CEI-P FEC and 802.3ap

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Agenda

- Overview of OIF CEI-P and its Firecode FEC
 - Frame Format overview
 - FEC performance
 - Implementation cost
 - Gates
 - Latency/Sync Time
- Relevance to 802.3ap
 - Fit with 802.3 layers
 - Potential Benefits of FEC
 - Options

OIF CEI-P Frame Format [1]

T ₀	64 Bit Payload Word O	T ₁	64 Bit Payload Word 1	T ₃	64 Bit Payload Word 2	S ₀
T ₃	64 Bit Payload Word 3	T_4	64 Bit Payload Word 4	T_5	64 Bit Payload Word 5	\square
Т ₆	64 Bit Payload Word 6	Т ₇	64 Bit Payload Word 7	T ₈	64 Bit Payload Word 8	S ₁
Т ₉	64 Bit Payload Word 9	T ₁₀	64 Bit Payload Word 10	T ₁₁	64 Bit Payload Word 11	
T ₁₂	64 Bit Payload Word 12	T ₁₃	64 Bit Payload Word 13	T ₁₄	64 Bit Payload Word 14	S ₂
T ₁₅	64 Bit Payload Word 15	T ₁₆	64 Bit Payload Word 16	T ₁₆	64 Bit Payload Word 17	
T ₁₈	64 Bit Payload Word 18	T ₁₉	64 Bit Payload Word 19	T ₂₀	64 Bit Payload Word 20	S ₃
T ₂₁	64 Bit Payload Word 21	T ₂₂	64 Bit Payload Word 22	T ₂₃	64 Bit Payload Word 23	
20	Overhead Bits : Oh[19:0]					•

Frame length = 1584 bits

- T_n= Transcode bit for 64bit payload word n Allows 65b64 or 10GBASE-R transcoding
- S_{0-3} = Supervisory channel Allows a telecom-style serial management channel
- Oh[19:0] = Firecode[19:0] xor State[2:0] Provides optional FEC
- All of packet is scrambled with a free-running $(X^{17} + X^{14} + 1)$ Scrambler
- Firecode generator polynomial $g(x) = (X^{13} + 1) (X^7 + X + 1)$
- Firecode is calculated over scrambled packet, then is itself scrambled
- CEI-P frame has exactly the same overhead as 64b66 encoding

CEI-P Firecode performance calculation

- The CEI-P firecode provides correction for a single 1 to 7bit burst error in each 1584 bit frame [2].
 - A Firecode was specifically chosen by CEI-P to address DFE error multiplication
- The corrected BER can be found by considering the probability of frames with 2 or more burst errors
 - The probability of n errors in a 1584 bit frame = (BERⁿ)*1584!/((1584-n)! * n!)
 - ... Ignoring n>2, probability of an uncorrectable frame = (BER²) (1584*1583)/2
- If a frame contains 2 or more separate bursts, then the errors are uncorrectable, and there is a 1 in 10 probability that 1-7 valid bits may be modified in error.
 - The unwanted correction is a 7 bit mask but only bits that are a 1 are altered, so on average 4 additional bits are corrupted per frame
 - Therefore, on average an uncorrectable frame contains 2+4/10 = 2.4 errors
- The corrected BER = (probability of uncorrectable frame)*(# of errors in frame)/1584
 - = (probability of Uncorrectable frame)*2.4/1584
 - $= (2.4/1584)^*(BER^2)^*(1584^{1583})/2$
 - $= 2.4*1583/2*(BER^2)$
- Corrected_BER ~= 1900*(BER²)

CEI-P Firecode FEC BER performance

- BER Corrected_BER 10-6 0.19*10-8 10-7 0.19*10-10 10-8 0.19*10-12 10-9 0.19*10-14 10-10 0.19*10-16 10-11 0.19*10-18 10-12 0.19*10-20 0.19*10-22 10-13
- For a thorough BER analysis that includes DFE/Firecode interactions, refer to Jim Hamstra's OIF contributions on the subject [3] & [4].
 - Results are at least an order of magnitude better than above.

CEI-P Firecode FEC costs

- Firecode protected frames can be corrected using simple error trapping [2].
 - Parallel implementations are practical and efficient at 10G data rates.
- Area for a 33bit wide datapath Transmit framer/coder : 3K gates @312.5Mhz Rx Sync : 4K gates @ 312.5Mhz Rx Error correction : 6K gates + 48x33 dual port RAM @ 312.5Mhz Rx deframer : 1K gates

Total Rx/Tx = 14K gates + 48x33 dual port RAM @ 312.5Mhz

- Latency 1584+ bits latency
- Sync time
 - ~500us worst case (without parallel sync engines)
 - Two 1584 frames to parse at all 1584 possible frame starts (= 2*1584² * 1E⁻¹⁰)
 - Presumes no errors

Relevance to 802.3ap

- The CEI-P FEC could be applied to 802.3ap
 - Re-use/reference CEI-P protocol
 - Create our own using similar techniques



Fit with 802.3 layering

- CEI-P or a similar frame could form a FEC sub-layer below the 10GBASE-R PCS (clause 49) in a similar manner to the WAN Interface Sublayer (Clause 50)
 - 10GBASE-R Sync bits would be collapsed into a single T bit
 - Provides full 10G data-rate, So no data-rate throttling would be needed, unlike WIS
 - May be possible to spoof existing BER monitor by corrupting Sync bits
- 10GBASE-R Latency : clause 49.2.15 states
 - "The sum of transmit delay contributed by a 10GBASE-R PCS shall be no more than 3584 BT."
 - This compares well with the frame delay required for error correction of 1584 BT.

TEXAS INSTRUMENTS

- It should be practical to meet this requirement with a combined 10GBASE-R PCS & FEC sub-layer.

Benefits to 802.3ap

- Ethernet BER objective of 10⁻¹² could be achieved with a 10⁻⁸ channel
- Channels meeting the Ethernet BER objective of 10⁻¹² could provide an effective BER of 10⁻²⁰
- Is this a way to reconcile the contradictory channel expectations within the TF ?
 - Ethernet system vendors with legacy backplanes happy with Ethernet BERs
 - Datacomm backplane users with 10⁻¹⁸ BER expectations
- CEI-P supervisory channel (S-bits) could be used as a back channel for adaptive Tx equalization
 - OIF have already discussed a scheme for this

Options

- Select FEC support by AN
 - Advertise sublayer support as an AN option
- Engage/disengage FEC error correction based on channel BER
 - Allows good channels to remove correction latency

References

 [1] Optical Internetworking Forum, oif2004.229.06, Common Electrical I/O -Protocol (CEI-P) Implementation Agreement, January 2005
[2] S. Lin and D. Costello, Error Control Coding : Fundamentals and Applications, Prentice Hall, Englewood Cliffs, New Jersey 1983
[3] Jim Hamstra, Effect of DFE Error propation on FEC, oif2003.267.02
[4] Jim Hamstra, Firecode Performance, oif2003.383.00



Summary

- A lightweight Firecode FEC can be implemented at low cost for 10G rates
- For BER's above 10⁻⁸ a Firecode FEC can provide substantial gains in payload BER
- A CEI-P like frame could be used as a WIS-like sublayer below 10GBASE-R



Backup foils

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Firecode Entropy Calculations

- The entropy of the error location is approx 10.63 bits + 6 more bits to extend burst to 7 bits = 16.63 bits - 2 -(20 - 16.63) = 10.34-1 which is probability of mis-correction in presence of arbitrary random error pattern.
 - log2 of 1584 = 10.63 bits to locate first errored bit, plus 6 additional error mask bits = 16.63 bits = total correction entropy
 - Total length of Fire Code polynomial = 20 bits 16.63 bits entropy = 3.37 bits = guard band against mis-correction
 - This is the combination of low probability of pointing beyond end of block for first error plus higher probability of more than 7 bit span in correction mask - both of the above indicate uncorrectable errors