LL-FC: Problem(s) Identification

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Goal: Prepare for Monterey Interim

- 1. Set objectives of an ideal LL-FC
- 2. Identify the issues re. PAUSE (deviations from 1)
 - agree on terminology
- 3. Select those problems that we can solve in '07
- 4. Later: Solution candidates
- 5. Set the LL-FC agenda for Monterey

Generic Objectives of LL-FC

- I. Correct by design:
 - > Doesn't encourage 'simple' and faulty implementations.
 - > Doesn't lead to inter-operability faults thru standard misinterpretation by vendor.
- Correctness defines the semantics of
 - > (a) lossy-lossless operation,
 - > (b) priority level / QoS support,
 - > (c) deadlock freedom.
- II. Efficiency options: Enables high performance, Bw- and power-efficient operation. Allows vendor differentiation thru options for
 - ✓ (a) e.g. load balancing/AR, reliable delivery;
 - ✓ (b) options against, e.g., HOL-blocking, hogging, transient HS congestion, persistent HS congestion.
- III. 3-way Compatibility:

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- > Backwards w/ legacy Ethernet; w/ established IETF protocols (IP, TCP, etc.).
- > Upwards w/ storage and cluster transports RDMA, MPI, iSCSI etc.
- > Forwards w/ new apps that may directly use the native capabilities of LL-FC.

Improved PAUSE: A Hollow Strawman

- 10GigE is a discontinuity in the Ethernet evolution
 - opportunity to address new needs and markets
 - however, improvements are needed
- Requirements of next-generation PAUSE
 - 1. Correct by design, not implementation
 - 1. Deadlock-free
 - No HOL₁- and, possibly reduced HOL₂-blocking
 Note: Do not try to address high-order HOL-blocking at link layer
 - 2. Configurable for both lossy and lossless operation
 - 3. QoS / 802.1p support
 - 4. Enables virtualization / 802.1q
 - 5. Beneficial or neutral to CM schemes (BCN, TCP, ...)
 - 6. Legacy PAUSE-compatible
 - 7. Simple to understand and implement by designers
 - 1. Min. no. of flow control domains: h/w queues and IDs in Ether-frame
 - 8. Compelling to use => always enabled...!

Boilerplate: Principles of LL-FC. Orthogonality and 5D Control

- LL-FC shall orthogonalize the following 4 dimensions
- Correctness

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- 1. C1: lossy lossless operation
- 2. C2: deadlock prevention and recovery
- 3. C3: priority classes / QoS service levels
- ➢ Performance
 - 4. P1: low-order HOL-blocking, resource hogging, transient congestion
 - 5. [P2: high-order HOL-blocking, persistent congestion (realm of CM).]

• If the solutions to the above 5 concerns are mutually exclusive (can not coexist in implementation or are not simultaneously operational), any such limitation/constrain will be explicitly stated, including the consequences thereof.

Deadlock Taxonomy: Deadlocks Possible in a Datacenter Interconnects

- The following types of dlocks may affect an Ethernet system:
 - > 1. (DLK1) circular dependency
 - \checkmark a) memory-2-memory circular dependency (inter-switch, CD = 1st order deadlock)
 - ✓ b) Load/Store (Rq/Reply) circular dependecy (transaction-induced deadlock)
 - > 2. (DLK2) priority blocking (PB=2nd order, improperly called deadlocks);
 - > 3. (DLK3) routing loop (RL=3rd order)
- Specific deadlock cases will be illustrated in Monterey

PAUSE Issues

- PAUSE-related issues interfere with BCN simulations
- Correctness
 - Deadlocks (some of them...)
 - \checkmark cycles in the routing graph (if multipath adaptivity is enabled)
 - multiple solutions exist
 - ✓ circular dependencies (in bidir fabrics)
 - > BCN can't help this => Solutions required
- Performance (to be elaborated in a future report)
 - Iow-order HOL-blocking and memory hogging
 - ✓ Non-selective PAUSE causes hogging, i.e., monopolization of common resources: e.g. shared memory may be monopolized by frames for a congested port (as shown <u>here</u>)
 - ✓ Consequences
 - best: reduced throughput
 - worst: unfairness, starvation, saturation tree, collapse
 - \checkmark properly tuned, BCN can address this problem