RPR MAC

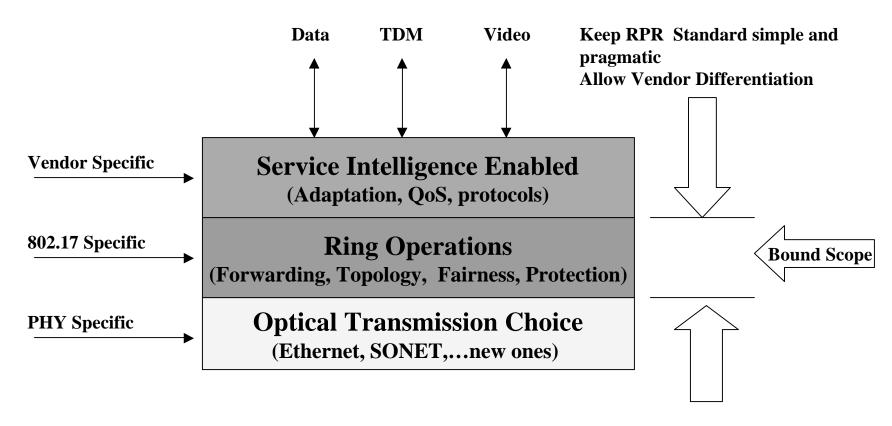
Distributed Cut-through Switching

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Agenda

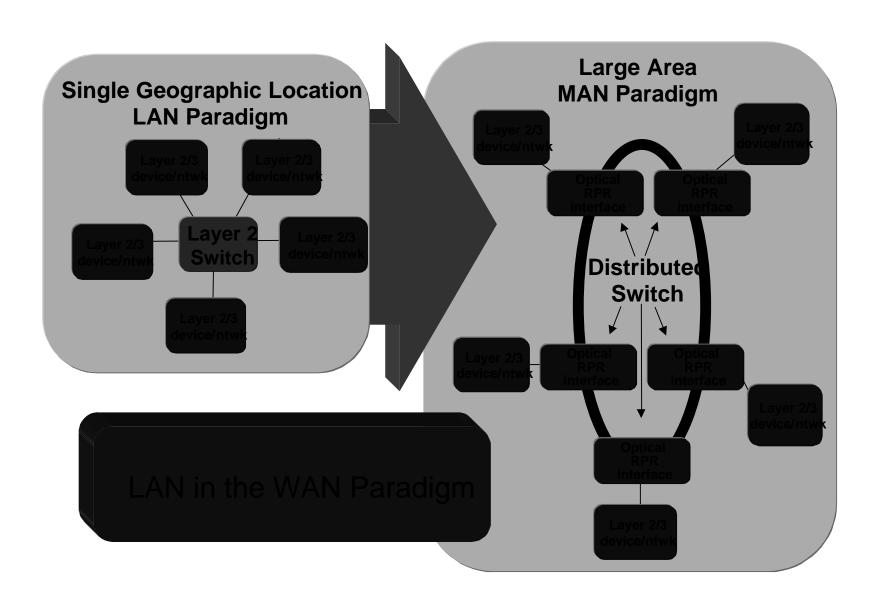
- What we see in the standard?
- What shouldn't be precluded?
- The Cut-Through concept and Architecture
- Distributed Switch Architecture
- Overall Control Plane
- Why Traffic engineering is requires links between traffic management and protection switching?

The Common Strategy

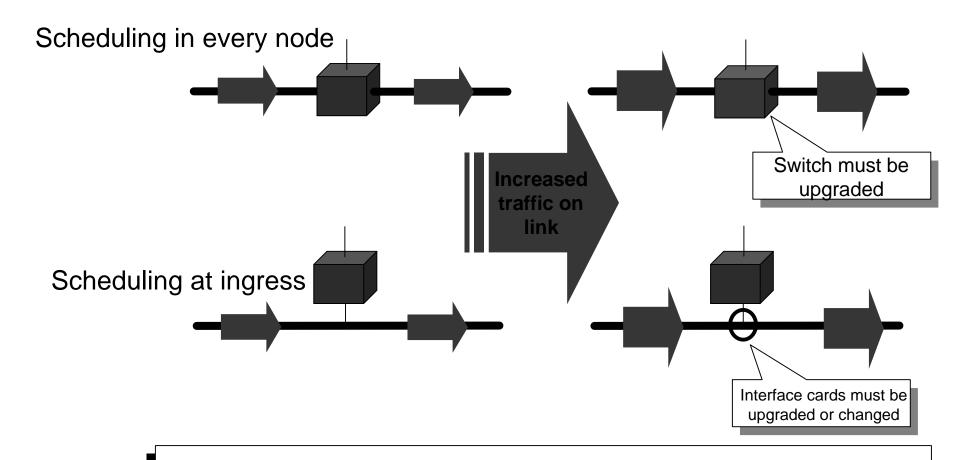


Keep Ring Operations PHY layer independent

RPR Conceptual View

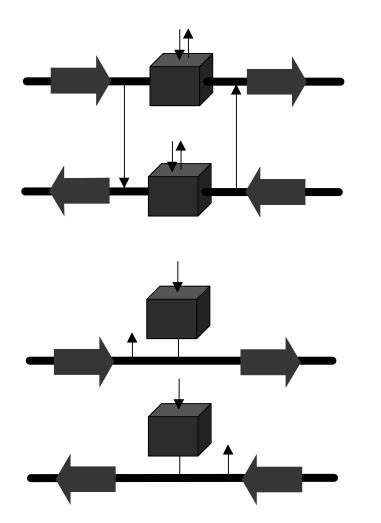


Scalability with Cut-through Switching in the Ring



A packet add-drop switch with ingress scheduling is the simplest RPR device able to scale in speed and size.

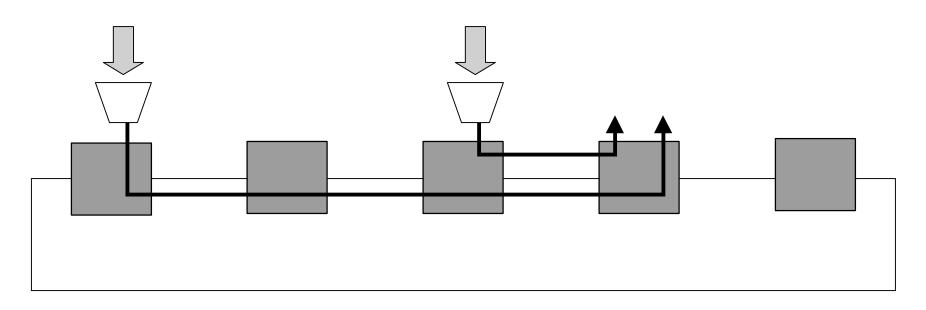
Functional Architecture



What we don't want:
A single MAC system

What is needed: A paralel MAC system

RPR: DISTRIBUTED traffic management

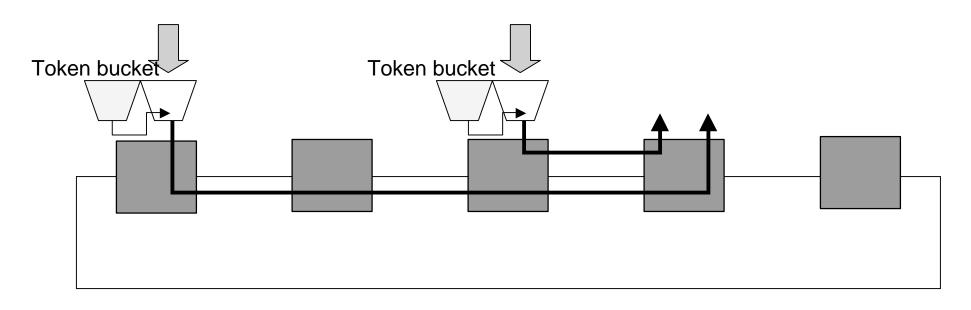


- Scheduling happens only at ingress
- No rescheduling at intermediate nodes!
- Asymmetric model: switched ports, ring ports
- Simple and inexpensive buffer management

Traffic Management Requirements

- Traffic Schemes are generally asymmetrical
- Sum of node traffic flows gives forecasted aggregated needs
- Possible bandwidth allocation between nodes done dynamically on the fly
- Support for IP DiffServe
- Simple provisioning
- Distributed Scheme that can scale

RPR: DISTRIBUTED token bucket



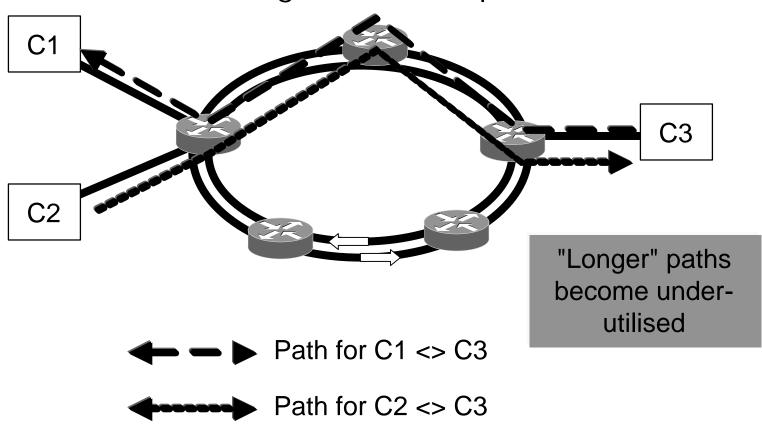
- Scheduling is achieved by a token bucket
- You can send when there is no traffic on the ring port and you are given a token
- A token correspond to a given bandwidth: 1Mbps
- Token Distribution is part of the control plane

RPR response to congestion

- Congestion in a node on the ring can be remedied immediately by reallocating spare bandwidth with the control plane though token distribution
- If a node is underdimensioned regarding its throughput
 - Another node can be introduced next to it on the ring
 - The backplane can be upgraded to double its processing capacity
- If the concerned ring is close to its maximum throughput
 - The ring can be segmented into two rings with maximum throughput in each
 - The link speed can be upgraded on-line by replacing the interface cards, thus increasing the ring throughput with a proportional factor

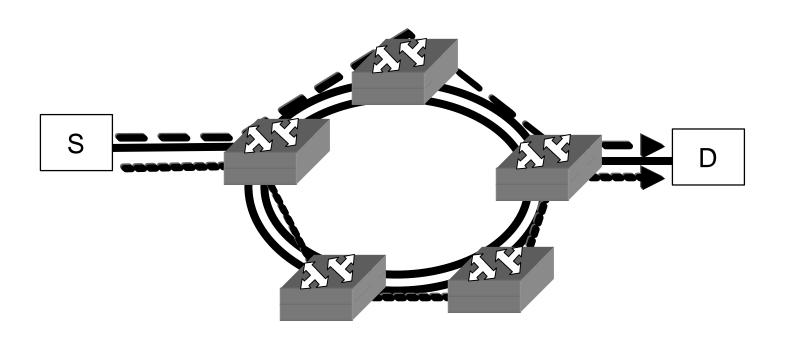
Today Bandwidth Bottlenecks

Today routing protocols create a single "shortest path"



RPR Ring Awareness

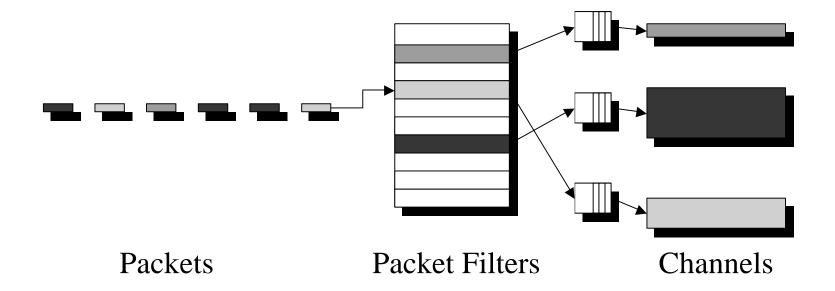
Source device determines the type of path on the basis of the service



— → Low delay (preferred for VoIP traffic)

Low bandwidth (preferred for HTTP)

Packets, Filters and Channels



- Packet filters are mapped to resource reservations (channel specifications)
- Channels created on demand (first packet arrives)

RPR Channels

- Three variations: Unicast, Multicast, Broadcast
- Created by a control message
- Control channel is bi-directional
- Switch nodes create channels on behalf of the sender
- A multi-hop channel must get OK from every switch along its path
- Portion of the bandwidth
- Synchronous between the sender and the receiver
- Virtually no loss of data

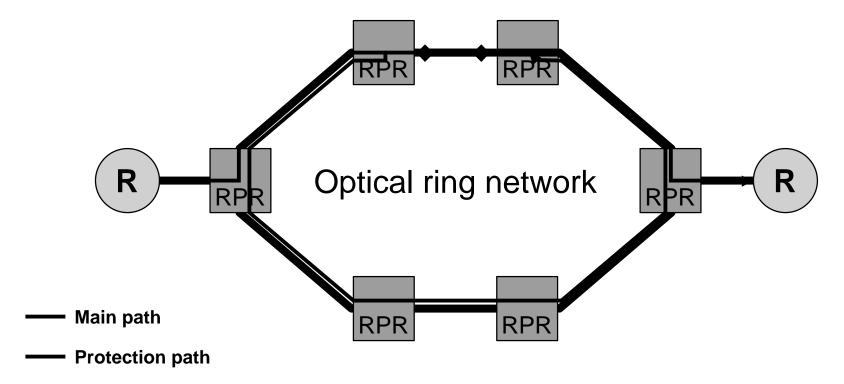
Control Plane v Data Plane

The data plane actually carries the information while the control plane sets up pathways through the data plane

RPR MAC and RPR BW both solve performance scalability problem by decoupling control and data planes

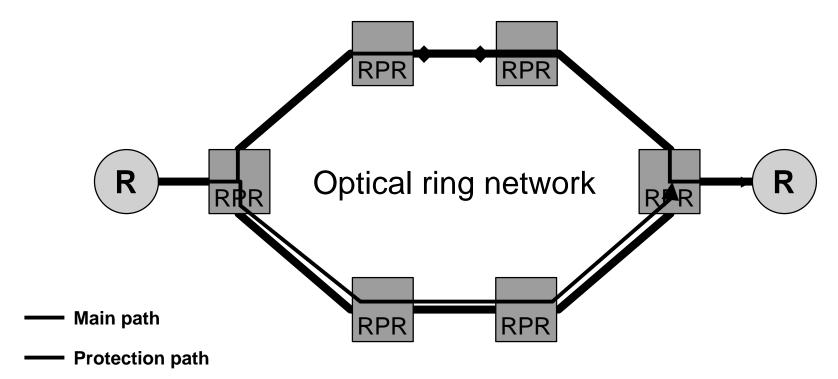
Wrapping Ring Protection

- Fast protection
- No control on traffic engineering
- Spans are used twice with limited SRP awareness



Steering Ring Protection

- Fast protection if pre-computed routes
- Full control on traffic engineering
- No need for dual control plane



Some Conclusions

- Cut Through Scheme is having some requirements on the MAC:
 - A control plane is needed to enable schedulers to work as one distributed switch
 - Control plane is also in charge of the steering of protection switching with communication of traffic management information
- It is a system approach on a ring and rings can be composed of different systems