



Canova Tech

The Art of Silicon Sculpting

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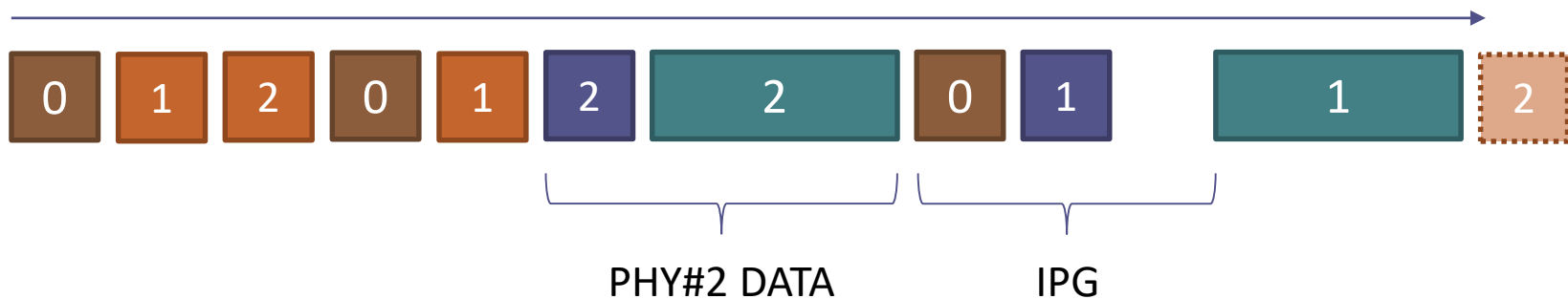
IEEE802.3cg TF

PHY-Level Collision Avoidance — Addendum #1

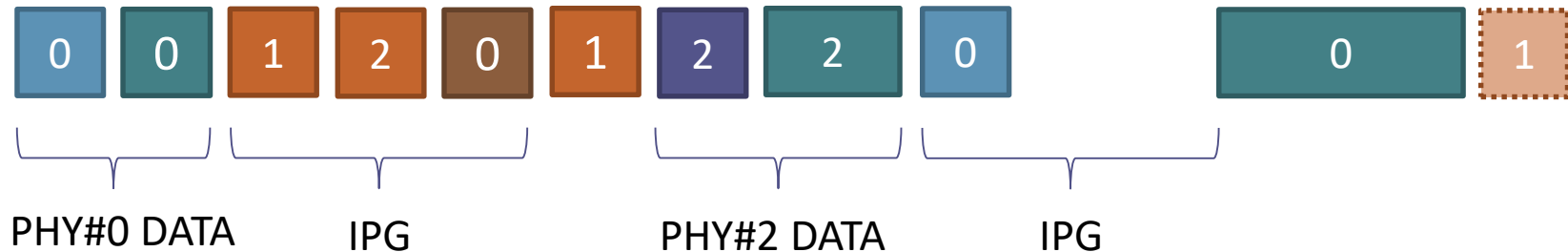
August 16th, 2017

Control signals

LINE



LINE



- **Control signals renamed since last pres.**
 - ACK → COMMIT
 - NACK → YIELD
- **Added new control signals**
 - COMMIT/S → same as COMMIT, sent by PHY #0 to allow synchronization
 - YIELD/S → same as YIELD, sent by PHY #0 to allow synchronization



COMMIT



YIELD



DATA



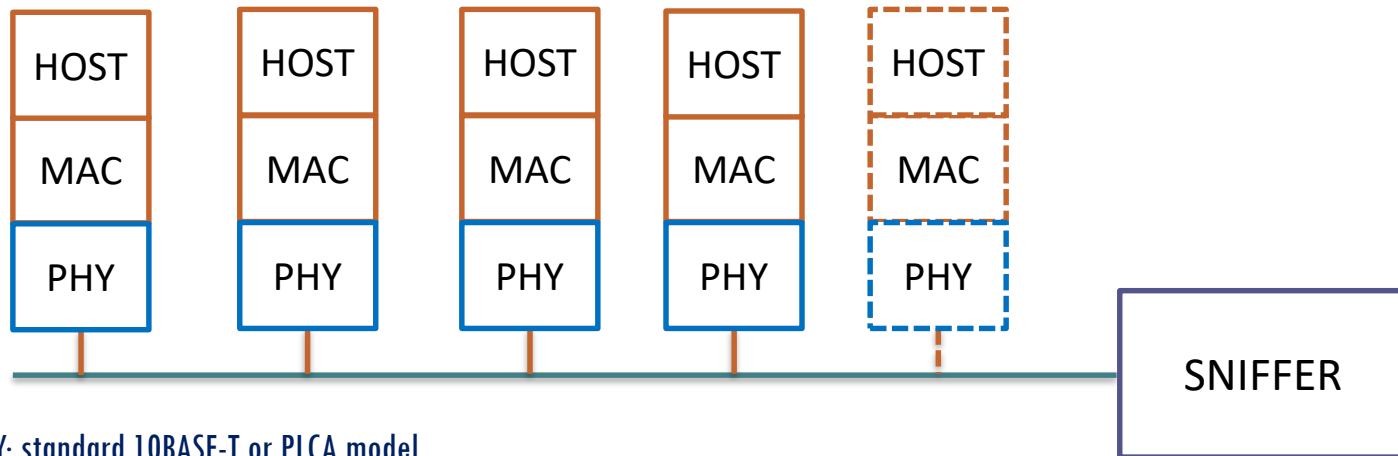
COMMIT/S



YIELD/S

- Packet Ordering
 - Node level: node “A” transmits packets #0, #1 and #2 at times $T_0 < T_1 < T_2$. Packets shall be received in the same order exactly (#0, #1, #2) with unspecified latency.
 - This is **guaranteed**
 - MAC procedures are synchronous → once a frame is waiting to be transmitted, the MAC does not process other frames until the current one is either transmitted or discarded because of excessive attempts.
 - » In PLCA proposal packets are NEVER discarded in such way as the packet is always transmitted within at maximum two attempts.
 - PHY is normally supposed to deliver the frame as soon as it's fed by the MAC.
 - » In PLCA proposal, the PHY shall buffer a *very short* portion of the packet in order to reduce latency.
 - » Not enough to allow misordering (would need to buffer at least two whole packets)
 - Network level: node “A” transmits packets #0, #1 and #2 while node “B” transmits packets #3, #4, #5 at times $T_0 < T_3 < T_1 < T_4 < T_2 < T_5$. Packets shall be received in the following order: #0, #3, #1, #4, #2, #5.
 - Not how ethernet works
 - Out of scope of PLCA proposal
 - Latency
 - Best case (min) occurs when the MAC initiate TX just before handshaking time (~ 0)
 - Worst case (max) occurs when the MAC initiate TX after handshaking time AND each other node commits to transmit a 1500 bytes packet ($\sim 7\text{ms}$ for a six node network)
 - Comparison with CSMA/CD (see next slides)

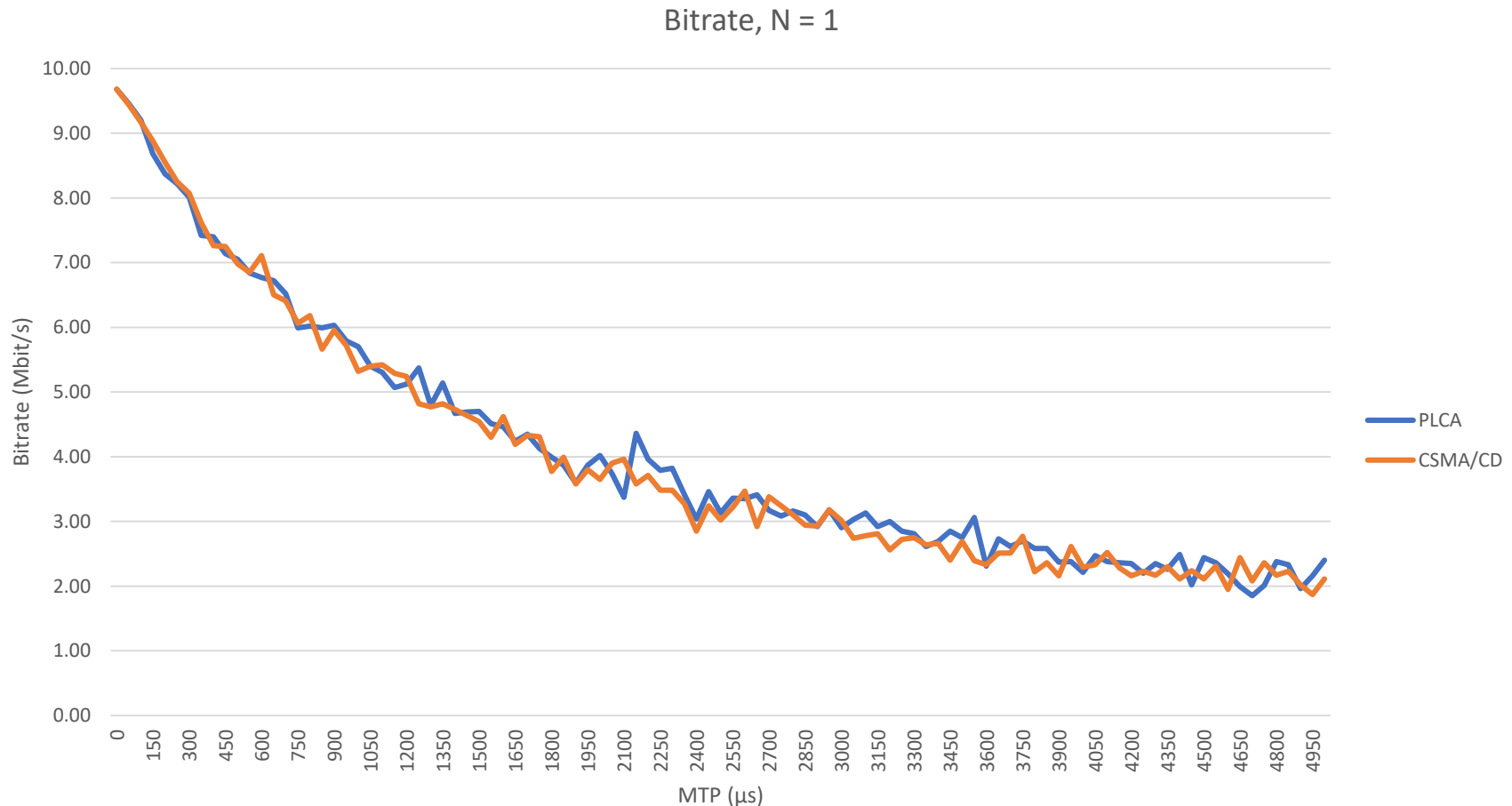
- Full digital simulation (verilog)
 - 4b/5b encoding + DME (25Mhz)
 - Use 5b codes S, R, T, H for COMMIT, YIELD, COMMIT/S, YIELD/S signaling



- PHY: standard 10BASE-T or PLCA model
- MAC: standard CSMA/CD capable MAC (802.3 clause 4)
 - host interface: DPRAM (one frame) + busy indication + size + trigger
 - PHY interface: MII (txd, txclk, txen, txer, rxd, rxclk, rxdv, rxer, **col**, **crs**)
- HOST: simple transmitter
 - wait for MAC BUSY = 0
 - wait random time between 0 and MTP (sim parameter, 0 = MAX speed)
 - write random data in DPRAM of size PKTSZ (sim. parameter $60 < \text{PKTSZ} < 1500$) or random size
- SNIFFER: measuring throughput, latency
 - throughput: number of received bytes (excluding FCS, PREAMBLE) / total simulation time
 - latency: time between BUSY = 1 and BUSY = 0 for each node

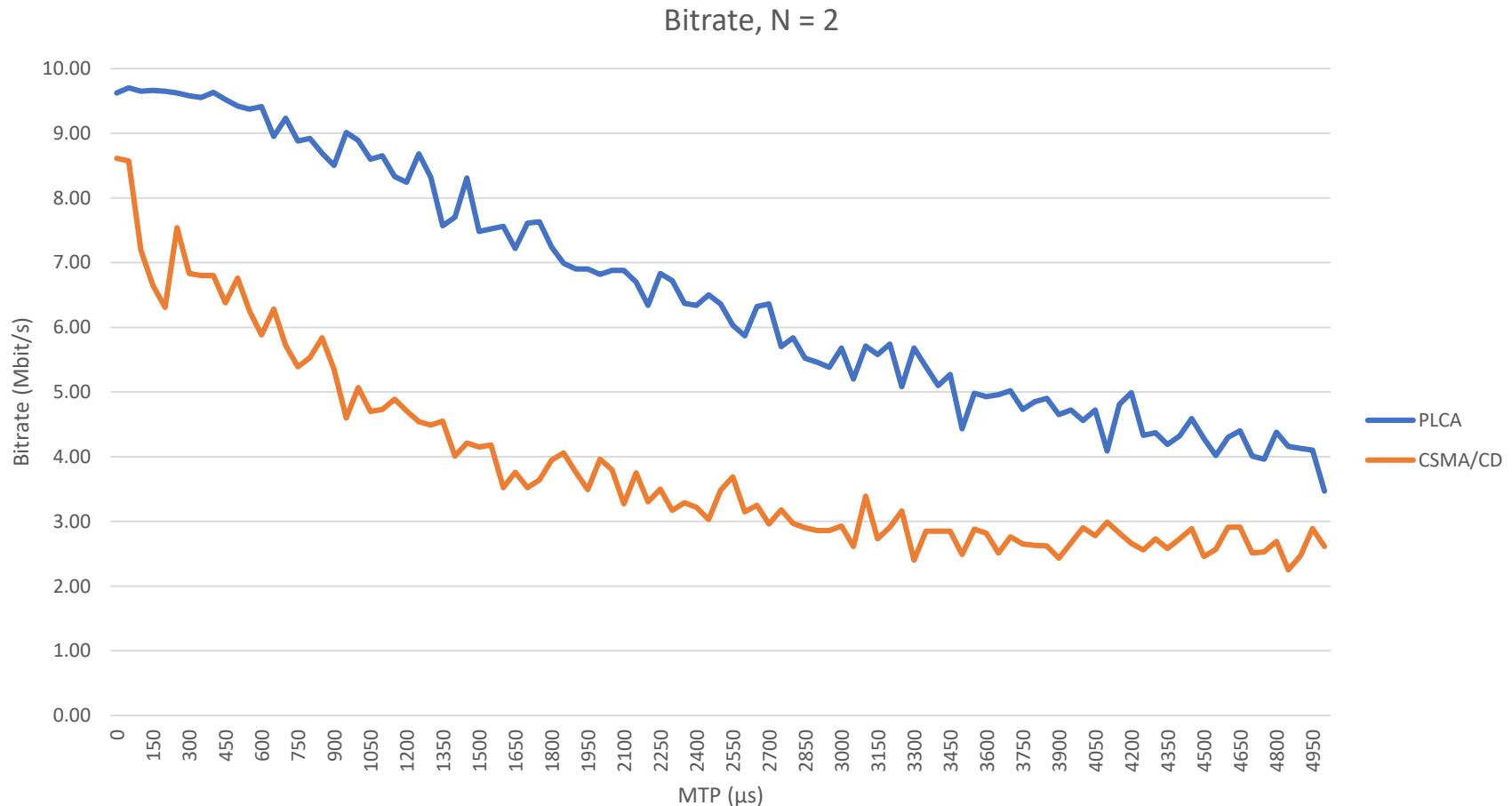
Throughput

- 100 pkts, random size (60, 1500), variable MTP and N (number of nodes)
- Comparison between 10base-T (simple CSMA/CD) and PLCA



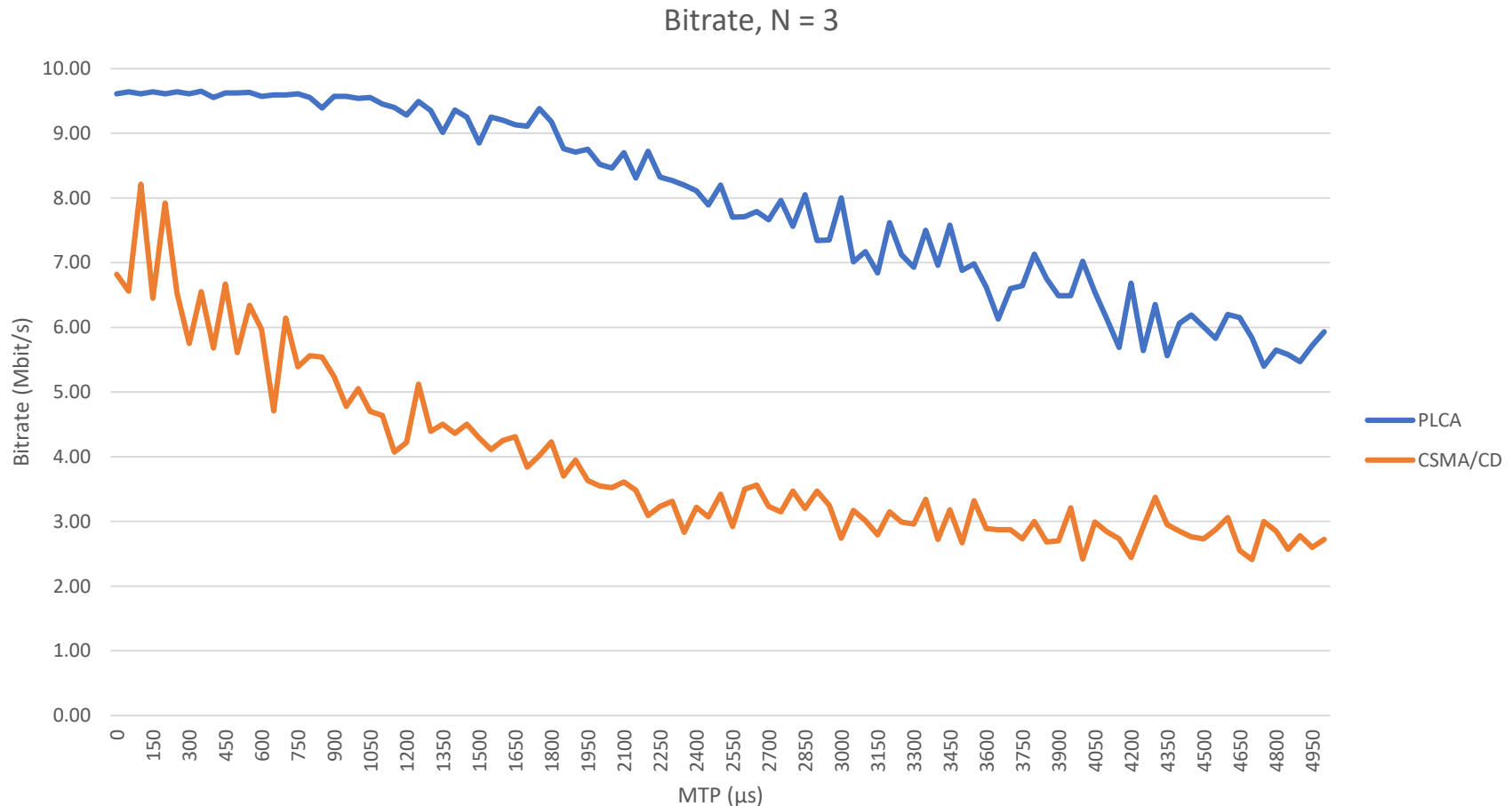
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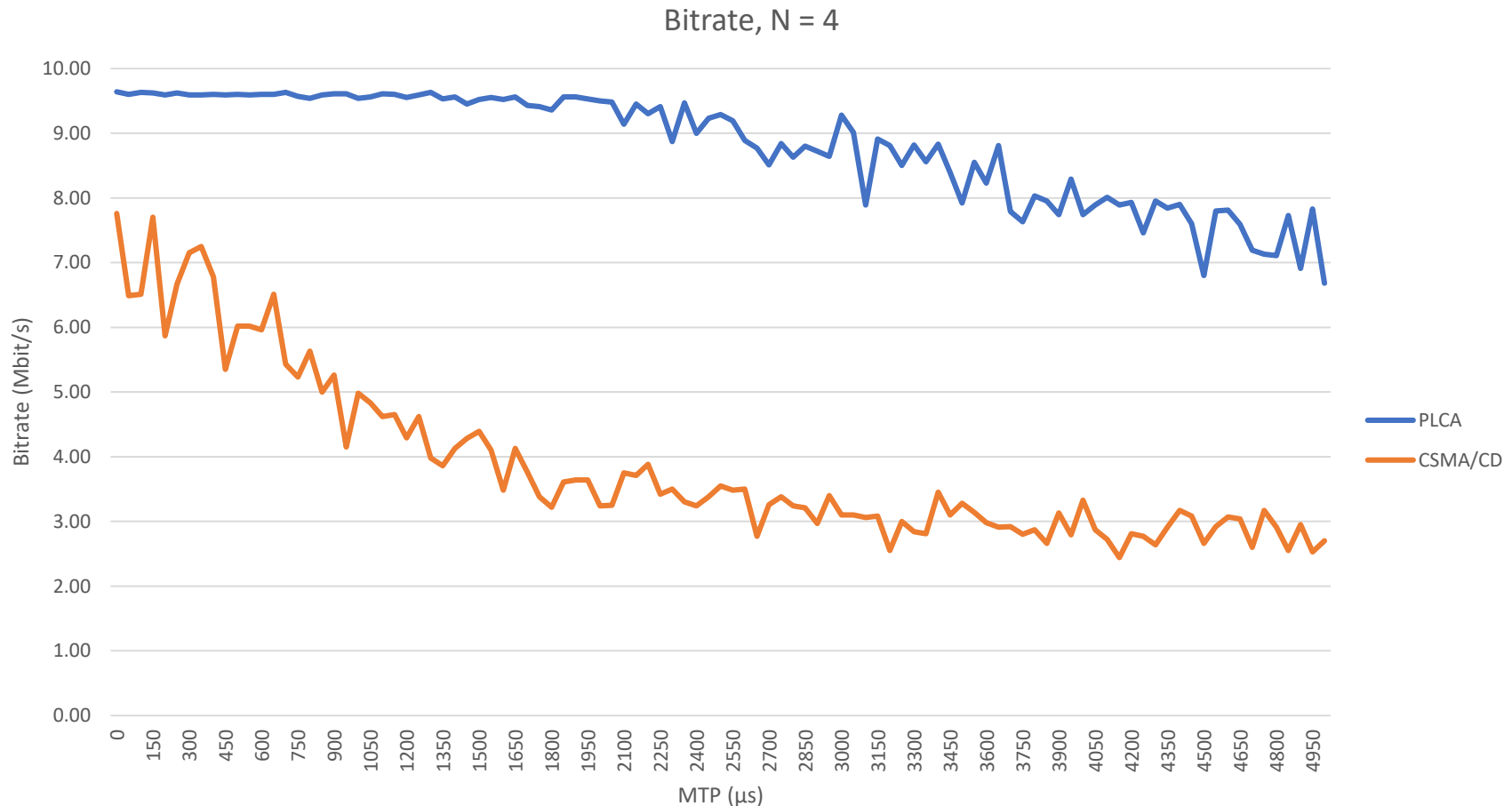
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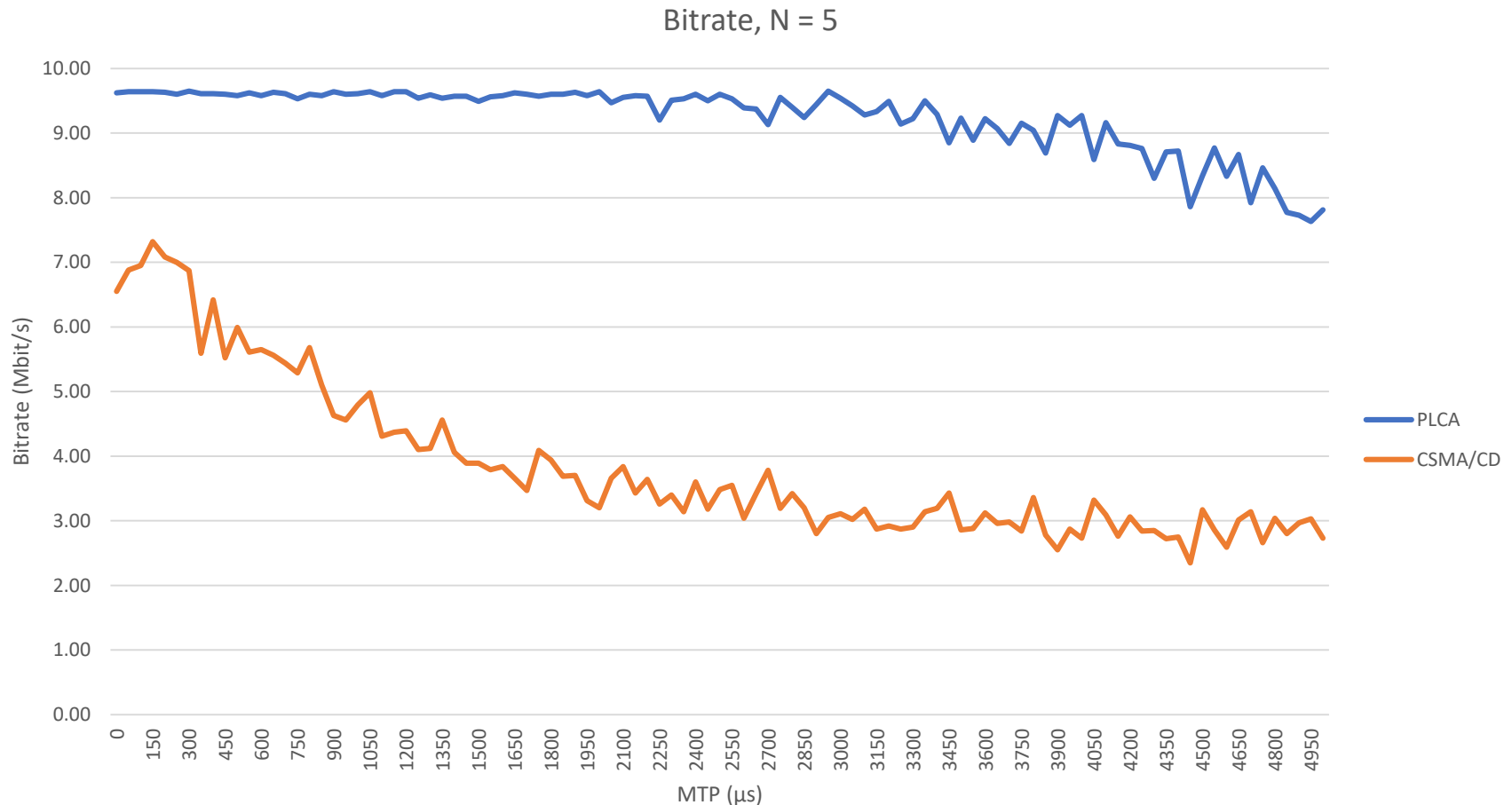
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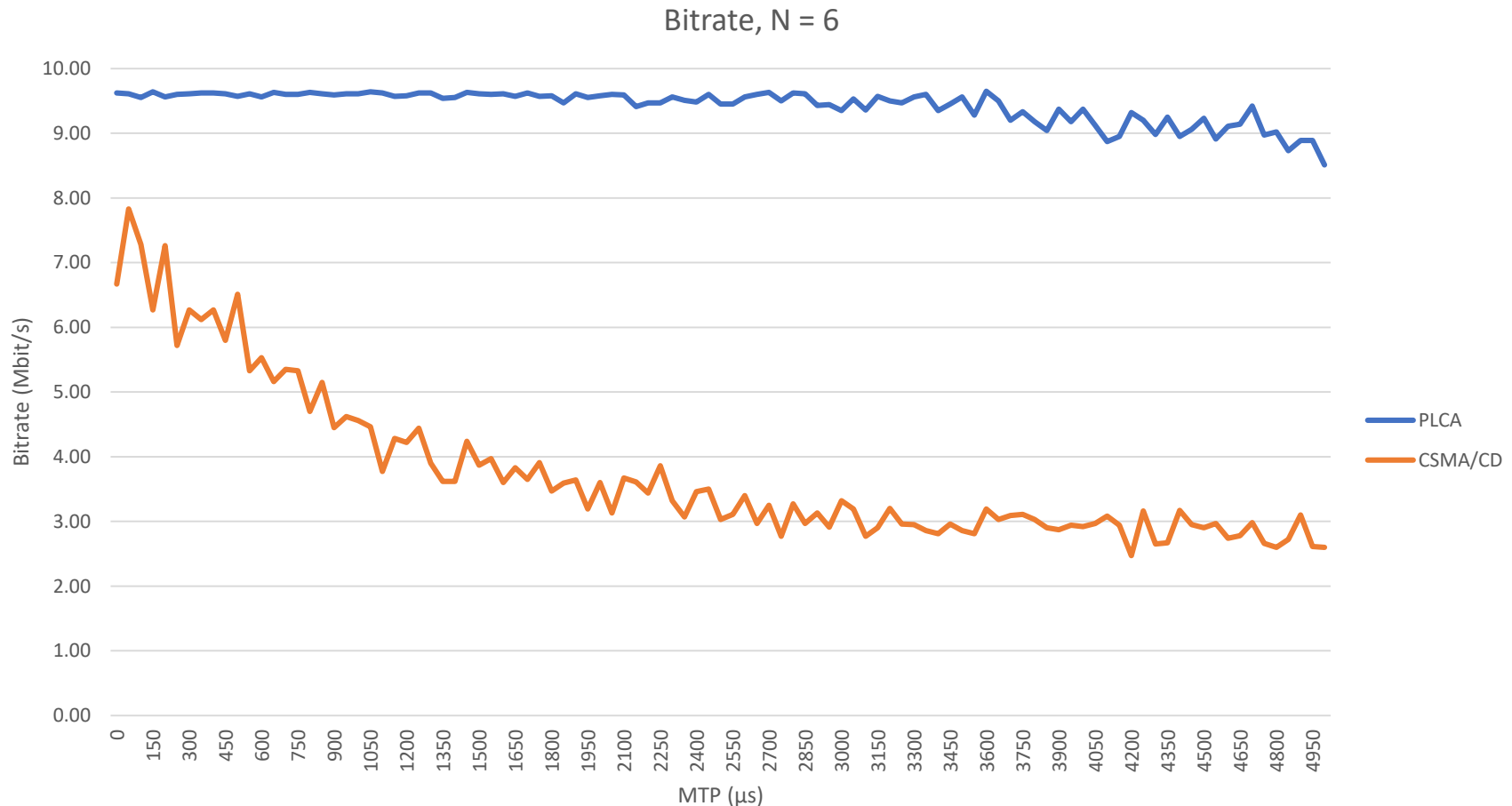
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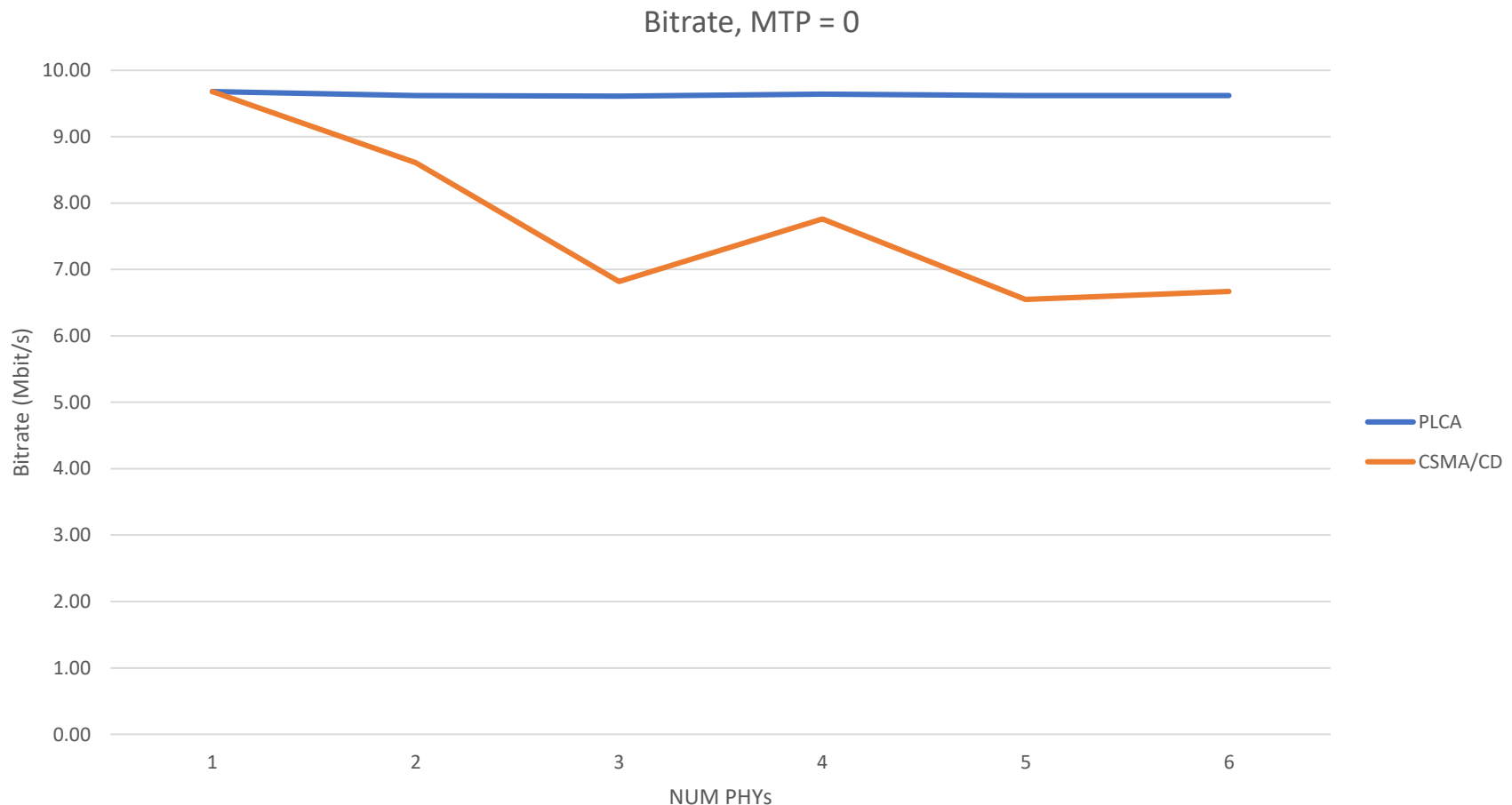
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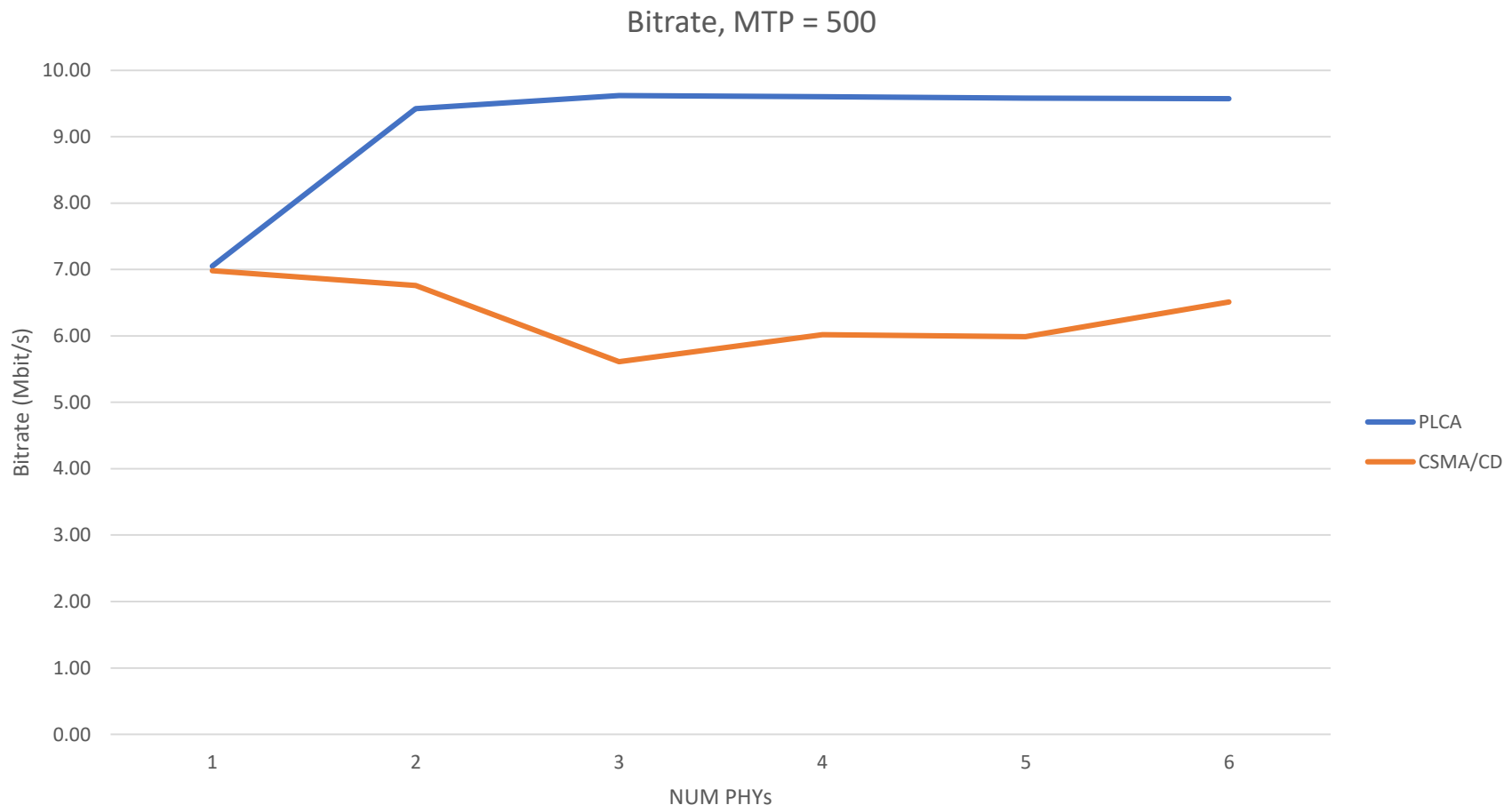
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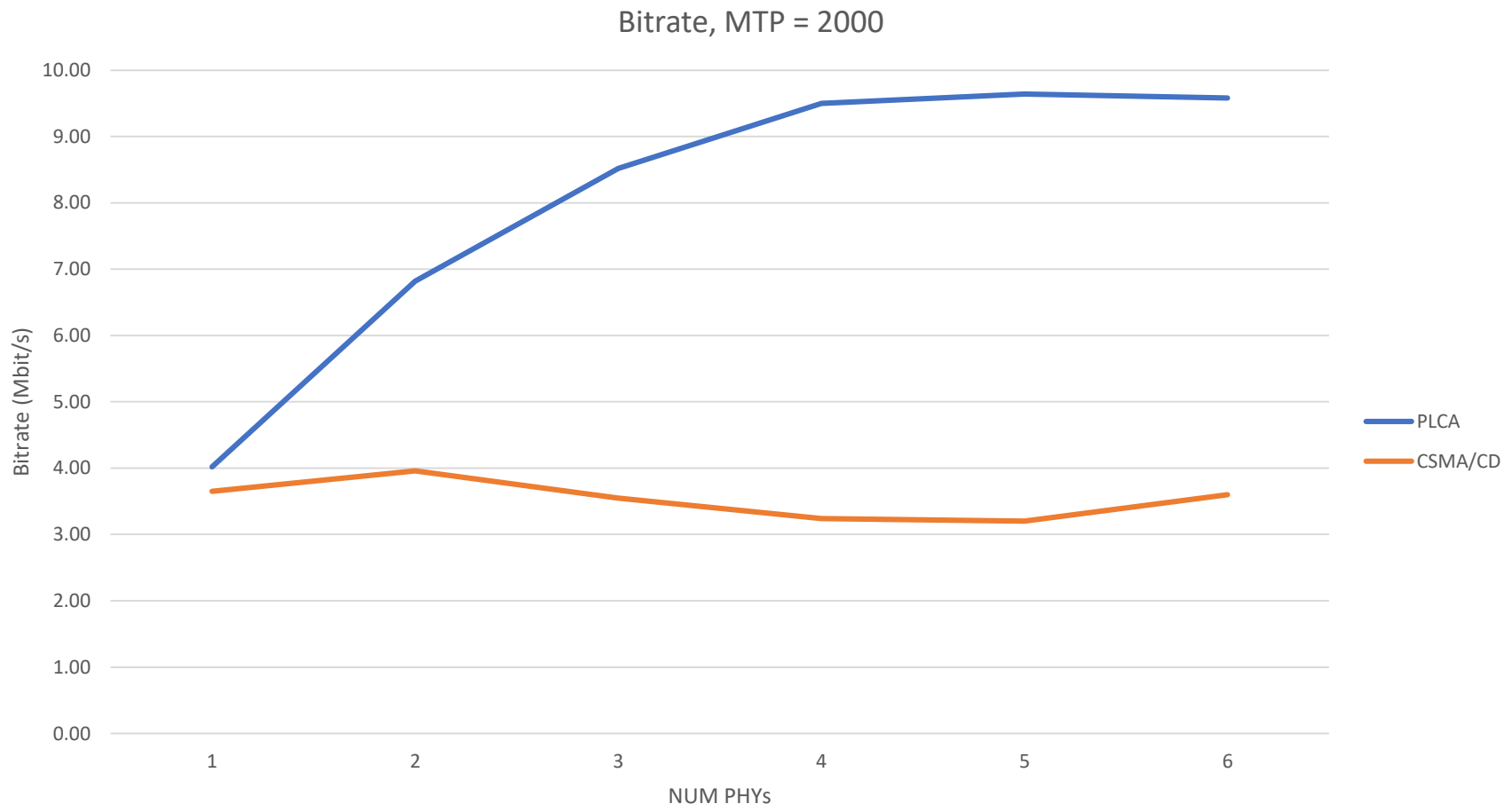
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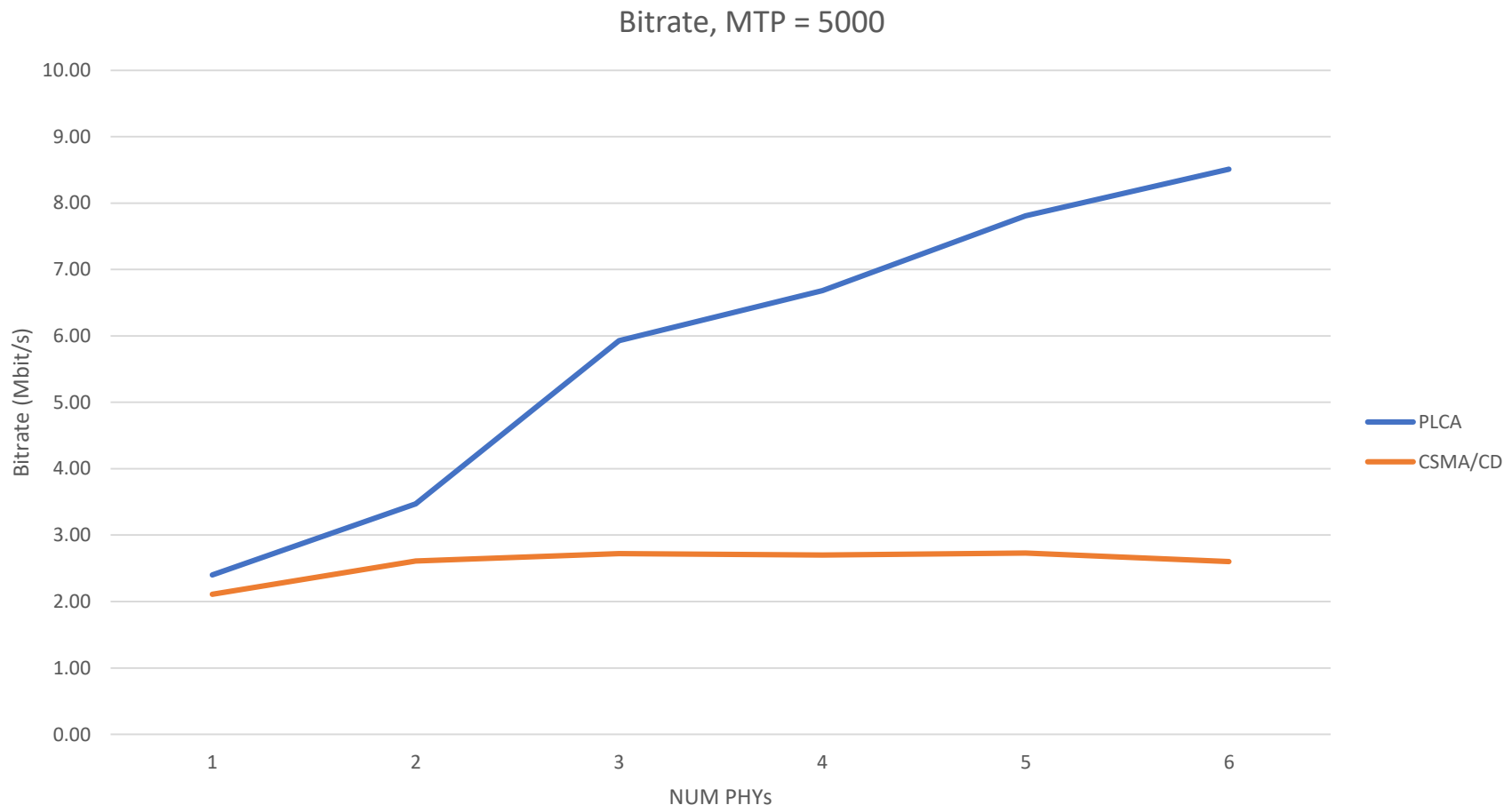
Throughput

- 100 pkts, random size (60, 1500), variable MTP and N (number of nodes)
- Comparison between 10base-T (simple CSMA/CD) and PLCA



Throughput

- 100 pkts, random size (60, 1500), variable MTP and N (number of nodes)
- Comparison between 10base-T (simple CSMA/CD) and PLCA



- 500 pkts, size = 60B, variable MTP, 6 nodes. Latencies in μs .
- Comparison between 10base-T (simple CSMA/CD) and PLCA

MTP	MAX_LAT	AVG_LAT	STDEV
0	57595.6	553.3	4826.0
500	59692.8	1034.2	4637.4
2000	29387.5	618.9	2298.2
5000	19645.4	264.0	1035.7

CSMA/CD

MTP	MAX_LAT	AVG_LAT	STDEV
0	443.4	441.1	26.2
500	546.4	186.4	90.7
2000	269.2	74.8	31.6
5000	223.7	64.0	17.8

PLCA

- Average latencies not so different at high loads
- Huge difference on MAX latency and STD deviation (due to collisions / backoff)
- PLCA is much more “deterministic”

- Still static configuration of the IDs
- How to handle “disabled” nodes?
 - COMMIT/YIELD is expected to be sent within a determined time (HS timer).
 - On timeout, next node will COMMIT/YIELD as appropriate
 - HS timer shall be long enough to accommodate for PHY latency but as short as possible not to degrade performance
 - Simulations performed with $HS_TIMER = 20\mu s$
 - Allows for $\sim 16\mu s$ of HS latency
 - Simulated with random missing nodes
 - Negligible impact on performance
- PHY with $ID = 0$ (master) use COMMIT/S and YIELD/S for handshaking.
 - Slave nodes synchronize on xxx/S signals before joining the HS mechanism
 - link_status shall not be signaled until first sync
 - Nodes receiving bad COMMIT/YIELD shall not transmit and re-synchronize
 - if this happens to the master, this one shall wait until there is no data on the line for a certain amount of time, then re-start transmitting HS.

- Simulations
 - More use-cases?
- Robustness
 - Simulate errored ACK/NACK handling
 - Corner cases study
- Optional
 - Auto-negotiation of IDs (instead of static config)?
 - not actively working on that
 - possible IDEA: keep a “virtual slot” at the end of the HS.
 - PHYs can send specific code + ID + CRC to get enumerated at that moment
 - monitor for collisions and backoff
 - On-the-fly election of new master for failover?

Thank You !