119.2.4.4 Alignment marker mapping and insertion

In order to support deskew and reordering of the 16 individual PCS lanes at the receive PCS, alignment markers are added periodically for each PCS lane. The alignment marker for each PCS lane is <u>composed of a unique 120 fixed 96</u>-bit block <u>interleaved with 24 pad bits to achieve alignment marker field positioning identical to that defined in 91.5.2.6</u>. <u>An</u> <u>alignment marker group is composed of the alignment markers for all PCS lanes plus an additional 136-bit pad to yield the</u> <u>equivalent of eight 257-bit blocks</u>. The alignment markers for all PCS lanes are inserted as a group, group is aligned to the beginning of <u>a two FEC</u> <u>messagesblock</u>, and interrupt any data transfer that is already in progress. <u>A 136-bit pad is</u>appended to the alignment markers to yield the equivalent of eight 257-bit blocks</u>. The pad bits shall be set to a free running PRBS9 pattern, defined by the polynomial $x^9 + x^5 + 1$. The initial value of the PRBS9 pattern generator may be any pattern other than all zeros. The pad <u>contents are ignored</u>shall not be checked on receive.

Room for the alignment markers group is created by the transmit PCS (see 119.2.4.1). Special properties of the alignment markers are group is that it is they are not scrambled, does not conform to the encoding rules as outlined in Figure 82-5 and is are not transcoded. This is possible because the alignment markers are group is added after encoding, transcoding, and scrambling, and removed before descrambling, transcoding, and 64B/66B decoding. The alignment markers are group is not scrambled, which allows in order to allow the receiver to directly search for and find the individual alignment markers, deskew the PCS lanes, and reassemble the aggregate stream before descrambling is performed. Additionally, Tthe fixed 96bit portion of the alignment markers is themselves are formed from a known pattern that is defined to be balanced and with many transitions and therefore scrambling is not necessary. The group of alignment markers shall be inserted so they appear every 163 840 257-bit blocks. The variable tx_scrambled_am is created by inserting the group of alignment markers in the variable tx_scrambled. Alignment marker mapping and repetition rate are shown in Figure 119–5 and Figure 119–6.

The format of the each PCS lane's alignment markers is shown in Figure 119–4. There is a portion that is common across all alignment markers (designated as CM_0 to CM_{75}), and then a unique portion per PCS lane (designated as UM_0 to UM_{65}). Common synchronization logic independent of the received PCS lane number can be used with the common p<u>ortionart</u> of the alignment marker.

The content of the *fixed 96-bit portion of the* alignment markers shall be as shown in Table 119–1. The contents depend on the PCS lane number and the octet number, with the first-64_48 bits being identical across all alignment markers to allow for common synchronization across lanes. The format shown in Table 119–1 *defines is*-how the alignment markers appear on the PCS lanes at the PMA service interface. In the FEC codewords, they appear in a permuted format due to the codeword interleaving that occurs before FEC codewords are distributed to PCS lanes.

The alignment marker mapping function *creates* a *set* of 16 alignment markers, *and in combination with an additional 136bit PRBS9 pad generates an alignment marker group.* Let $am_x < 119:0 > be the alignment marker for PCS lane x, x=0 to 15, where bit 0 is the first bit transmitted. The alignment markers shall be mapped to am_mapped<1919:0> in a manner that yields the same result as the following process.$

For x=0 to 15, am_x<119:0> is constructed as follows.

a) am_x < 23:0 > is set to CM₀, CM₁, and CM₂, as shown in Figure 119-4 (bits 23:0) using the values in Table 119-1 for PCS lane number x.

<u>b) *if even(x)*</u>

_____x<31:24>={PRBS9<2*x+99:2*x+98>, PRBS9<6*x+5:6*x>}

else

<u>am_x<31:24>={PRBS9<2*x+95:2*x+94>, PRBS9<6*x+5:6*x>}</u> As shown in Figure 119-4 (bits 31:24) is an 8-bit pad value of PRBS9 pattern bits, where bit 6*x is the first PRBS9 bit output of the 8-bit pad.

c) am_x<55:32> is set to CM₃, CM₄, and CM₅, as shown in Figure 119-4 (bits 55:32) using the values in Table 119-1 for

PCS lane number x.

d) <u>if even(x)</u>

_____am_x<63:56>={PRBS9<4*x+195:4*x+192>, PRBS9<4*x+135:4*x+132>}

<u>else</u>

_____am_x<63:56>={PRBS9<4*x+195:4*x+192>, PRBS9<4*x+127:4*x+124>}

As shown in Figure 119-4 (bits 63:56) is an 8-bit pad value of PRBS9 pattern bits, where bit 4*x+128 is the first PRBS9 bit output of the 8-bit pad.

e) am_x<87:64> is set to UM_0 , UM_1 , and UM_2 , as shown in Figure 119-4 (bits 87:64) using the values in Table 119-1 for PCS lane number x.

f) if even(x)

______am__x<95:88>={PRBS9<6*x+299:6*x+294>, PRBS9<2*x+257:2*x+256>} ______else

 $\underline{\text{am}_x<95:88>=\{PRBS9<6^*x+287:6^*x+282>, PRBS9<2^*x+257:2^*x+256>\}}$ As shown in Figure 119-4 (bits 95:88) is an 8-bit pad value of PRBS9 pattern bits, where bit 2*x+256 is the first PRBS9 bit output of the 8-bit pad.

g) am_x<119:96> is set to UM_3 , UM_4 , and UM_5 , as shown in Figure 119-4 (bits 119:96) using the values in Table 119-1 for PCS lane number x.

As an example, the: is sent as (left most bit sent first) lane marker for 400GBASE-R lane number 0 variable am_0 is sent as (left most bit sent first):

10000011 00010110 10000100 00101111 0111100 01111001 01111011 11010000 TBD 01011001 01010010 01100100 <PRBS9(0:5), PRBS9(98:99)> 10100110 10101101 10011011 <PRBS9(132:135), PRBS9(192:196)> 01111001 11010111 11100100 <PRBS9(256:257), PRBS9(294:299)> 10000110 00101000 00011011

The variable am_mapped is then derived from 10-bit interleaving the group of 16 alignment markers am_x per the following procedure.

For all k=0 to 11 For all j=0 to 7

<u>if even(k)</u>

else

 $\frac{\text{am}_{mapped} < 160^{*} \text{k} + 20^{*} \text{j} + 9:160^{*} \text{k} + 20^{*} \text{j} > = \text{am}_{2^{*} \text{j}} < 10^{*} \text{k} + 9:10^{*} \text{k} > \text{am}_{mapped} < 160^{*} \text{k} + 20^{*} \text{j} + 19:160^{*} \text{k} + 20^{*} \text{j} + 10 > = \text{am}_{2^{*} \text{j}} + 1 < 10^{*} \text{k} + 9:10^{*} \text{k} > \text{am}_{2^{*} \text{j}} + 1 < 10^{*} \text{k} + 9:10^{*} \text{k} > \text{am}_{2^{*} \text{j}} + 1 < 10^{*} \text{k} + 9:10^{*} \text{k} > 10^{*} \text{k} = \text{am}_{2^{*} \text{j}} + 1 < 10^{*} \text{k} + 9:10^{*} \text{k} > 10^{*} \text{k} = 10^{*} \text{k} + 9:10^{*} \text{k} > 10^{*} \text{k} = 10^{*} \text{$

The additional 136-bit pad is appended to variable am_mapped as follows.

<u>am_mapped<2055:1920> = PRBS9<519:384></u>

In this expression, PRBS9<384> is the first PRBS9 bit output of the 136-bit pad.

The alignment marker *group am_mapped*<2055:0> shall be inserted so *it* appears every 163 840 257-bit blocks. The variable tx_scrambled_am<10279:0> is constructed in one of two ways. Let the set of vectors tx_scrambled_i<256:0> represent consecutive values of tx_scrambled<256:0>. For a *10280-bit* block with *an* alignment marker *group* inserted:

 $\frac{\text{tx_scrambled_am<2055:0>} = \text{am_mapped<2055:0>}}{\text{for all i=0 to 31}}$

 $\underline{\text{tx_scrambled_am < 257*i + 2312: 257*i + 2056 > = tx_scrambled_i < 256:0 > }}$

For a 10280-bit block without an alignment marker group:

for all i=0 to 39

 $tx_scrambled_am < 257*i + 256:257*i > = tx_scrambled_i < 256:0 >$

Alignment marker mapping and repetition rate are shown in Figure 119–5 and Figure 119–6.

{CM0, CM1, CM2} Pad {CM3, CM4, CM5} Pad {UM0, UM1, UM2} Pad {UM3, UM4, UM5}	Bit Position:	0 23	24 31	32 55	56 63	64 87	88 95	5 96 119
		{CM0, CM1, CM2}	Pad	{CM3, CM4, CM5}	Pad	{UM0, UM1, UM2}	Pad	{UM3, UM4, UM5}

Common Marker = {CM0, CM1, CM2, CM3, CM4, CM5}; Unique Marker = {UM0, UM1, UM2, UM3, UM4, UM5}

Figure 119-4 – Alignment marker format

<Replaces current Figure 119-4>

Table 119-1 – *Encodings for fixed 96-bit portion of* 400GBASE-R Alignment marker encodings

PCS	Encoding														
lane number	{CM0,CM1,CM2,CM3,CM4,CM5,UM0,UM1,UM2,UM3,UM4,UM5}														
0	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x9E,	0xEB,	0x27,	0x61,	0 x14,	0xD8			
1	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x50,	0x74,	0x88,	0xAF,	0x8B,	0x77			
2	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0xB4,	0x B7,	0xEA,	0x4B,	0x48,	0x1 5			
3	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0xE4,	0xFB,	0xF1,	0x1B,	0x04,	0x0E			
4	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0xDC,	0x58,	0xEE,	0x23,	0xA7,	0x11			
5	0x9A,	0x4A,	0x26,	<mark>0x65</mark> ,	0xB5,	0xD9,	0xBD,	0 xA9,	0xBF,	0x42,	0x56,	0x40			
6	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x97,	0x67,	0x77,	0x68 ,	0x98,	0x88			
7	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x24,	0x35,	0xA5,	0xDB,	0xCA,	0x5A			
8	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x57,	0x64,	0x51,	0xA8,	0x9B,	0xAE			
9	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x28,	0xF9,	0x3E,	0xD7,	0x06,	0xC1			
10	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0xCB,	0xD1,	0xAD,	0x34,	0x2E,	0x52			
11	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x5E,	0x1E,	0x38,	0xA1,	0xE1,	0xC7			
12	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x19,	0x98,	0x F9,	0xE6,	0x67,	0x06			
13	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x84,	0xEC,	0x20,	0x7B,	0x13,	0xDF			
14	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x13,	0 xA4,	0xED,	0xEC,	0x5B,	0x12			
15	0x9A,	0x4A,	0x26,	0x65,	0xB5,	0xD9,	0x3F,	0x8 A,	0xBE,	0xC0,	0x75,	0x41			

<Replaces current Table 119-1>

PCS			Ree	d-Sol	omon	symb	ol ind	lex, <i>k</i>	(10-b	it sym	bols)		
lane, i	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0					am_0	l					119	
1						am_1							
2						am_2							
3						am_3							
4	am_4												
5	am_5												
6	am_6												
7 8	am_7 am_8												
9						am 9							
10													
11	am_10 am_11												
12						am_12							
13					i	am_13	3						
14					i	am_14	4						
15					á	am_15	5						
		Fig	ure 1	19–5	—Ali	ignm	ent n	narke	er Ma	ppin	g to F	CS I	anes
am_0												1_0]
am_1												<u>1</u>	-
am_2											<u> </u>	1_2	-
am_3 am_4												1_3 1_4	-
am_4												1_4 1_5	-
am_0												1_0 1_6	-
am 7												1_0 1_7	-
am_8												1_8	1
am_9												 1_9	1
am_10												10	1
am_11	1										am	11	1
am_12	2										am	12	
	3											13	1

am_14

am_15

163 840 257-bit blocks between AM insertions

am_14

am_15

Figure 119-6—Alignment marker insertion period

119.2.4.5 Pre-FEC Distribution

Two Reed-Solomon FEC codewords are interleaved before data is distributed to the PCS lanes to improve error correctioncapability. Data is distributed to two 5140-bit message blocks (m_A-and m_B-are both arrays of 514 10-bit symbols) byperforming a 10-bit round robin distribution of the tx_scrambled_am<256:0> data as follows. In order to improve error_ correction capability, each set of two consecutive Reed-Solomon FEC codewords is interleaved before being distributed to_ form the PCS lanes. To enable this interleaving, the Pre-FEC Distribution function receives a 10280-bit block tx_scrambled_ am, and performs a 10-bit symbol round robin distribution to form two 514-symbol FEC messages, which are subsequently_ each encoded by the RS FEC. The following describes the 10-bit round robin distribution process.

For all j=0 to 39, tx_temp<10279:0> shall be constructed as follows:

_____tx_temp<(257j+256):(257j)> = tx_scrambled_am_j<256:0>

For all i=0 to 513, mA<513:0> and mB<513:0> shall be constructed as follows: $m_A < (513-i) > = \frac{tx_temp_tx_scrambled_am}{(20_*i+9):(20_*i)}$ $m_B < (513-i) > = \frac{tx_temp_tx_scrambled_am}{(20_*i+19):(20_*i+10)}$