

# Proposed Text for 2.5GBASE-X PCS/PMA

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To make editing instructions easier to understand for the purposes of assembling the initial Framemaker file the following conventions are used below:

**Table 125-1** – Highlight yellow means it cross references something in the 802.3cb draft specification.

**1.2.3** – Green text means it references a Clause in IEEE 802.3 that does not appear in the 802.3cb draft specification.

**Insert these 2 rows in table 1-23** – Red text means editing instructions for person assembling the initial Framemaker file and should not appear in the document.

Underlined blue text – New text inserted in existing clauses. Note that text in new clauses are not underlined.

~~Cross out purple text~~ – Text that should be deleted in existing clauses.

***Editor's Note: Insert following rows in Table 1-23*** – Bold italic font look like editing instructions for the person assembling the Framemaker file, but they are instructions for the IEEE editor and are actually part of the of the 802.3cb draft specification as bold italic font. Leave these as is in the text and do not delete or take any action on these. Only act on the editing instructions in **red**.

Ignore this page with table of contents. It is included here to make sure all the headings in this document are formatted properly as headings. The final assembly of all the clauses in Framemaker should build the table of content in that document

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## 200 Physical Coding Sublayer (PCS) and Physical Medium Attachment (PMA) sublayer for 8B/10B, type 2.5GBASE-X

### 200.1 Overview

#### 200.1.1 Scope

This clause specifies the Physical Coding Sublayer (PCS) and the Physical Medium Attachment (PMA) sublayer that are common to a family of 2.5 Gb/s Physical Layer implementations known as 2.5GBASE-X. The 2.5GBASE-X PCS and 2.5GBASE-X PMA are sublayers of the 2.5 Gb/s BASE-X PHY listed in [Table 125-1](#). The term 2.5GBASE-X is used when referring generally to Physical Layers using the PCS and PMA defined in this clause.

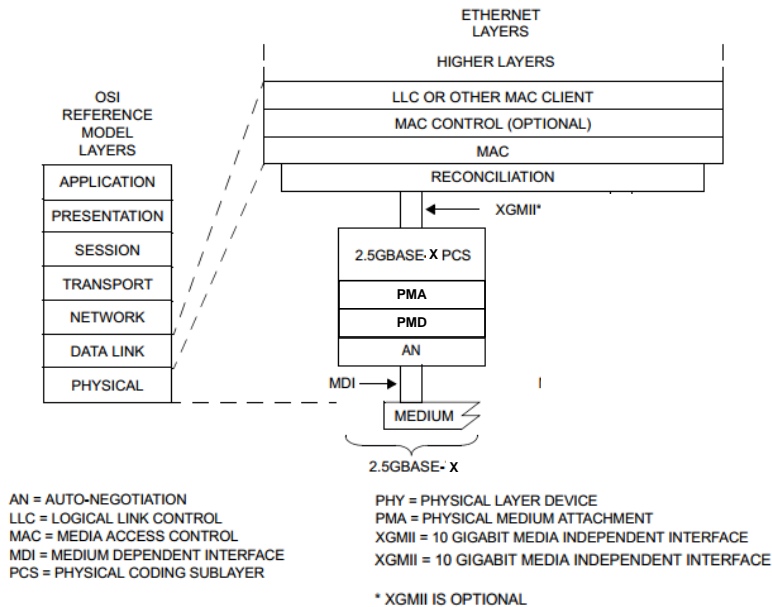
#### 200.1.2 Relationship of 2.5GBASE-X to other standards

[Figure 125-1](#) depicts the relationships among the 2.5GBASE-X sublayers, the Ethernet MAC and reconciliation layers, and the higher layers. The 2.5GBASE-X service interface is the XGMII, which is defined in [Clause 46](#).

#### 200.1.3 Summary of 2.5GBASE-X sublayers

[Figure 200-1](#) shows the relationship of the 2.5GBASE-X PCS sublayer with other sublayers to the ISO Open System Interconnection (OSI) reference model.

FM editor's note: Copy figure 125-1 from 802.3bz and draw diagram as sketched below:



**Figure 200-1** 2.5GBASE-X PCS and PMA relationship to the ISO/IEC Open Systems Interconnection (OSI) reference model and IEEE 802.3 Ethernet model

### **200.1.3.1 Physical Coding Sublayer (PCS)**

The PCS service interface is the XGMII, which is defined in Clause 46 running at 2.5Gb/s. The 2.5GBASE-X PCS provides all services required by the XGMII including Encoding (decoding) of XGMII data octets to (from) ten-bit code-groups (8B/10B) for communication with the underlying PMA.

### **200.1.3.2 Physical Medium Attachment (PMA) sublayer**

The PMA provides a medium-independent means for the PCS to support the use of serial-bit oriented physical media. The 2.5GBASE-X PMA performs the following functions:

- a) Mapping of transmit and receive code-groups between the PCS and PMA via the PMA Service Interface
- b) Serialization (deserialization) of code-groups for transmission (reception) on the underlying serial PMD
- c) Recovery of clock from the 8B/10B-coded data supplied by the PMD
- d) Mapping of transmit and receive bits between the PMA and PMD via the PMD Service Interface

### **200.1.4 Inter-sublayer interfaces**

The upper interface of the PCS may connect to the Reconciliation Sublayer through the XGMII. The lower interface of the PCS connects to the PMA sublayer to support a PMD. The 2.5GBASE-X PCS has a nominal rate at the PMA service interface of 3.125 Gb/s, which provides capacity for the MAC data rate of 2.5 Gb/s.

It is important to note that, while this specification defines interfaces in terms of bits, octets, and frames, implementations may choose other data-path widths for implementation convenience.

200.1.5 Functional Block Diagram

Figure 200-2 provides a functional block diagram of the 2.5GBASE-X PHY.

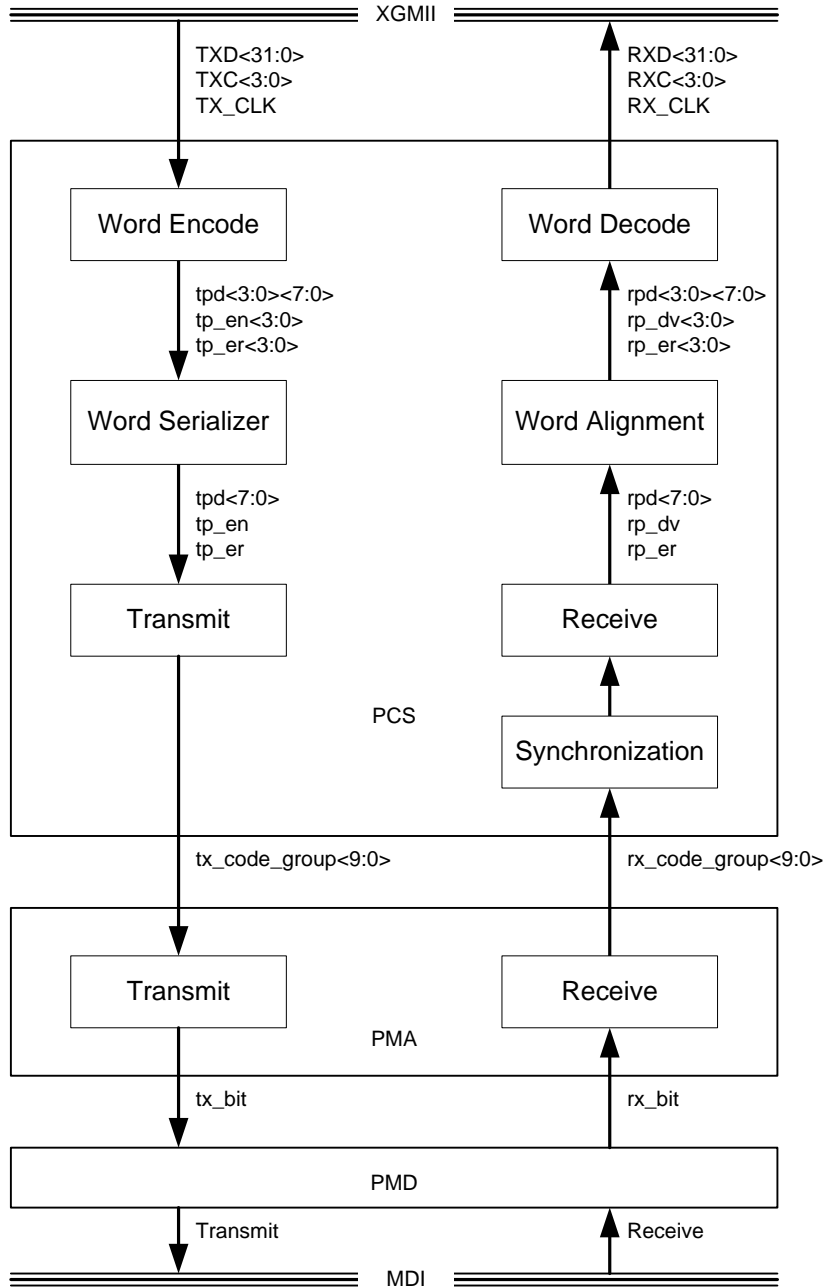


Figure 200-2 Functional block diagram

## 200.2 Physical Coding Sublayer (PCS)

### 200.2.1 PCS service interface (XGMII)

The PCS service interface allows the 2.5GBASE-X PCS to transfer information to and from a PCS client. A PCS client is generally the Reconciliation Sublayer. The PCS Interface is precisely defined as the 10 Gigabit Media Independent Interface (XGMII) scaled to run at 2.5Gb/s in Clause 46.

### 200.2.2 Functions within the PCS

The PCS includes the Word Encode, Word Serializer, Transmit, Synchronization, Receive, Word Alignment, and Word Decode processes for 2.5GBASE-X. The PCS shields the RS (and MAC) from the specific nature of the underlying channel.

When communicating with the XGMII, the PCS uses, in each direction, 32 data signals (TXD <31:0> and RXD <31:0>), four control signals (TXC <3:0> and RXC <3:0>), and a clock (TX\_CLK and RX\_CLK).

When communicating with the PMA, the PCS uses a ten-bit wide, synchronous data path, which conveys ten-bit code-groups. At the PMA Service Interface, code-group alignment and MAC packet delimiting are made possible by embedding special non-data code-groups in the transmitted code-group stream. The PCS provides the functions necessary to map packets between the XGMII format and the PMA service interface format.

The Word Encode process continuously generates four 2.5GPll symbols based upon the TXD <31:0> and TXC <3:0> signals on the XGMII, sending them to the Word Serializer process.

The Word Serializer process takes the four 2.5GPll symbols and outputs them one 2.5GPll symbol at a time to the PCS Transmit Process.

The PCS Transmit process continuously generates code-groups based upon the tpd<7:0>, tp\_en, and tp\_er signals on the 2.5GPll, sending them immediately to the PMA Service Interface via the PMA\_UNITDATA.request primitive.

The PCS Synchronization process continuously accepts code-groups via the PMA\_UNITDATA.indication primitive and conveys received code-groups to the PCS Receive process via the SYNC\_UNITDATA.indicate primitive. The PCS Synchronization process sets the sync\_status flag to indicate whether the PMA is functioning dependably.

The PCS Receive process continuously accepts code-groups via the SYNC\_UNITDATA.indicate primitive. The PCS Receive process monitors these code-groups and generates rpd<7:0>, rp\_dv, and rp\_er on the 2.5GPll.

The Word Alignment process queues the received 2.5GPll symbols and aligns them in groups of four 2.5GPll symbols. Symbols may be deleted or idle symbols added in order to do the alignment.

The Word Decode process continuously accepts the four 2.5GPll symbols from the Word Alignment process and generates RXD <31:0> and RXC <3:0> on the XGMII.

All PCS processes are described in detail in the state diagrams in 200.2.6.2.

### 200.2.3 Use of code-groups

The transmission code used by the PCS, referred to as 8B/10B, is identical to that specified in Clause 36. The PCS maps XGMII characters into an intermediate 2.5GPll which is then mapped to 10-bit code-groups, and vice versa, using the 8B/10B block coding scheme. Implicit in the definition of a code-group is an establishment of code-group boundaries by a PCS Synchronization process. The 8B/10B transmission code as well as the rules by which the PCS ENCODE and DECODE functions generate, manipulate, and interpret code-groups are specified in 200.2.5. The XGMII to 2.5GPll mapping and vice versa are specified in 200.2.4. Code-groups and the 2.5GPll are unobservable and have no meaning outside the PCS.

## 200.2.4 XGMII to 2.5GPll mapping

The Word Encode, Word Serializer, Word Alignment, and Word Decode processes together define the XGMII to 2.5GPll mapping. This mapping function formats and serializes/de-serializes the data between the four XGMII lanes and the single lane of the PCS transmit/receive process.

### 200.2.4.1 2.5Gb/s PCS Internal Interface (2.5GPll)

The 2.5Gb/s PCS Internal Interface (2.5GPll) is a logical interface that is internal to the 2.5GBASE-X PCS and exists solely for the purposes of defining the 2.5GBASE-X PCS functionality. Physical implementation of the 2.5GPll is optional and is not exposed outside of the PCS.

The 2.5GPll consists of the following variables: `tp_en`, `tp_er`, `tpd<7:0>`, `rp_dv`, `rp_er`, `rpd<7:0>` and its encoding is similar but not identical to the GMII in clause 35. The permissible encodings are shown in [Table 200-1](#) and [Table 200-2](#).

A 2.5GPll symbol is defined to be a set of `tp_en`, `tp_er`, `tpd<7:0>` variables or `rp_dv`, `rp_er`, `rpd<7:0>` variables.

The Word Encode and Word Decode processes maps between the four XGMII lanes to four 2.5GPll symbols. The four 2.5GPll symbols are indexed as `tp_en<3:0>`, `tp_er<3:0>`, `tpd<3:0><7:0>`, `rp_dv<3:0>`, `rp_er<3:0>`, `rpd<3:0><7:0>`. The nominal rate of operation is 12.8ns +/- 0.01%

The Word Serializer and Word Alignment processes serializes/de-serializes four 2.5GPll symbols to/from four consecutive single 2.5GPll symbols. The nominal rate of operation of the single 2.5GPll symbol is 3.2ns +/- 0.01%.

**Table 200-1** Permissible encodings of `tpd<7:0>`, `tp_en`, `tp_er`

<code>tp_en</code>	<code>tp_er</code>	<code>tpd&lt;7:0&gt;</code>	Description	Abbreviation
0	0	0x00 to 0xFF	Normal inter-frame	Idle
0	1	0x00	Reserved	
0	1	0x01	Assert LPI	LPI
0	1	0x02 to 0x9B	Reserved	
0	1	0x9C	Sequence	Seq
0	1	0x9D to 0xFF	Reserved	
1	0	00 to FF	Data	Data x
1	1	0x00 to 0xFF	Transmit Error	Err

**Table 200-2** Permissible encodings of `rpd<7:0>`, `rp_dv`, `rp_er`

<code>rp_dv</code>	<code>rp_er</code>	<code>rpd&lt;7:0&gt;</code>	Description	Abbreviation
0	0	0x00 to 0xFF	Normal inter-frame	idle
0	1	0x00	Reserved	
0	1	0x01	Assert LPI	LPI
0	1	0x02 to 0x0D	Reserved	
0	1	0x0E	False carrier indication	
0	1	0x0F	Carrier Extend (odd byte packets)	CE
0	1	0x10 to 0x9B	Reserved	
0	1	0x9C	Sequence	Seq
0	1	0x9D to 0xFF	Reserved	
1	0	00 to FF	Data	Data x
1	1	0x00 to 0xFF	Receive Error	Err



**200.2.4.2 Word Encode**

The Word Encode process maps the four XGMII lanes onto four 2.5GPPII symbols as shown in **Table 200-3**. The XGMII encoding is specified in **Table 46-3**. The 2.5GPPII encoding is specified in **Table 200-1**. The mapping of the sequence ordered set is dependent on the current state of the wencode\_state variable as shown in column 5. The state of wencode\_state is updated once the mapping occurs per the last column.

**Table 200-3 Word Encode Mapping**

XGMII						2.5GPPII				
Lane 0	Lane 1	Lane 2	Lane 3	wencode_state		2.5GPPII<0>	2.5GPPII<1>	2.5GPPII<2>	2.5GPPII<3>	wencode_state
Data A/Err	Data B/Err	Data C/Err	Data D/Err	X	➔	Data A/Err	Data B/Err	Data C/Err	Data D/Err	DATA
Idle	Idle	Idle	Idle	X	➔	Idle	Idle	Idle	Idle	IDLE
LPI	LPI	LPI	LPI	X	➔	LPI	LPI	LPI	LPI	DATA
SOP	Data A/Err	Data B/Err	Data C/Err	X	➔	0x55 Data	Data A/Err	Data B/Err	Data C/Err	DATA
Terminate	Idle	Idle	Idle	X	➔	Idle	Idle	Idle	Idle	IDLE
Data A/Err	Terminate	Idle	Idle	X	➔	Data A/Err	Idle	Idle	Idle	DATA
Data A/Err	Data B/Err	Terminate	Idle	X	➔	Data A/Err	Data B/Err	Idle	Idle	DATA
Data A/Err	Data B/Err	Data C/Err	Terminate	X	➔	Data A/Err	Data B/Err	Data C/Err	Idle	DATA
Sequence	Data X	Data Y	Data Z	IDLE	➔	Seq	Data S0	Seq	Data S1	SEQ
Sequence	Data X	Data Y	Data Z	SEQ	➔	Seq	Prev Data S2	Seq	Prev Data S3	IDLE
Sequence	Data X	Data Y	Data Z	DATA	➔	Idle	Idle	Idle	Idle	IDLE
else					➔	Err	Err	Err	Err	DATA

A sequence ordered set |Q| appears to the PMA as /K28.5/W/K28.5/W/K28.5/W/K28.5/W/. On the 2.5GPPII this appears as Seq, Data S0, Seq Data S1, Seq, Data S2, Seq, Data S3. The same sequence ordered-set is assumed to be repeating over multiple XGMII cycles. Every other consecutive Sequence ordered-set on the XGMII is ignored. If a Sequence ordered-set occurs over odd number of cycles on the XGMII then the final one will be truncated as Seq, Data S0, Seq Data S1 and Seq, Data S2, Seq, Data S3 are not sent.

The 24 bit Data X, Data Y, and Data Z from the sequence ordered set is mapped to Data S0, Data S1, Data S2, Data S3 as shown in Equation (200-1).

**(200-1)**

$$\begin{aligned}
 S_0\langle 7 \rangle &= S_3\langle 7 \rangle = 0, S_1\langle 7 \rangle = S_2\langle 7 \rangle = 1 \\
 S_0\langle 5:0 \rangle &= \text{Data X}\langle 5:0 \rangle \\
 S_1\langle 5:0 \rangle &= \text{Data Y}\langle 3:0 \rangle, \text{Data X}\langle 7:6 \rangle \\
 S_2\langle 5:0 \rangle &= \text{Data Z}\langle 1:0 \rangle, \text{Data Y}\langle 7:4 \rangle \\
 S_3\langle 5:0 \rangle &= \text{Data Z}\langle 7:2 \rangle \\
 S_n\langle 6 \rangle &= S_n\langle 7 \rangle \text{ if } S_n\langle 2 \rangle = 0 \\
 S_n\langle 6 \rangle &= S_n\langle 5 \rangle \text{ if } S_n\langle 2 \rangle = 1
 \end{aligned}$$

The signal ordered-set |Fsig| uses the same equation except  $S_2\langle 7 \rangle$  is set to 0.

Since sequence ordered-set can be sent back to back it is necessary to determine the boundaries of the ordered-set.  $\{S_0\langle 7 \rangle, S_1\langle 7 \rangle, S_2\langle 7 \rangle, S_3\langle 7 \rangle\}$  can be used to determine the boundary.  $\{S_0\langle 7 \rangle, S_1\langle 7 \rangle\}$  will always be 01, while the 01 combination will never occur over  $\{S_1\langle 7 \rangle, S_2\langle 7 \rangle\}$ ,  $\{S_2\langle 7 \rangle, S_3\langle 7 \rangle\}$ , or  $\{S_3\langle 7 \rangle, S_0\langle 7 \rangle\}$ .

Only 128 combinations of  $S_n\langle 7:0 \rangle$  are possible. When encoded to their 10-bit equivalent these 128 code-groups are defined to be in the set of |W|.

### 200.2.4.3 Word Serializer

The Word Serializer process takes the four 2.5GPll symbols ( $tp\_en<3:0>$ ,  $tp\_er<3:0>$ ,  $tpd<3:0><7:0>$ ) from by the Word Encoder and presents one symbol at a time ( $tp\_en$ ,  $tp\_er$ ,  $tpd<7:0>$ ) to the PCS transmit process. Index 0 is presented first and index 3 is presented last.

The Word Serializer process shall be synchronized to the PCS transmit process such that the 2.5GPll index 0 and 2 symbols are presented to the PCS transmit process which will result in the corresponding ordered set to be output to the PMA when the variable  $tx\_even$  is TRUE and index 1 and 3 variables when  $tx\_even$  is FALSE.

### 200.2.4.4 Word Alignment

The Word Alignment process de-serializes a sequence of 2.5GPll symbols ( $rp\_dv$ ,  $rp\_er$ ,  $rp d<7:0>$ ) from the PCS receive process to four 2.5GPll symbols ( $rp\_dv<3:0>$ ,  $rp\_er<3:0>$ ,  $rp d<3:0><7:0>$ ). Index 0 is the earliest to arrive and index 3 is the latest.

The alignment process will insert 2.5GPll idle symbols, or delete 2.5GPll symbols from the sequence of 2.5GPll symbols received to achieve the following conditions.

- a) A transition of a 2.5GPll idle symbol to a data or error symbol shall place the data or error symbol on index 0.
- b) A transition of a 2.5GPll idle symbol to a LPI symbol shall place the LPI symbol on either index 0 or 2.
- c) A transition of a 2.5GPll LPI symbol to an idle symbol shall place the idle symbol on either index 0 or 2.
- d) The start of |Q| or |Fsig| set shall always occur on index 0. Clause 200.2.4.2 describes how the start of |Q| and |Fsig| can be determined.

The Word Alignment process maintains a Deficit Idle Count (DIC) that represent the cumulative count of 2.5GMll symbols added or deleted from the sequence of received 2.5GPll symbols. The DIC is incremented by one for each 2.5GPll symbol deleted and decremented by one for each 2.5GPll idle symbol added. The DIC shall be bounded to a minimum of 0 and a maximum of 3.

Note that in a properly behaved system 2.5GPll symbol deletion should only occur at most once at the beginning of link, and afterwards no further insertions or deletions are required. In order to interoperate with the application described in Annex 200B, additional symbol insertions and deletions may be required during normal operation.

The only symbol that may be inserted is a 2.5GPll idle symbol. However any 2.5GPll symbol may be deleted. Usually this will either be a 2.5GPll idle or LPI symbols, though in pathological error conditions (i.e. unterminated packet followed immediately with sequence ordered-set) some other symbol may be deleted.

**200.2.4.5 Word Decode**

The Word Decode process maps the four 2.5GPll symbols onto the four XGMll lanes as shown in **Table 200-4**. The XGMll encoding is specified in **Table 46-4**. The 2.5GPll encoding is specified in **Table 200-2**. The mapping is dependent on the current state of the wdecode\_state and next\_seq\_s2\_s3 variables as shown in columns 5 and 6. The state of wdecode\_state is updated once the mapping occurs per the last column.

The 24 bit Data X, Data Y, and Data Z from the sequence ordered is reconstructed from Data S0, Data S1, Data S2, Data S3 according to Equation **(200-2)**.

**(200-2)**

```

if (S0<7>, S1<7>, S2<7>, S3<7> = 0110) then output
  XGMll = Sequence, Data X, Data Y, Data Z where
    Data X<7:0> = S1<1:0>, S0<5:0>
    Data Y<7:0> = S2<3:0>, S1<5:2>
    Data Z<7:0> = S3<5:0>, S2<5:4>
else
  XGMll = Idle, Idle, Idle, Idle
    
```

**Table 200-4 Word Decode Mapping**

2.5GPll				XGMll							
2.5GPll<0>	2.5GPll<1>	2.5GPll<2>	2.5GPll<3>	wdecode_state	next_seq_s2s3		Lane 0	Lane 1	Lane 2	Lane 3	wdecode_state
Data A/Err	Data B/Err	Data C/Err	Data D/Err	!IDLE	X	➔	Data A/Err	Data B/Err	Data C/Err	Data D/Err	DATA
Data *	Data A/Err	Data B/Err	Data C/Err	IDLE	X	➔	SOP	Data A/Err	Data B/Err	Data C/Err	DATA
Idle	Idle	Idle	Idle	!DATA	X	➔	Idle	Idle	Idle	Idle	IDLE
Idle	Idle	Idle	Idle	DATA	X	➔	Terminate	Idle	Idle	Idle	IDLE
Data A/Err	Idle or CE	Idle	Idle	DATA	X	➔	Data A/Err	Terminate	Idle	Idle	IDLE
Data A/Err	Data B/Err	Idle	Idle	DATA	X	➔	Data A/Err	Data B/Err	Terminate	Idle	IDLE
Data A/Err	Data B/Err	Data C/Err	Idle or CE	DATA	X	➔	Data A/Err	Data B/Err	Data C/Err	Terminate	IDLE
LPI	LPI	LPI	LPI	X	X	➔	LPI	LPI	LPI	LPI	IDLE
Idle	Idle	LPI	LPI	X	X	➔	LPI	LPI	LPI	LPI	IDLE
LPI	LPI	Idle	Idle	X	X	➔	Idle	Idle	Idle	Idle	IDLE
Seq	Data S0	Seq	Data S1	X	TRUE	➔	Sequence	Data X	Data Y	Data Z	SEQ
Seq	Data S0	Seq	Data S1	X	FALSE	➔	Idle	Idle	Idle	Idle	IDLE
Seq	Data S2	Seq	Data S3	SEQ	X	➔	Sequence	Data X	Data Y	Data Z	IDLE
else						➔	Err	Err	Err	Err	ERR

**200.2.5 8B/10B transmission code**

The transmission code used by the PCS, referred to as 8B/10B, is identical to that specified in Clause 36. In addition to the requirements in this clause, a 2.5GBASE-X PCS shall also meet the 8B/10B transmission code requirements specified in 36.2.4.1 through 36.2.4.6, 36.2.4.8, and 36.2.4.9. The relationship of code-group bit positions to PMA and other PCS constructs is illustrated in Figure 200-3.

FM editor's note: Copy Figure 36-3 and modify as shown below

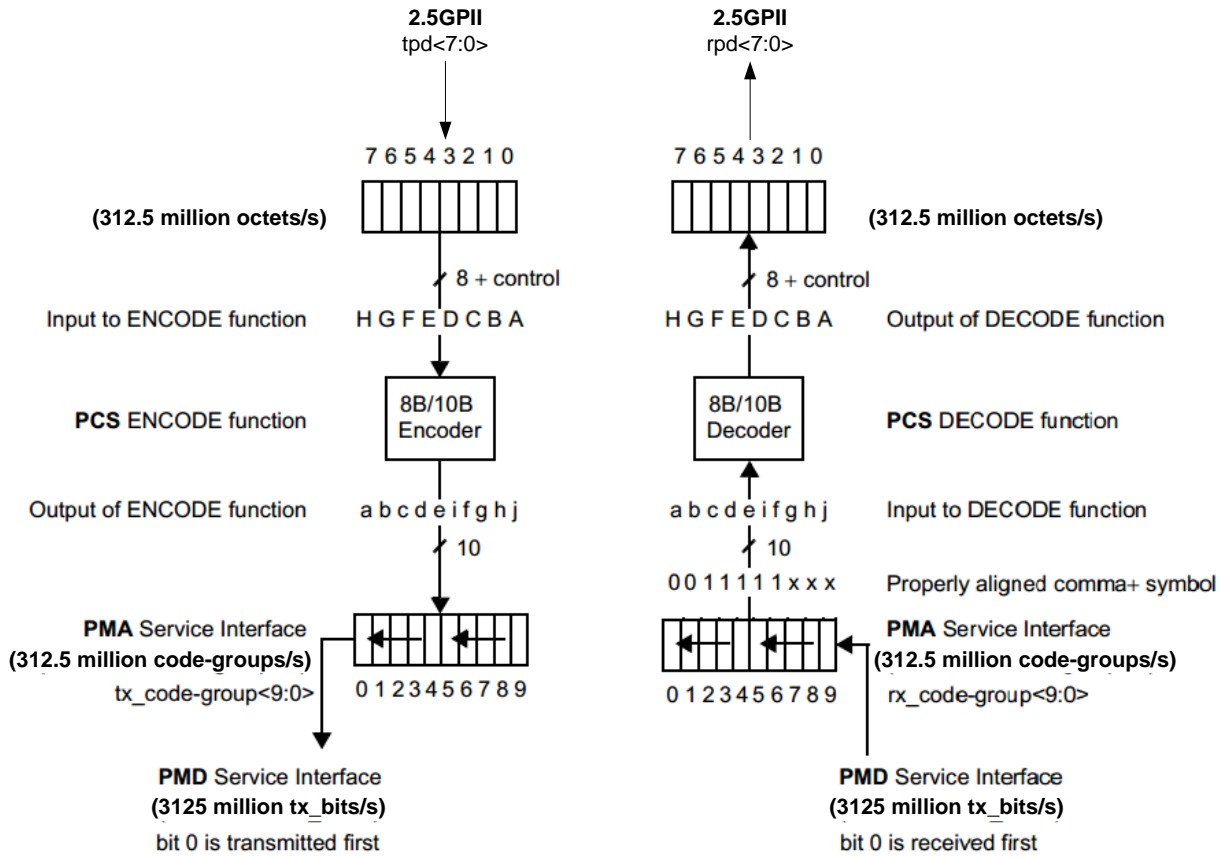


Figure 200-3 PCS 8B/10B reference diagram

**200.2.5.1 Notation conventions**

The 8B/10B transmission code uses letter notation for describing the bits of an unencoded information octet and a single control variable according to clause 36.2.4.1.

**200.2.5.2 Transmission order**

The code-group bit transmission order is illustrated in Figure 200-3 and defined in clause 36.2.4.2.

**200.2.5.3 Generating and checking Code-groups**

Valid code-groups are defined in clause 36.2.4.3.  
 The running disparity rules are defined in clause 36.2.4.4.  
 The code-group generation is defined in clause 36.2.4.5.  
 The check on the validity of received code-groups is defined in 36.2.4.6.

### 200.2.5.4 Ordered\_sets

Table 200-5 lists the defined ordered\_sets, consisting of a single special code-group or combinations of special and data code-groups. Ordered\_sets which include /K28.5/ provide the ability to obtain bit and code-group synchronization and establish ordered\_set alignment (see 36.2.4.9 and 200.3.2.4). Ordered\_sets provide for the delineation of a packet and synchronization between the transmitter and receiver circuits at opposite ends of a link. Certain PHYs include an option (see 78.3) to transmit or receive /LI/, /LI1/ and /LI2/ to support Energy-Efficient Ethernet (see Clause 78).

Ordered\_sets are specified according to the following rules:

- a) Ordered\_sets consist of either one, two, or four code-groups.
- b) The first code-group of all ordered\_sets is always a special code-group.
- c) The second code-group of all multi-code-group ordered\_sets is always a data code-group. The second code-group is used to distinguish the ordered set from all other ordered sets.

Table 200-5 Defined ordered\_sets

Code	Ordered_Set	Number of Code-Groups	Encoding
/I/	IDLE		Correcting /I1/, Preserving /I2/
/I1/	IDLE 1	2	/K28.5/D5.6/
/I2/	IDLE 2	2	/K28.5/D16.2/
	<b>Encapsulation</b>		
/R/	Carrier_Extend	1	/K23.7/
/S/	Start_of_Packet	1	/K27.7/
/T/	End_of_Packet	1	/K29.7/
/V/	Error_Propagation	1	/K30.7/
/LI/	<b>LPI</b>		Correcting /LI1/, Preserving /LI2/
/LI1/	LPI 1	2	/K28.5/D6.5/
/LI2/	LPI 2	2	/K28.5/D26.4/
	<b>Link Status</b>		
/Q/	Sequence ordered_set	8	/K28.5/W0/K28.5/W1/K28.5/W2/K28.5/W3/ (note a)
	<b>Reserved</b>		
/Fsig/	Signal ordered_set	8	/K28.5/W0/K28.5/W1/K28.5/W2/K28.5/W3/ (note a)

<sup>a</sup> /W0/, /W1/, /W2/, /W3/ are the 10-bit /Dx.y/ version of S0, S1, S2, S3 as defined per clause 200.2.4.2 and will never have a value of /D5.6/, /D16.2/, /D6.5/, /D26.4/.

### 200.2.5.5 Comma considerations

The comma considerations are described in clauses 36.2.4.8 and 36.2.4.9.

### 200.2.5.6 Sequence (/Q/)

A sequence ordered\_set is used to convey various link status such as local fault or remote fault. The 24 bit data of the sequence ordered\_set on the XGMII are mapped to S0, S1, S2, S3 per clause 200.2.4.2, and /W0/, /W1/, /W2/, /W3/ are the 8B/10B mapped version.

S0, S1, S2, S3 are constructed such that each /Wn/ can be mapped to only 128 out of 256 possible /Dx.y/ code-groups. /W/ is defined to be the set of 128 /Dx.y/ code-groups that can appear in a sequence ordered\_set. /W/ does not contain /D5.6/, /D16.2/, /D6.5/, /D26.4/.

It is possible for the transmitter to send out a truncated sequence ordered\_set that appears as /K28.5/W0/K28.5/W1/. When a truncated sequence ordered\_set is received, the Word Decode process will convert it to idles.

#### **200.2.5.7 Data (/D/)**

A data code-group, when not used to distinguish or convey information for a defined ordered set, conveys one octet of arbitrary data between the XGMII and the PCS. The sequence of data code-groups is arbitrary, where any data code-group can be followed by any other data code-group. Data code-groups are coded and decoded but not interpreted by the PCS. Successful decoding of the data code-groups depends on proper receipt of the Start\_of\_Packet delimiter, as defined in 200.2.5.10 and the checking of validity, as defined in 36.2.4.6.

#### **200.2.5.8 Idle (/I/)**

IDLE ordered sets (/I/) are transmitted continuously and repetitively whenever the XGMII is idle. /I/ provides a continuous fill pattern to establish and maintain clock synchronization. /I/ is emitted from, and interpreted by, the PCS. /I/ consists of one or more consecutively transmitted /I1/ or /I2/ ordered sets, as defined in Table 200-5.

The /I1/ ordered set is defined such that the running disparity at the end of the transmitted /I1/ is opposite that of the beginning running disparity. The /I2/ ordered set is defined such that the running disparity at the end of the transmitted /I2/ is the same as the beginning running disparity. The first /I/ following a packet or sequence ordered set restores the current positive or negative running disparity to a negative value. All subsequent /I/s are /I2/ to ensure negative ending running disparity.

Distinct carrier events are separated by /I/s.

A received ordered set that consists of two code-groups, the first of which is /K28.5/ and the second of which is a data code-group other than any of the 128 possible /M/, /D21.5/, or /D2.2/ (or /D6.5/ or /D26.4/ to support EEE capability), is treated as an /I/ ordered set.

#### **200.2.5.9 Low Power Idle (/LI/)**

LPI is transmitted in the same manner as IDLE. LPI ordered sets (/LI/) are transmitted continuously and repetitively whenever the XGMII is indicating LPI.

#### **200.2.5.10 Start\_of\_Packet (SPD) delimiter**

A Start\_of\_Packet delimiter (SPD) is used to delineate the starting boundary of a data transmission sequence.

Upon initiation of packet transmission, the PCS replaces the first octet of the MAC preamble with SPD. Upon initiation of packet reception, the PCS replaces the received SPD delimiter with the data octet value associated with the first preamble octet. A SPD delimiter consists of the code-group /S/ as defined in Table 200-5.

#### **200.2.5.11 End\_of\_Packet delimiter (EPD)**

An End\_of\_Packet delimiter (EPD) is used to delineate the ending boundary of a packet. The EPD is transmitted by the PCS following the last data octet comprising the FCS of the MAC packet. On reception, EPD is interpreted by the PCS as terminating a packet. An EPD always consists of the code-groups /T/R/. If /R/ is transmitted in an even-numbered code-group position, the PCS appends a single additional /R/ to the code-group stream to ensure that the subsequent /I/ is aligned on an even numbered code-group boundary. An /I/ always follows the conclusion of EPD. The /T/R/ or /T/R/R/ occupies part of the region considered by the MAC to be the IPG.

The code-group /T/ and /R/ are defined in Table 200-5.

### 200.2.5.12 Error\_Propagation (/V/)

Error\_Propagation (/V/) indicates that the PCS client wishes to indicate a transmission error to its peer entity. /V/ is emitted from the PCS when the XGMII indicates transmit error propagation, or when the XGMII presents a mapping that is undefined in [Table 200-3](#) to the Word Encode process. The code-group /V/ is defined in [Table 200-5](#).

The presence of Error\_Propagation or any invalid code-group on the medium denotes an error condition. Invalid code-groups are not intentionally transmitted onto the medium.

## 200.2.6 Detailed functions and state diagrams

This sub-clause starts with [Clause 36.2.5](#) and mark up the differences.

The body of this subclause is comprised of state diagrams, including the associated definitions of variables, constants, and functions. Should there be a discrepancy between a state diagram and descriptive text, the state diagram prevails. The notation used in the state diagrams follows the conventions of [21.5](#). State diagram timers follow the conventions of [14.2.3.2](#).

### 200.2.6.1 State variables

#### 200.2.6.1.1 Notation conventions

/x/

Denotes the constant code-group specified in [200.2.6.1.2](#) (valid code-groups must follow the rules of running disparity as per [200.2.5.3](#)).

[x/]

Denotes the latched received value of the constant code-group (/x/) specified in [200.2.6.1.2](#) and conveyed by the SYNC\_UNITDATA.indicate message described in [200.2.6.1.6](#).

#### 200.2.6.1.2 Constants

/COMMA/

The set of special code-groups which include a comma as specified in [36.2.4.9](#) and listed in [Table 36-2](#).

/D/

The set of 256 code-groups corresponding to valid data, as specified in [200.2.5.7](#).

/Dx.y/

One of the set of 256 code-groups corresponding to valid data, as specified in [200.2.5.7](#).

//

The IDLE ordered set group, comprising either the /I1/ or /I2/ ordered sets, as specified in [200.2.5.8](#).

/INVALID/

The set of invalid data or special code-groups, as specified in [36.2.4.6](#).

/Kx.y/

One of the set of 12 code-groups corresponding to valid special code-groups, as specified in [Table 36-2](#).

/Q/

Sequence ordered set as specified in [200.2.5.6](#). A properly formed sequence order set appears as /K28.5W/K28.5W/K28.5W/K28.5W/. A truncated sequence order set appears as /K28.5W/K28.5W/.

/R/

The code-group used as either: End\_of\_Packet delimiter part 2; End\_of\_Packet delimiter part 3.

/S/

The code-group corresponding to the Start\_of\_Packet delimiter (SPD) as specified in 200.2.5.10.

/T/

The code-group used for the End\_of\_Packet delimiter part 1.

/V/

The Error\_Propagation code-group, as specified in 200.2.5.12.

/W/

The set of 128 code-groups that is generated by ENCODE(s<7:0>) where for all 128 possible values of x<6:0> s<7> = x<6>, s<5:0> = x<5:0>, s<6> is set to x<5> when x<2> = 1 and s<6> is set to x<6> when x<2> = 0  
Note that /W/ is a subset of /D/

PL\_LIMIT

The number of 2.5GPll symbols to preload in the Word Alignment process after release from the RESET state. The number of symbols to preload is implementation dependent and should be minimized to reduce latency. The minimum number must be 3 to account for the deficit idle counting in clause 200.2.4.4.

The following constant is used only for the EEE capability:

/LI/

The LP\_IDLE ordered set group, comprising either the /LI1/ or /LI2/ ordered sets, as specified in 200.2.5.9.

### 200.2.6.1.3 Variables

assert\_seq

Alias used for sequence ordered set, consisting of the following terms:  
tp\_en=0 \* tp\_er=1 \* (tpd<7:0>=0x9C)

cgbad

Alias for the following terms: ((rx\_code-group∈/INVALID/) + (rx\_code-group=/COMMA/ \*rx\_even=TRUE)) \* PMA\_UNITDATA.indication

cggood

Alias for the following terms: !((rx\_code-group∈/INVALID/) + (rx\_code-group=/COMMA/ \*rx\_even=TRUE)) \* PMA\_UNITDATA.indication

EVEN

The latched state of the rx\_even variable, when rx\_even=TRUE, as conveyed by the SYNC\_UNITDATA.indicate message described in 200.2.6.1.6.

mr\_loopback

A Boolean that indicates the enabling and disabling of data being looped back through the PHY. Loopback of data through the PHY is enabled when Control register bit 3.0.14 is set to one.

Values: FALSE; Loopback through the PHY is disabled.  
TRUE; Loopback through the PHY is enabled.

mr\_main\_reset

Controls the resetting of the PCS via Control Register bit 3.0.15.

Values: FALSE; Do not reset the PCS.  
TRUE; Reset the PCS.



ODD

The latched state of the rx\_even variable, when rx\_even=FALSE, as conveyed by the SYNC\_UNITDATA.indicate message described in [200.2.6.1.6](#).

power\_on

Condition that is true until such time as the power supply for the device that contains the PCS has reached the operating region. The condition is also true when the device has low power mode set via Control register bit 3.0.11.

Values: FALSE; The device is completely powered (default).  
TRUE; The device has not been completely powered.

NOTE—Power\_on evaluates to its default value in each state where it is not explicitly set.

rpdc7:0>

Single lane 2.5GPII receive data from the PCS receive process.

rpdcx><7:0>

2.5GPII receive data from the Word Alignment process. x= 0, 1, 2, 3 for the four sets of 2.5GPII.

rp\_dv

Single lane 2.5GPII receive data valid from the PCS receive process.

Values: 0 or 1

rp\_dvx>

2.5GPII receive data valid from the Word Alignment process. x= 0, 1, 2, 3 for the four sets of 2.5GPII.

Values: 0 or 1

rp\_er

Single lane 2.5GPII receive error from the PCS receive process.

Values: 0 or 1

rp\_erx>

2.5GPII receive error from the Word Alignment process. x= 0, 1, 2, 3 for the four sets of 2.5GPII.

Values: 0 or 1

**Editor's Note: TBD should point to PMD section that describes PMD\_UNITDATA.indication**

rx\_bit

A binary parameter conveyed by the PMD\_UNITDATA.indication service primitive, as specified in [TBD](#), to the PMA.

Values: ZERO; Data bit is a logical zero.  
ONE; Data bit is a logical one

rx\_code-group<9:0>

A 10-bit vector represented by the most recently received code-group from the PMA. The element rx\_code-group<0> is the least recently received (oldest) rx\_bit; rx\_code-group<9> is the most recently received rx\_bit (newest). When code-group alignment has been achieved, this vector contains precisely one code-group.

rx\_even

A Boolean set by the PCS Synchronization process to designate received code-groups as either even- or odd-numbered code-groups as specified in [200.2.5.2](#).

Values: TRUE; Even-numbered code-group being received.  
FALSE; Odd-numbered code-group being received.

RXC<3:0>

The RXC<3:0> signal of the XGMII as specified in Clause 46. Set by the Word Decode process.

RXD<31:0>

The RXD<31:0> signal of the XGMII as specified in Clause 46. Set by the Word Decode process.

signal\_detect

A Boolean set by the PMD continuously via the PMD\_SIGNAL.indication(signal\_detect) message to indicate the status of the incoming link signal.

Values: FAIL; A signal is not present on the link.  
OK; A signal is present on the link.

sync\_status

A parameter set by the PCS Synchronization process to reflect the status of the link as viewed by the receiver. The values of the parameter are defined for code\_sync\_status. The equation for this parameter is  $\text{sync\_status} = \text{code\_sync\_status} + \text{rx\_lpi\_active}$

NOTE—If EEE is not supported, the variable rx\_lpi\_active is always false, and this variable is identical to code\_sync\_status controlled by the synchronization state diagram.

tpd<7:0>

Single lane 2.5GPII transmit data passed to the PCS transmit process.

tpd<x><7:0>

2.5GPII transmit data from the Word Encoder process. x= 0, 1, 2, 3 for the four sets of 2.5GPII.

tp\_en

Single lane 2.5GPII transmit enable passed to the PCS transmit process.

Values: 0 or 1

tp\_en<x>

2.5GPII transmit enable from the Word Encoder process. x= 0, 1, 2, 3 for the four sets of 2.5GPII.

Values: 0 or 1

tp\_er

Single lane 2.5GPII transmit error passed to the PCS transmit process.

Values: 0 or 1

tp\_er<x>

2.5GPII transmit error from the Word Encoder process. x= 0, 1, 2, 3 for the four sets of 2.5GPII.

Values: 0 or 1

**Editor's Note: TBD should point to PMD section that describes PMD\_UNITDATA.request**

tx\_bit

A binary parameter used to convey data from the PMA to the PMD via the PMD\_UNITDATA.request service primitive as specified in **TBD**.

Values: ZERO; Data bit is a logical zero.  
ONE; Data bit is a logical one.

tx\_code-group<9:0>

A vector of bits representing one code-group, as specified in Tables 36-1a to 36-1e or 36-2, which has been prepared for transmission by the PCS Transmit process. This vector is conveyed to the PMA as the parameter of a PMD\_UNITDATA.request(tx\_bit) service primitive. The element tx\_codegroup<0> is the first tx\_bit transmitted; tx\_code-group<9> is the last tx\_bit transmitted.

tx\_disparity

A Boolean set by the PCS Transmit process to indicate the running disparity at the end of code-group transmission as a binary value. Running disparity is described in 36.2.4.3.

Values: POSITIVE  
NEGATIVE

tx\_even

A Boolean set by the PCS Transmit process to designate transmitted code-groups as either even or odd-numbered code-groups as specified in 200.2.5.2.

Values: TRUE; Even-numbered code-group being transmitted.  
FALSE; Odd-numbered code-group being transmitted.

tx\_o\_set

One of the following defined ordered sets: /T/, /R/, /I/, /S/, /V/, /LI/, /K28.5/, or the code-group /D/.

TXC<3:0>

The TXC<3:0> signal of the XGMII as specified in Clause 46.

TXD<31:0>

The TXD<31:0> signal of the XGMII as specified in Clause 46.

wdecode\_state

Word Decoder State used by the Word Decode process to properly assemble the next XGMII value to output.

Values: DATA, IDLE, SEQ

wencode\_state

Word Encoder State used by the Word Encode process to properly assemble the Sequence ordered-set.

Values: DATA, ERR, IDLE, SEQ

The following variables are used only for the EEE capability:

assert\_lpidle

Alias used for the optional LPI function, consisting of the following terms:

$(tp\_en=0 * tp\_er=1 * (tpd<7:0>=0x01))$

code\_sync\_status

A parameter set by the PCS Synchronization process to reflect the status of the link as viewed by the receiver.

Values: FAIL; The receiver is not synchronized to code-group boundaries.  
OK; The receiver is synchronized to code-group boundaries.

idle\_d

Alias for the following terms:

$SUDI( ![D21.5/] * ![D2.2/] )$

that uses an alternate form to support the EEE capability:

$SUDI( ![D21.5/] * ![D2.2/] * ![D6.5/] * ![D26.4/] )$

rx\_lpi\_active

A Boolean variable that is set to TRUE when the receiver is in a low power state and set to FALSE when it is in an active state and capable of receiving data.

rx\_quiet

A Boolean variable set to TRUE while in the RX\_QUIET state and set to FALSE otherwise.

**Editor's Note: TBD should point to PMD section that describes transmit disable in EEE**

tx\_quiet

A Boolean variable set to TRUE when the transmitter is in the TX\_QUIET state and set to FALSE otherwise. When set to TRUE, the PMD will disable the transmitter as described in **TBD**.

#### 200.2.6.1.4 Functions

carrier\_detect

In the PCS Receive process, this function uses for input the latched code-group ([/x/]) and latched rx\_even (EVEN/ODD) parameters of the SYNC\_UNITDATA.indicate message from the PCS Synchronization process. When SYNC\_UNITDATA.indicate message indicates EVEN, the carrier\_detect function detects carrier when either:

- a) A two or more bit difference between [/x/] and both /K28.5/ encodings exists (see Table 36–2); or
- b) A two to nine bit difference between [/x/] and the expected /K28.5/ (based on current running disparity) exists.

Values: TRUE; Carrier is detected.  
FALSE; Carrier is not detected.

check\_end

Prescient End\_of\_Packet function used by the PCS Receive process. The check\_end function returns the current and next two code-groups in rx\_code-group<9:0>.

DECODE ([/x/])

In the PCS Receive process, this function takes as its argument the latched value of rx\_codegroup<9:0> ([/x/]) and the current running disparity, and returns the corresponding 2.5GPll rpd<7:0> per Table 36–1a–e. DECODE also updates the current running disparity per the running disparity rules outlined in 36.2.4.4.

ENCODE(x)

In the PCS Transmit process, this function takes as its argument (x), where x is a 2.5GPll tpd<7:0> octet, and the current running disparity, and returns the corresponding ten-bit code-group per Table 36–1a–e. ENCODE also updates the current running disparity per Table 36–1a–e.

NEXTSEQ()

Prescient function used by the Word Decode process that returns whether the next four 2.5GPll symbols presented to the Word Decode process is of the form Sequence, Data, Sequence, Data.

Values: TRUE; Next four 2.5GPll symbols are Sequence, Data, Sequence, Data  
FALSE; Next four 2.5GPll symbols are not Sequence, Data, Sequence, Data

signal\_detectCHANGE

In the PCS Synchronization process, this function monitors the signal\_detect variable for a state change. The function is set upon state change detection.

Values: TRUE; A signal\_detect variable state change has been detected.  
FALSE; A signal\_detect variable state change has not been detected (default).

NOTE—Signal\_detectCHANGE is set by this function definition; it is not set explicitly in the state diagrams. Signal\_detectCHANGE evaluates to its default value upon state entry.

SINSERT(x)

Add a single 2.5GPll symbol to the end of a queue that stores the 2.5GPll symbol presented by the receive process. The variable x is the 2.5GPll symbol (rp\_dv, rp\_er, rpd<7:0>). If x is a null set then all content in the queue is emptied. The depth of the queue is implementation dependent.

VOID(x)

x ∈ /D/, /T/, /R/, /K28.5/. Substitutes /V/ on a per code-group basis as requested by the 2.5GPll.

```
If [tp_en=0 * tp_er=1 * tpd<7:0>≠(0000 1111)],
    then return /V/;
Else if [tp_en=1 * tp_er=1],
    then return /V/;
Else return x.
```

WALIGN()

In the PCS Word Alignment process, this function performs the alignment according to clause 200.2.4.4. Four 2.5GPll symbols (rp\_dv<3:0>, rp\_er<3:0>, rpd<3:0><7:0>) are returned by this function. The SINSERT(x) function adds 2.5GPll symbols to the queue. The WALIGN() functions removes one to seven 2.5GPll symbols from the front of the queue every time it is called.

When no 2.5GPll symbols are inserted or deleted, this function will return the first four 2.5GPll symbols in the queue and remove them from the queue.

When X 2.5GPll idles symbols are inserted then X 2.5GPll idles symbols and the first 4 – X 2.5GPll symbols are returned by the function, and the first 4 - X 2.5GPll symbols are removed from the queue.

When X 2.5GPll symbols are deleted then the first X 2.5GPll symbols are removed from the queue and the next four in the queue are returned by the function and then removed from the queue.

WDECODE(x, y, z)

In the PCS Word Decode process, this function performs the mapping according to clause 200.2.4.5. The variable x is four sets of 2.5GPll variables rp\_dv<3:0>, rp\_er<3:0>, rpd<3:0><7:0>, the variable y is the current state of the wdecode\_state variable, and the variable z indicates whether the next four 2.5GPll variables are the final four 2.5GPll symbols of the |Q| or |Fsig| ordered-set.

The output the XGMll RXC<3:0>, RXD<31:0>, wdecode\_state

WENCODE(x, y)

In the PCS Word Encode process, this function performs the mapping according to clause 200.2.4.2. The variable x is the XGMll TXC<3:0>, TXD<31:0>, and the variable y is the current state of the wencode\_state variable.

The output is four sets of 2.5GPll variables and the updated state of the wencode\_state variable and is ordered as follows: tp\_en<3:0>, tp\_er<3:0>, tpd<3:0><7:0>, wencode\_state

### 200.2.6.1.5 Counters

good\_cgs

Count of consecutive valid code-groups received.

plcnt

Count of number the number of 2.5GPll symbols to preload in the Word Alignment process after release from the RESET state. The number of symbols to preload is implementation dependent and should be minimized to reduce latency.

The following counter is used only for the EEE capability:

wake\_error\_counter

A counter that is incremented each time that the LPI receive state diagram enters the RX\_WTF state indicating that a wake time fault has been detected. The counter is reflected in register 3.22 (see 45.2.3.10).

### 200.2.6.1.6 Messages

PMA\_UNITDATA.indication(rx\_code-group<9:0>)

A signal sent by the PMA Receive process conveying the next code-group received over the medium (see 200.3.1.2).

PMA\_UNITDATA.request(tx\_code-group<9:0>)

A signal sent to the PMA Transmit process conveying the next code-group ready for transmission over the medium (see 200.3.1.1).

PMD\_SIGNAL.indication(signal\_detect)

A signal sent by the PMD to indicate the status of the signal being received on the MDI.

PUDI

Alias for PMA\_UNITDATA.indication(rx\_code-group<9:0>).

PUDR

Alias for PMA\_UNITDATA.request(tx\_code-group<9:0>).

SUDI

Alias for SYNC\_UNITDATA.indicate(parameters).

SYNC\_UNITDATA.indicate(parameters)

A signal sent by the PCS Synchronization process to the PCS Receive process conveying the following parameters:

Parameters:     [/x/]; the latched value of the indicated code-group (/x/);  
                  EVEN/ODD; The latched state of the rx\_even variable;

Value:    EVEN; Passed when the latched state of rx\_even=TRUE.  
          ODD; Passed when the latched state of rx\_even=FALSE.

TX\_OSET.indicate

A signal sent to the PCS Transmit ordered set process from the PCS Transmit code-group process signifying the completion of transmission of one ordered set.

The following messages are used only for the EEE capability:

**PMD\_RXQUIET.request(rx\_quiet)**

A signal sent by the PCS/PMA LPI receive state diagram to the PMD. Note that this message is ignored by devices that do not support EEE capability.

Values: TRUE: The receiver is in a quiet state and is not expecting incoming data.  
FALSE: The receiver is ready to receive data.

**PMD\_TXQUIET.request(tx\_quiet)**

A signal sent by the PCS/PMA LPI transmit state diagram to the PMD. Note that this message is ignored by devices that do not support the optional LPI mechanism.

Values: TRUE: The transmitter is in a quiet state and may cease to transmit a signal on the medium.  
FALSE: The transmitter is ready to transmit data.

### 200.2.6.1.7 Timers

**cg\_timer**

A continuous free-running timer.

Values: The condition `cg_timer_done` becomes true upon timer expiration.

Restart when: immediately after expiration; restarting the timer resets the condition `cg_timer_done`.

Duration: 3.2 ns nominal.

The following timers are used only for the EEE capability:

**rx\_tq\_timer**

This timer is started when the PCS receiver enters the `START_TQ_TIMER` state. The timer terminal count is set to  $T_{QR}$ . When the timer reaches terminal count, it will set the `rx_tq_timer_done = TRUE`.

**rx\_tw\_timer**

This timer is started when the PCS receiver enters the `RX_WAKE` state. The timer terminal count shall not exceed the maximum value of  $T_{WR}$  in [Table 200-7](#). When the timer reaches terminal count, it will set the `rx_tw_timer_done = TRUE`.

**rx\_wf\_timer**

This timer is started when the PCS receiver enters the `RX_WTF` state, indicating that the receiver has encountered a wake time fault. The `rx_wf_timer` allows the receiver an additional period in which to synchronize or return to the quiescent state before a link failure is indicated. The timer terminal count is set to  $T_{WTF}$ . When the timer reaches terminal count, it will set the `rx_wf_timer_done = TRUE`.

**tx\_ts\_timer**

This timer is started when the PCS transmitter enters the `TX_SLEEP` state. The timer terminal count is set to  $T_{SL}$ . When the timer reaches terminal count, it will set the `tx_ts_timer_done = TRUE`.

**tx\_tq\_timer**

This timer is started when the PCS transmitter enters the `TX_QUIET` state. The timer terminal count is set to  $T_{QL}$ . When the timer reaches terminal count, it will set the `tx_tq_timer_done = TRUE`.

**tx\_tr\_timer**

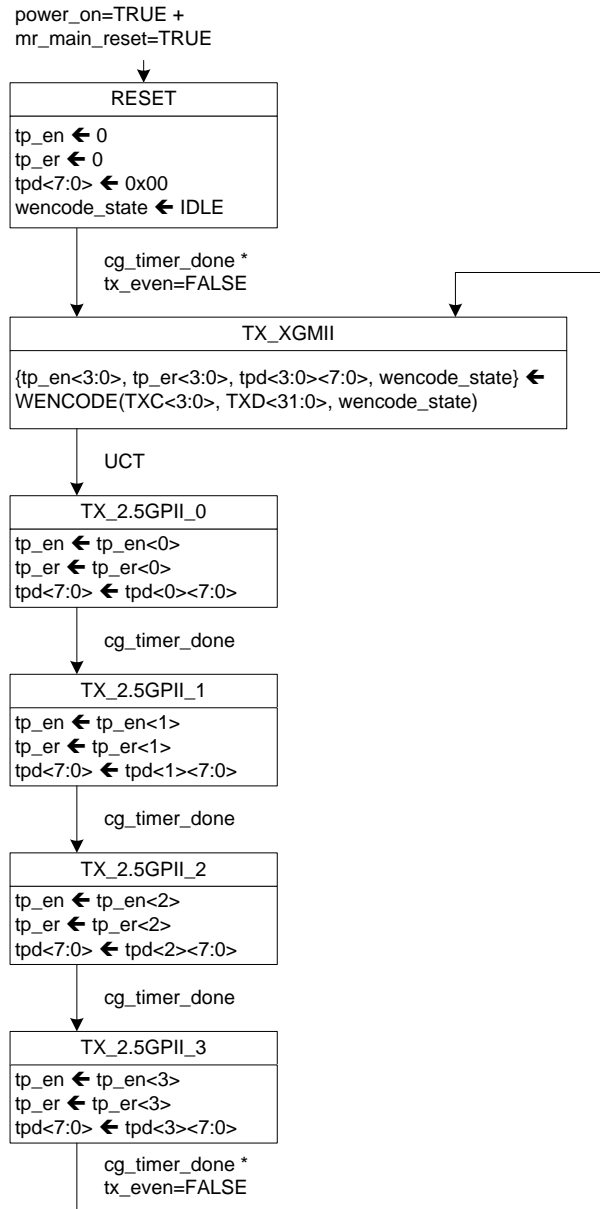
This timer is started when the PCS transmitter enters the `TX_REFRESH` state. The timer terminal count is set to  $T_{UL}$ . When the timer reaches terminal count, it will set the `tx_tr_timer_done = TRUE`.

**200.2.6.2 State diagrams**

**200.2.6.2.1 Word Encode and Serializer**

The Word Encode process and Word Serializer process are merged into one state diagram depicted in **Figure 200-4**, including compliance with the associated state variables as specified in **200.2.6.1**.

The Word Encode process continuously maps the four XGMII lanes to four 2.5GPPII symbols via the WENCODE function in the TX\_XGMII state. The four 2.5GPPII symbols are then serialized and output one at a time by the Word Serializer process. The presentation of the 2.5GPPII symbols are synchronized to the PCS transmit process such that the 2.5GPPII index 0 and 2 symbols are presented to the PCS transmit process which will result in the corresponding ordered set to be output to the PMA when the variable tx\_even is TRUE and index 1 and 3 variables when tx\_even is FALSE.



**Figure 200-4** PCS Word Encode and Serializer state diagram



### 200.2.6.2.2 Transmit

The PCS Transmit process is depicted in two state diagrams: PCS Transmit ordered\_set and PCS Transmit code-group. The PCS shall implement its Transmit process as depicted in Figure 200-5 and Figure 200-6, including compliance with the associated state variables as specified in 200.2.6.1.

The Transmit ordered\_set process continuously sources ordered\_sets to the Transmit code-group process. Upon the assertion of tp\_en by the 2.5GPll, the SPD ordered\_set is sourced. Following the SPD, /D/ code-groups are sourced until tp\_en is deasserted. Following the de-assertion of tp\_en, EPD ordered\_sets are sourced. If tp\_en and tp\_er are both deasserted, the /R/ ordered\_set may be sourced, after which the sourcing of /I/ is resumed. If, while tp\_en is asserted, the tp\_er signal is asserted, the /V/ ordered\_set is sourced except when the SPD ordered set is selected for sourcing. If the 2.5GPll indicates sequence then /Q/ ordered\_sets are sourced. If the optional EEE is enabled and the 2.5GPll indicates low power idles then /LI/ ordered\_sets are sourced.

The Transmit code-group process continuously sources tx\_code-group<9:0> to the PMA based on the ordered\_sets sourced to it by the Transmit ordered\_set process. The Transmit code-group process determines the proper code-group to source based on even/odd-numbered code-group alignment, running disparity requirements, and ordered\_set format.

FM editor's note: This diagram is modified from Figure 36-5. Note that TX\_EN, TX\_ER, TXD has to be change to tp\_en, tp\_er, tpd.

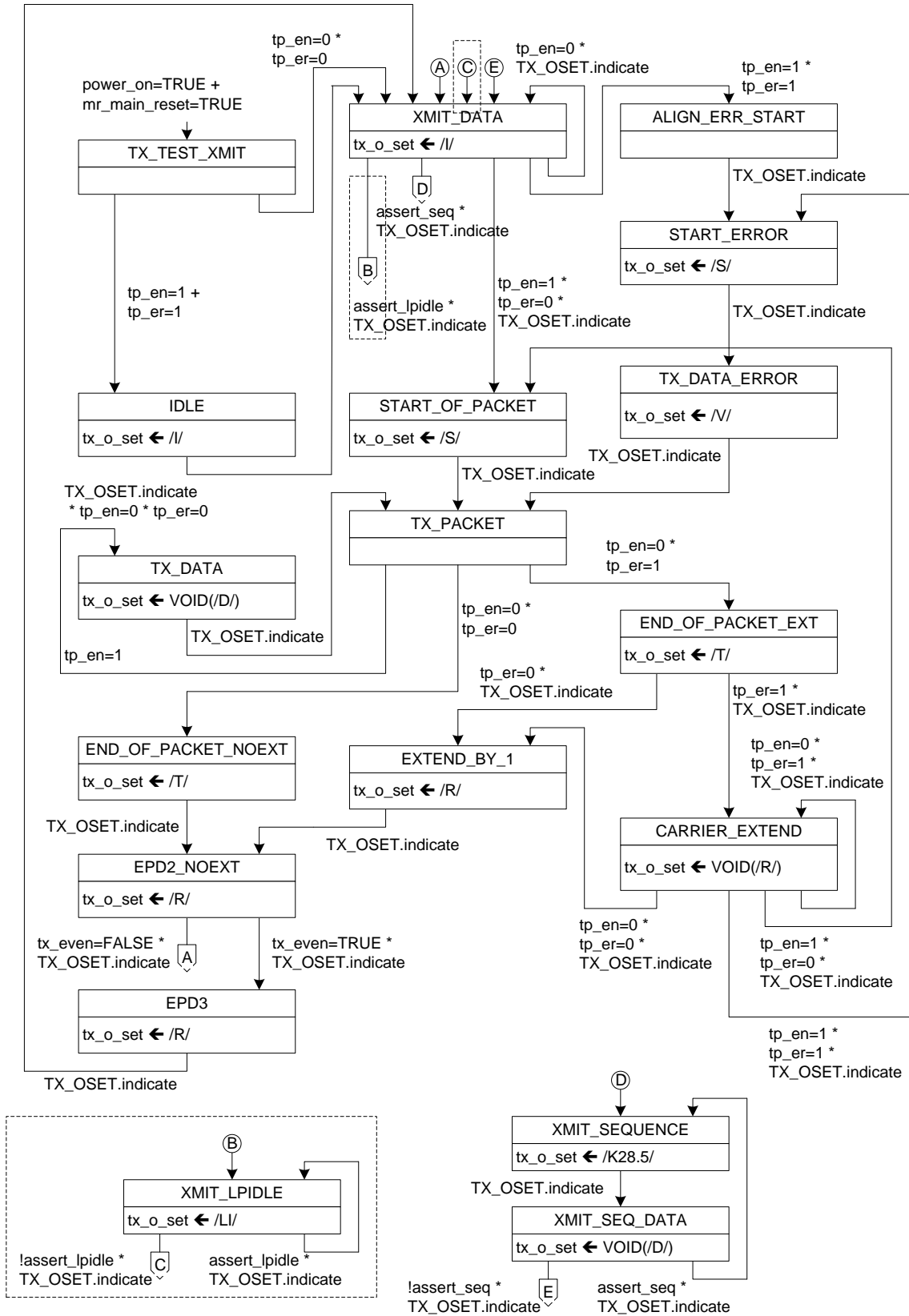


Figure 200-5 PCS transmit ordered\_set state diagram

FM editor's note: This diagram is modified from Figure 36-6

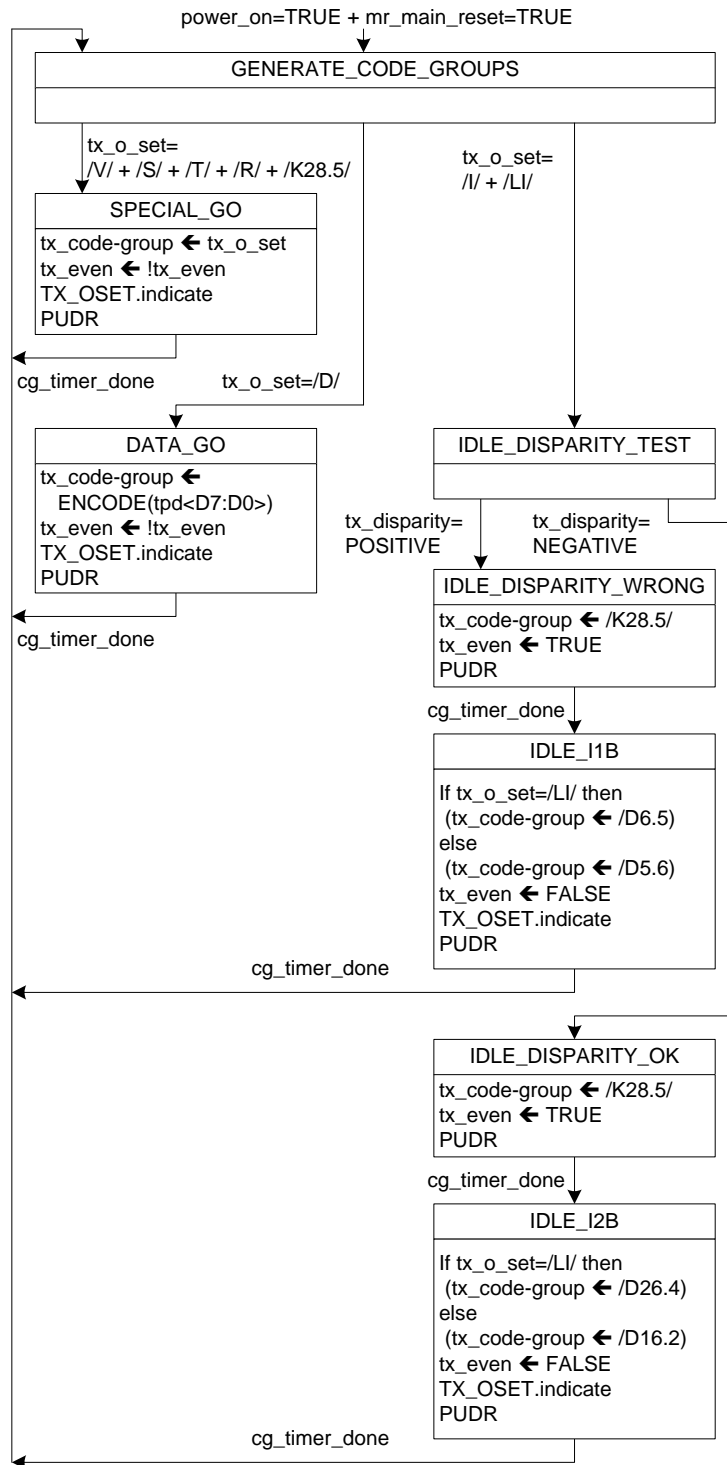


Figure 200-6 PCS transmit code-group state diagram

### 200.2.6.2.3 Synchronization

The PCS shall implement the Synchronization process as depicted in [Figure 200-7](#) including compliance with the associated state variables as specified in [200.2.6.1](#). The Synchronization process is responsible for determining whether the underlying receive channel is ready for operation. Failure of the underlying channel typically causes the PMA client to suspend normal actions.

A receiver that is in the LOSS\_OF\_SYNC state and that has acquired bit synchronization attempts to acquire code-group synchronization via the Synchronization process. Code-group synchronization is acquired by the detection of three ordered\_sets containing commas in their leftmost bit positions without intervening invalid code-group errors. Upon acquisition of code-group synchronization, the receiver enters the SYNC\_ACQUIRED\_1 state. Acquisition of synchronization ensures the alignment of multi-code-group ordered\_sets to even-numbered code-group boundaries.

Once synchronization is acquired, the Synchronization process tests received code-groups in sets of four code-groups and employs multiple sub-states, effecting hysteresis, to move between the SYNC\_ACQUIRED\_1 and LOSS\_OF\_SYNC states.

For EEE capability the relationship between sync\_status and code\_sync\_status is given by [Figure 200-10](#) PCS receive state diagram, part c (only required for the optional EEE capability); otherwise sync\_status is identical to code\_sync\_status.

FM editor's note: This diagram is identical to Figure 36-9

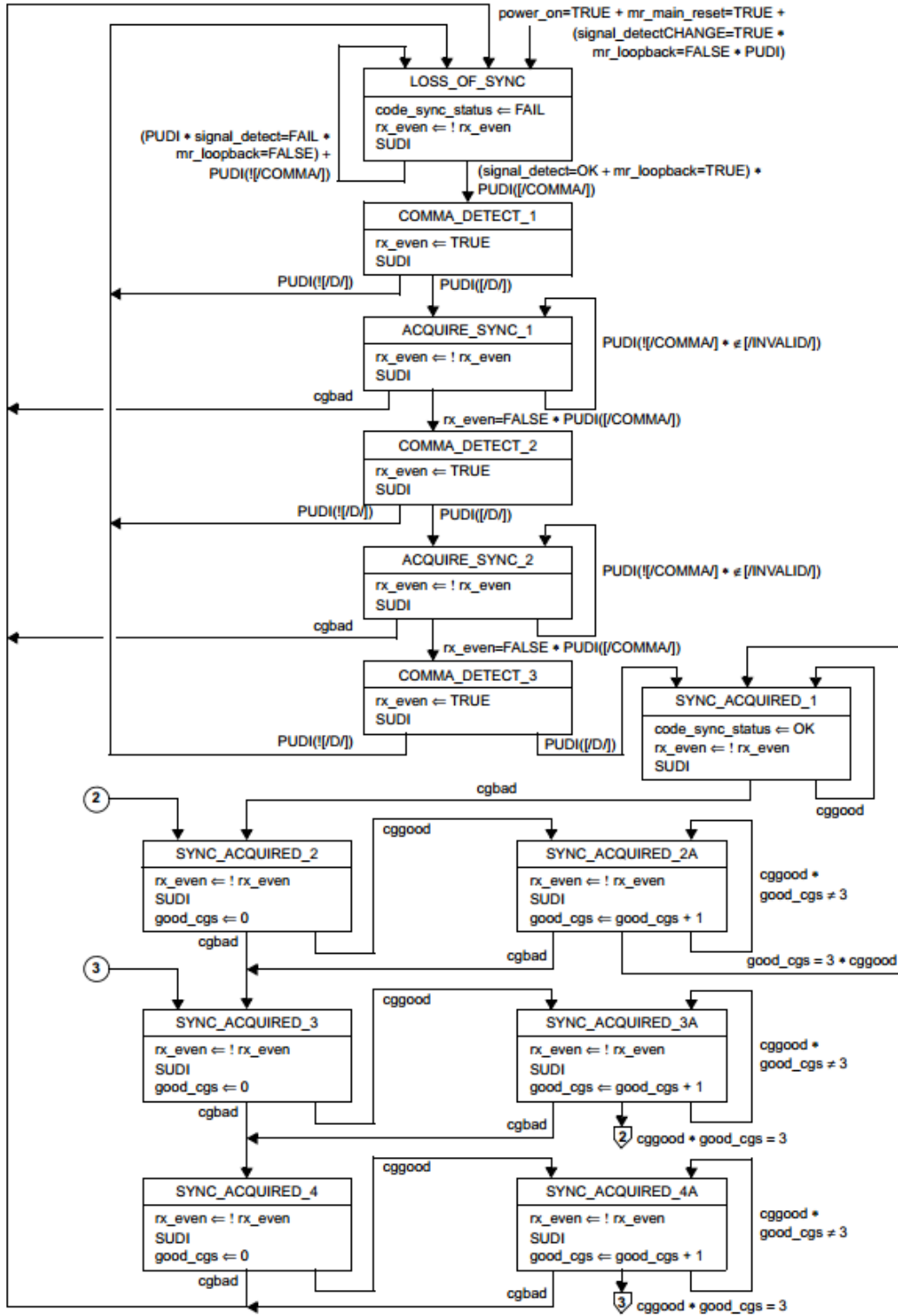


Figure 200-7 Synchronization state diagram

#### **200.2.6.2.4 Receive**

The PCS shall implement its Receive process as depicted in Figure 200-8 and Figure 200-9, including compliance with the associated state variables as specified in 200.2.6.1. Figure 200-10 shall be implemented if the optional EEE is enabled.

The PCS Receive process continuously passes  $rp_d<7:0>$  and sets the  $rp\_dv$  and  $rp\_er$  signals to the 2.5GPll based on the received code-group from the PMA.

FM editor's note: This diagram is modified from Figure 36-7a. Note that RX\_DV, RX\_ER, RXD has to be change to rp\_dv, rp\_er, rpd.

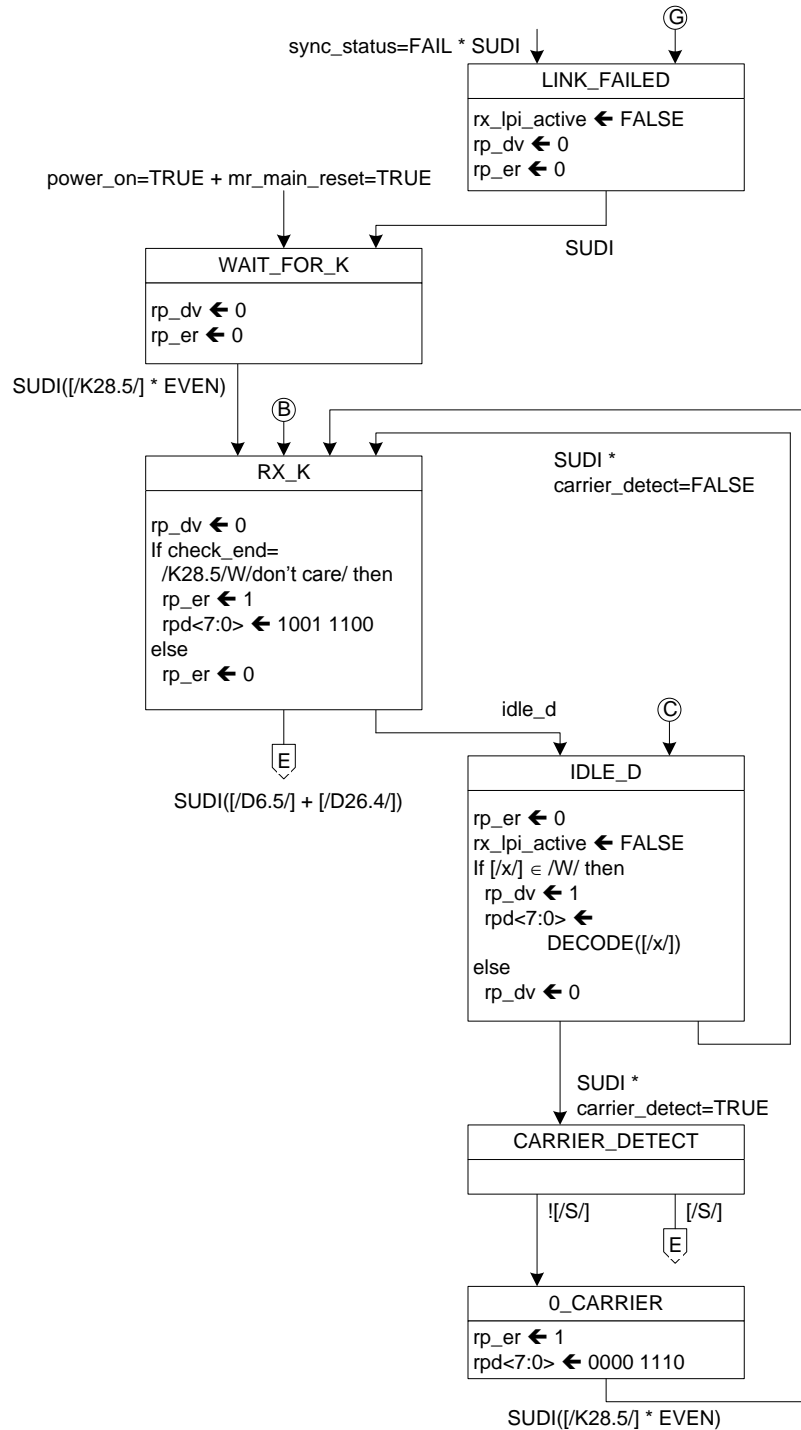


Figure 200-8 PCS receive state diagram, part a

FM editor's note: This diagram is modified from Figure 36-7b. Note that RX\_DV, RX\_ER, RXD has to be change to rp\_dv, rp\_er, rpd.

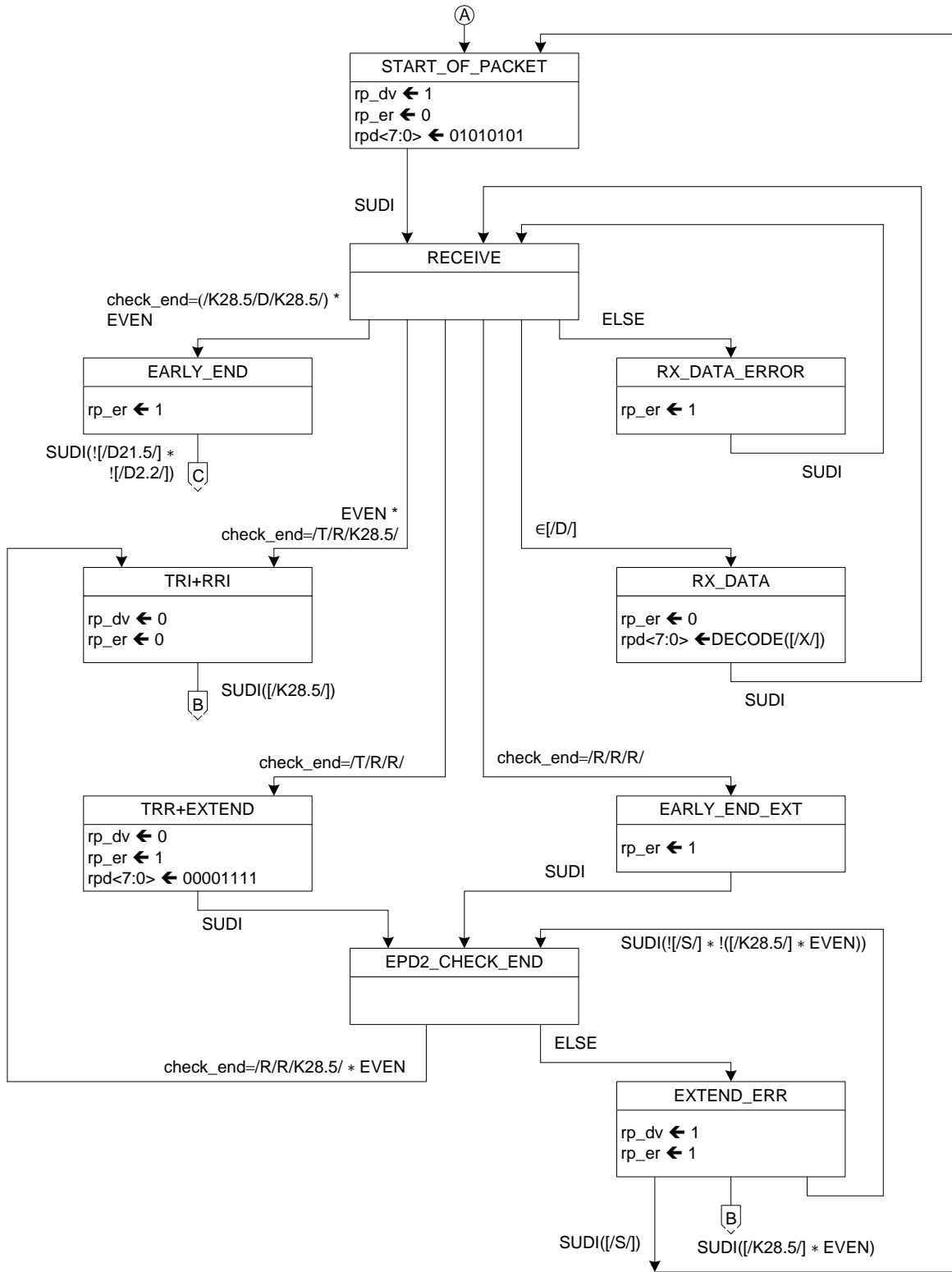


Figure 200-9 PCS receive state diagram, part b



FM editor's note: This diagram is modified from Figure 36-7c. Note that RX\_DV, RX\_ER, RXD has to be change to rp\_dv, rp\_er, rpd.

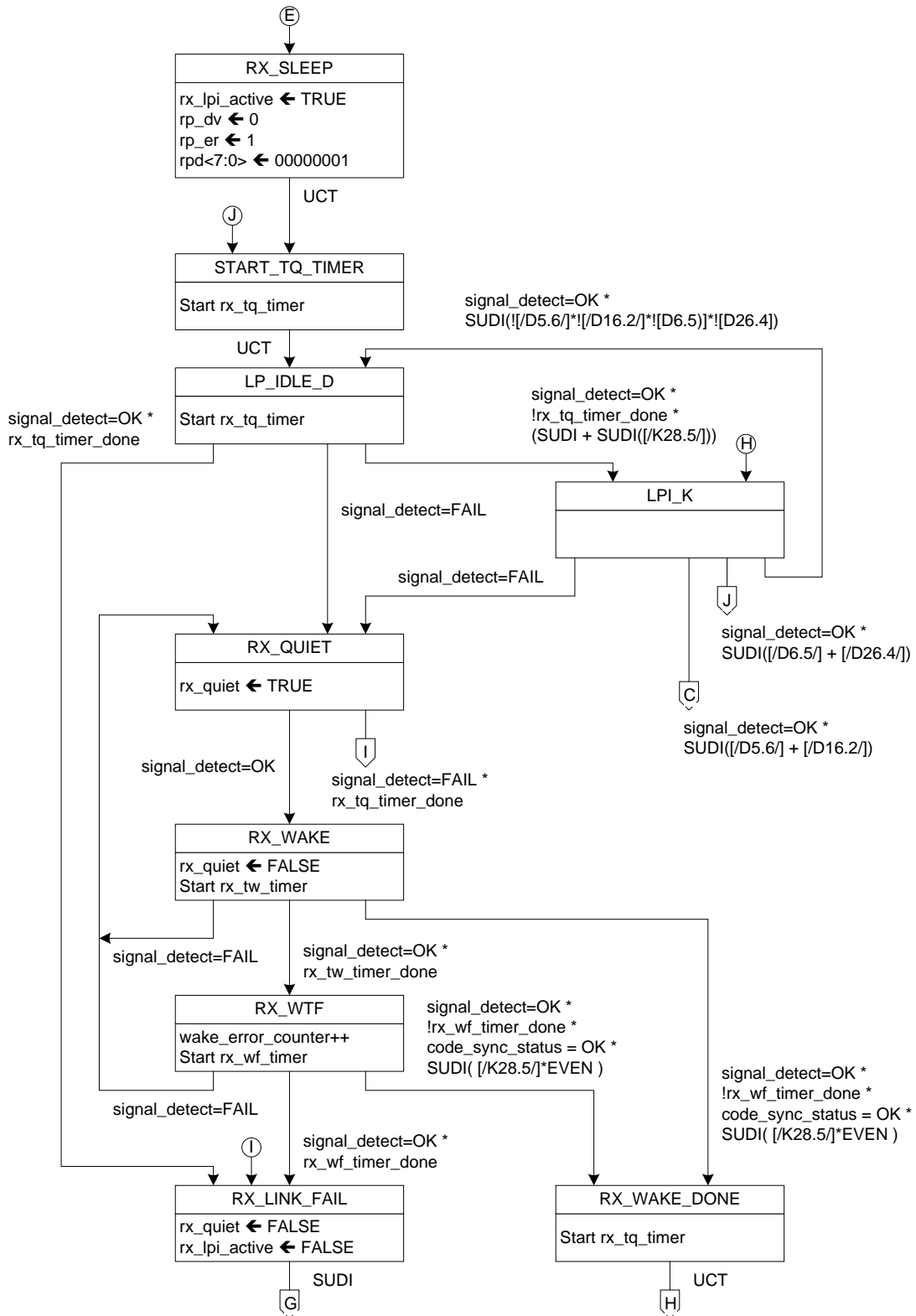
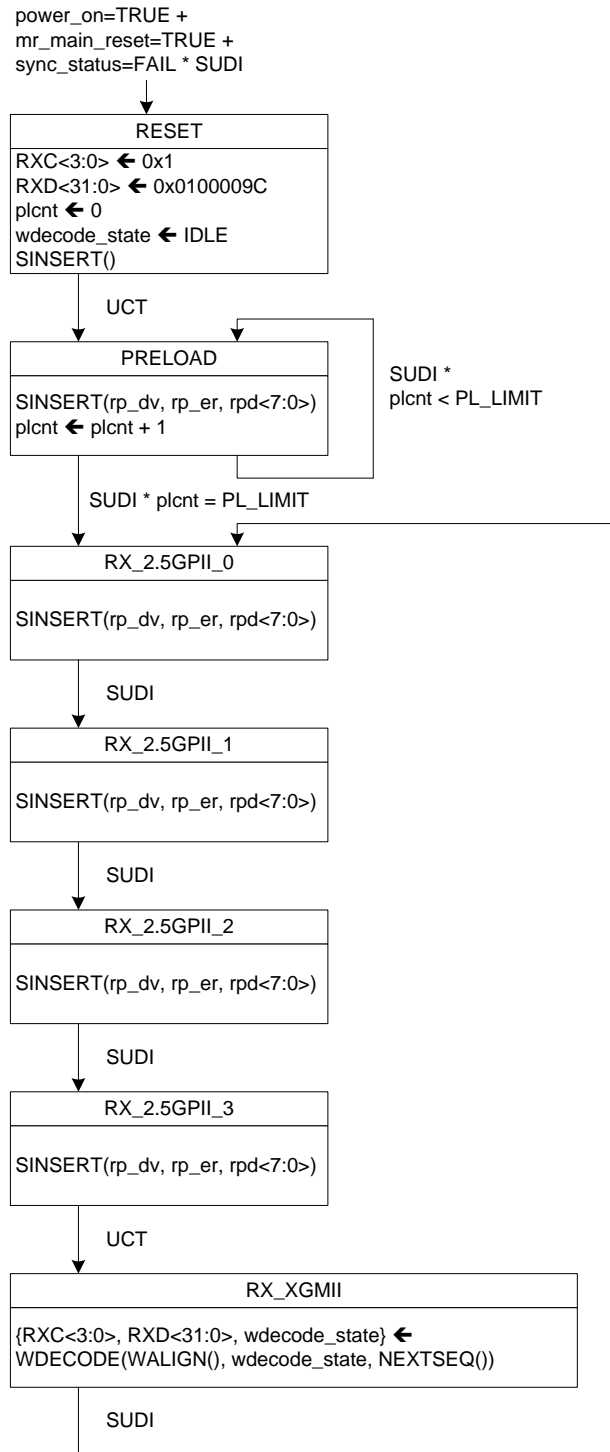


Figure 200-10 PCS receive state diagram, part c (only required for the optional EEE capability)

#### **200.2.6.2.5 Word Alignment and Decode**

The Word Alignment process and Word decode process are merged into one state diagram depicted in [Figure 200-11](#), including compliance with the associated state variables as specified in [200.2.6.1](#).

The Word Alignment process continuously queues the incoming 2.5GPll symbols from the PCS receive process and outputs four 2.5GPll symbols at a time using the WALIGN function. Symbols may be dropped or idles symbols added by the WALIGN function. The Word Decode process continuously maps the four 2.5GPll symbols and presents them on the XGMll using the WDECODE function in the RX\_XGMll state.

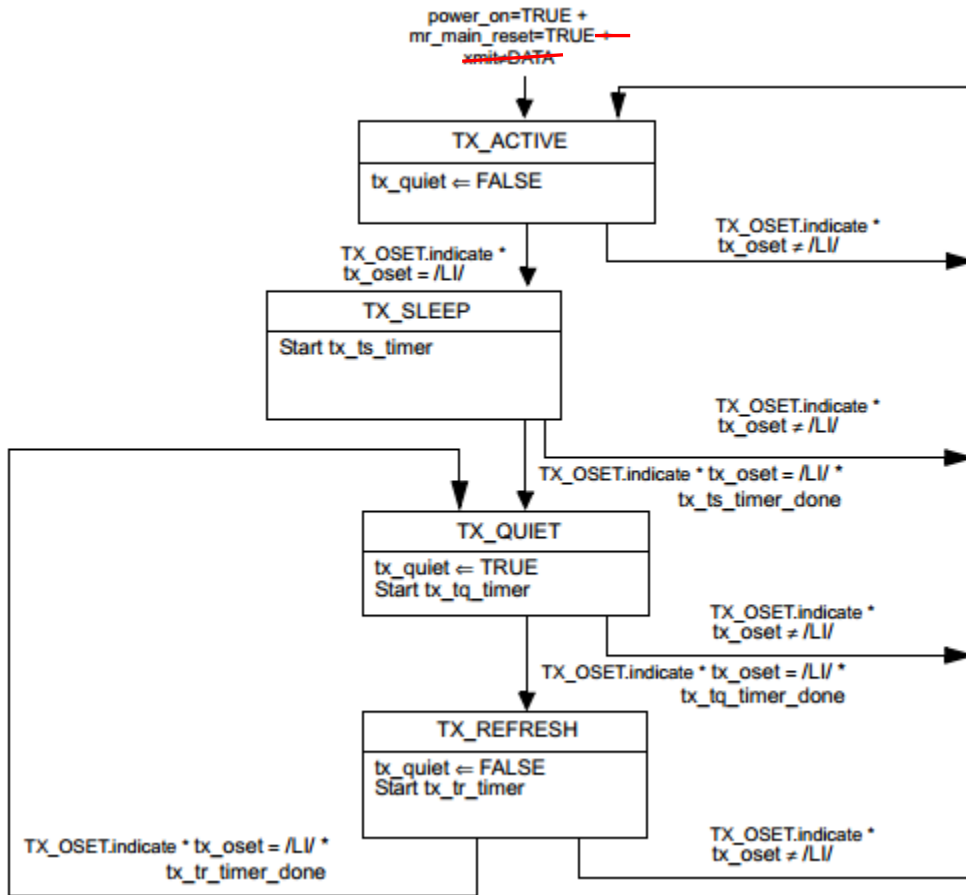


**Figure 200-11** Word Alignment and Decode state diagram

**200.2.6.2.6 LPI state diagram**

A PCS that supports the EEE capability shall implement the LPI transmit process as shown in **Figure 200-12**. The transmit LPI state diagram controls tx\_quiet, which disables the transmitter when true.

FM editor's note: This diagram is identical to Figure 36-10 except "+ xmit !=DATA" deleted



**Figure 200-12** LPI Transmit state diagram

The timer values for these state diagrams are shown in [Table 200-6](#) for transmit and [Table 200-7](#) for receive.

**Table 200-6** Transmitter LPI timing parameters

FM editor's note – copy table 36-8 as is.

Parameter	Description	Min	Max	Units
T <sub>SL</sub>	Local Sleep Time from entering the TX_SLEEP state to when tx_quiet is set to TRUE	19.9	20.1	μs
T <sub>QL</sub>	Local Quiet Time from when tx_quiet is set to TRUE to entry into the TX_REFRESH state	2.5	2.6	ms
T <sub>UL</sub>	Local Refresh Time from entry into the TX_REFRESH state to entry into the TX_QUIET state	19.9	20.1	μs

**Table 200-7** Receiver LPI timing parameters

FM editor's note – copy table 36-9 as is.

Parameter	Description	Min	Max	Units
T <sub>QR</sub>	The time the receiver waits for signal detect to be set to OK while in the LP_IDLE_D, LPI_K and RX_QUIET states before asserting a rx_fault	3	4	ms
T <sub>WR</sub>	Time the receiver waits in the RX_WAKE state before indicating a wake time fault (WTF)		11	μs
T <sub>WTF</sub>	Wake time fault recovery time		1	ms

#### 200.2.6.2.7 LPI status and management

For EEE capability, the PCS indicates to the management system that LPI is currently active in the receive and transmit directions using the status variables shown in [Table 36-10](#).

## **200.3 Physical Medium Attachment (PMA) sublayer**

### **200.3.1 Service Interface**

The PMA provides a Service Interface to the PCS. These services are described in an abstract manner and do not imply any particular implementation. The PMA Service Interface supports the exchange of code-groups between PCS entities. The PMA converts code-groups into bits and passes these to the PMD, and vice versa. It also generates an additional status indication for use by its client.

The following primitives are defined:

PMA\_UNITDATA.request(tx\_code-group<9:0>)

PMA\_UNITDATA.indication(rx\_code-group<9:0>)

#### **200.3.1.1 PMA\_UNITDATA.request**

This primitive defines the transfer of data (in the form of code-groups) from the PCS to the PMA. PMA\_UNITDATA.request is generated by the PCS Transmit process.

##### **200.3.1.1.1 Semantics of the service primitive**

PMA\_UNITDATA.request(tx\_code-group<9:0>)

The data conveyed by PMA\_UNITDATA.request is the tx\_code-group<9:0> parameter defined in [200.2.6.1.3](#).

##### **200.3.1.1.2 When generated**

The PCS continuously sends, at a nominal rate of 312.5 MHz, tx\_codegroup<9:0> to the PMA.

##### **200.3.1.1.3 Effect of receipt**

Upon receipt of this primitive, the PMA generates a series of ten PMD\_UNITDATA.request primitives, requesting transmission of the indicated tx\_bit to the PMD.

#### **200.3.1.2 PMA\_UNITDATA.indication**

This primitive defines the transfer of data (in the form of code-groups) from the PMA to the PCS. PMA\_UNITDATA.indication is used by the PCS Synchronization process.

##### **200.3.1.2.1 Semantics of the service primitive**

PMA\_UNITDATA.indication(rx\_code-group<9:0>)

The data conveyed by PMA\_UNITDATA.indication is the rx\_code-group<9:0> parameter defined in [200.2.6.1.3](#).

##### **200.3.1.2.2 When generated**

The PMA continuously sends one rx\_code-group<9:0> to the PCS corresponding to the receipt of each code-group aligned set of ten PMD\_UNITDATA.indication primitives received from the PMD. The nominal rate of the PMA\_UNITDATA.indication primitive is 312.5 MHz, as governed by the recovered bit clock.

##### **200.3.1.2.3 Effect of receipt**

The effect of receipt of this primitive by the client is unspecified by the PMA sublayer.

### 200.3.2 Functions within the PMA

Figure 200-2 and Figure 200-3 depicts the mapping of the four octet-wide data path of the XGMII to the ten-bit-wide code-groups of the PMA Service Interface, and on to the serial PMD Service Interface. The PMA comprises the PMA Transmit and PMA Receive processes for 2.5GBASE-X.

The PMA Transmit process serializes tx\_code-groups into tx\_bits and passes them to the PMD for transmission on the underlying medium, according to Figure 200-3. Similarly, the PMA Receive process deserializes rx\_bits received from the PMD according to Figure 200-3. The PMA continuously conveys ten bit code-groups to the PCS, independent of code-group alignment. After code-group alignment is achieved, based on comma detection, the PCS converts code-groups into XGMII data octets, according to 200.2.6.2.4 and 200.2.6.2.5.

The proper alignment of a comma used for code-group synchronization is depicted in Figure 200-3.

#### 200.3.2.1 Data delay

The PMA maps a nonaligned one-bit data path from the PMD to an aligned, ten-bit-wide data path to the PCS, on the receive side. Logically, received bits must be buffered to facilitate proper code-group alignment. These functions necessitate an internal PMA delay of at least ten bit times. In practice, code-group alignment may necessitate even longer delays of the incoming rx\_bit stream.

#### 200.3.2.2 PMA transmit function

The PMA Transmit function passes data unaltered (except for serializing) from the PCS directly to the PMD. Upon receipt of a PMA\_UNITDATA.request primitive, the PMA Transmit function shall serialize the ten bits of the tx\_code-group<9:0> parameter and transmit them to the PMD in the form of ten successive PMD\_UNITDATA.request primitives, with tx\_code-group<0> transmitted first, and tx\_code-group<9> transmitted last.

#### 200.3.2.3 PMA receive function

The PMA Receive function passes data unaltered (except for deserializing and possible code-group slipping upon code-group alignment) from the PMD directly to the PCS. Upon receipt of ten successive PMD\_UNITDATA.indication primitives, the PMA shall assemble the ten received rx\_bits into a single tenbit value and pass that value to the PCS as the rx\_code-group<9:0> parameter of the primitive PMA\_UNITDATA.indication, with the first received bit installed in rx\_code-group<0> and the last received bit installed in rx\_code-group<9>. An exception to this operation is specified in 200.3.2.4.

#### 200.3.2.4 Code-group alignment

In the event the PMA sublayer detects a comma+ within the incoming rx\_bit stream, it may realign its current code-group boundary, if necessary, to that of the received comma+ as shown in Figure 36–3. This process is referred to in this document as code-group alignment. During the code-group alignment process, the PMA sublayer may delete or modify up to four, but shall delete or modify no more than four, ten-bit code-groups in order to align the correct receive clock and code-group containing the comma+. This process is referred to as code-group slipping.

In addition, the PMA sublayer is permitted to realign the current code-group boundary upon receipt of a comma-pattern.

### 200.3.3 Loopback mode

Loopback mode shall be provided, as specified in this subclause, by the transmitter and receiver of a device as a test function to the device. When Loopback mode is selected, transmission requests passed to the transmitter are shunted directly to the receiver, overriding any signal detected by the receiver on its attached link. A device is explicitly placed in Loopback mode (i.e., Loopback mode is not the normal mode of operation of a device). The method of implementing Loopback mode is not defined by this standard.

NOTE—Loopback mode may be implemented either in the parallel or the serial circuitry of a device.

#### 200.3.3.1 Receiver considerations

Entry into or exit from Loopback mode may result in a temporary loss of synchronization.

### **200.3.3.2 Transmitter considerations**

While in Loopback mode, the transmitter output is not defined.

### **200.3.4 Test functions**

A limited set of test functions may be provided as an implementation option for testing of the transmitter function or for testing of an attached receiver.

Some test functions that are not defined by this standard may be provided by certain implementations. Compliance with the standard is not affected by the provision or exclusion of such functions by an implementation. Random jitter test patterns for 2.5GBASE-X are specified in Annex 200A.

A typical test function is the ability to transmit invalid code-groups within an otherwise valid PHY bit stream. Certain invalid PHY bit streams may cause a receiver to lose word and/or bit synchronization. See ANSI X3.230-1994 [B21] (FC-PH), subclause 5.4, for a more detailed discussion of receiver and transmitter behavior under various test conditions.

### **200.4 Compatibility considerations**

There is no requirement for a compliant device to implement or expose any of the interfaces specified for the PCS or PMA. Implementations of a XGMII shall comply with the requirements as specified in Clause 46.

### **200.5 Delay constraints**

Predictable operation of the MAC Control PAUSE operation (Clause 31, Annex 31B) demands that there be an upper bound on the propagation delays through the network. This implies that MAC, MAC Control sublayer, and PHY implementers must conform to certain delay maxima, and that network planners and administrators conform to constraints regarding the cable topology and concatenation of devices.

The sum of transmit and receive delay contributed by the 2.5GBASE-X PCS and PMA shall be no more than TBD BT.

The reference point for all MDI measurements is the 50% point of the mid-cell transition corresponding to the reference bit, as measured at the MDI.

### **200.6 Environmental specifications**

*Editor's Note: Need input on whether this is adequate*

All equipment subject to this clause shall conform to the requirements of 71.8.

### **200.7 Protocol implementation conformance statement (PICS) proforma for Clause 200, Physical Coding Sublayer (PCS) and Physical Medium Attachment (PMA) sublayer, type 2.5GBASE-X**

*Editor's Note: To be inserted later by PICS editor after text stabilizes*



**Annex 200A**

(informative)

**Jitter test patterns**

This annex defines test patterns for 2.5GBASE-X PMDs. The patterns described in Annex 36A can be used for 2.5GBASE-X except the nominal bit rate is 2.5 times faster and any references to the GMII applies to the XGMII.

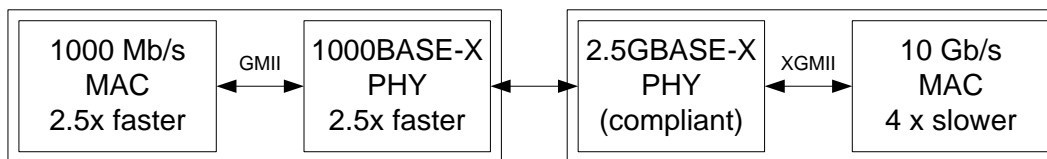
**Annex 200B**

(informative)

**Compatibility of 2.5GBASE-X PCS/PMA with 1000BASE-X PCS/PMA running 2.5 times faster**

This annex discusses the restrictions when operating 2.5GBASE-X PCS/PMA to a 1000BASE-X PCS/PMA link partner running 2.5 times faster. Compatibility of the PMD is outside the scope of this annex. In this annex when 1000BASE-X PCS/PMA is referred to, the 2.5 times speed up is implied.

The 2.5GBASE-X PCS/PMA is specified to be compatible with a link partner running 1000BASE-X PCS/PMA running 2.5 times faster as shown in **Figure 200B-1** with some restrictions.



**Figure 200B-1** 2.5X speed 1000BASE-X to 2.5GBASE-X Link

2.5GBASE-X is defined to operate only in full duplex. Hence the 1000BASE-X PCS can only operate in full duplex. If the half duplex carrier extend or carrier extend with errors are sent by the 1000BASE-X PCS, the 2.5GBASE-X PCS will convert it to receive errors. Note that a /T/R/ or a /T/R/R/ at the end of packet will be correctly converted as idles.

The 2.5GBASE-X PCS does not support clause 37 Auto-Negotiation. Hence, the 1000BASE-X PCS is expected to have its clause 37 Auto-Negotiation functionality disabled so that the /C/ ordered\_set will not be transmitted. If a 2.5GBASE-X PCS receives /C/ ordered\_set then undefined behavior may occur.

Since the 2.5GBASE-X PCS is attached to a MAC that can send out sequence ordered\_set (/Q/), a compliant 1000BASE-X PCS will interpret each /Q/ ordered\_set as four /I/ ordered set. Unlike /I/, /LI/, and /C/ ordered\_sets, there is no concept of correcting vs preserving the running disparity when /Q/ ordered\_sets are generated.

Unlike the GMII, there is no false carrier defined in the XGMII. The 2.5GBASE-X PCS Receive process can detect false carrier, but these will be converted to receive error by the PCS Word Decode process.

It is permissible for a compliant 1000BASE-X PCS transmit process to truncated the first byte of preamble in order to align the start of packet on the EVEN boundary. The implication of this is the 2.5Gb/s MAC has to be able to accept a seven byte preamble on the XGMII.

Since the start of packet has to align to lane 0 of the XGMII the 2.5Gb/s MAC may delete idle bytes in order to do the alignment. The 1G MAC running 2.5 time faster has to be able to handle IPG shrinkage.