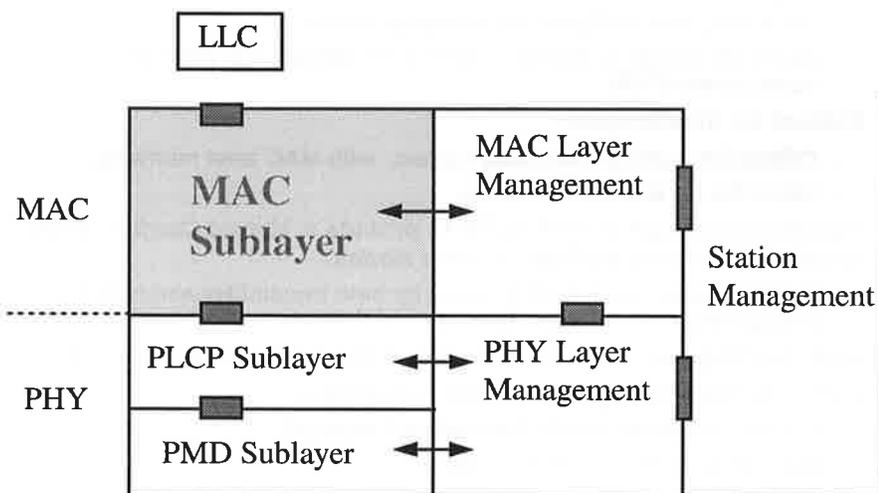


802.11 Tutorial

802.11 MAC Entity: MAC Basic Access Mechanism Privacy and Access Control

802.11 Protocol Entities



Main Requirements

- **Single MAC to support multiple PHYs.**
 - Support single and multiple channel PHYs.
 - and PHYs with different *Medium Sense* characteristics
- **Should allow overlap of multiple networks in the same area and channel space.**
 - Need to be able to share the medium.
 - Allow re-use of the same medium.
- **Need to be *Robust for Interference*.**
 - Microwave interferers
 - Other un-licensed spectrum users
 - Co-channel interference
- **Need mechanisms to deal with *Hidden Nodes*.**
- **Need provisions for *Time Bounded Services*.**
- **Need provisions for *Privacy and Access Control*.**

Copyright ©1996 IEEE. All rights reserved. This contains parts from an unapproved draft, subject to change

3

Basic Access Protocol Features

- **Use Distributed Coordination Function (DCF) for efficient medium sharing without overlap restrictions.**
 - Use CSMA with Collision Avoidance derivative.
 - Based on *Carrier Sense* function in PHY called ***Clear Channel Assessment (CCA)***.
- **Robust for interference.**
 - ***CSMA/CA + ACK*** for unicast frames, with MAC level recovery.
 - CSMA/CA for Broadcast frames.
- **Parameterized use of RTS / CTS to provide a *Virtual Carrier Sense* function to protect against *Hidden Nodes*.**
 - ***Duration*** information is distributed by both transmitter and receiver through separate RTS and CTS Control Frames.
- **Includes fragmentation to cope with different PHY characteristics.**
- **Frame formats to support the access scheme**
 - For Infrastructure and Ad-Hoc Network support
 - and ***Wireless Distribution System***.

Copyright ©1996 IEEE. All rights reserved. This contains parts from an unapproved draft, subject to change

4

802.11 Tutorial
March 96

CSMA/CA Explained

- **Reduce collision probability where mostly needed.**
 - Stations are waiting for medium to become free.
 - Select Random Backoff after a Defer, resolving contention to avoid collisions.
- **Efficient Backoff algorithm stable at high loads.**
 - Exponential Backoff window increases for retransmissions.
 - Backoff timer elapses only when medium is idle.
- **Implement different fixed priority levels.**
 - To allow immediate responses and PCF coexistence.

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 5

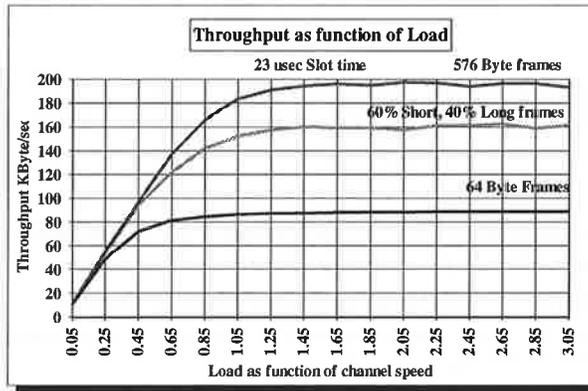
802.11 Tutorial
March 96

CSMA/CA + ACK protocol

- **Defer access based on Carrier Sense.**
 - CCA from PHY and *Virtual Carrier Sense* state.
- **Direct access when medium is sensed free longer then DIFS, otherwise defer and backoff.**
- **Receiver of directed frames to return an ACK immediately when CRC correct.**
 - When no ACK received then retransmit frame after a random backoff (up to maximum limit).

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 6

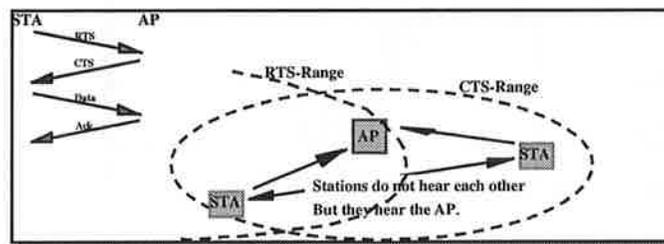
Throughput Efficiency



- **Efficient and stable throughput.**
 - **Stable throughput at overload conditions.**
 - **To support “Bursty Traffic” characteristics.**

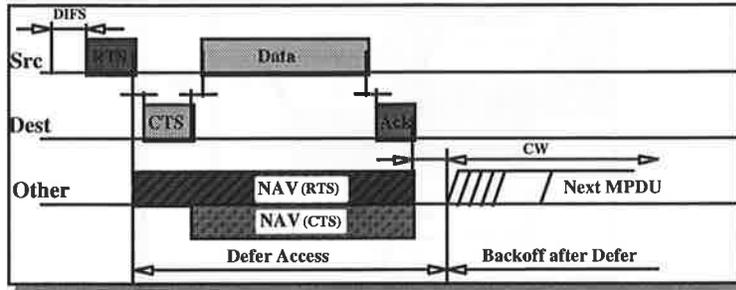
“Hidden Node” Problem

- **Transmitters contending for the medium may not “Hear each other” as shown below.**



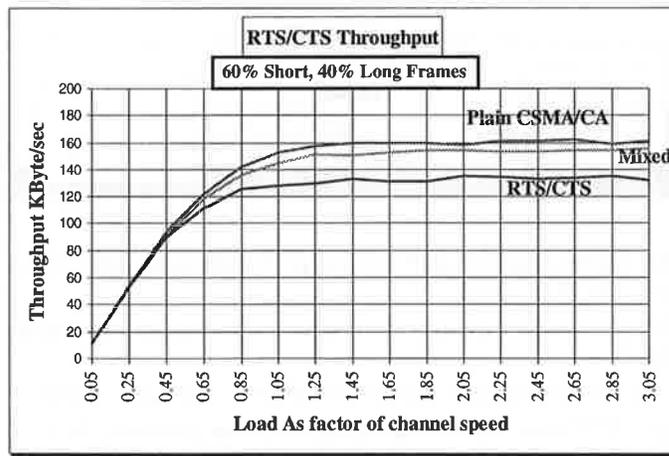
- **Separate Control frame exchange (RTS / CTS) between transmitter and receiver will *Reserve the Medium* for subsequent data access.**
 - **Duration is distributed around both Tx and Rx station.**

“Hidden Node” Provisions



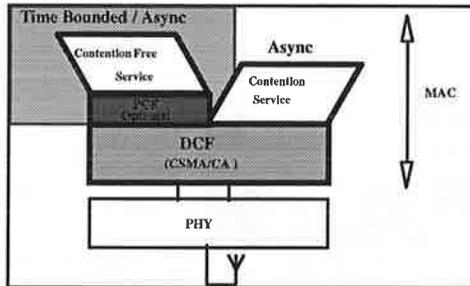
- Duration field in RTS and CTS frames distribute **Medium Reservation** information which is stored in a **Net Allocation Vector (NAV)**.
- Defer on either NAV or "CCA" indicating **Medium Busy**.
- Use of RTS / CTS is optional but must be implemented.
- Use is controlled by a **RTS_Threshold** parameter per station.
 - To limit overhead for short frames.

RTS/CTS Overhead Impact



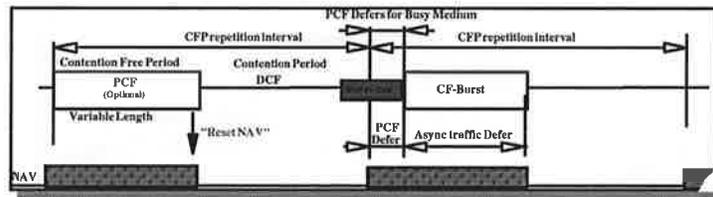
Good mixed Throughput (long inbound frames) efficiency.

Optional Point Coordination Function (PCF)



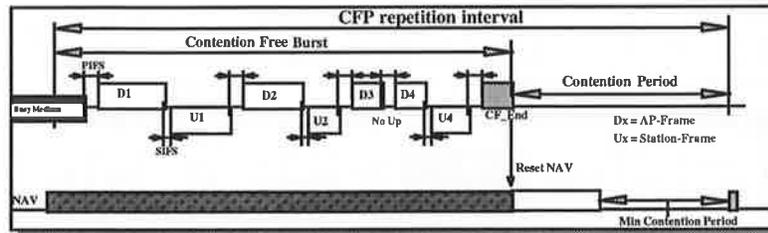
- Contention Free Service uses Point Coordination Function (PCF) on a DCF Foundation.
 - PCF can provide lower *transfer delay* variations to support **Time Bounded Services**.
 - Async Data, Voice or mixed implementations possible.
 - Point Coordinator resides in AP.
- Coexistence between Contention and optional Contention Free does not burden the implementation.

Contention Free operation



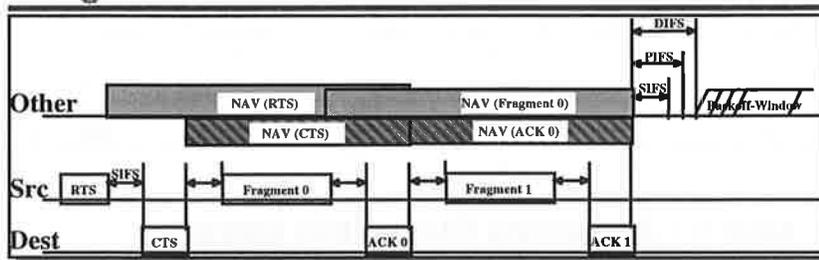
- Alternating **Contention Free** and **Contention** operation under PCF control.
- NAV prevents **Contention** traffic until reset by the last PCF transfer.
 - So variable length **Contention Free** period per interval.
- Both PCF and DCF defer to each other causing PCF Burst start variations.

PCF Burst



- CF-Burst by Polling bit in CF-Down frame.
- Immediate response by Station on a CF_Poll.
- Stations to maintain NAV to protect CF-traffic.
- Responses can be variable length.
- "Reset NAV" by last (CF_End) frame from AP.
- "ACK Previous Frame" bit in Header.

Fragmentation



- Burst of Fragments which are individually acknowledged.
 - For Unicast frames only.
- Random backoff and retransmission of failing fragment when no ACK is returned.
- Duration information in data fragments and Ack frames causes NAV to be set, for medium reservation mechanism.

802.11 Tutorial March 96

Frame Formats

| | | | | | | | | | | |
|-----------------------------|-------------|---------|--------|---------|------------------|--------|------------|-----------|-----|------|
| Bytes: | | | | | | | | | | |
| 2 | 2 | 6 | 6 | 6 | 2 | 6 | 0-2312 | 4 | | |
| Frame Control | Duration ID | Addr 1 | Addr 2 | Addr 3 | Sequence Control | Addr 4 | Frame Body | CRC | | |
| ← 802.11 MAC Header → | | | | | | | | | | |
| Bits: 2 2 4 1 1 1 1 1 1 1 1 | | | | | | | | | | |
| Protocol Version | Type | SubType | To DS | From DS | More Frag | Retry | Pwr Mgt | More Data | WEP | Rsvd |
| Frame Control Field | | | | | | | | | | |

- **MAC Header format differs per Type:**
 - Control Frames (several fields are omitted)
 - Management Frames
 - Data Frames
- **Includes Sequence Control Field for filtering of duplicate caused by ACK mechanism.**

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 15

802.11 Tutorial March 96

Address Field Description

| To DS | From DS | Address 1 | Address 2 | Address 3 | Address 4 |
|-------|---------|-----------|-----------|-----------|-----------|
| 0 | 0 | DA | SA | BSSID | N/A |
| 0 | 1 | DA | BSSID | SA | N/A |
| 1 | 0 | BSSID | SA | DA | N/A |
| 1 | 1 | RA | TA | DA | SA |

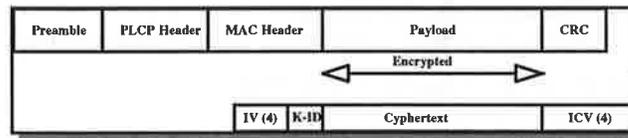
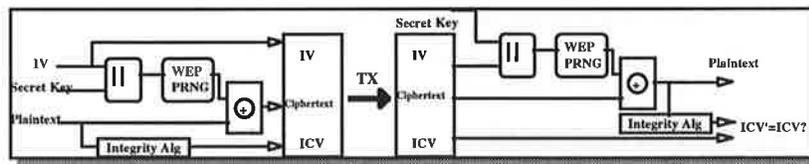
- **Addr 1 = All stations filter on this address.**
- **Addr 2 = Transmitter Address (TA)**
 - Identifies transmitter to address the ACK frame to.
- **Addr 3 = Dependent on *To* and *From DS* bits.**
- **Addr 4 = Only needed to identify the original source of WDS (*Wireless Distribution System*) frames.**

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 16

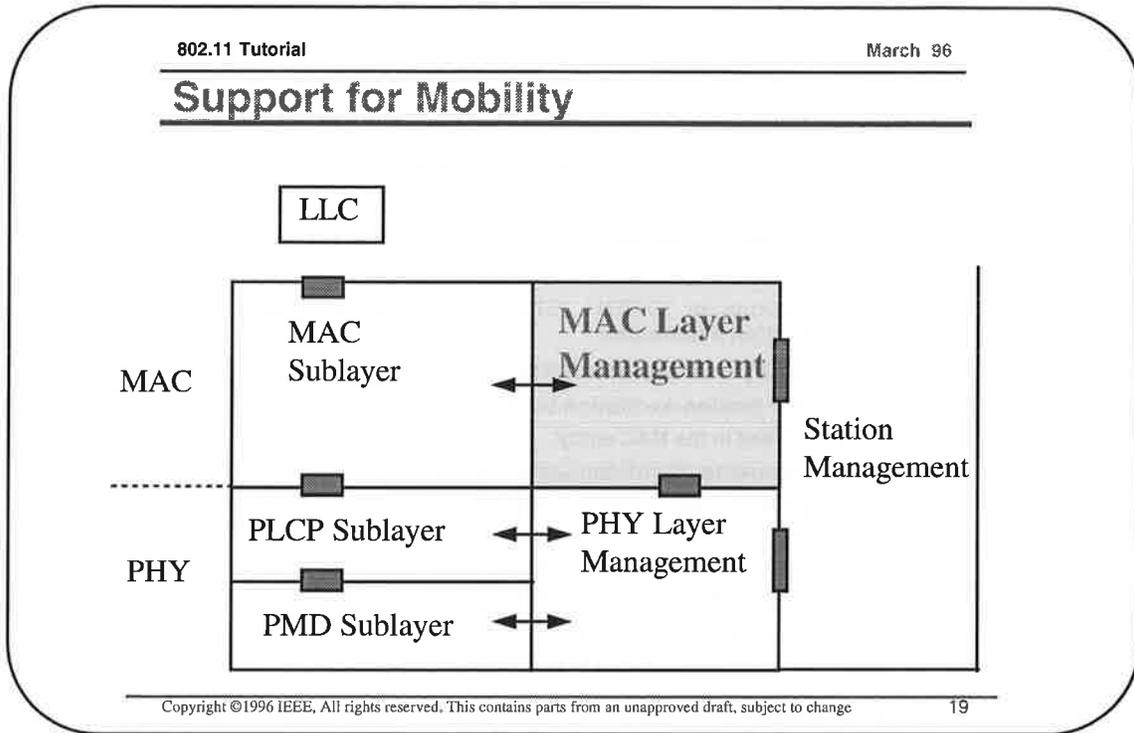
Privacy and Access Control

- **Goal of 802.11 is to provide “Wired Equivalent Privacy” (WEP)**
 - Usable worldwide
- **802.11 provides for an Authentication mechanism**
 - To aid in access control.
 - Has provisions for “OPEN”, “Shared Key” or proprietary authentication extensions.
- **Optional (WEP) Privacy mechanism defined by 802.11.**
 - Limited for Station-to-Station traffic, so not “end to end”.
 - » Embedded in the MAC entity.
 - Only implements “Confidentiality” function.
 - Uses RC4 PRNG algorithm based on:
 - » a 40 bit secret key (No Key distribution standardized)
 - » and a 24 bit IV that is send with the data.
 - » includes an ICV to allow integrity check.
 - Only payload of Data frames are encrypted.
 - » Encryption on per MPDU basis.

Privacy Mechanism



- **WEP bit in Frame Control Field indicates WEP used.**
 - Each frame can have a new IV, or IV can be reused for a limited time.
 - If integrity check fails then frame is ACKed but discarded.



- 802.11 Tutorial March 96
-
- ## MAC Management Layer
-
- **Synchronization**
 - finding and staying with a WLAN
 - Synchronization functions
 - » TSF Timer, Beacon Generation
 - **Power Management**
 - sleeping without missing any messages
 - Power Management functions
 - » periodic sleep, frame buffering, Traffic Indication Map
 - **Association and Reassociation**
 - Joining a network
 - Roaming, moving from one AP to another
 - Scanning
 - **Management Information Base**
- Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 20

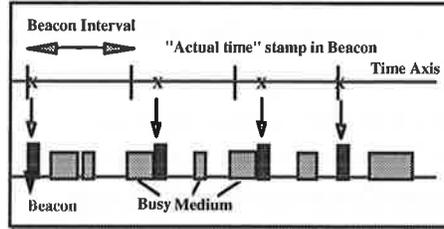
Synchronization in 802.11

- **Timing Synchronization Function (TSF)**
- **Used for Power Management**
 - Beacons sent at well known intervals
 - All station timers in BSS are synchronized
- **Used for Point Coordination Timing**
 - TSF Timer used to predict start of Contention Free burst
- **Used for Hop Timing for FH PHY**
 - TSF Timer used to time Dwell Interval
 - All Stations are synchronized, so they hop at same time.

Synchronization Approach

- **All stations maintain a local timer.**
- **Timing Synchronization Function**
 - keeps timers from all stations in synch
 - AP controls timing in infrastructure networks
 - distributed function for Independent BSS
- **Timing conveyed by periodic Beacon transmissions**
 - Beacons contain Timestamp for the entire BSS
 - Timestamp from Beacons used to calibrate local clocks
 - not required to hear every Beacon to stay in synch
 - Beacons contain other management information
 - » also used for Power Management, Roaming

Infrastructure Beacon Generation



- APs send Beacons in infrastructure networks.
- Beacons scheduled at Beacon Interval.
- Transmission may be delayed by CSMA deferral.
 - subsequent transmissions at expected Beacon Interval
 - not relative to last Beacon transmission
 - next Beacon sent at Target Beacon Transmission Time
- Timestamp contains timer value at transmit time.

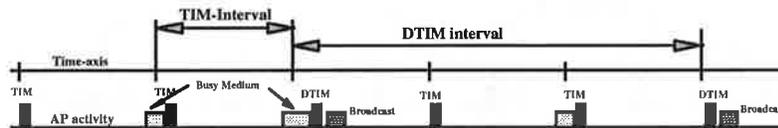
Power Management

- Mobile devices are battery powered.
 - *Power Management* is important for mobility.
- Current LAN protocols assume stations are always ready to receive.
 - Idle receive state dominates LAN adapter power consumption over time.
- How can we power off during idle periods, yet maintain an active session?
- 802.11 Power Management Protocol:
 - allows transceiver to be off as much as possible
 - is transparent to existing protocols
 - is flexible to support different applications
 - » possible to trade off throughput for battery life

Power Management Approach

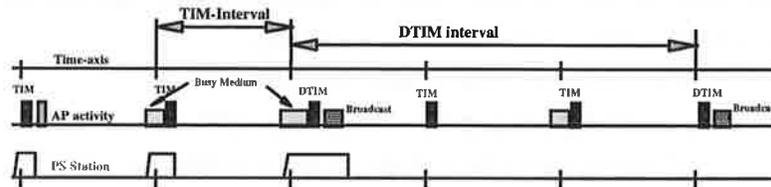
- **Allow idle stations to go to sleep**
 - station's power save mode stored in AP
- **APs buffer packets for sleeping stations.**
 - AP announces which stations have frames buffered
 - Traffic Indication Map (TIM) sent with every Beacon
- **Power Saving stations wake up periodically**
 - listen for Beacons
- **TSF assures AP and Power Save stations are synchronized**
 - stations will wake up to hear a Beacon
 - TSF timer keeps running when stations are sleeping
 - synchronization allows extreme low power operation
- **Independent BSS also have Power Management**
 - similar in concept, distributed approach

Infrastructure Power Management



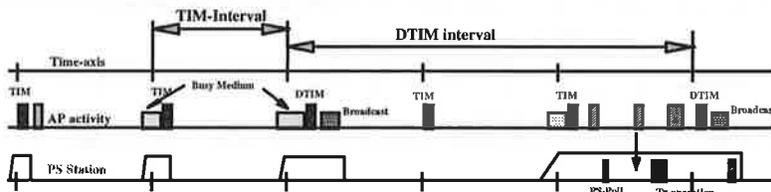
- **Broadcast frames are also buffered in AP.**
 - all broadcasts/multicasts are buffered
 - broadcasts/multicasts are only sent after DTIM
 - DTIM interval is a multiple of TIM interval

Infrastructure Power Management



- **Broadcast frames are also buffered in AP.**
 - all broadcasts/multicasts are buffered
 - broadcasts/multicasts are only sent after DTIM
 - DTIM interval is a multiple of TIM interval
- **Stations wake up prior to an expected (D)TIM.**

Infrastructure Power Management



- **Broadcast frames are also buffered in AP.**
 - all broadcasts/multicasts are buffered
 - broadcasts/multicasts are only sent after DTIM
 - DTIM interval is a multiple of TIM interval
- **Stations wake up prior to an expected (D)TIM.**
- **If TIM indicates frame buffered**
 - station sends PS-Poll and stays awake to receive data
 - else station sleeps again

802.11 Tutorial March 96

Wireless LAN Infrastructure Network

Station 1 Station 2 Access Point B Station 5 Station 6

Access Point A Station 4 Access Point C Station 7

- **Each Station is Associated with a particular AP**
 - Stations 1, 2, and 3 are associated with Access Point A
 - Stations 4 and 5 are associated with Access Point B
 - Stations 6 and 7 are associated with Access Point C

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 29

802.11 Tutorial March 96

Roaming

Station 1 Station 2 Access Point B Station 5 Station 6

Access Point A Station 4 Access Point C Station 7

- **Mobile stations may move...**

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 30

802.11 Tutorial March 96

Roaming

• Mobile stations may move...
– beyond the coverage area of their Access Point

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 31

802.11 Tutorial March 96

Roaming

• Mobile stations may move...
– beyond the coverage area of their Access Point
– but within range of another Access Point

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 32

802.11 Tutorial March 96

Roaming

- **Mobile stations may move...**
 - beyond the coverage area of their Access Point
 - but within range of another Access Point
- **Reassociation allows station to continue operation**

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 33

802.11 Tutorial March 96

Roaming Approach

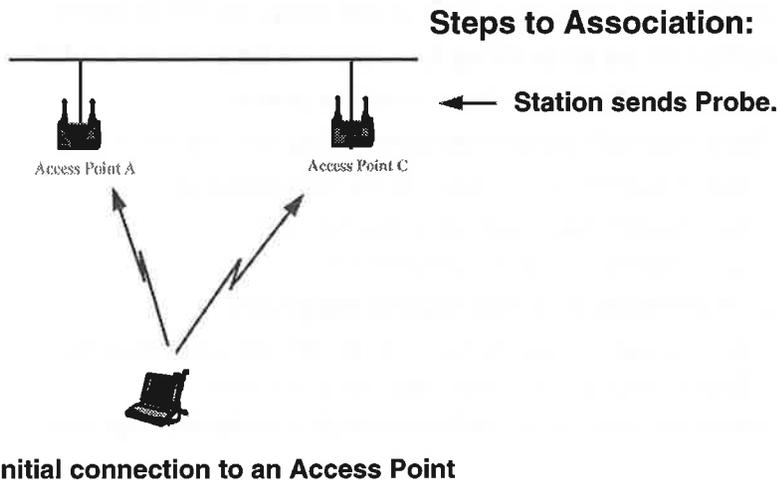
- **Station decides that link to its current AP is poor**
- **Station uses scanning function to find another AP**
 - or uses information from previous scans
- **Station sends Reassociation Request to new AP**
- **If Reassociation Response is successful**
 - then station has roamed to the new AP
 - else station scans for another AP
- **If AP accepts Reassociation Request**
 - AP indicates Reassociation to the Distribution System
 - Distribution System information is updated
 - normally old AP is notified through Distribution System

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 34

Scanning

- **Scanning required for many functions.**
 - finding and joining a network
 - finding a new AP while roaming
 - initializing an Independent BSS (ad hoc) network
- **802.11 MAC uses a common mechanism for all PHY.**
 - single or multi channel
 - passive or active scanning
- **Passive Scanning**
 - Find networks simply by listening for Beacons
- **Active Scanning**
 - On each channel
 - » Send a Probe, Wait for a Probe Response
- **Beacon or Probe Response contains information necessary to join new network.**

Active Scanning Example

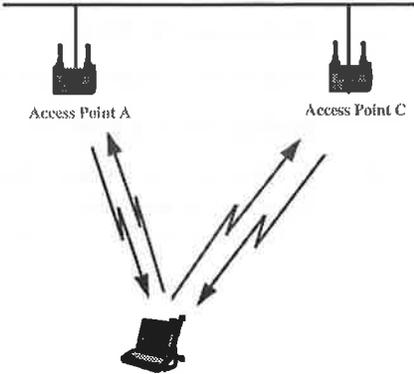


802.11 Tutorial March 96

Active Scanning Example

Steps to Association:

- ← Station sends Probe.
- APs send Probe Response.



Initial connection to an Access Point

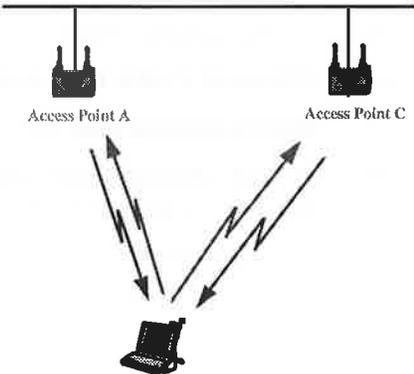
Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 37

802.11 Tutorial March 96

Active Scanning Example

Steps to Association:

- ← Station sends Probe.
- APs send Probe Response.
- Station selects best AP.**



Initial connection to an Access Point

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 38

802.11 Tutorial March 96

Active Scanning Example

Access Point A Access Point C

Steps to Association:

- ← Station sends Probe.
- APs send Probe Response.
- Station selects best AP.
- ← Station sends Association Request to selected AP.

Initial connection to an Access Point

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 39

802.11 Tutorial March 96

Active Scanning Example

Access Point A Access Point C

Steps to Association:

- ← Station sends Probe.
- APs send Probe Response.
- Station selects best AP.
- ← Station sends Association Request to selected AP.
- AP sends Association Response.

Initial connection to an Access Point

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 40

802.11 Tutorial March 96

Active Scanning Example

Initial connection to an Access Point
- ReAssociation follows a similar process

Steps to Association:

- ← Station sends Probe.
- APs send Probe Response.
- Station selects best AP.
- ← Station sends Association Request to selected AP.
- AP sends Association Response.

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 41

802.11 Tutorial March 96

MAC Management Frames

- **Beacon**
 - Timestamp, Beacon Interval, Capabilities, ESSID, Supported Rates, parameters
 - Traffic Indication Map
- **Probe**
 - ESSID, Capabilities, Supported Rates
- **Probe Response**
 - Timestamp, Beacon Interval, Capabilities, ESSID, Supported Rates, parameters
 - same for Beacon except for TIM
- **Association Request**
 - Capability, Listen Interval, ESSID, Supported Rates
- **Association Response**
 - Capability, Status Code, Station ID, Supported Rates

Copyright ©1996 IEEE, All rights reserved. This contains parts from an unapproved draft, subject to change 42

More MAC Management Frames

- **Reassociation Request**
 - Capability, Listen Interval, ESSID, Supported Rates, Current AP Address
- **Reassociation Response**
 - Capability, Status Code, Station ID, Supported Rates
- **Disassociation**
 - Reason code
- **Authentication**
 - Algorithm, Sequence, Status, Challenge Text
- **Deauthentication**
 - Reason

802.11 MAC

the end...