

Design alternatives of Simple Reservation Protocol

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Design Choice: Number of Signaling Pass

□ Two-pass signaling in RSVP

- An important function of PATH messages is to set up the reverse path for forwarding RESV messages, since in IP networks the routing can be asymmetric.
- Other functions of PATH messages include non-RSVP cloud marking and gathering end-to-end QoS information.

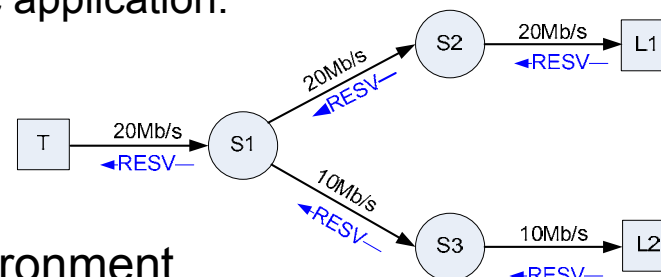
□ ResE

- In ResE, all nodes share a same view of spanning tree. The routing is symmetric.
- However, one-pass scheme has deficiency:
 - A reservation message will be flooded if its DA aren't existing in bridge's FDB.
 - Flooding will cause inefficient resource usage and misleading reservation results.
- Three-pass scheme
 - Listener would “ping” the talker before reservation.
 - Talker responses to listener ping messages.
 - Bridges along the path will learn the talker's address.
 - This scheme incurs additional round-trip delay

Design Choice: Reservation Initiator

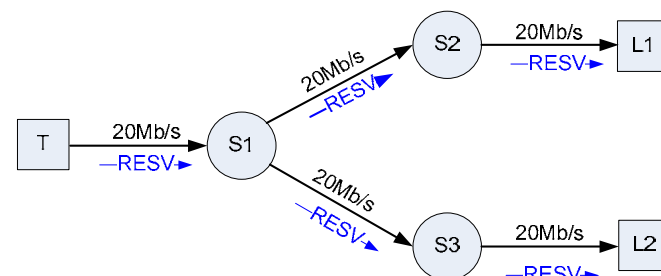
❑ Listener imitated

- Essential for heterogeneous listener environment
 - However, heterogeneous reservation is not useful unless the switches can intelligently discard frames for each specific application.



❑ Talker initiated

- Applicable for homogeneous listener environment
 - It would significantly reduce the complexity in signaling
 - Messages follow the data path
 - » But routing for the data path should exist before reservation.
 - No need for merge of reservation messages from different receivers



Design Choice: Reservation Policy

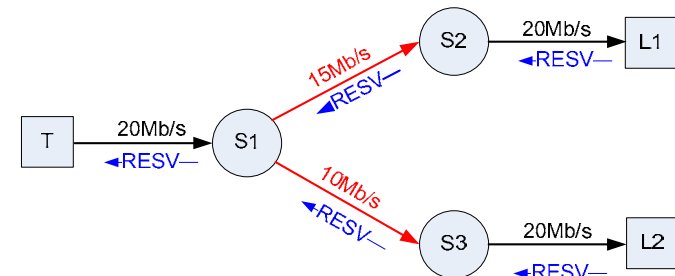
❑ Strictly guaranteed reservation

- If the admission control failed with a receiver's resource requirement somewhere, then the request will be denied. The receiver should retry his reservation.

❑ Flexible reservation

- Even the admission control failed with a receiver's resource requirement somewhere, SRP would first reserve the maximum allowable resource locally, and then continue with

- the original resource requirement (Greedy style)
 - Signaling is simple, but may incur more wasted resources
- updated resource requirement (Tight style)
 - With receiver initiated scheme, this scheme will cause “virtually heterogeneous” reservation.



- Data frames with flexible reservation should be marked with special flag. Intermediate bridges use this flag to discard frames that exceed the committed bandwidth.
- Flexible reservations may be updated to strictly guaranteed reservations upon reservation refreshing.
- Should provide notifications to related receivers about their flexible reservations.

Example 1: Three-Pass Listener Initiated SRP

□ Three-pass signaling

- Receiver first sends a “ping” message to desirable talker.
- Talker responds to listener ping message
 - Talker address will be learnt by intermediate bridges
- Listener sends reservation message to talker upon receiving talker ping response
- Listener sends teardown message to talker to release reservation

□ Listener initiated reservation

- Listener sends reservation message to talker upon receiving talker ping response.
 - Reservation messages should never be flooded
 - Reservation messages are merged in the intermediate nodes. A reservation request message travels only as far as the first bridge at which another receiver's reservation already exists for the same session.

Example 1: Three-Pass Listener Initiated SRP (cont.)

- Bridges execute admission control on the port receiving reservation messages
 - Set up reservation state if the admission control is successful
 - Send back Error message if the admission control is failed
- Each node periodically refreshes its reservation.
 - Otherwise, reservation states will age out and send Error messages downstream

□Implicit ACK and explicit NACK

- Application data frames are used as implicit acknowledgement
- Error messages are used as explicit negative acknowledgement

□Flexible reservation

- Can be achieved using greedy reservation scheme

Example 2: Two-Pass Talker Initiated SRP

❑ Two-pass signaling: use GMRP as the first pass

- Listeners use GMRP registration to join the isochronous multicast group.
- Listeners use GMRP deregistration to leave the isochronous multicast group. There is no need for dedicate teardown messages.
 - Upon a GID_Leave.indication event for an isochronous group address on a given port, corresponding reservation states and reserved resources (if any) will be released.

❑ Talker initiated reservation: only one type of message is needed

- Talker sends reservation message towards listeners in the isochronous group.
 - Carry a flag to indicate whether the resource reservation is successful upstream.
- Reservation messages traverse end-to-end even admission control is failed.

❑ For each incoming reservation message, intermediate bridge will:

- Get the outbound ports list for this message from group registration entries
- Execute admission control on each outbound port, and update the reservation states in bridge
 - Update the flag according to the execution result of admission control

Example 2: Two-Pass Talker Initiated SRP (cont.)

- Upon a GID_join.indication event for an isochronous group address on a given port, intermediate bridge will send a reservation message out of this port if corresponding reservation state has been set up in this bridge.
 - This will disassociate the set-up delay from the talkers' refresh timer value
- Talker periodically refreshes its reservation.
 - Otherwise, reservation states will age out and reservation messages with failure flag will be sent downstream

□ Explicit ACK and NACK

- When a listener receives the reservation message, it knows whether the reservation is successful according the corresponding flag.

□ Flexible reservation

- Reservation messages carry the “bottlenecked” bandwidth instead of requested bandwidth.
- Reservation messages serve as a notification to receivers about the flexible reservation.

Comparison

□ Example 1: Three-Pass Listener Initiated SRP

▪ Pros:

- Admission control would be done on Resv message receiving port. No need for FDB retrieval.
- Error messages clean up unnecessary reservation automatically.
- Extendable for heterogeneous reservation scenarios

▪ Cons:

- No explicit ACK
- Three-pass signaling incurs additional set-up delay
- Several new messages and processing method should be defined:
 - Ping, PingResponse, Resv, Error, Teardown; Most are new to Ethernet.

Comparison (cont.)

□ Example 2: Two-Pass Talker Initiated SRP

▪ Pros:

- Explicit ACK/NACK
- Only one new message will be defined. Signaling processing is simpler.
 - Resv; GMRP is reutilized.
- Listeners can be non-ResE devices in some simplified scenarios.
- Be natural for realizing flexible reservation.

▪ Cons:

- Admission control should be done on Resv message outbound port.
- Upon failed reservation, listeners should explicit deregister to avoid possible wasted reservation.

Thank you!