## P802.1CQ/D0.6 Preview v03

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thanks to Antonio de la Oliva for review and constructive comments

## Summary

- P802.1CQ/D0.5 was reviewed in TG Ballot.
- Comment resolution was completed in November.
- Address Blocks were introduced for address claiming
- In March, Editor presented "Block Address Registration and Claiming (BARC)"
- cq-marks-BARC-0321-v00.pdf
- address blocks used for registrar-managed addresses as well
- Address Registration and Claiming (ARC)
- address blocks, and also claiming address ranges using MAAP
- presented also to IEEE 1722 Working Group
- Main issues raised in March concerned VLAN operation
- This contribution previews P802.1CQ/D0.6
- refinements and details since March presentation
- discussion on improved VLAN support
- v01 presented to TSN at May 802.1 Interim; this version (v03) adds detail


## BARC assigns MAC Addresses in Address Blocks

1) Address Blocks (ABs) are sets of local addresses.
2) An AB includes equal-sized unicast and multicast address subblocks.
3) No BARC address falls within more than one AB.
4) An Address Block Designation (ABD) is a CABA or a RABI.
5) Claimable AB Address (CABA) is claimable by a Claimant without using a Registrar.

- identifies Claimable Address Blocks (CABs) holding Claimable Addresses (CAs)
- CABA is a multicast MAC address, not in any AB, and used as a DA.

6) RABI

- identifies a Registrable Address Block (RAB) holding Registrable Addresses (RAs)
- Registrable Address Block Indicators (RABIs): held in inventory of a Registrar
- may be assigned to Claimants
- may be claimed by Registrants

7) A large set of Temporary Unicast Addresses (TUAs) is specified

- useful for initial discovery by Claimant lacking a unicast address


## MAC Address Categorization

| determinable via inspection: | Expanded name | I/G | indicates, by inspection |
| :---: | :---: | :---: | :---: |
| CA <br> [ICA=unicast <br> GCA=multicast] | claimable address, in claimable address block (CAB) | U,M | CABA, CAB Size, CAB (including all other CAs in CAB) |
| CABA | CAB Address | M | CAB Size, CAB |
| RA [IRA=unicast GRA=multicast] | registrable address, in registrable address block (RAB) | U,M | BABI Size <br> [Basic ABI Size] |
| TUA | temporary unicast address | U | note: $\sim 6.9 \mathrm{E} 10$ to choose among |

## Address Block Designation (ABD) Categorization

| ABD | Address Block Designation <br> (CABA or RABI) | not an <br> address | AB <br> (including all RAs in AB) |
| :--- | :--- | :---: | :--- |
| RABI | RAB Identifier | not an <br> address | RAB Size, BABI Size, RAB, <br> MABI Size [Multiple ABI Size] |
| CABA | CAB address | $M$ | CAB Size, CAB |

## BARC MAC Address Structure

| N11 | $\mathrm{r} \mathrm{i}_{\mathrm{j}} \mathrm{j}_{\mathrm{k}}$ |
| :---: | :---: |
| N10 | 111 m |
| N9 |  |
| N8 |  |
| N7 |  |
| N6 |  |
| N5 |  |
| N6 |  |
| N5 |  |
| N2 |  |
| N1 |  |
| N0 |  |
| $\begin{aligned} & 12 \text { nibbles } \\ & \text { per 48-bit } \\ & \text { address } \end{aligned}$ |  |

for registrable addresses, $r=1$; for claimable addresses, $r=0$
m is the usual multicast (I/G) bit; 111 is local "SAI" range per IEEE Std 802c
0000 for claimable addresses

- address block includes subblocks of
- $16^{j k}$ claimable addresses, or
- $16^{j k}$ registrable addresses (or aggregated into larger blocks)
- for claimable addresses, $i$ distinguishes
- Claimable Addresses (CAs) from
- CABAs
- identifiers that are also used as addresses
- see Appendix for details

|  | $r$ | $i$ | $j k$ | $m$ |
| :---: | :---: | :---: | :---: | :---: |
| CA | 0 | 1 | CAB | I/G |
| CABA | 0 | 0 | Size | 1 |
| TUA | 0 | 0 | 0 | 0 |
| RA | 1 | RABI <br> Option | BABI <br> Size | $\mathrm{I} / \mathrm{G}$ |

Claimant of $C_{A B A}^{X A B}$ $\begin{aligned} & \text { listens to } \mathrm{CABA}_{X} \\
& \text { multicast address }\end{aligned}$

| $\mathrm{CABA}_{5}$ |
| :---: | :---: |


(1) CABA $_{1}$ : DISCOVER state

LAN
Claiming (simplified)


## Registrar

- Claimant need not be aware of Registrar when initiating a claim.
- Registrar maintains an inventory of RABIs.
- a protocol specifies how Registrars acquire RABIs.
- set of RABs is disjoint from the set of CABs
- AB is either claimable (CAB) or registrable (RAB); not both
- Registrar listens for all messages to a CABA.
$-r=0, i=0, m=1$, i.e. DA begins $00^{* *}-1111$
- [MMRP NumberOfValues field is 13 bits]
- Registrar can respond to a DISCOVER with an offer of a RABI in its inventory.
- The offer can also defend the DISCOVER's CABA.
- Registrar confirms registration of request for offered RABI.
- Pre-claim Inquiry lets Claimant reach Registrar or Advisor.
- Client can learn of Registrars and received Claim proposals.


# Operation with Registrars 



## Inquiry to (anticipated) Registrar or Advisor



## BARC Design

- A BARC architecture follows, with details including state machines.
- additional details in Appendix
- BARC (Block Address Registration and Claiming) is put into the broader context of Address Registration and Claiming (ARC), which supports both:
- address blocks (ABs), identified by Address Block Identifiers (ABIs)
- address ranges (ARs), excluding addresses specified by BARC
- ARC is the general protocol
- BARC handles ABI Registration and CABA Claiming
- existing MAAP handles AR Claiming


## ARC Architecture - ARC Claimant



## BARCPDU Summary

| field name | purpose | content |
| :--- | :--- | :--- |


| DA | dest addresss | DA |
| :---: | :---: | :--- |
| SA | source address |  |
| E | Ethertype | [tbd; could be 22F0 (MAAP Ethertype)] |
| t | subtype | [tbd, per 1722 WG; see IEEE 1722 Table 6] |


|  | State 1 | D (Discover), C (Claimed), V (Vacant), <br> R (Registered), I (Inquiry), <br> P (Proposal), A (address), <br> RD, RC, RV, RX, N(null) |
| :--- | :---: | :--- |
| S1 | Identifier 2 | 48-bit address or ABI |
| S2 | State 2 | O (Offered), A (address), N (null) |
| I2 | Identifier 2 | 48-bit address or ABI |
| S3 | State 3 | A (address), T (token) |
| I3 | Identifier 3 | 48-bit address or token |

AVTPDU Summary

| field name | purpose | content |
| :---: | :--- | :--- |


| DA | dest addresss | $91:$ E0:F0:00:FF:00 for MAAP multicast |
| :---: | :---: | :--- |
| E | Ethertype | 22F0 (MAAP Ethertype) |
| t | subtype | FE per IEEE 1722 Table 6 |



## ABD Claimant/Registrar Procedure



## RABI Registration



## Renewal and Withdrawal of a Registration



## Inquiry followed by Registrar Offer



## Inquiry followed by Advisor Proposal



## ABD Claimant: ABD State Machine



## ARC Claimant: cBARCPDU_in


\{Note: square-bracketed parameters are sometimes absent.\}

## ARC Claimant: cBARCPDU_out




## AR Claimant Procedure



## BARC Registrar: AVTPDU Processor



## AR State Transition Table

|  | State |  |  |
| :--- | :--- | :--- | :--- |
|  | VACANT (V) | DISCOVERY (D) | ACQUIRED (A) |
| Event | sMAAP(Begin(AR,sa)!) <br> DISCOVERY |  |  |
| rMAAP(AR:Defend)! |  | Outcome(A,AR) <br> ACQUIRED |  |
| rMAAP(AR:Initial)! |  | Outcome(F,AR) <br> VACANT | Outcome(X)[AR] <br> VACANT |

rMAAP(AR:State!) invokes an event at the state machine when the MAAP state changes to State sMAAP(Action!) invokes Action! event at MAAP state machine

BARC Architecture - Registrar


## BARC Registrar Application: Disco Process




## BARC Registrar: rBARCPDU_in



## BARC Registrar: rBARCPDU_out



## Registrar: RABI State Machine



## VLANs

- All address assignments are specific to the VLAN (or VLANs) in which messaging is communicated and under which the assignment was completed.
- Usage of any address may need to be limited to the VLAN or VLANs under which it was obtained.
- Due to the possibility that the same unicast address may be assigned in different VLANs, Independent VLAN Learning (IVL) may be needed in bridges, per IEEE Std 802.1Q Annex F (F.1.2).
- This requirement could be relaxed in some cases
- e.g. when assigned unicast addresses are declared via MMRP (instead of learning)
- This issue is common to claiming protocols generally.
- Some approaches follow.


IEEE Std 802-2014 says "Local MAC addresses need to be unique on a LAN or bridged LAN unless the bridges support VLANs with independent learning."

With IVL, each VLAN has an independent forwarding table.
-but IVL is not always possible
BARC claiming on each VLAN is independent
a duplicate address may occur in more than one VLAN; that is not harmful if managed carefully
A claimant with multiple VLANs needs to claim in each VLAN.
Claimed address is usable only in claimed VLAN:
Claimant needs to bind address to VLAN
For convenience, Claimant may claim the same address in each of its VLANs
-Still, requires multiple claim messages and multiple forwarding table entries.
-Device needing many VLANs should consider an EUI


With SVL, VLANs share a forwarding table.
BARC claiming on each VLAN is independent
an address could become a duplicate, existing in more than one VLAN
forwarding table is limited to one entry per address, so duplication is catastrophic.


Network is configured with Registrar on all active VLANs on which BARC is used.
BARC claim from any VLAN is delivered to Registrar.
-Offer delivered on Claimant's VLAN
Registrar ensures that registered address is unique across all (or perhaps only some) of its VLANs.
-SVL or IVL will work
Registrar needs to remember over which VLANs the address was assigned.
-should be retained in State Machine


SVL is used for Asymmetric VLAN (IEEE Std 802.1Q Annex F.1.3)
Registrar can assign address to be unique across all VLANs available to the Registrar.

## Summary

- Claimants operate with or without Registrars.
- Multiple registrars are supported, holding claims of disjoint RABIs.
- The block discretization provides:
- a vast set of addresses to a LAN
- a large set of temporary unicast addresses
- operational efficiency and simplicity
- both unicast and multicast addresses to Claimant
- unicast and multicast subblocks share the same range, except for the $m$ bit
- could be exploited
- devices needing both unicast and multicast addresses need make only one claim
- Could integrate with MMRP to limit propagation and eliminate learning of unicast $A B$ content.
- MMRP needs to efficiently handle address ranges
- BARP could be specified as alternative MRP application (e.g. would understand an ABD)


## Appendix 1

- additional details on BARC addresses and identifiers


## BARC Address Parsing



## CABA and CA, CAB Size 0-3



## CABA CAB

 X8

| X 7 | X 7 |
| :---: | :---: |
| X 6 | X 6 |
| X 5 | X 5 |
| X 4 | X 4 |
| X 3 | X 3 |
| X 2 | X 2 |
| X 1 | X 1 |
| X 0 | X 0 |


| 0 | 1 | 0 | 0 |
| :--- | :--- | :--- | :--- |
| 1 | 1 | 1 | $*$ |
| 0 | 0 | 0 | 0 |
| X 8 |  |  |  |
| X 7 |  |  |  |
| X 6 |  |  |  |
| X 5 |  |  |  |
| X 4 |  |  |  |
| X 3 |  |  |  |
| X 2 |  |  |  |
| X 1 |  |  |  |
| X 0 |  |  |  |

CAB Size C=1

## CABA CAB



| 0 | 1 | 0 | 1 |
| :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | $\star$ |
| 0 | 0 | 0 | 0 |
| X 8 |  |  |  |
| X 7 |  |  |  |
| X 6 |  |  |  |
| X 5 |  |  |  |
| X 4 |  |  |  |
| X 3 |  |  |  |
| X 2 |  |  |  |
| X 1 |  |  |  |
| $\|c\|$ |  |  |  |

CAB Size C=2
CABA CAB

| 0 | 0 | 1 | 0 | 0 | 1 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 |  |  |  |  |  |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 |

CAB Size C=3
CABA CAB

| 0 | 01 | 1 | 0 | 1 | 1 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 11 | 1 | 1 | 1 | 1 | * |
| 0 | 00 | 0 | 0 | 0 | 0 | 0 |
|  | X8 |  |  |  | X8 |  |
|  | X7 |  |  |  | X |  |
|  | X6 |  |  |  | $\times 6$ |  |
|  | X5 |  |  |  | X |  |
|  | X4 |  |  |  | X4 |  |
|  | X3 |  |  |  | X3 |  |
|  | 0 |  |  |  | * |  |
|  | 0 |  |  |  | * |  |
|  | 0 |  |  |  | * |  |

2 contiguous subblocks per CABA (one unicast, one multicast)

- 6.9E10 Size 0 CABAs
- 1 CA/subblock
4.3E9 Size 1 CABAs 16 CAs/subblock
2.7E8 Size 2 CABAs
-256 CA/subblock
1.7E7 Size 3 CABAs 4096 CAs/subblock


## CABA/CAB Math


$C(X)=(X \& 0 \times 300000000000) / 0 \times 100000000000$ [extracts CAB Size $C$ when $X$ is CABA or CA]
Cmask $(C)=\sim\left(0 \times 410000000000+0 \times 10^{* *} C-1\right)$ [CABA mask, per Size; used to create CABA from CA]
CABA(CA) = CA\&Cmask(C(CA))
Example: [Note: Underlining on the middle four nibbles is shown only as a reading aid.]

- CA = 0x6F0123456789 = 0110-1111-0000-0001-0010-0011-0100-0101-0110-0111-1000-1001
- C(CA) $=0 \times 200000000000 / 0 \times 100000000000=2$
- Cmask $(0 \times 2)=\sim(0 \times 410000000000+0 \times 0100-1)=\sim(0 \times 4100000000 F F)=0 \times B E F F F F F F F F 00$
- CABA(CA) $=$ CA\& $0 \times B E F F$ FFFFFF $00=0 \times 2$ E0123456700
$A C A$ is within $C A B A_{x}$ if and only if $C A B A(C A)=C A B A_{x}$
-this requires identical CAB Size
The CAB of $C A B A_{x}$ is the set of all CAs that satisfy $C A B A(C A)=C A B A_{x}$
Lowest ICA in CABA: CABA|0x400000000000 (example: $0 \times 6 E 01 \underline{23456700)}$
Lowest GCA in CABA: CABA|0x410000000000 (example: $0 \times 6 F 01 \underline{23456700)}$
Highest ICA in CABA: (CABA|0x4000000000000) + 0x10**C(CABA) - 1 (example: $0 \times 6 E 01234567 F F)$
Highest GCA in CABA: (CABA|0x4100000000000) + 0x10**C(CABA)-1 (example: 0x6F01234567FF)


## ABD (CABA/RABI) Format

ABD
CABA


RABI


## RABI Aggregation

Address

CAB


RAB


## -RAB Size $R=B+M$

- for 48-bit addresses, set $a=0$; then $R=j k+b c d \leq 3+7=10$, matching the 10 available nibbles N0-N9
- could use ijk as the BABI Size, and/or the full abcd as the MABI Size; e.g., for 64-bit addresses
- A RABI may aggregate other RABls.
-A RABI of RAB Size $R$ and MABI Size $M$ can be disaggregated into:
-16 RABIs of RAB Size $R-1$ (MABI Size $M-1$ ), or
-16² RABIs of RAB Size $R-2$ (MABI Size $M-1$ ), or
-16n ${ }^{n}$ RABIs of RAB Size $R-n$ (MABI Size $M-n$ ), or
- .. $16^{M}$ RABIs of RAB Size $B$ (MABI Size 0), or
-A RABI of RAB Size $B$ (MABI Size 0) cannot be disaggregated.
-An RA appears in one and only RABI of each $M$.


## RABI and RA, MABI Size 0, BABI Size 0-3

| BABI Size 0 |  |  | BABI Size 1 |  |  |  | BABI Size 2 |  |  |  | BABI Size 3 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| RABI | RAB |  | RABI |  | RAB |  | RABI |  | RAB |  | RABI |  |  | RAB |  |
| 1 i l 000 | 1 i | 0 0 | 1 | O 1 | 1 i | 0 | 1 i | i 10 | 1 i | 10 | 1 |  | 1 | 1 i | 11 |
| 0 0 000 | 11 | 1 * | 0 | 0 0 0 | 11 | 1 |  | 00 | 11 | $1{ }^{*}$ |  | 0 | 0 | 11 | 1 |
| X9 |  | X9 |  | X9 |  | X9 |  | X9 | X9 |  |  | 9 |  | X |  |
| X8 |  | X8 |  | X8 |  | X8 |  | X8 | X8 |  |  | X8 |  | X |  |
| X7 |  | X7 |  | X7 |  | X7 |  | X7 | X7 |  |  | X7 |  | X |  |
| X6 |  | X6 |  | X6 |  | $\times 6$ |  | X6 | X6 |  |  | X6 |  | X | X |
| X5 |  | X5 |  | X5 |  | X5 |  | X5 | X 5 |  |  | X5 |  | X | ¢ |
| X4 |  | X4 |  | X4 |  | X4 |  | X4 | X |  |  | X4 |  | X | 4 |
| X3 |  | X3 |  | X3 |  | X3 |  | X3 | X3 |  |  | X3 |  | X | X |
| X2 |  | X2 |  | X2 |  | X2 |  | X2 | X |  |  | 0 |  | * |  |
| X1 |  | X1 |  | X1 |  | X1 |  | 0 | * |  |  | 0 |  | * | * |
| X0 |  | X0 |  | 0 |  | * |  | 0 | * |  |  | 0 |  | * | * |

2 contiguous subblocks per RABI (one unicast, one multicast)
-1.1 E 12 Size 0 RABIs
-1 RA/subblock
-6.9E10 Size 1 RABIs - 16 RAs/subblock

2.7E8 Size 3 RABIs 4096 RAs/subblock

## Aggregation Example: BABI Size 3, various MABI Sizes

## BABI Size 3



## Example:

## Hierarchical RABI Addressing with common Registrar

Held by Registrar

MABI Size 2
RABI RAB

| 1 | i |  |  | i | i | 1 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 |  |  |  | 1 | 1 | * |
| X9 |  |  |  | X9 |  |  |  |
| X8 |  |  |  | X8 |  |  |  |
| X7 |  |  |  | X7 |  |  |  |
| X6 |  |  |  | X6 |  |  |  |
| X5 |  |  |  | X5 |  |  |  |
| 0 |  |  |  | \# |  |  |  |
| 0 |  |  |  | \# |  |  |  |
| 0 |  |  |  | * |  |  |  |
| 0 |  |  |  | * |  |  |  |
| 0 |  |  |  |  | * |  |  |

## Assigned to Bridge 1 by Registrar



RABI RAB


Assigned by Registrar to bridges, end stations, etc. connected to Bridge 1


Bridge serves as Advisor to connected station; proposes PRABI, and Registrar Address,
in response to an Inquiry.

## Example: Hierarchical CABA Addressing

Claimed by Bridge 1

| CAB Size C=3 |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| CABA |  |  | CAB |  |  |
| 0 | 01 | 1 | 0 | 1 | 11 |
| 1 | 1 | 1 | 1 | 11 | 1 |
| 0 | 00 | 0 | 0 | 00 | 0 |
|  | X8 |  |  | X8 |  |
|  | X7 |  |  | X7 |  |
|  | X6 |  |  | X6 |  |
|  | X5 |  |  | X5 |  |
|  | X4 |  |  | X4 |  |
|  | X3 |  |  | X3 |  |
|  | 0 |  |  | * |  |
|  | 0 |  |  | * |  |
|  | 0 |  |  | * |  |

Claimed by bridges, end stations, etc. connected to Bridge 1
CAB Size C=2
CAB Size C=2

CABA CAB CABA CAB

| 0 | 01 | 0 | 0 | 1 |  | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 11 | 1 | 1 | 1 | 1 |  |
| 0 | 00 | 0 | 0 | 0 | 0 | 0 |
|  | X8 |  |  | X8 |  |  |
|  | X7 |  |  | X7 |  |  |
|  | X6 |  |  | X6 | - |  |
|  | X5 |  |  | X5 |  |  |
|  | X4 |  |  | X4 | 4 |  |
|  | X3 |  |  | X3 | 3 |  |
|  | 0 |  |  | 0 | 0 |  |
|  | 0 |  |  |  |  |  |
|  | 0 |  |  | * |  |  |

Bridge serves as Advisor to connected stations; proposes CABA in response to an Inquiry.

## Claimant as Registrar



RABI registered as a Claimant could be disaggregated and reassigned by a Registrar function managed jointly with the Claimant.
A RABI in the "REGISTERED" state of the ABD Claimant State Machine could be considered to be in the Inventory of the RABI Registrar State Machine (along with Claimed RABIs) and could be disaggregated, Offered and Registered by that Registrar.

## Hierarchical RABI Disaggregation with Tiered Registrars

Assigned to Bridge 1
by Registrar

| Used by Bridge 1 |
| :---: |
| directly |

MABI Size 1

| RABI | RAB |  |
| :---: | :---: | :---: |
| 1 i 1 1 <br>     | 1 i 1 |  |
| 0 0 0 1 <br>     | 1 1 1 <br>    |  |
| X9 | X9 |  |
| X8 | X8 | common RAB |
| X7 | X7 | prefix |
| X6 | X6 |  |
| X5 | X5 | $\nabla$ |
| 8-15 | 8-15 |  |
| 0 | \# |  |
| 0 | * |  |
| 0 | * |  |
| 0 | * |  |

Reassigned by Bridge 1 as secondary Registrar to connected bridges, end stations, etc.

| MABI Size 1 |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| RABI |  |  | RAB |  |  |
| 1 | i 1 |  | 1 | 1 | 1 |
| 00 | 00 |  | 1 |  | * |
|  | X9 |  |  | X9 |  |
|  | X8 |  |  | X8 |  |
|  | X7 |  |  | X7 |  |
|  | X6 |  |  | X6 |  |
|  | X5 |  |  | X5 |  |
|  | 0 |  |  | 0 |  |
|  | 0 |  |  | \# |  |
|  | 0 |  |  | * |  |
|  | 0 |  |  | * |  |
|  | 0 |  |  | * |  |


| MABI Size 1 |  |  |
| :---: | :---: | :---: |
| RABI | RAB |  |
| 1 i 1 1 | 1 i | 11 |
| 0 0 0 1 | 111 | 1 |
| X9 | X9 |  |
| X8 | X8 |  |
| X7 | X7 |  |
| X6 | X6 |  |
| X5 | X5 |  |
| 1 | 1 |  |
| 0 | \# |  |
| 0 | * |  |
| 0 | * |  |
| 0 | * |  |

## RABI

| N11 |
| :--- |
| N10 |



- $j k$ indicates the BABI Size $B$ -bcd indicates the MABI Size $M$

| $\&$ | bitwise AND |
| :---: | :---: |
| $\sim$ | bitwise NOT |
| $/$ | divide |
| $==$ | binary equality |

$i(X)=(X \& 0 x 040000000000) / 0 x 400000000000$ [RABI Option, when $X$ is RABI or RA]
$B(X)=(X \& 0 \times 300000000000) / 0 \times 100000000000$ [BABI Size, when $X$ is RABI or RA]
$M(X)=(X \& 0 \times 070000000000) / 0 x 010000000000$ [MABI Size when $X$ is a RABI]
$R(X)=B(X)+M(X)[R A B$ Size, when $X$ is a RABI]
Rmask $(\mathrm{N})=\sim\left(0 \times 0 \mathrm{~F} 0000000000+0 \times 10^{* *} \mathrm{~N}-1\right)$ [mask, used below]
An RA is within RABIx if and only if $\operatorname{RA\& Rmask}(R(R A B I x))=R A B I x \& R m a s k(R(R A B \mid x))$
-this requires identical RABI Option and BABI Size
The RAB of RABIx is the set of all RAs that satisfy RA\&Rmask(R(RABIx)) = RABIx\&Rmask(R(RABIx))
Example:

- RABI $=0 x F 20123400000=111-0010-0000-0001-0010-0011-0100-0000-0000-0000-0000-0000$
- $B($ RABI $)=0 \times 300000000000 / 0 \times 100000000000=3$
- $M($ RABI $)=0 \times 020000000000 / 0 \times 010000000000=2$
- $\mathrm{R}(\mathrm{RABI})=3+2=5$
- Rmask(5) $=\sim\left(0 \times 0 F 0000000000+0 \times 10^{* *} 5-1\right)=\sim(0 \times 0 F 00000 F F F F F)=0 \times F 0 F F F F F 00000$
-RABI\&Rmask(5) = 0xF00123400000
- RA $=0 x F E 0123456789=1011-0010-0000-0001-0010-0011-0100-0101-0110-0111-1000-1001$
-RA\&Rmask(5) $=0 \times F 00123400000=$ RABI\&Rmask(5)
- so RA is within RABI


RABIcheck(RABI1,RABI2) determines whether RABI RABI1 overlaps RABI2.
RABIcheck(R1,R2) = R1\&Rmask(R(R2))
RABI1 shares RAs with RABI2 if and only if:
RABlcheck(RABI1, RABI2) = RABlcheck(RABI2, RABI1)
-Note: This can be ruled out by inspection if the two N11 nibbles differ
Example:

- RABI2 $=0 x$ F501000000000 $=1111-0101-0000-0001-0000-0000-0000-0000-0000-0000-0000-0000$
- $\mathrm{B}(\mathrm{RABI} 2)=0 \times 300000000000 / 0 \times 100000000000=3$
- $M($ RABI2 $)=0 \times 070000000000 / 0 \times 100000000000=5$
- $R($ RABI 2$)=3+5=8$
- Rmask(8) $=0 \times 50$ FF00000000
- RABI1 = 0xF20123400000
-RABI1\&Rmask(8) $=$ 0xF001000000000
- RABI2\&Rmask(5) = 0xF001000000000
- so RABI1 and RABI2 have RAs in common

Lowest IRA in RABI: RABI \& 0xF0FFFFFFFFFF + 0x0E0000000000
Example: 0xFE01234000000)
Lowest GRA in RABI: RABI \& 0xF0FFFFFFFFFF + 0x0F00000000000
Example: 0xFF0123400000)
Highest IRA in RABI: RABI \& 0xF0FFFFFFFFFF + 0x0E0000000000 + 0x10**(R(RABI)) - 1
Example: 0xFE01234FFFFF)
Highest GRA in RABI: RABI \& 0xF0FFFFFFFFFF + 0x0F0000000000 + 0x10**(R(RABI)) - 1
Example: 0xFF01234FFFFF)

|  | RABI |  |  | RA |  |  | RABI Math |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| N11 | 1 | i | j k | 1 | i | k | $\cdot j k$ indicates the BABI Size $B$ $\cdot b c d$ indicates the MABI Size $M$ | \& | bitwise AND |
|  |  |  |  |  |  |  |  | $\sim$ | bitwise NOT |
| N10 | 0 | b | c d | 1 | 1 | 1 * |  | 1 | divide |
|  |  |  |  |  |  |  |  | == | binary equality |

An RA exists in one and only one RABI (called $\mathrm{RABI}_{M}$ ) of each MABI Size $M$. Given the MABI Size $M$, what is $\mathrm{RABI}_{M}$ ?

A RABI exists in one and only one aggregated RABI called RABI) of each larger MABI Size.
Given the MABI Size $M$, what is RABI $_{M}$ ?
$\operatorname{RABI}_{M}\left(\mathrm{X}, M_{n}\right)=\mathrm{X} \& \operatorname{Rmask}\left(M_{n}+\mathrm{B}(\mathrm{X})\right)+M_{n}(\mathrm{X})^{*} 0 \times 010000000000\left[\mathrm{X}\right.$ is an RA, or a RABI with $\left.\mathrm{M}[\mathrm{X}]<M_{n}\right]$
Example:

- RABI1 = 0xF20123400000 B(RABI1)=3; M(RABI1)=2
- $M_{n}=5 ; M_{n}+\mathrm{B}=8$
- RABI1\&Rmask(8) $=0 \times 520123400000 \& 0 x F 0 F F 00000000=0 x F 00100000000$
-RABIm(RABI1,5) $=0 \times 500100000000+0 x 050000000000=0 x F 50100000000$
Example:
- RA $=0 \times F E 0123456789 \quad[B(R A)=3]$
- $M_{n}=5 ; M_{n}+B=8$
- RA\&Rmask(8) $=0 \times F E 0123456789 \& 0 x F 0 F F 00000000=0 x F 00100000000$
$\cdot \operatorname{RABIm}(\mathrm{RA}, 8)=0 x F 00100000000+0 x 050000000000=0 x F 50100000000$
Example:
- RA $=0 x F E 0123456789 \quad[B(R A)=3]$
- $M_{n}=2 ; M_{n}+B=5$
- RA\&Rmask(5) $=0 \times$ xFE0123456789\&0xF0FFFFF00000 $=0 \times 500123400000$
$\cdot \operatorname{RABIm}(R A, 5)=0 x F 00123400000+0 x 020000000000=0 x F 20123400000$


## Temporary Unicast Address (TUA)

- For temporary use
- device without a source address selects a random temporary unicast address for initial discovery only
- protocol then assigns at least one persistent unicast address
- simultaneous duplicate temporary addresses may lead to message loss in some circumstances
- network learns route to source as initial message crosses the network
- before response is returned, another initial message with duplicate source address
crosses the path and rewrites the route
- unlikely to be disastrous
- loss of initial message will be corrected eventually
- nevertheless, need to consider the likelihood of duplication
- Temporary address range includes 9 full nibbles of 16 values each ( $0-F)$
$-16^{9}=68,719,476,736(=N)$ temporary addresses in the pool
- chance of no duplicates with $k$ randomly selected addresses is approximated $\exp \left(-k^{*}(k-1) /\left(2^{*} N\right)\right)$
- with $k=1000$ devices simultaneously using a temporary address, chance of no duplicates is $\sim 0.99988$
- can add $j k$ bits to the pool if more entropy is needed

| 0 | 0 | 0 |
| :---: | :---: | :---: |
| 1 | 1 | 0 |
| 0 | 0 | 0 |
| * |  |  |
|  | * |  |
|  | * |  |
|  | * |  |
|  | * |  |
|  | * |  |
|  | * |  |
|  | * |  |
|  | * |  |

## null CABA $\left(\mathrm{CABA}_{0}\right)$



## Proposed RABI (PRABI) and null RABI/PRABI

- Usable only as the content of a BARC Inquiry or BARC Proposal message.
- Indicates a set of RABIs characterized by RABI Option $i$, BABI Size B $j k$, and MABI Size M bcd.
- If the PRABI has the form of a RABI (with 0 in the $\mathrm{B}+\mathrm{M}$ least significant nibbles), then the PRABI set is only that RABI.
- Non-zero values in the lower $\mathrm{B}+\mathrm{M}$ nibbles can signify that some bits of the higher nibbles are "don't care."

PRABI

proposed RABIs


0
0
0
0
0

- As a PRABI in a BARC Inquiry or Proposal message, indicates a set of RABIs characterized by RABI Option i, BABI Size B $j k$, and MABI Size M bcd, without expressing a preference for the RABI values of the 0 nibbles.
- As a RABI in a BARC Offer message, indicates to Claimant that no RABI is offered.
null RABI/ null PRABI

| 1 | i 1 | 1 |
| :---: | :---: | :---: |
| 0 | c | c |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |
|  | 0 |  |

- For example, the lower B+B nibbles of the PRABI could form a bitmask of the higher nibbles.

Proposed RABIs are those that satisfy:
RABI | Pmask = (PRABI\&~(0x10**(R(PRABI)) - 1)) | Pmask
where $\operatorname{Pmask}(\operatorname{PRABI})=10^{\star *}(\mathrm{R}(\mathrm{PRABI}))^{\star} \operatorname{PRABI\& }\left(0 \times 10^{\star \star}(\mathrm{Rcap})-1\right)$
and $R c a p=\min (R(P R A B I), 10-R(P R A B I))$

## Appendix 2

- additional procedural details


## BARC Address Propagation with MMRP

The ARC Claimant Application AddABD Process includes "declare with MMRP". This entails declaring, to MMRP (when available), MMRP attributes, using an MMRPDU per IEEE Std 802.1Q § 10.12.1.6:

- The multicast address represented by the CABA

FirstValue field = CABA/NumberOfValues=1

- The unicast address subblock indicated by the ABD (CABA or RABI)

FirstValue field $=$ first address in unicast subblock
NumberOfValues = $166^{\text {Size }}$ per ABD Size

- limited to Size up to 3: $16^{3}=4096$, and MRP provides 13 bits of NumberOfValues $\left(2^{13}=8192\right)$

The ARC Claimant Application AddABD Process includes ("select CABA"); this selection should consider any local MMRP registration database to avoid selecting a registered CABA.

Unicast MMRP declaration can be useful because:
(1) A one-step declaration covers a contiguous range of self-assigned unicast addresses.
(2) Eliminates flooding for all the unicast addresses in the assignment.
(3) Eliminates the need for learning of each unicast address when used.
(4) Precludes erroneous re-learning of an address when a false duplicate is used elsewhere in the network.

- Could be a way to control duplication.
- Security issues to study.

BARC could alternatively specify "BARP," a new MRP application. This could entail the following changes:
(a) the BARP application would be enabled to Join and Leave with the ABD as the declared attribute
(b) the BARP application would be specified to understand the semantics of the ABD and extract from it the indicated ABD multicast address and the indicated unicast address set, then use it to populate the FDB
(c) In the BARC Claimant ABD State Machine, the CABA claim ["sBARC(CABA:C)]" might not be needed, since the a BARP declaration could convey the claim to the CABA as well as the declaration of interest in receiving at the CABA multicast address

BARP might be better suited to specification within IEEE Std 802.1Q instead of 802.1CQ.

## ARC Claimant Application Process: BARC Management



## ARC Claimant Application Process: Drop Claim



## ARC Claimant Application Process: MAAP Management



## BARC Registrar Application: ClaimCheck Process



## BARC Registrar Application: RABI Claiming



## RABI Claimant Application Process: BARC Management



