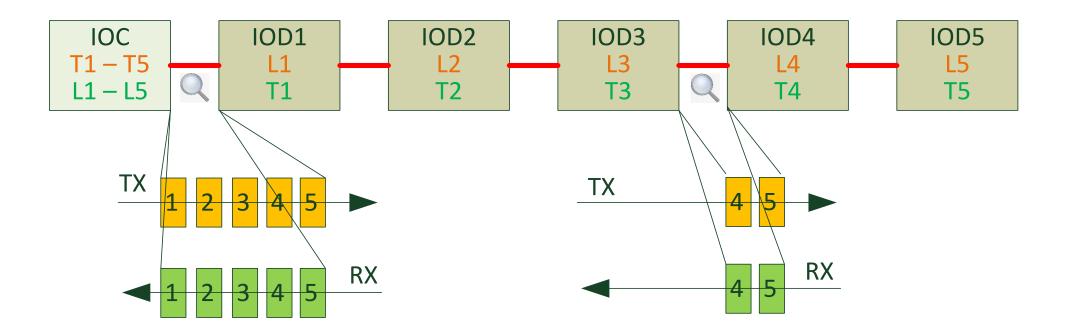
**Stream Class based Scheduling: Lessons learned** 

- $\rightarrow$  No benefit from frame-based scheduling combined with deadline model
- $\rightarrow$  Allows decoupling of application from frame scheduling
- → One application starting time at deadline, independent e.g. from seamless transmission
- $\rightarrow$  In use for more than 30 years in Industrial Automation
- → Integrated in hardware of technology modules (e.g. drives, counters, positioning, IOs, ...)
- → Fulfills all TSN-IA use cases including required topologies and network cycle times,

except "Drives without common application cycle"

**Stream Class based Scheduling: Example** 

- IOC Talker sends frames in convoy using "longest path first" strategy
- Each IOD Talker sends frame at begin of network cycle



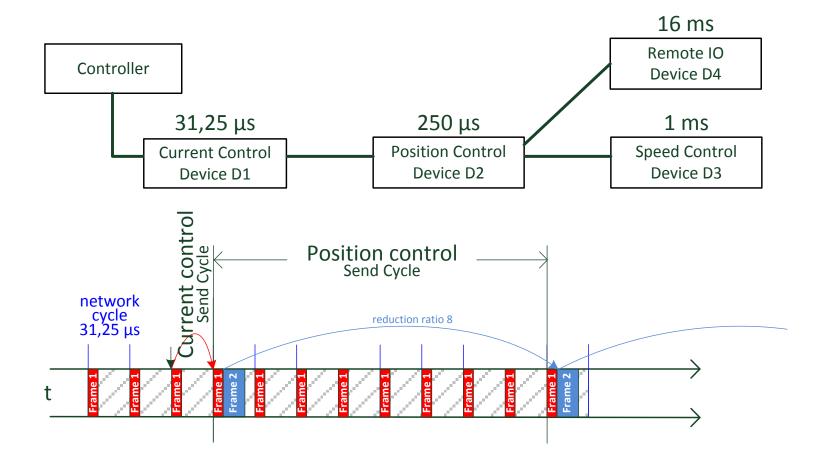
## **Thank You!**



# **Questions?**

#### IEC/IEEE 60802 JP TSN Industrial Profile: Stream Class based Scheduling

"Talker frames are assigned to send cycles, which are multiples of the network cycle time"



#### IEC/IEEE 60802 JP TSN Industrial Profile: Stream Class based Scheduling

"Transmission of frames as a convoy starts at network cycle Start"

