

State-of-the-Art Fault-Tolerant Clock Synchronization for Ultra-high Reliable Systems

Wilfried Steiner

- Byzantine Fault-Tolerance
- SAE AS6802 Byzantine Fault-Tolerant (Clock) Synchronization
- Protocol Verification for Ultrahigh Reliable Systems
- Conclusion

ТГГесһ

- Byzantine Fault-Tolerance
- SAE AS6802 Byzantine Fault-Tolerant (Clock) Synchronization
- Protocol Verification for Ultrahigh Reliable Systems
- Conclusion

Byzantine Fault-Tolerance

(3k+1) clocks are necessary to tolerate the Byzantine Failure of k clocks (e.g., $k=2 \rightarrow 7$ clocks)

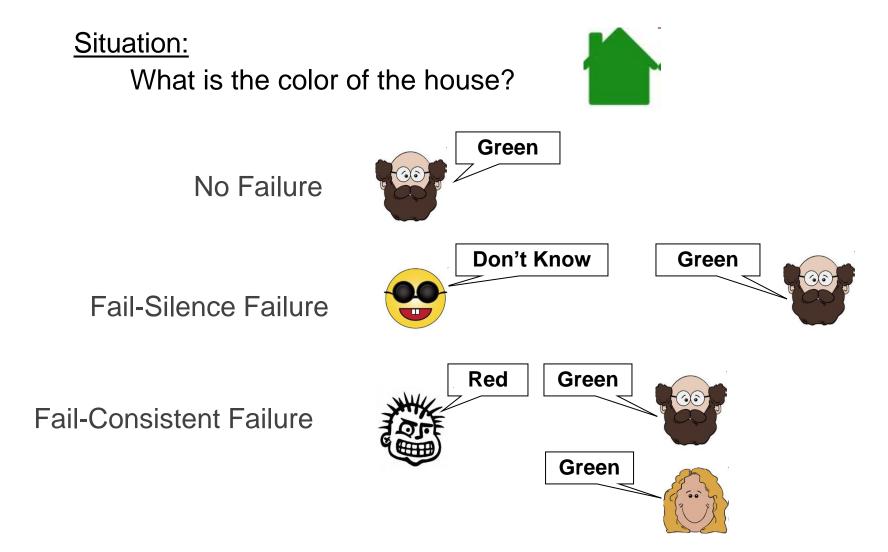
- Ultrahigh reliable systems require a system failure rate to be in the order of 10⁻⁹ failures/h or lower.
- Only a distributed fault-tolerant computer system with well-defined fault-containment units (FCUs) achieves the ultrahigh reliability requirement.
- FCUs may fail in a Byzantine failure mode and the distributed computer system must be designed to mitigate this failure mode.
- In order to tolerate k Byzantine failures n=3k+1 nodes are needed.
- The concept to Byzantine fault tolerance has been introduced in: Lamport, L., Shostak, R., Pease, M. (1982). *The Byzantine Generals Problem*. Comm. ACM TOPLAS. Vol. 4 (3). (pp.382-401).



https://www.ieee802.org/1/files/public/docs2012/new-avbwsteiner-fault-tolerant-clock-synchronization-0112-v01.pdf

Tlech

Byzantine Fault-Tolerance (cont.)



ТГГесһ

Byzantine Fault-Tolerance (cont.)

Situation:

What is the color of the house?



Static Situation – one Truth

Situation:

What is the color of the ball ?

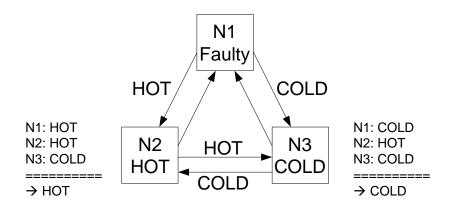


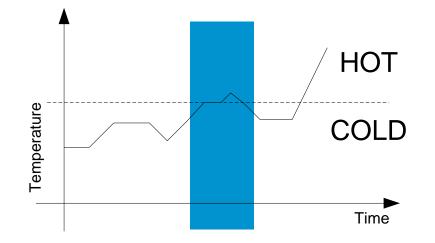
Dynamic Situation – >one Truth

Byzantine Fault-Tolerance (cont.)

A distributed system that measures the temperature of a vessel shall raise an alarm when the temperature exceeds a certain threshold. The system shall tolerate the arbitrary failure of one node. How many nodes are required?

How many messages are required?





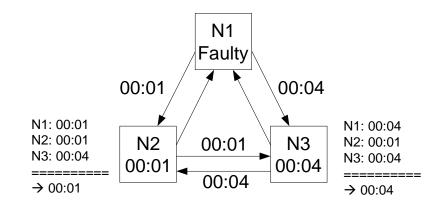
In general, three nodes are insufficient to tolerate the arbitrary failure of a single node. The two correct nodes are not always able to agree on a value. A decent body of scientific literature exists that address this problem of dependable systems, in particular dependable communication.

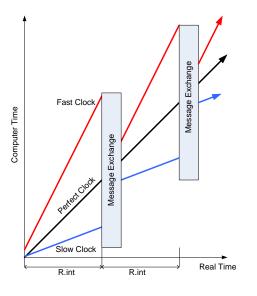
Byzantine Fault-Tolerance (cont.)

A distributed system in which all nodes are equipped with local clocks, all clocks shall become and remain synchronized.

The system shall tolerate the arbitrary failure of one node.

How many nodes are required? How many messages are required?



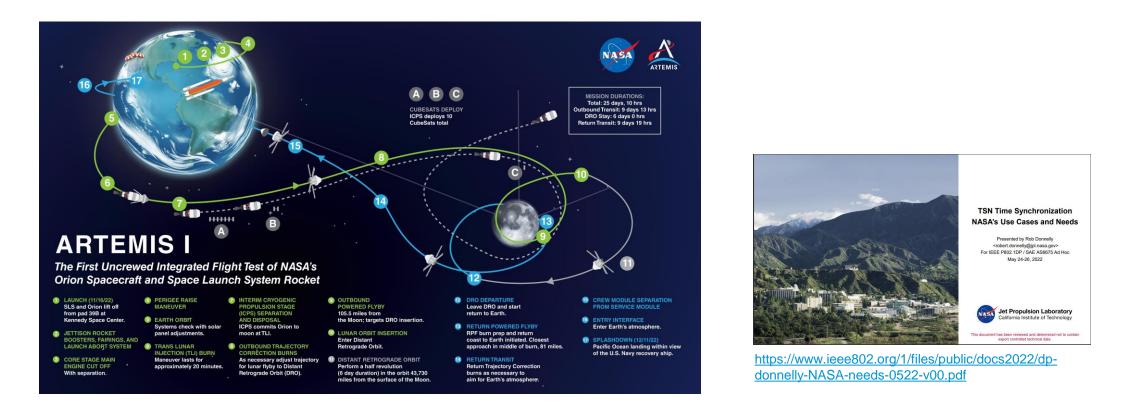


In general, three nodes are insufficient to tolerate the arbitrary failure of a single node. The two correct nodes are not always able to bring their clocks into close agreement. A decent body of scientific literature exists that address this problem of fault-tolerant clock synchronization.

- Byzantine Fault-Tolerance
- SAE AS6802 Byzantine Fault-Tolerant (Clock) Synchronization
- Protocol Verification for Ultrahigh Reliable Systems
- Conclusion

SAE AS6802^{*)}: Current State-of-the-Art in Fault-Tolerant (Clock) Synchronization Protocols for Avionics

First mission in 2014, last mission in 2022, next mission in 2024 (first crewed flight)

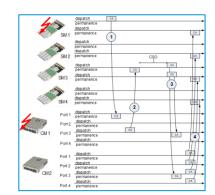


^{*)} Standard-relevant patent families: EP2297885, EP2297886

ТГГесһ

SAE AS6082 standardizes

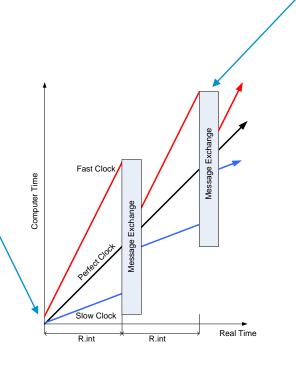
Fault-Tolerant Startup and Restart Protocol



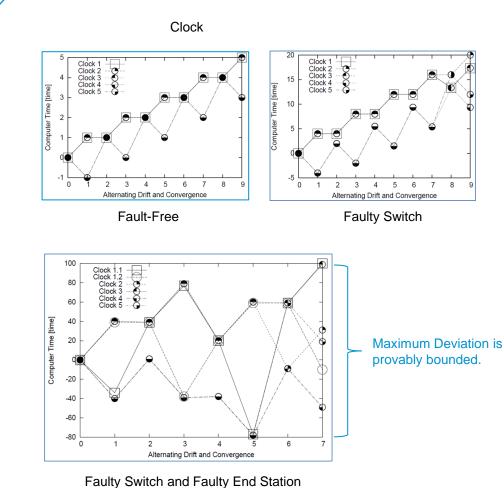
Synchronization Masters (SM) and Compression Masters (CM) interaction

	SM 1	SM 2	SM 3	SM 4	SM 5		SM 1	SM 2	SM 3	SM 4	SM 5
A:	Flood	Unsync	Unsync	Unsync	Flood	-l:	Tent	Sync	Tent	Sync	Sync
	CA CA				CA						
					UA ↓	J:	Sync	Unsync	Sync	Unsync	Unsync
B:	Flood	Unsync	Unsync	Unsync	Unsync		- /				
C:	Wait	Unsync	Unsync	Unsync	Wait	K:	Integrate	Unsync	Sync	Unsync	Unsync
	ÎN				ÎN ÎN			čs čs			
		IN ↓	IN IN	IN IN	IN ¥		Integrate	Wait	Wait	Wait	Wait
D:	Tent	Unsync	Unsync	Unsync	Tent	L.	integrate	wan	wan	Wan	Wait
E:	Unsync	Unsync	Sync	Sync	Unsync	м	Unsync	Wait	Wait	Wait	Wait
		čs čs	↓ ↓ N N	in in			čs	ÎN ÎN	ÎN ÎN	ÎN ÎN	ÎN ÎN
	CS	IN IN	IN IN	IN IN	IN IN			CS	CS	IN	IN J.
F:	Unsync	Unsync	Sync	Sync	Unsync	N:	Unsync	Wait	Wait	Tent	Tent
G:	Flood	Sync	Sync	Sync	Sync	0	Unsync	Flood	Flood	Tent	Tent
	ČA CA		CA				ÇA	CA CA CA	CA CA CA	ÇA	ÇA
H:	Flood	Sync	Sync	Sync	Sync	P:	Unsync	Flood	Flood	Unsync	Unsync
Е	Tent	Sync	Tent	Sync	Sync	Q	Unsync	Sync	Sync	Sync	Sync

Startup w/ faulty Switch and Faulty End Station



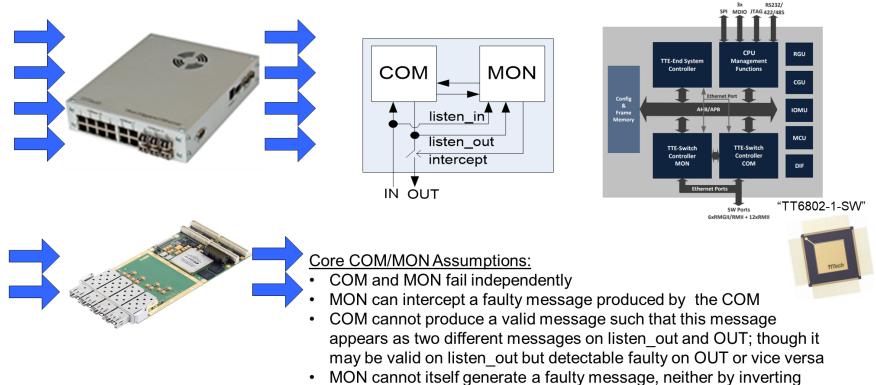
Fault-Tolerant Clock Synchronization Protocol



T[**F**ech

SAE AS6802 does not standardize

how to design devices to a sufficient quality level, e.g., Commander/Monitor (COM/MON) Structures



listen_out to an output, nor by toggling the intercept signal

- Byzantine Fault-Tolerance
- SAE AS6802 Byzantine Fault-Tolerant (Clock) Synchronization
- Protocol Verification for Ultrahigh Reliable Systems
- Conclusion

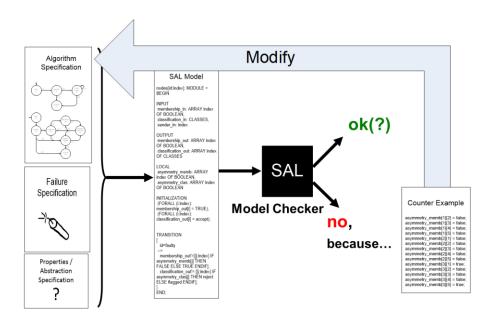
Formal Verification of SAE AS6802 Protocols

Formal Verification (model-checking or theorem proving) is state-of-the-art for ultrahigh reliable systems.

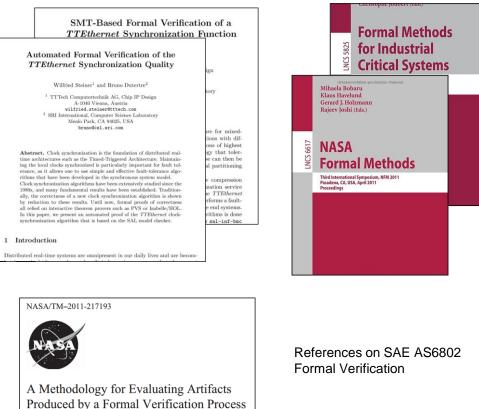
Radu I Siminiceanu

Paul S. Miner and Suzette Person Langley Research Center, Hampton, Virginia

National Institute of Aerospace, Hampton, Virginia



Formal Verification w/ Model Checking



- Byzantine Fault-Tolerance
- SAE AS6802 Byzantine Fault-Tolerant (Clock) Synchronization
- Protocol Verification for Ultrahigh Reliable Systems
- Conclusion

Conclusions

Let's not re-invent the wheel but build on proven solutions.

- Byzantine fault tolerance is state-of-the-art in ultra-high reliable systems.
- SAE AS6802 (Time-Triggered Ethernet) tolerates Byzantine faults and is the state-of-the-art in fault-tolerant (clock) synchronization for ultra-high reliable systems such as avionics.
- It is state-of-the-art to use formal methods like model checking and theorem proving to verify protocol correctness. All SAE AS6802 protocols have been formally verified.
- SAE AS6802 has been standardized in 2011 and concluded a revision in 2023. It has been proven in use for ultra-high reliable systems.
- Proposal 1: work towards using SAE AS6802 synchronization protocols as default FTM.
 - Maybe prepare an SAE AS6802 "Appendix D" for "Time-Triggered Ethernet Realization on IEEE802.1 (TSN)" (today Appendix B is for IEEE 802.3 and Appendix C is for ARINC 664-p7).
- Proposal 2: let's define a "slimmed-down" version of gPTP to simplify a chip-only implementation.
 - E.g., make peer-delay-measurements and clock rate measurements optional.

Trech