



Institute of Computer Science Chair of Communication Networks Prof. Dr. Tobias Hoßfeld



Latency Model and Example Reservation Flow in RAP IEEE 802.1 Meeting, January 2023

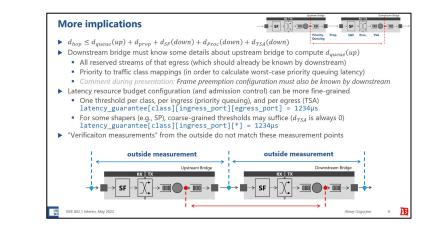
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Overview for this Presentation

- Recap: last presentation (measurement points for latency models)
 - cf. <u>dd-grigorjew-measurement-points-0522-v02.pdf</u>
 - What are the measurement points (and resulting delay segments)?
 - What are the reasons for this change?
 - Some implications
- More implications
 - Visualization of delay segments
 - What happens at the Listener?
 - Suggestion: merge two configurations per delay segment
 - What happens with different Shapers?
 - Suggestion: communicate the behavior at the Priority Transmission Selection Queue
- Example reservation flow
 - Very simple scenario (2 switches, 2 streams)
 - Clarify general procedure
 - Clarify the implications of delay segments





if we have time SW 1 SW 2

Recap: Last Presentation

MEASUREMENT POINTS FOR LATENCY MODELS



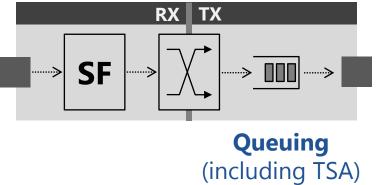


Extended delay model, including transmission selection algorithm

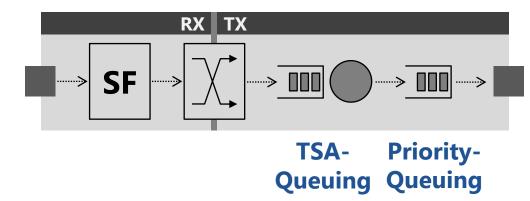
- Split "queuing" latency of formal latency models into...
 - Transmission Selection Algorithm (TSA)
 - Priority-Queuing, where only the eligible frames interfere

Previous model:

WÜ



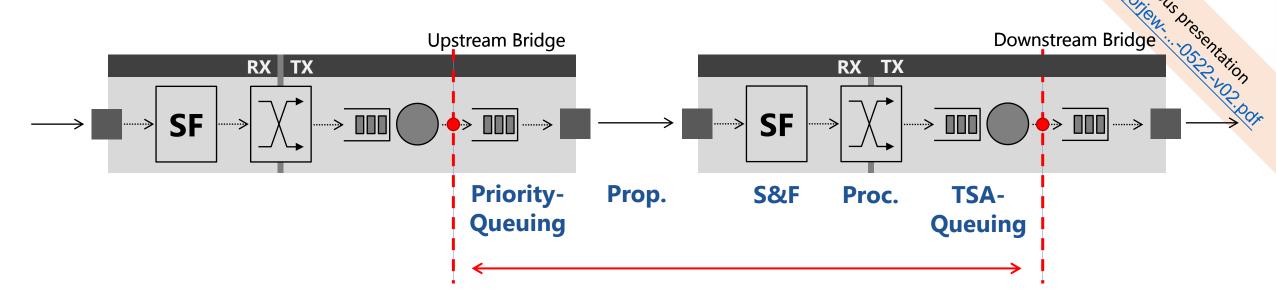
Extended model:



- Add measurement point during queuing when frame becomes eligible for transmission
 - SP: Immediately after enqueuing
 - CBSA: When credits >= 0, the head of the queue becomes eligible for transmission
 - ATS: When the defined eligibility time for that frame is reached (cf. Qcr)
 - CQF: When queues swap roles (receive \rightarrow send), all frames in the send queue become eligible

previous presentation

Suggestion: Use ATS measurement points for all shapers in RAP



- Suggestion: Use the ATS measurement points for all TSAs & latency models in RAP
- Per-hop latency is given by...
 - Queuing after eligibility time was reached (upstream)
 - Propagation

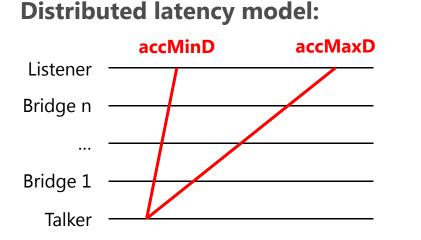
WÜ

- Store-and-Forward (downstream)
- Processing (downstream)
- Queuing until eligibility time is reached (downstream)
- Comment during presentation: PHY can often introduce a delay after priority queuing. The simple suggestion is to account for it as part of the upstream processing delay, even if it technically occurs after the measurement point.

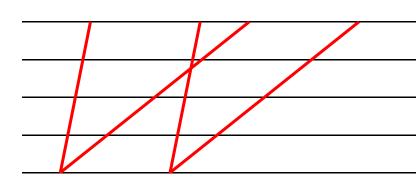
// queuing for priority transmission selection

// queuing for transmission selection algorithm

Why is shaper-to-shaper latency beneficial?



CQF (edge to edge measurement):



CQF (shaper to shaper):

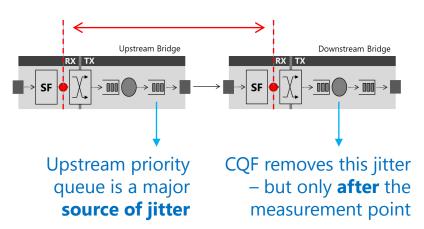
dd-grigorjew-strict-priority-latency-0320-v02.pdf

Generally:

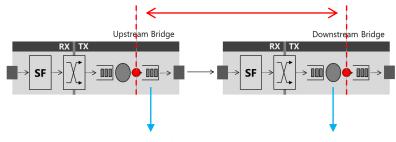
WÜ

- Minimum delay and maximum delay accumulated per hop
- Accumulating bursts are calculated based on (accMaxD - accMinD)
- A lower latency variance is better for downstream delay computation

Fully-received to fully-received:



Shaper to shaper:



All sources of jitter can be removed; Well-defined traffic pattern, as intended by the TSA, is measured directly after the TSA.

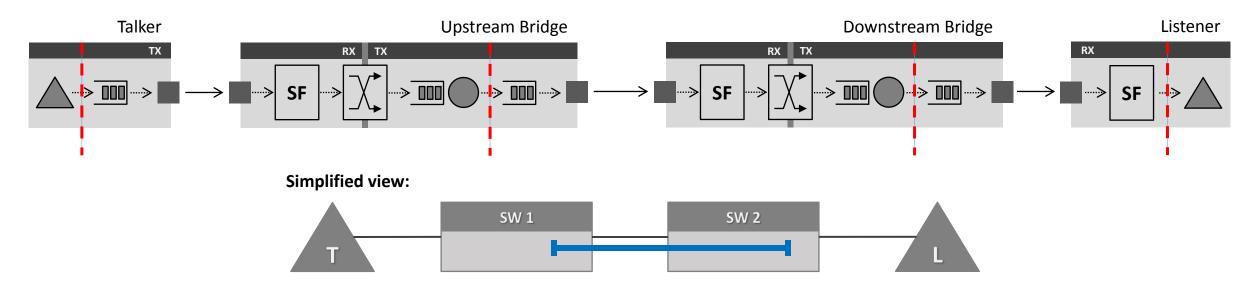
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Visualization and new Suggestions

MORE IMPLICATIONS

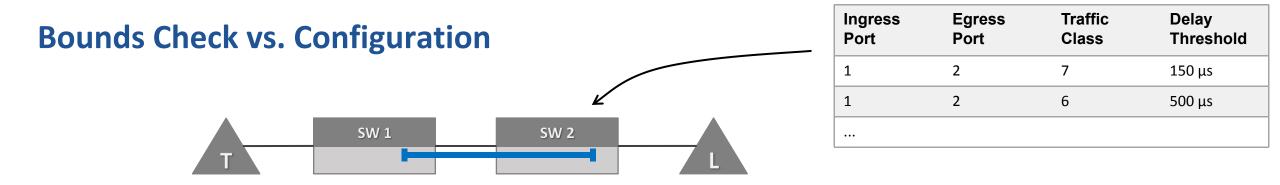


Full End-to-end Path with Delay Segments



One delay segment includes

- TX of upstream bridge (SW 1)
- RX of downstream bridge (SW 2)
- TX of downstream bridge (SW 2)
- Downstream bridge (SW 2) performs the bounds check during reservation
- But where does the configuration (delay threshold) come from? SW 1 or SW 2?
- General problem: on any path with n bridges (2 bridges), we have n+1 delay segments (3 delay segments)



- Initial suggestion: SW 2 performs bounds check and contains the delay threshold config
 - But: we don't really like the fact that SW 1 has no say, although it is involved in the delay segment

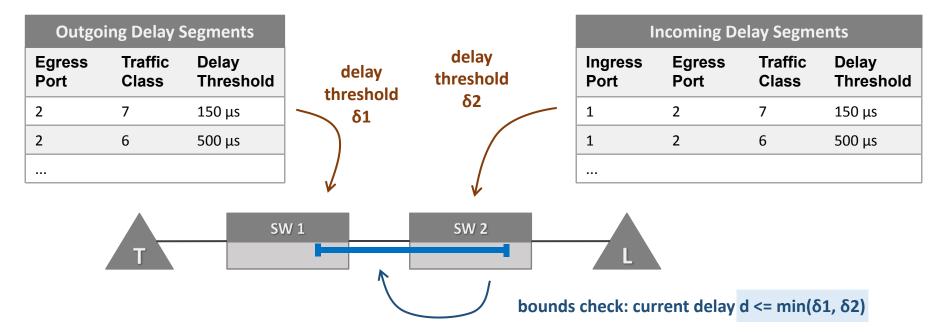
In addition: what happens at the Listener?



It can perform bounds checks

- But we don't really want to configure that aspect in our end devices
- (Config sources can be: default configuration, profile, CLI, Network Management System)

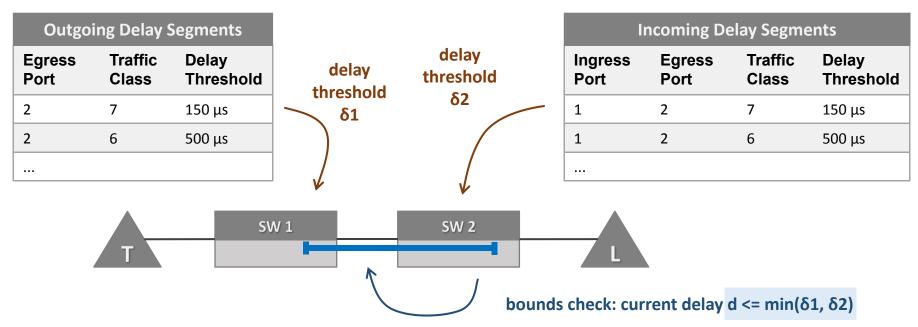
Suggestion: Both Devices Suggest a Delay Threshold



Suggestion: split threshold configuration for each delay segment into two configs

- Upstream bridge (SW 1) has one config for each egress port and traffic class
- Downstream bridge (SW 2) has one config **for each in ingress/egress port pair** and traffic class
- Each bridge will have two delay config tables: one for outgoing delay segments, one for incoming segments
- Upstream (SW 1) communicates the outgoing $\delta 1$ with the downstream neighbor (SW 2)
 - Downstream aggregates both configurations and selects the minimum of both for bounds checking

Implications of Having two Delay Threshold Tables



► The listener no longer needs a delay threshold configuration

- It can simply use $\delta 1$ of upstream (SW 2 in that case)
- It can still specify its own δ2 where necessary (e.g., routers are listeners from layer 2 RAP point of view)
- When optimizing a network's configuration (e.g., via NMS), simply use the same value for $\delta 1$ and $\delta 2$
 - It is the same delay segment after all
- Upstream (SW 1) could specify "don't care" in order to prevent unnecessary resource constraints
 - Technically, it still **needs** a valid outgoing $\delta 1$ config in case an end device connects to that port



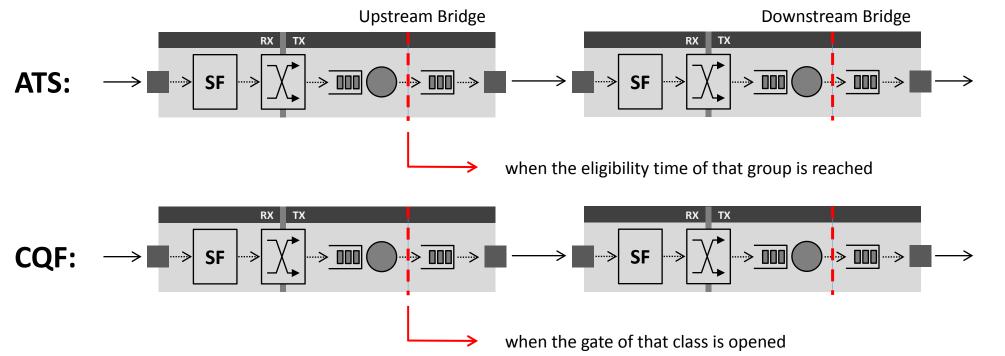
A Closer Look at Delay Segments with Different Shapers

- Recap old presentation: delay segments begin when the frame "becomes eligible for transmission"
- More specifically, we want delay segments to be tied to the events that change the shaper's state
- ► This ensures that the shaper has the intended effect on the latency model
- ► This is simple for ATS and CQF:

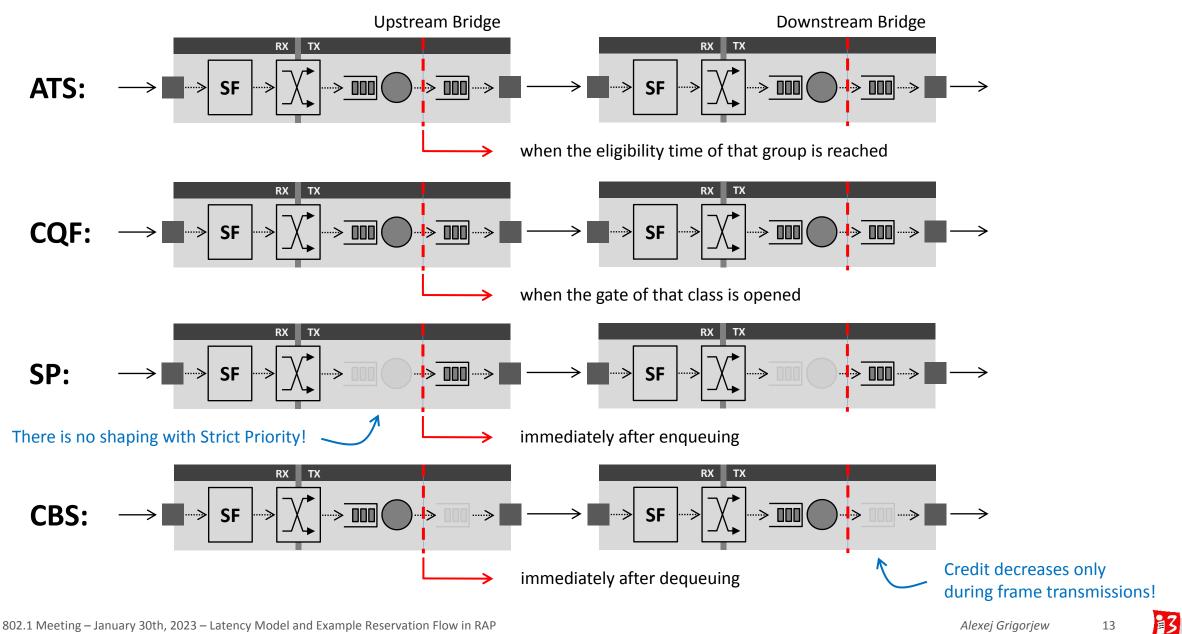
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CQF (shaper to shaper):



Not All Shapers Use All Delay Segments



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TAs, LAs, Bounds Checks, Example Values

EXAMPLE RESERVATION PROCESS

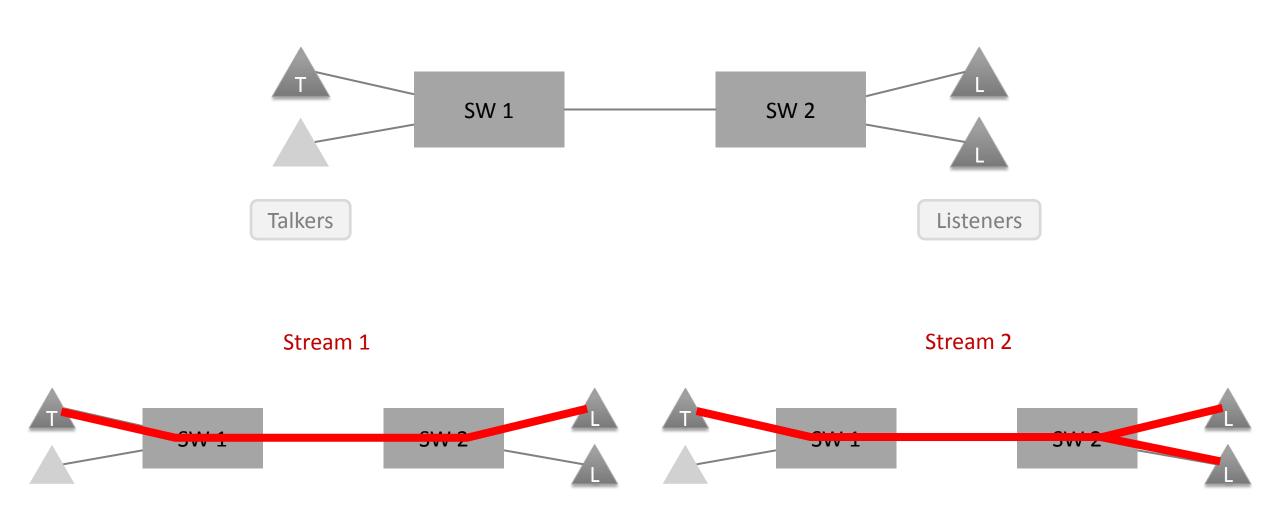


Disclaimer

- Just a simple example!
- Many things are simplified
- Some things are only suggestions
- Some things are subject to change in the standard
- See this as a means for easy introduction
- Please do not cling to the details

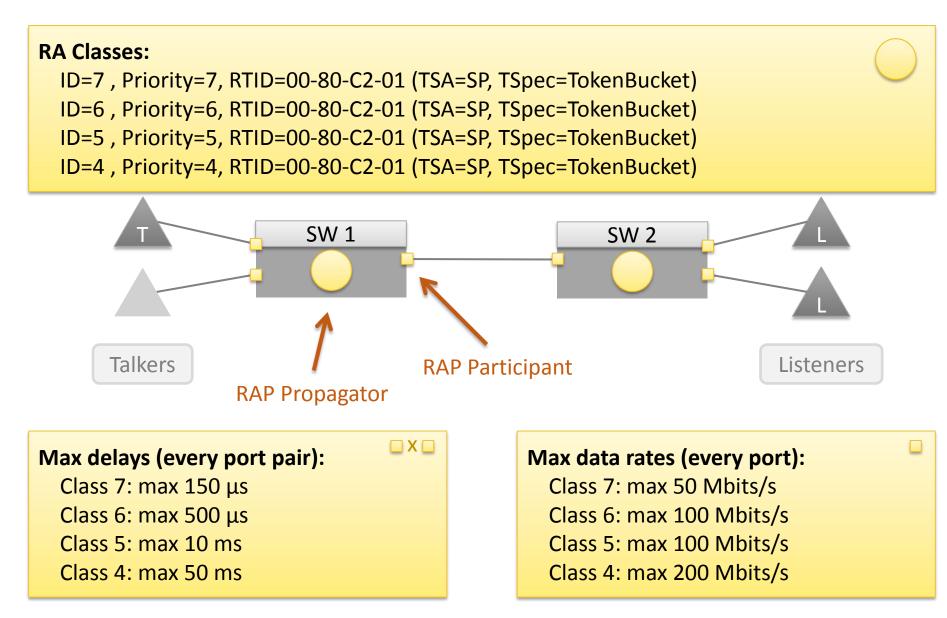
E3

Example Topology Overview



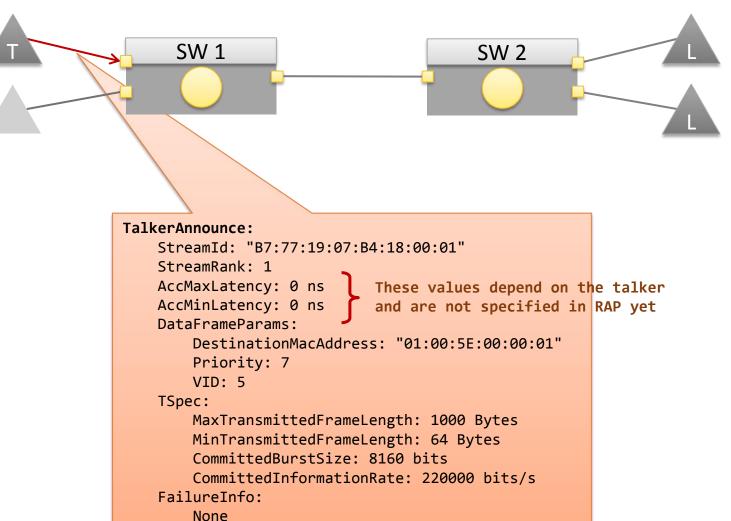
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Configuration





Stream 1

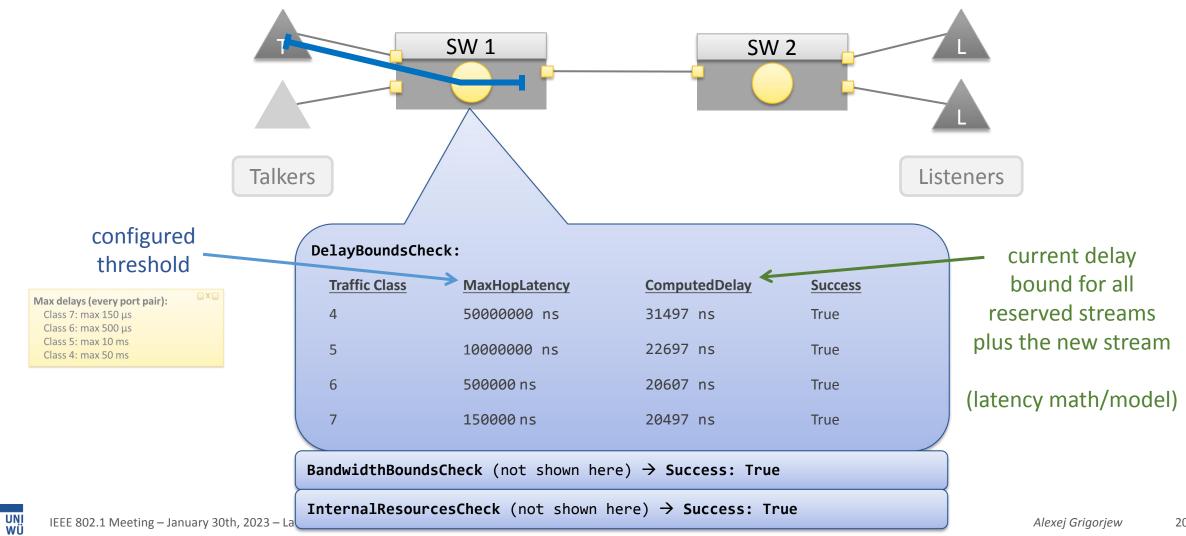


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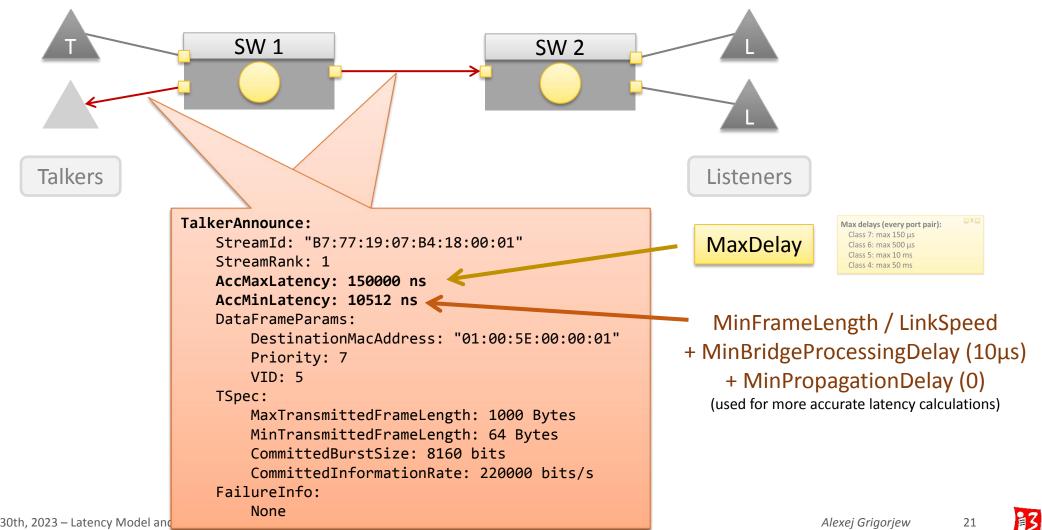
i3

Bounds Check on SW 1 (EgressPort: SW 2)



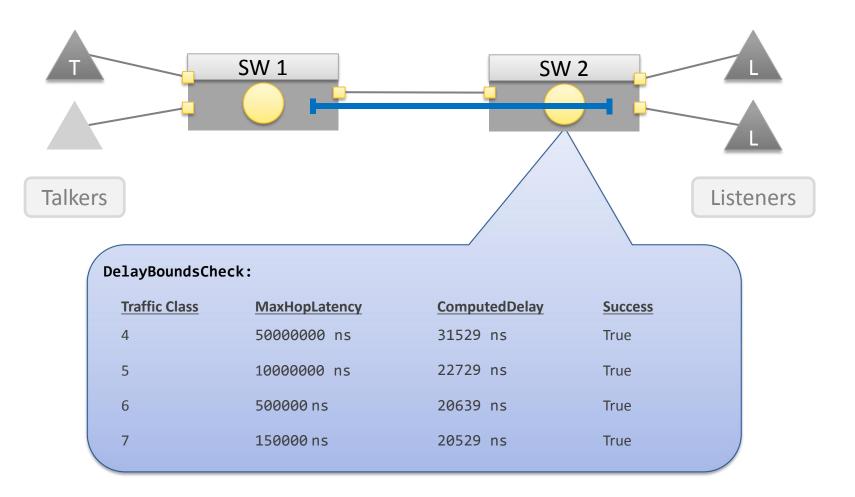
E3

Adjusted TA is Propagated to other Ports



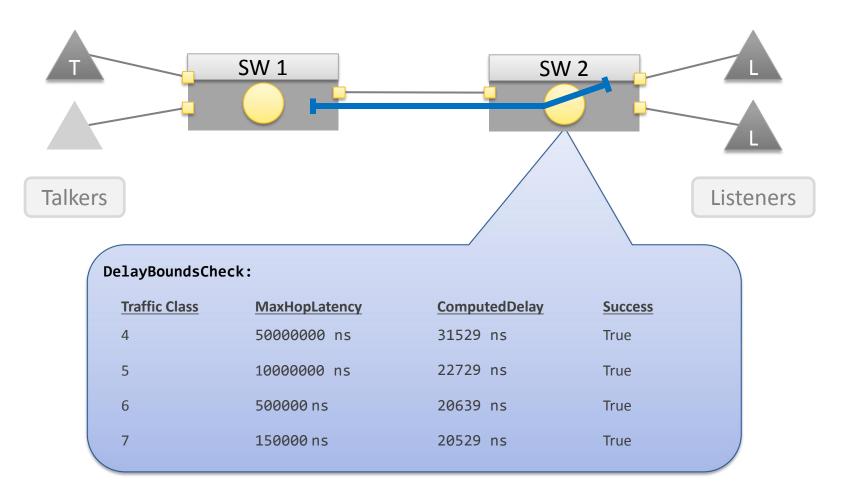
UN

Bounds Check on SW 2 (EgressPort: Listener 2)



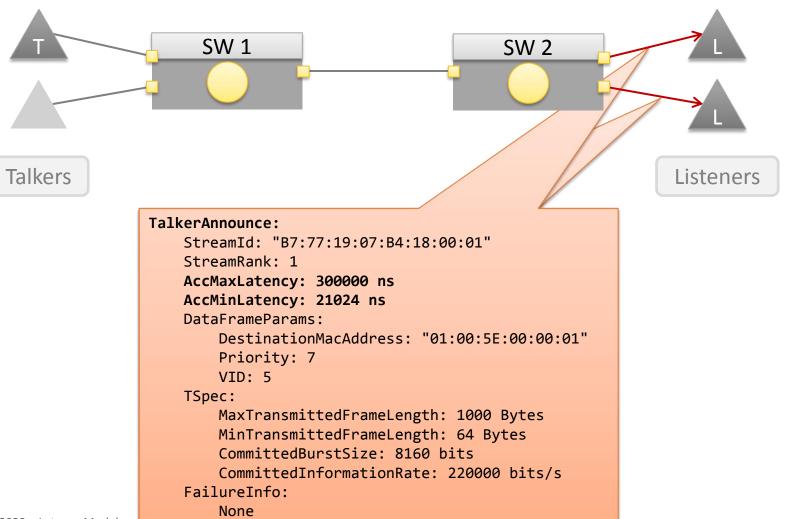
UNI WÜ 3

Bounds Check on SW 2 (EgressPort: Listener 1)



UNI WÜ 3

Adjusted TA is Propagated to other Ports

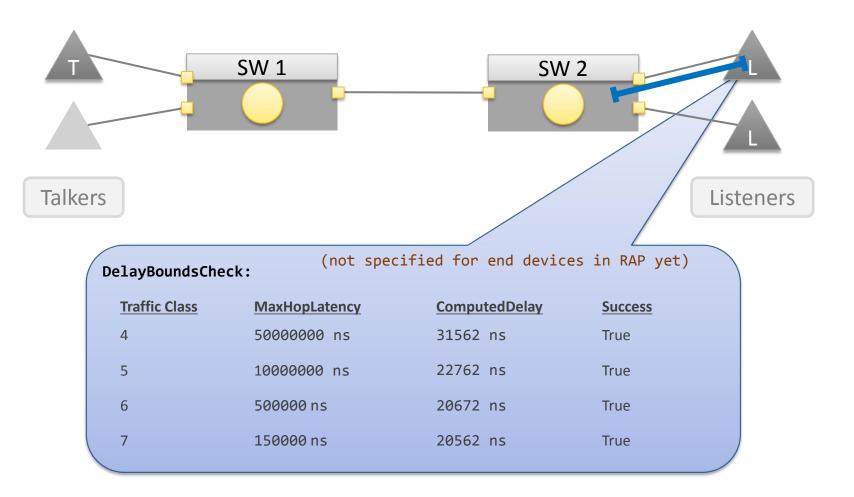


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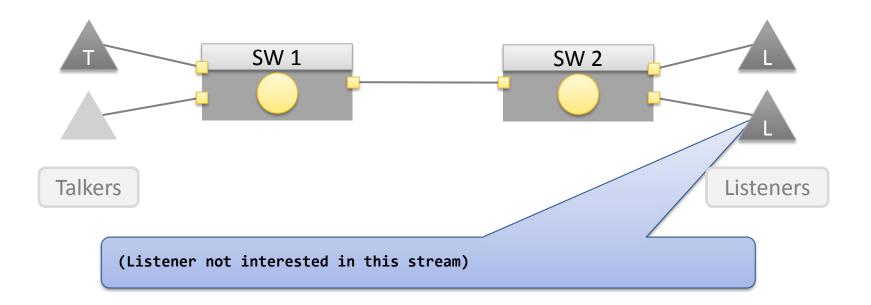
i3

Bounds Check on Listener 1



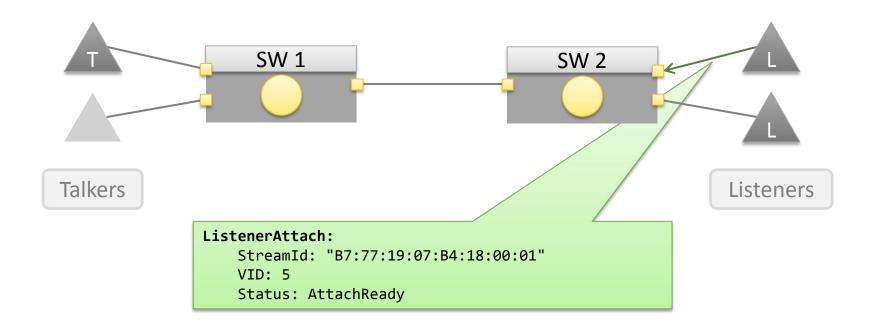


Listener 2 is not Attaching



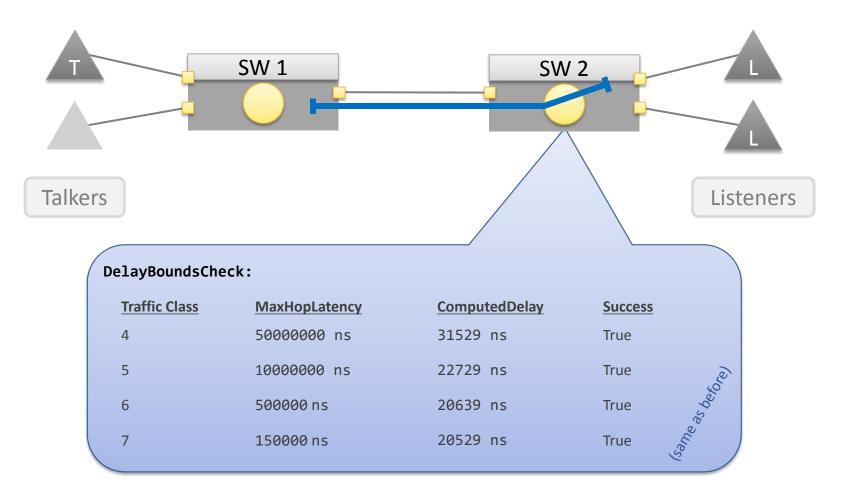


Listener 1 sends LA



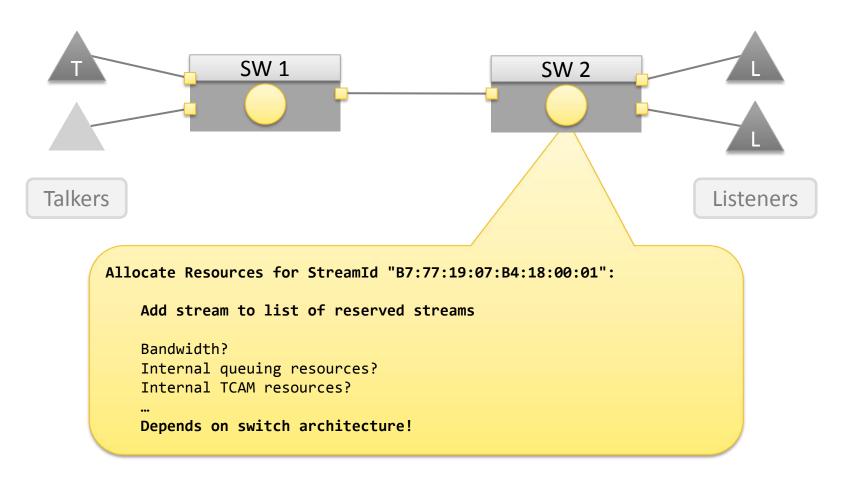


Bounds Check on SW 2 (EgressPort: Listener 1)



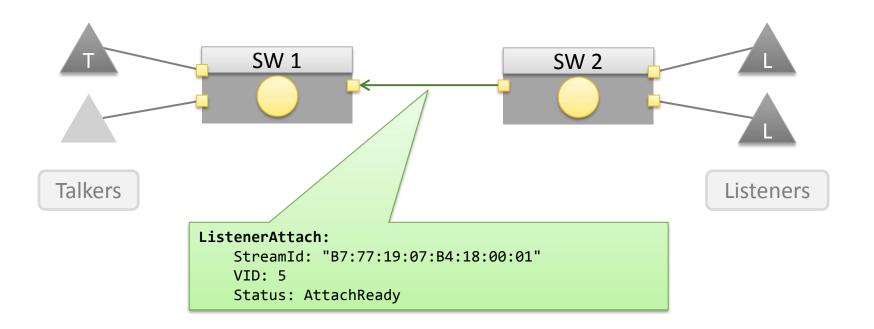


Reservation on SW 2 Successful



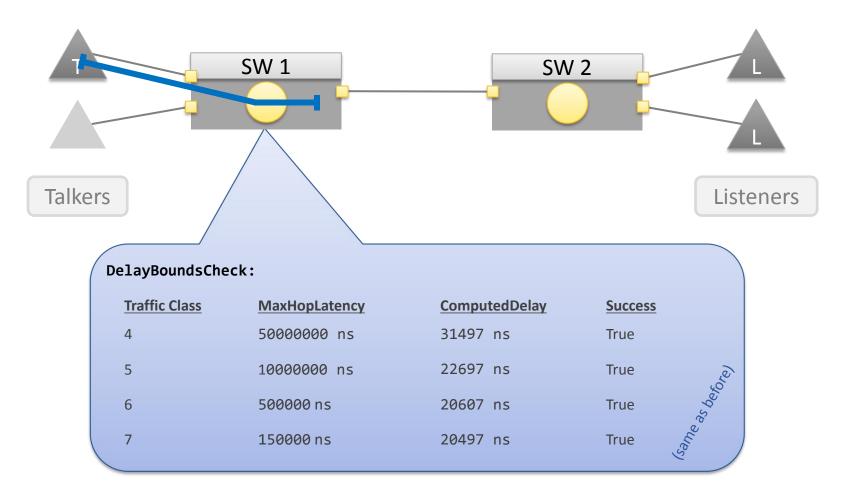
UNI WÜ 13

SW 2 Forwards the LA to SW 1



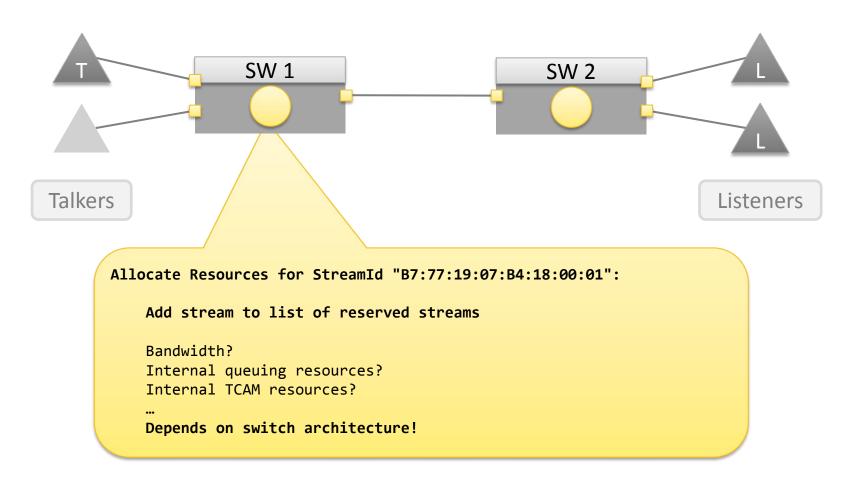


Bounds Check on SW 1 (EgressPort: SW 2)



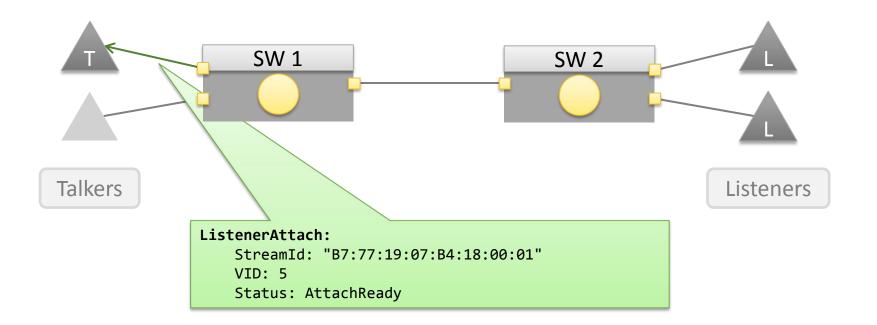


Reservation on SW 1 Successful



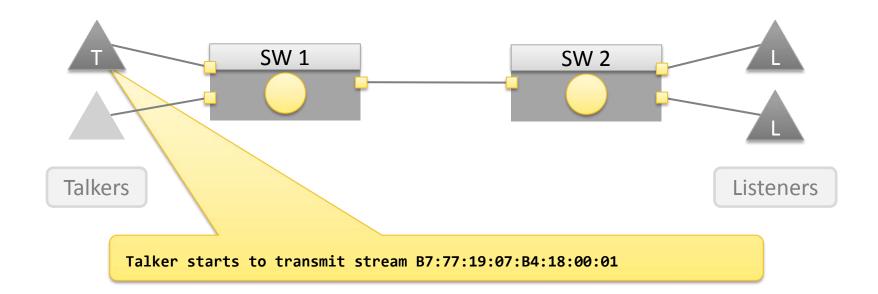
UNI WÜ 3

SW 1 Forwards the LA to the Talker



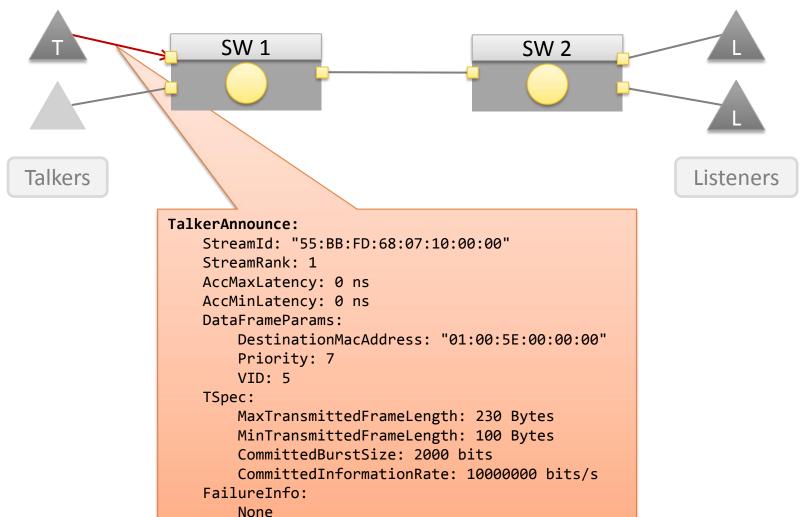


Talker Received Successful LA and Starts to Transmit



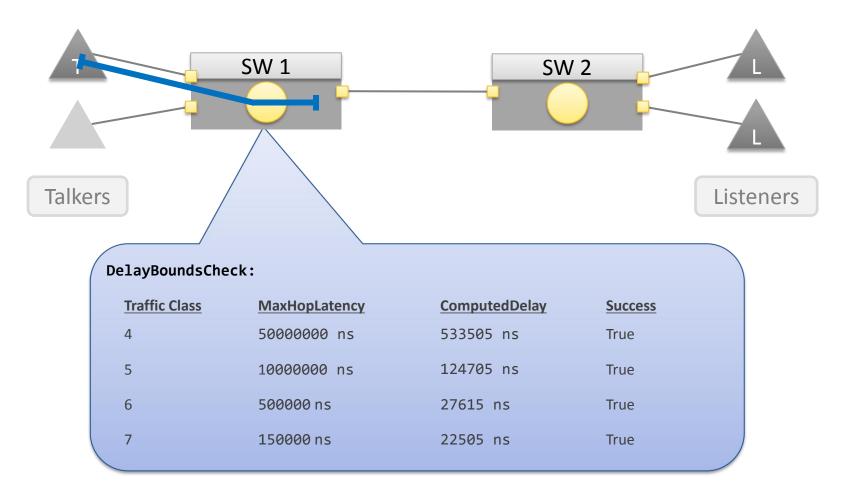


Stream 2



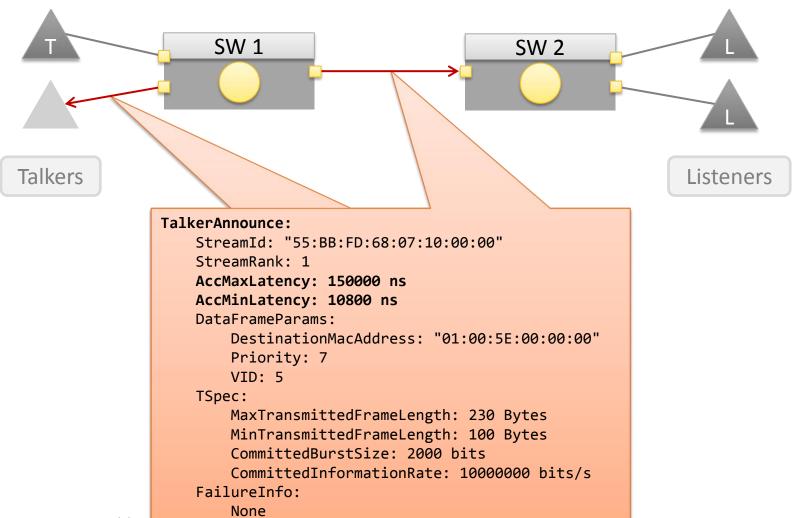
UNI WÜ 13

Bounds Check on SW 1 (EgressPort: SW 2)



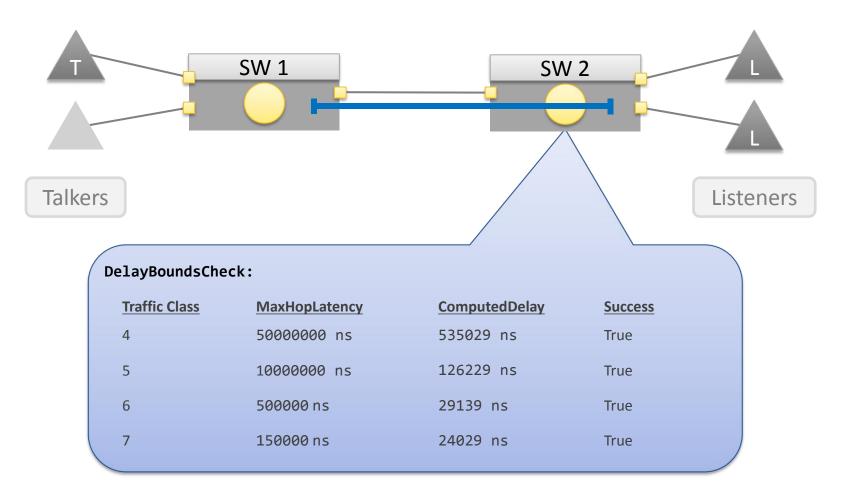


Adjusted TA is Propagated to other Ports



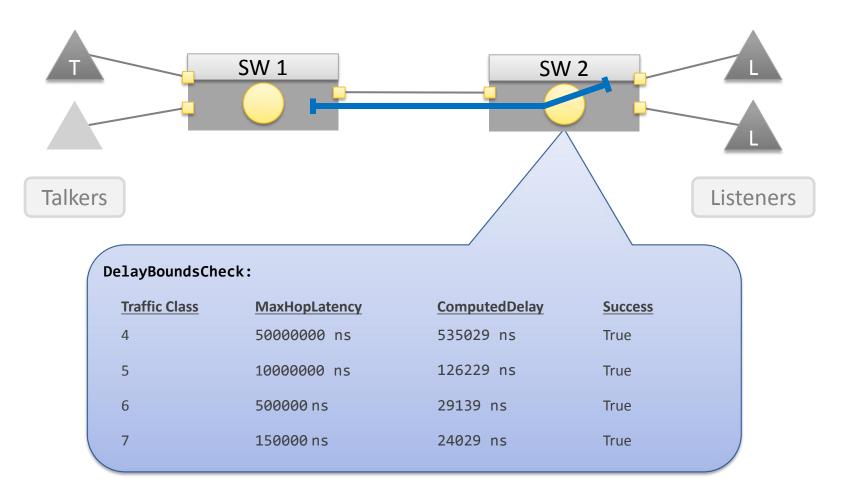
UNI WÜ **i**3

Bounds Check on SW 2 (EgressPort: Listener 2)



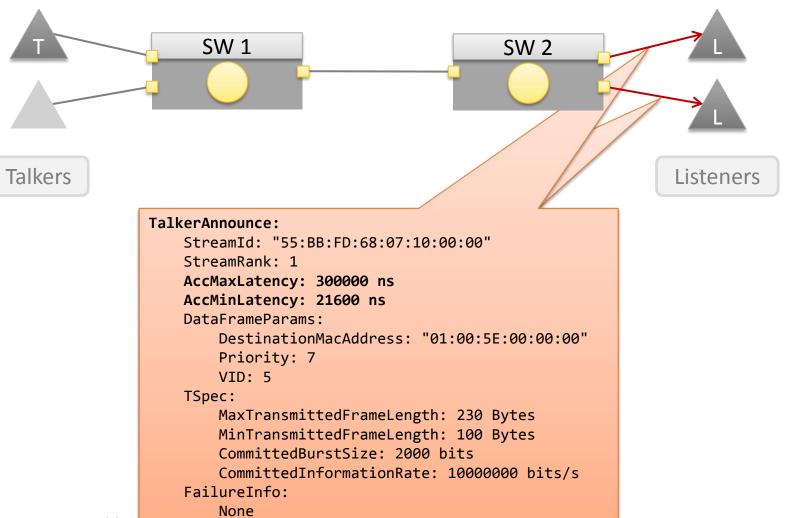


Bounds Check on SW 2 (EgressPort: Listener 1)





Adjusted TA is Propagated to other Ports

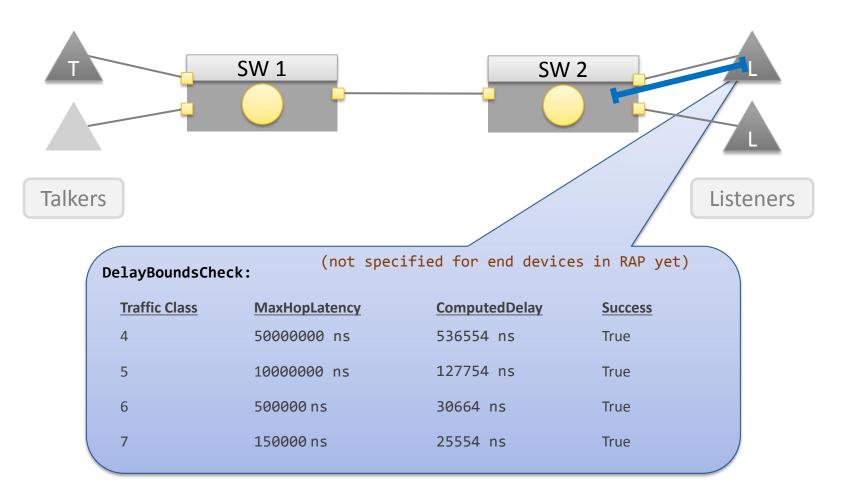


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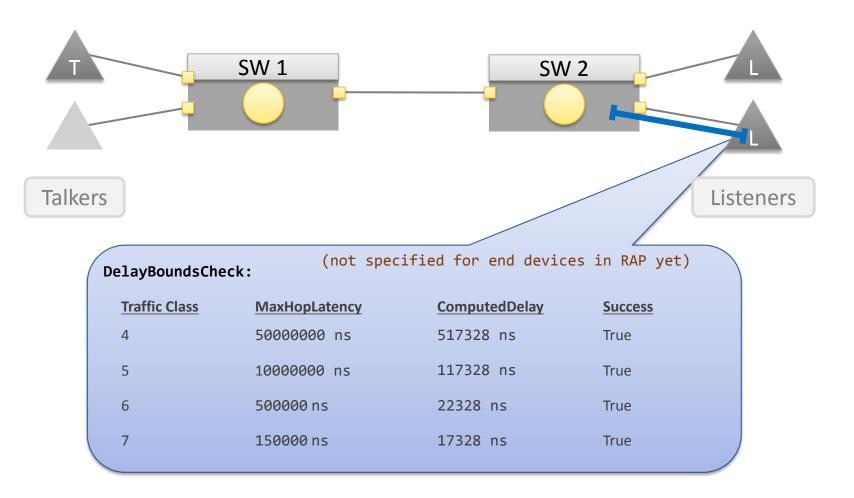
i3

Bounds Check on Listener 1



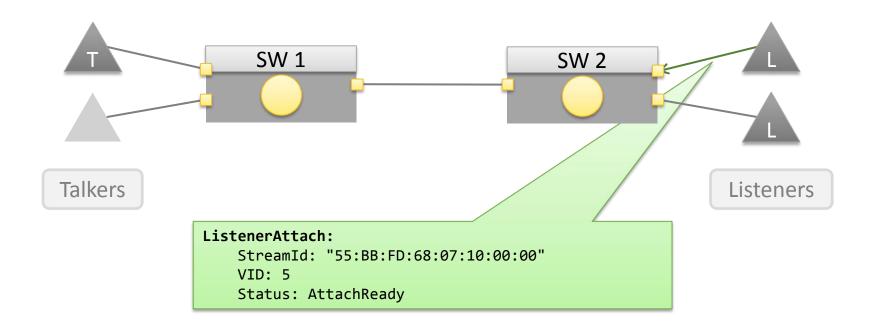
UNI WÜ

Bounds Check on Listener 2



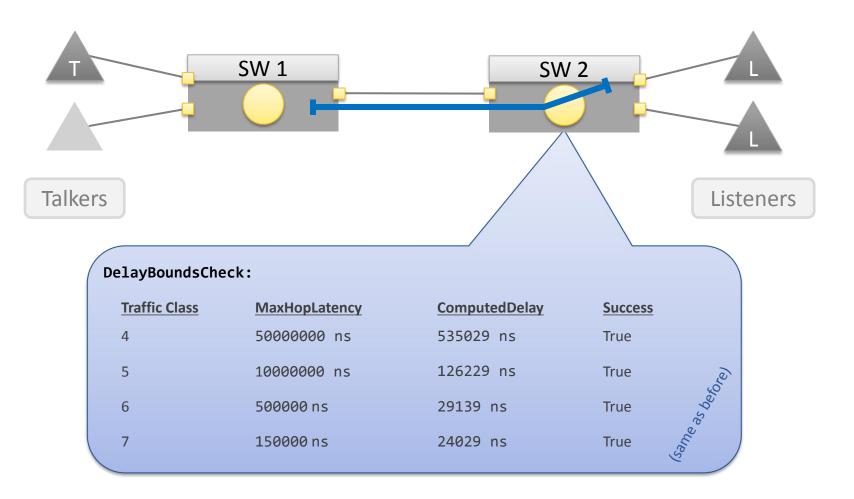
UNI WÜ

Listener 1 sends LA



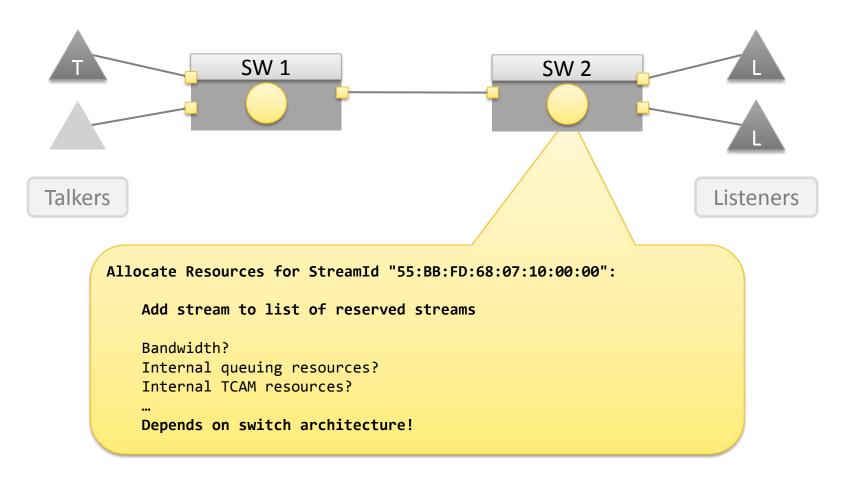


Bounds Check on SW 2 (EgressPort: Listener 1)



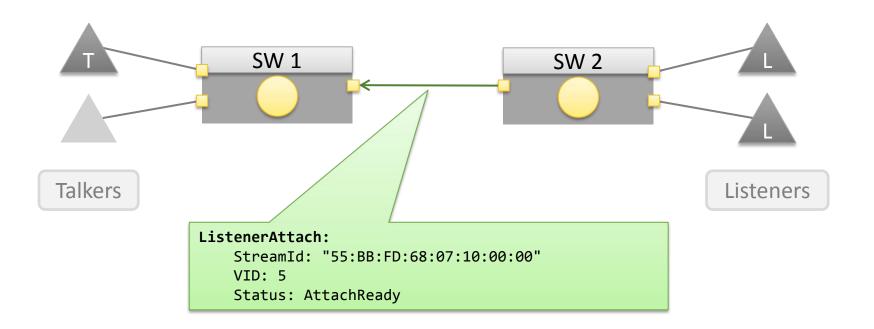
UNI WÜ

Reservation on SW 2 (Port 1) Successful



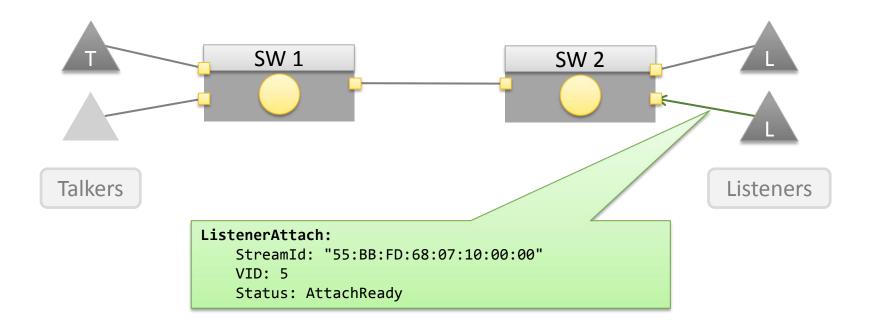
UNI WÜ **B**3

SW 2 Forwards the LA to SW 1



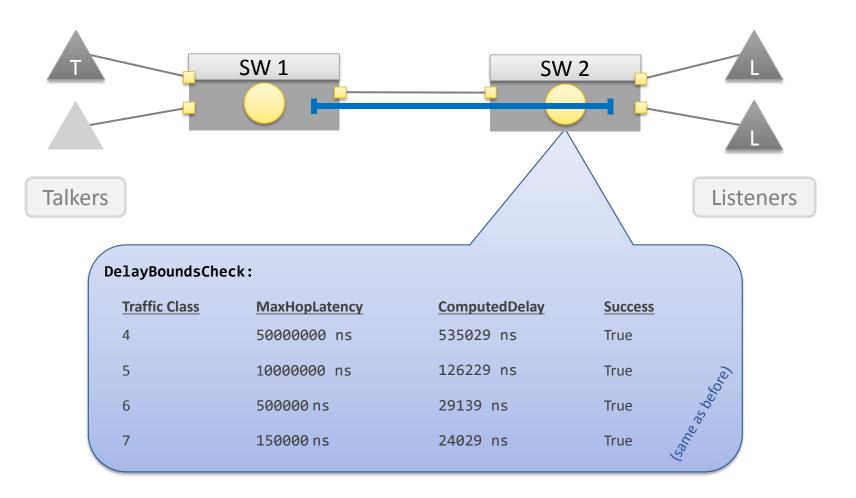


Listener 2 sends LA



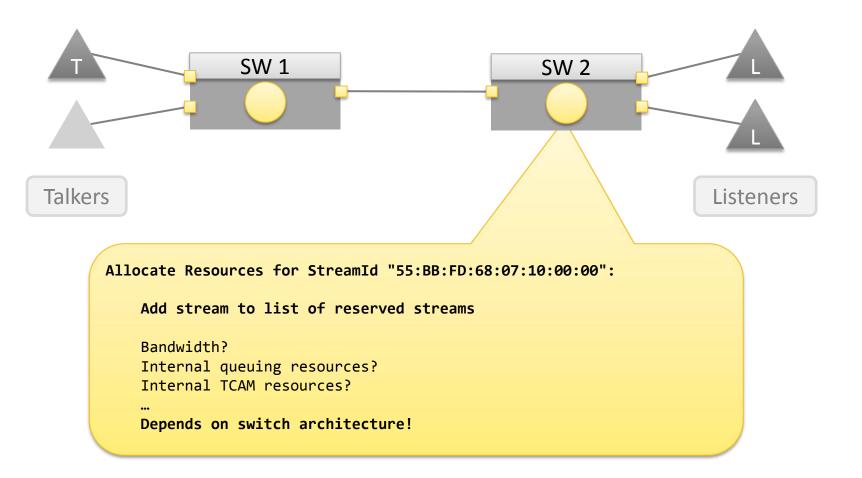


Bounds Check on SW 2 (EgressPort: Listener 2)



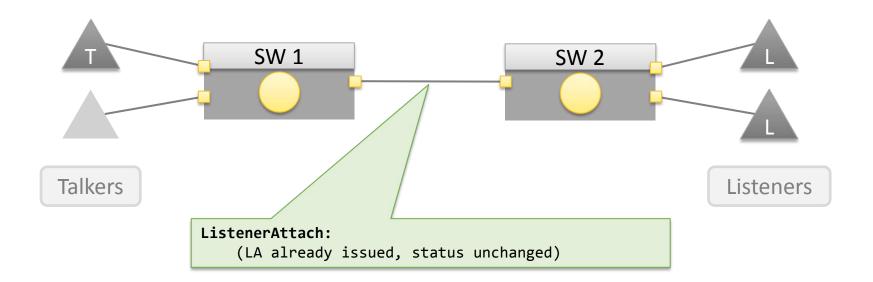


Reservation on SW 2 (Port 2) Successful



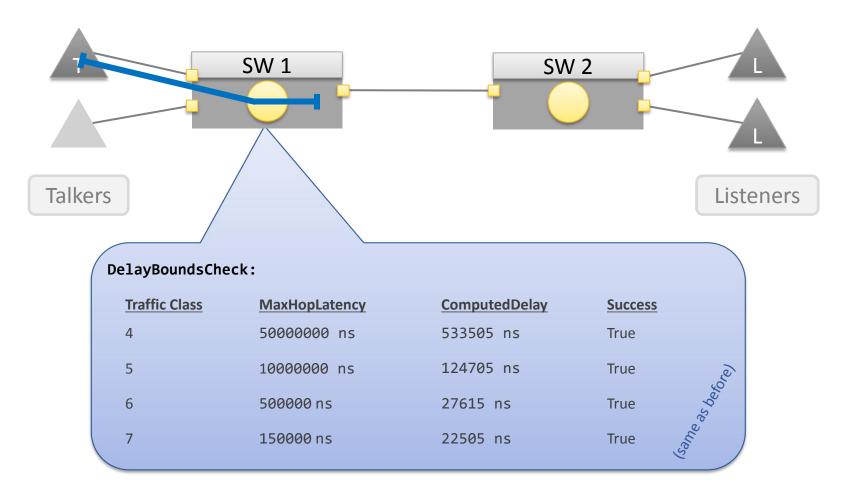
UNI WÜ **B**3

SW 2 does Nothing, Existing LA for this Stream Unchanged



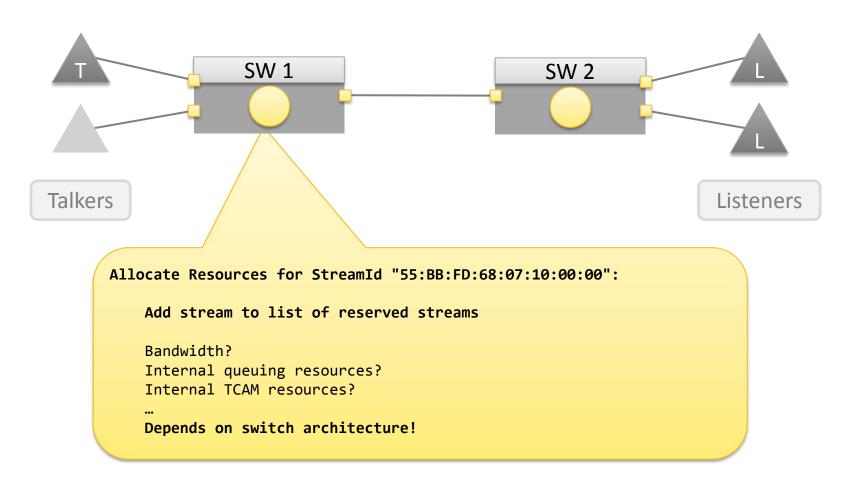


Bounds Check on SW 1 (EgressPort: SW 2)



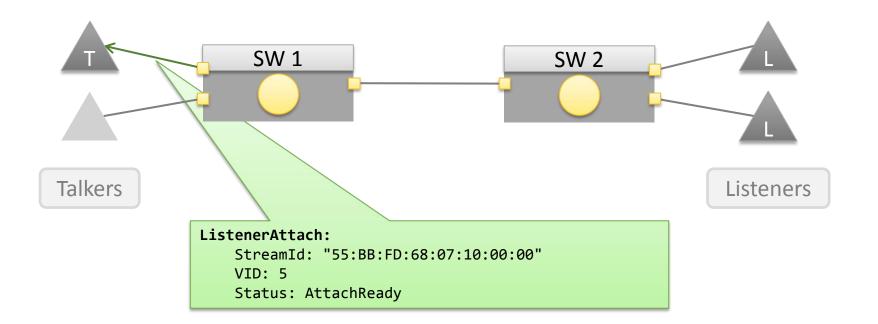
UNI WÜ

Reservation on SW 1 Successful



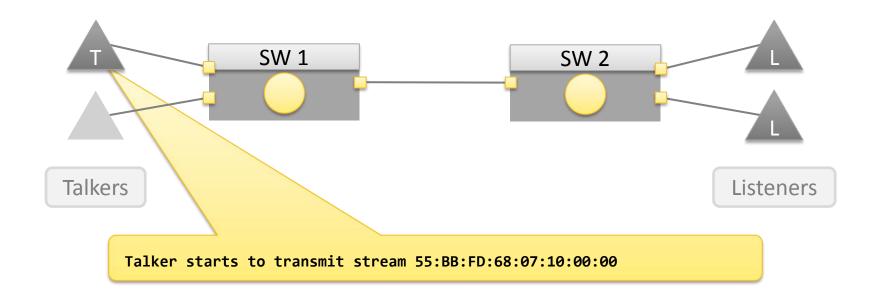
UNI WÜ

SW 1 Forwards the LA to the Talker





Talker Received Successful LA and Starts to Transmit





THANK YOU!

Questions, comments, suggestions?



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