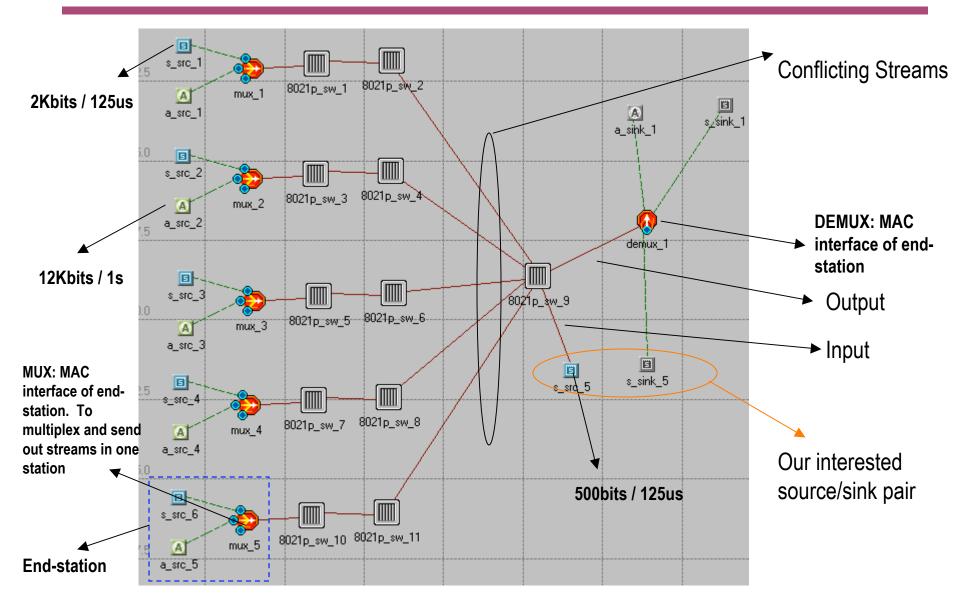
# On the worst case, and pacing

Felix Feifei Feng feng.fei@samsung.com

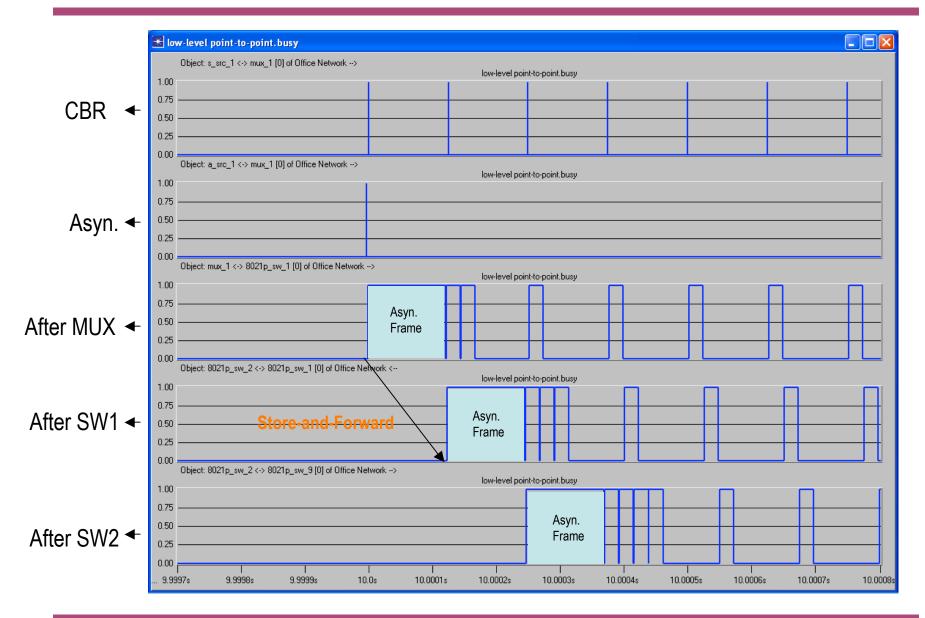
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### Scenario 1: Using 802.1p switches

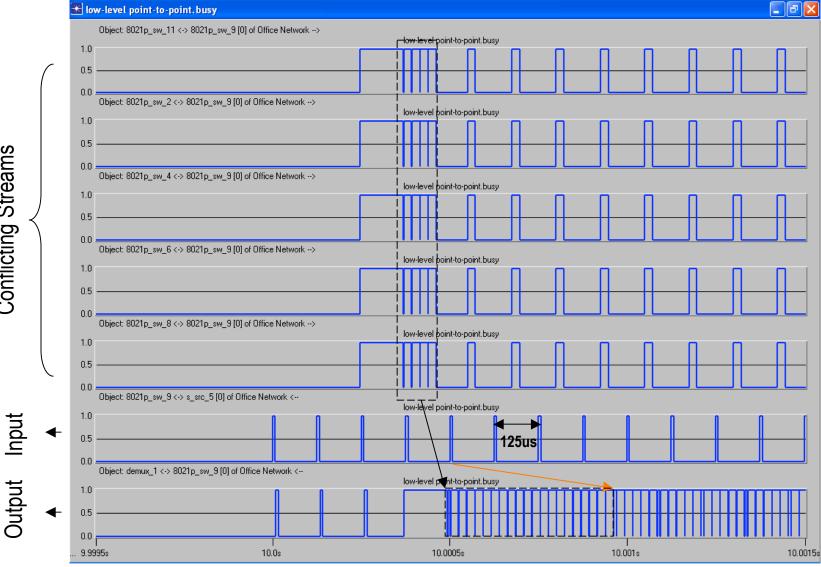


## **Distortion of CBR traffics**



## **Contention**

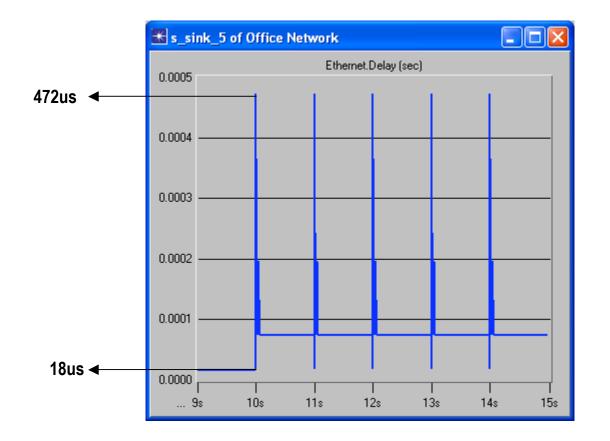
🛣 low-level point-to-point.busy



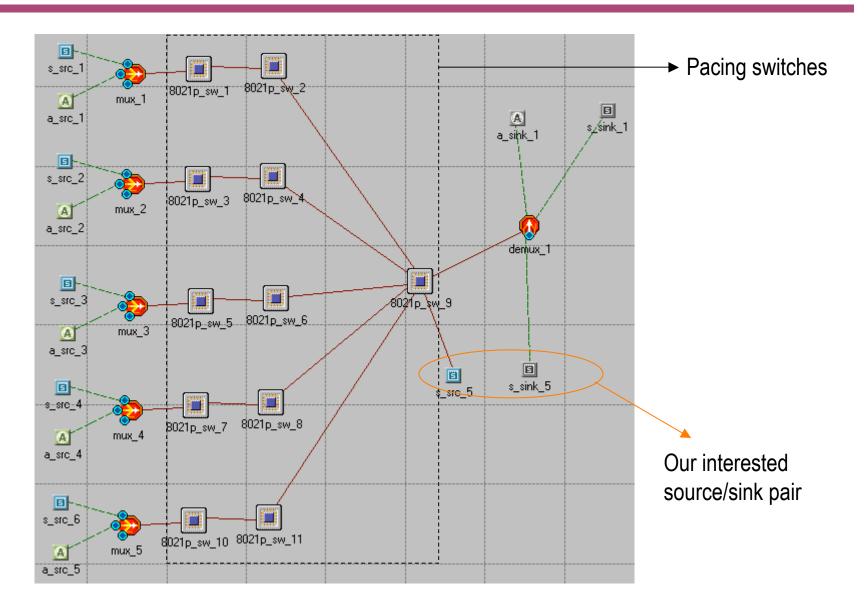
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Conflicting Streams

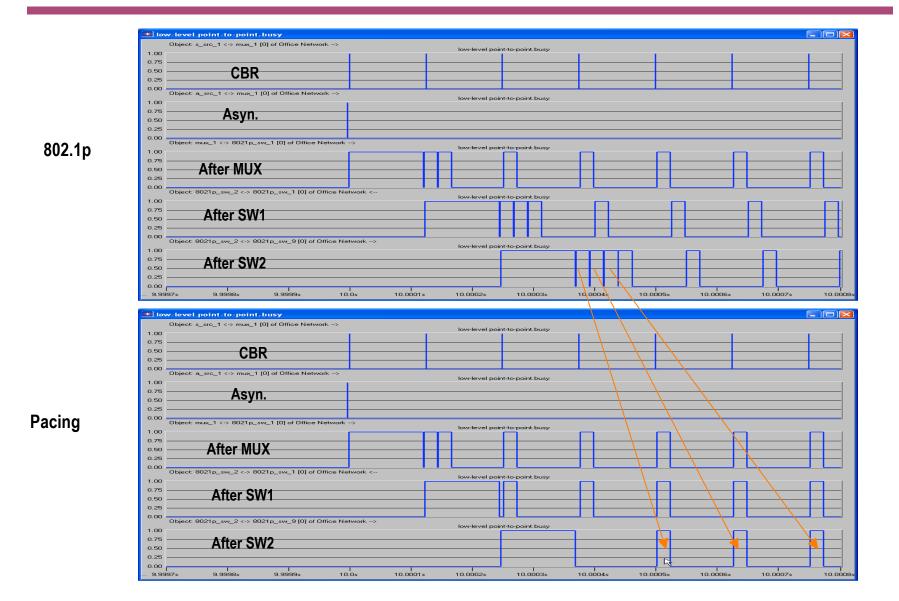
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### Scenario 2: Using Pacing-based Switches



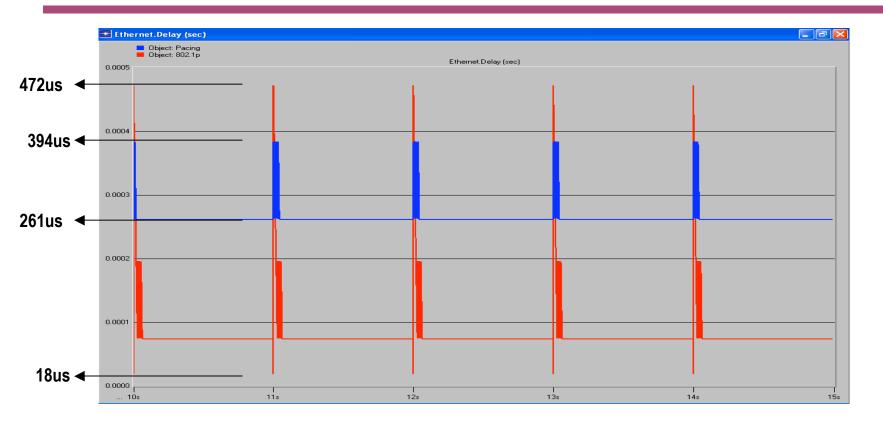
### Pacing Avoids the Traffic Distortion



# Pacing

#### 🛣 low-level point-to-point.busy Object: 8021p\_sw\_11 <-> 8021p\_sw\_9 [0] of Office Network --> low-level point-to-point.busy 1.0 0.5 — 0.0 💻 Object: 8021p\_sw\_2 <-> 8021p\_sw\_9 [0] of Office Network --> low-level point-to-point.busy 1.0 0.5 -----low-level point-to-point.busy 1.0 0.5 — 0.0 Object: 8021p\_sw\_6 <-> 8021p\_sw\_9 [0] of Office Network --> low-level point-to-point.busy 1.0 0.5 — low-level point-to-point.busy 1.0 -0.5 —— 0.0 🚃 Object: 8021p\_sw\_9 <-> s\_src\_5 [0] of Office Network <-low-level point-to-point.busy 1.0 0.5 — low-level point-te-point.busy 1.0 250us 0.5 ----0.0 9.9995s 10.0s 10.0005s 10.001s 10.0015s

# **Comparison of Delay Results**



In this case, pacing can decrease both the delay and the delay variation. But the difference is margina I.

### However...

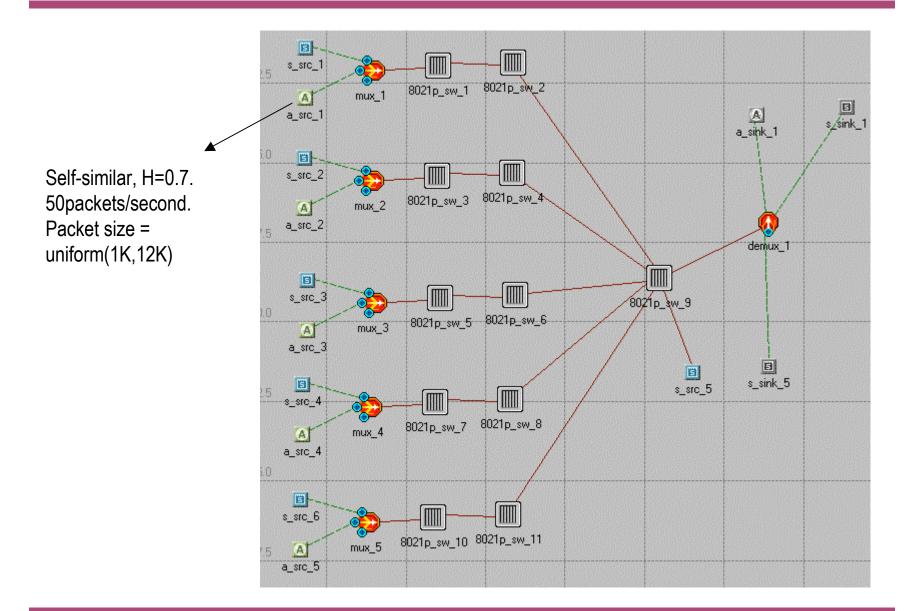
### This scenario is just an artificial case

- •We made all five conflicting CBR streams centralized. So when they come to switch-9, all five streams get bunched bursts and conflict with our interested stream at the same time. We made this case by:
  - Issuing a maximum size conflicting asynchronous packet to each conflicting CBR stream at the same time.
  - •The conflicting CBR stream traverses through several store-and-forward switches, which makes the CBR packets bunched together
  - •The link is almost fully loaded.

#### In a realistic network

This kind of bunched conflicting may rarely occur

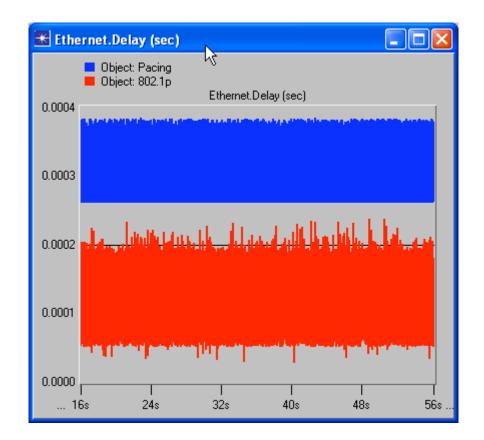
### Scenario 3: A 'More Realistic' Scenario



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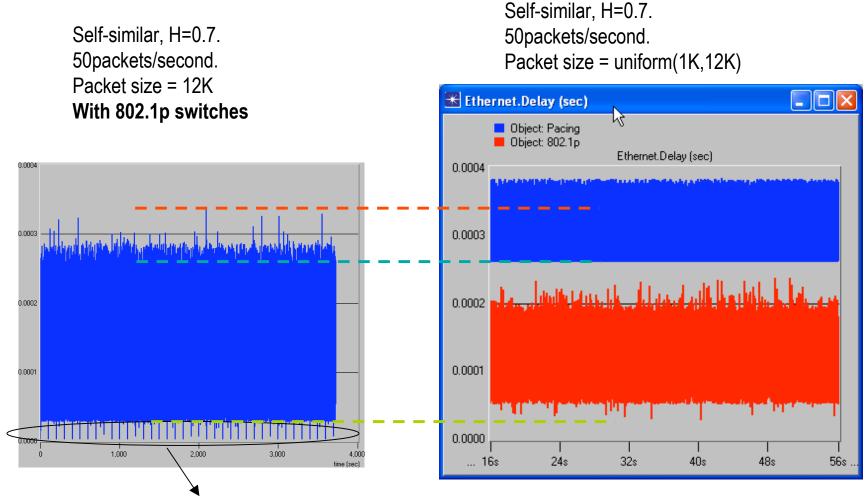
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# **Delay Results**



In this scenario, over 60 seconds, we didn't notice the occurrence of the worst case. Here pacing scheme has larger delay but smaller delay variation.

### Additional tests: MTU Packets, Self-Similar Arrival

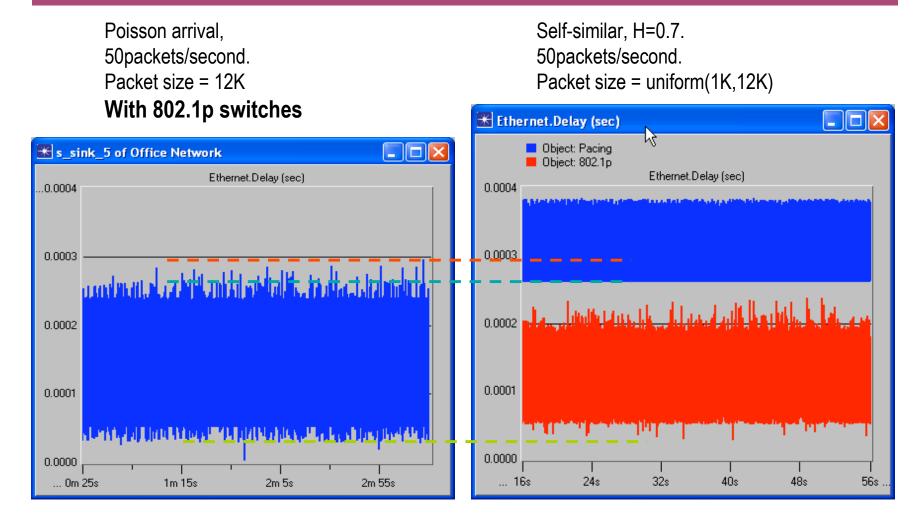


These small delays are because that I make sink nodes broadcast their addresses every 100 seconds.

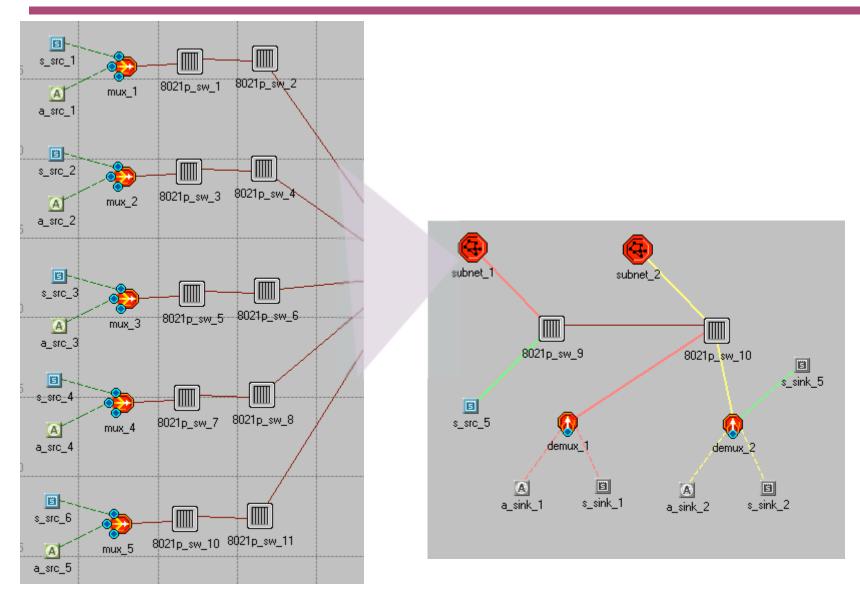
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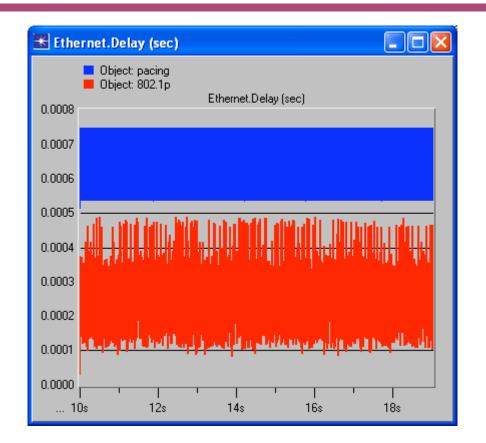
# Additional tests: MTU Packets, Poisson Arrival



### Scenario 4: Multi-hop Scenario



# **Delay Results**



Within the observation period, pacing scheme still has larger delay but smaller delay variation.

- •With pacing method, jitter is not accumulated along the multi-hop path.
- •With 802.1p method, jitter could be accumulated along the multi-hop path. But as long as timing synchronization is implemented, this jitter can be removed using reasonable size of buffer.

# Summary

□In a worst case, 802.1p has larger delay and delay variation than pacing scheme.

- The worst case delay of 802.1p is related to the number of incoming conflicting streams, the extent of distortion of those conflicting streams, and the cycle size of CBR traffic
- In our simulated scenarios, the performance differences are marginal

□In a realistic scenario, 802.1p shows a smaller delay, but larger delay variation (jitter).

 As long as timing synchronization is implemented, this jitter can be removed using reasonable size of buffer. Note that pacing scheme do need some kinds of timing synchronization.