

802.1aj Two port MAC Relay status

ADVANCE

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Two port MAC Relay

- ▶ Industry recognises that full 802.1 bridges are sometimes unnecessarily complex
- ▶ TPMR (802.1aj) attempts to provide a simpler relay function than a VLAN bridge
- ▶ PAR granted December 2004
 - ▶ Initial draft 0.0 May 2005
 - ▶ Draft 1.0 July 2005
 - ▶ Draft 1.1 August 2005
 - ▶ Draft 1.2 November 2005
 - ▶ Draft 1.3 May 2006
 - ▶ Draft 1.4 June 2006
 - ▶ Draft 2.0 January 2007
 - ▶ Draft 2.1 May 2007
 - ▶ Draft 2.2 October 2007 (Working group ballot)
- ▶ Interest from other standards bodies including MEF and DSL Forum
- ▶ This presentation represents a personal view of the status following the November 2007 meeting of 802.1 and of how a TPMR might be used.



TPMR topics

- ▶ Zero configuration option – should work out-of-the-box
- ▶ Topologies
- ▶ Link maintenance
- ▶ Discovery
- ▶ Management
- ▶ A TPMR is a bridge
- ▶ Forwarding
- ▶ MAC types
- ▶ Loopback
- ▶ Link status propagation



Topologies

- ▶ TPMRs can be deployed singly, or in a chain
- ▶ A typical application might be as a demarcation device (NID)
- ▶ A TPMR has exactly two ports
 - ▶ Each port can be Ethernet or any MAC or emulated MAC which supports the 802.1 Internal Sublayer Service
- ▶ Protection is not supported in the draft standard
- ▶ Management using SNMP over Ethernet
 - ▶ Envisaged to be from an intelligent device which proxies the TPMR's managed objects into its own MIB



Link maintenance (not in 802.1aj)

- ▶ For Ethernet links, 802.3ah EFM OAM may be employed
 - ▶ This provides an indication of link up/down
 - ▶ Ethernet MAC link down indication is notoriously unreliable
- ▶ 802.3ah also provides
 - ▶ Link status change information
 - ▶ Link statistics including errored seconds etc.
 - ▶ Managed object access, which is NOT used in TPMR
- ▶ E-LMI (MEF UNI Phase 2) was considered, but is not suited as a link maintenance protocol
 - ▶ Intended for CE to retrieve status and service attributes from the network
 - ▶ Includes UNI and per-EVC configuration and status information
- ▶ Other MACs and emulated MACs can use their own protocol



Discovery

- ▶ A mechanism is required to allow discovery of TPMRs, so that the managing device knows what to manage
- ▶ Mandatory CFM (802.1ag) is the primary discovery method
 - ▶ At least a level 0 MIP is required in TPMR
 - ▶ Attached bridge or station can use Linktrace to discover connectivity of attached TPMR chain
 - ▶ All TPMRs in a chain can be found, but a method is needed to know when the end of the chain has been reached
 - ▶ CFM tells you what kind of device it is (uncertain if TPMR defined)
- ▶ LLDP (802.1ab) may be used for further probing
 - ▶ LLDP support is optional
- ▶ Ethernet EFM OAM (802.3ah) could have been chosen for Ethernet links, but is harder to use for chain discovery.



Management

- ▶ SNMP over Ethernet, without IP, is mandatory
 - ▶ SNMP over IP was rejected because of the desire to avoid IP address management and NMS interaction with individual TPMRs
 - ▶ 802.3ah EFM/OAM was rejected because of concerns over scalability to a chain and lack of "Set" capability
 - ▶ CORBA was considered too much of a stretch given that nothing else in 802.1 uses it
- ▶ Management is required to be supported on at least one of the data ports on the TPMR
- ▶ SNMP over Ethernet is specified in RFC4789
 - ▶ Untagged frames are used
 - ▶ Management VLAN option not yet discussed much
- ▶ SNMP over other transports is not precluded
 - ▶ For example, traditional SNMP over IP is allowed



Management, continued

- ▶ Discovery is used first to find what to manage
- ▶ How remote management is done is not specified, but perhaps
 - ▶ Retrieved objects are incorporated into the managing device's MIB
 - ▶ Incorporation into Interface MIB objects is a possibility
- ▶ Which ports can be used to manage the device?
 - ▶ A management block is provided to prevent access from the customer port
 - ▶ This block can be turned on and off by management
 - ▶ Unspecified issues:
 - ▶ Is access provided by authentication, to allow a device which is installed the wrong way round to be "recovered" remotely?
 - ▶ In a device with different port types, which port then?



A TPMR is a kind of 802.1 Bridge

- ▶ Only two ports
- ▶ No MAC address learning
- ▶ No VLAN tagging, but can be priority aware
- ▶ No Spanning Tree
 - ▶ BPDUs require special treatment (see later slide)



Forwarding

- ▶ General idea is to be transparent to protocols the TPMR does not implement
 - ▶ But some protocols are filtered out by the MAC, e.g. Pause
 - ▶ Transparent to BPDUs
 - ▶ Transparent to LACP (despite the layering violation)
 - ▶ One reserved address will be terminated by the TPMR and used for the LAN Status Propagation Protocol
- ▶ No modification of user data frames (e.g., tagging)
- ▶ Multiple queues are optional
 - ▶ Extract priority from Q-tag and 802.1ae LinkSec tag
 - ▶ Only for integrity-protected frames – unencrypted
 - ▶ Recognise L2 control protocols and place in fastest queue
 - ▶ Typically BPDUs
- ▶ Otherwise like 802.1d/Q
 - ▶ Note that MRP (802.1ak) needs special handling in a Q-bridge



Loopback (not part of 802.1aj)

- ▶ Per-link loopback on Ethernet with 802.3ah EFM/OAM
 - ▶ Invoked by SNMP to previous hop
 - ▶ Beware that EFM OAM loopback discards returned frames
- ▶ Multi-hop CFM-based (802.1ag) loopback
 - ▶ Uses a special loopback frame
 - ▶ Can contain arbitrary data inside a TLV
 - ▶ Non-intrusive, in that user data continues
 - ▶ Stateless
 - ▶ Issues include VLAN-non-awareness
- ▶ Stateful per-VLAN loopback is not supported
 - ▶ CFM rejected this idea as not sufficiently useful at resolving data-driven errors
 - ▶ Could be provided using an EFM/OAM extension invoked from the previous hop using SNMP



LAN Status Propagation

- ▶ 802.1aj incorporates a link loss forwarding function called LAN Status Propagation
- ▶ Where TPMRs are placed in between existing bridges, they could interfere with protection and restoration processes
- ▶ LAN Status Propagation ensures that changes in connectivity are signalled to the bridges at the ends of a chain of TPMRs
- ▶ The signalling is compatible with existing 802.1 bridges

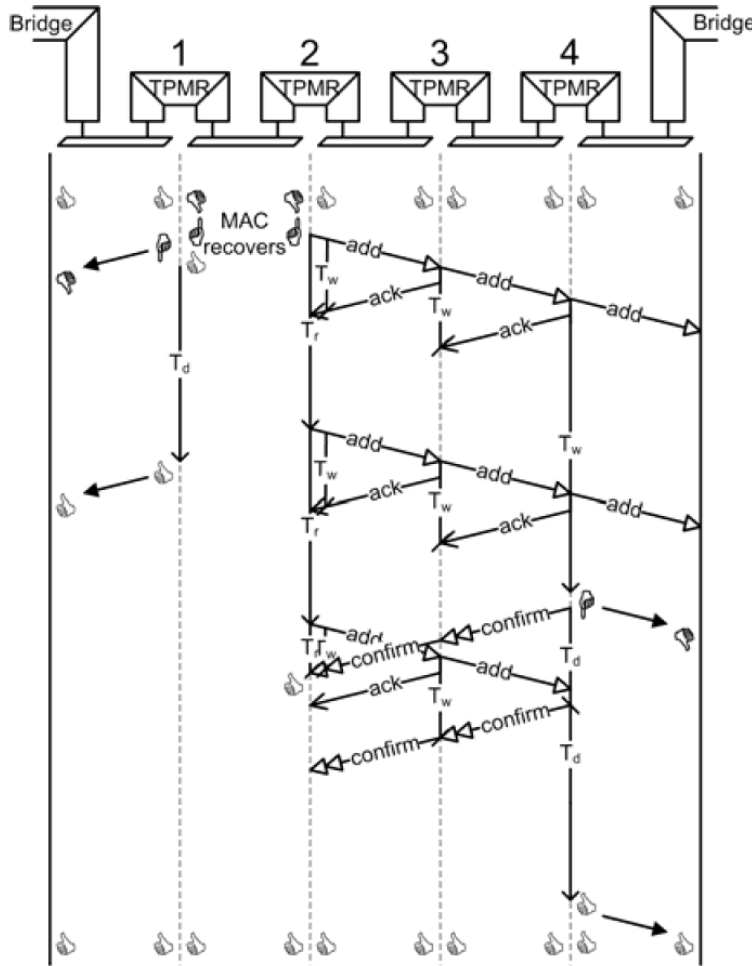


LAN Status Propagation Protocol

- ▶ The protocol is aimed at links over which RSTP may be running
 - ▶ It's important that when the connectivity is announced, it is available
- ▶ Another aim is to keep links up when possible, to allow management traffic to a maximal subset of TPMRs in a chain
- ▶ Basic idea:
 - ▶ For loss: send a loss message; if unacknowledged, blip the link;
 - ▶ For add: send an add message; if unacknowledged, blip the link;
 - ▶ Blipping the link alerts the next layer that connectivity has changed
- ▶ The protocol communicates only changes, not states
 - ▶ Link down is not periodically sent
- ▶ Optimisations are possible when a TPMR already knows that its neighbour **doesn't** speak the LAN Status Propagation Protocol (LSPP)



New Connectivity (general case)



Key



OperUp



OperDown but locally
MAC_Enabled



MAC_Enabled FALSE



MAC_Operational, but
reporting OperDown
to clients


Timers

- ▶ T_d = blip time (longish)
- ▶ T_r = retransmit time
- ▶ T_w = ACK wait time

Figure 23-5—New connectivity with MAC status notification



New Connectivity (general case)

- ▶ In the more general case, TPMRs don't know whether their link partners speak LSPP and must try it with timeouts
- ▶ To illustrate this, consider TPMR2 and its rightward neighbours
- ▶ MAC recovers in TPMR1 and TPMR2: MAC_Operational detected
- ▶ Both TPMRs enter  to avoid creating loops
- ▶ As TPMRs 2-4 don't know if their neighbours speak LSPP:
 - ▶ TPMR2 sends "add" to its right neighbour and starts T_r and T_w
 - ▶ TPMRs 3 and 4 forward the "add", ack leftward, and start T_w
 - ▶ Received acks cancel T_w in TPMRs 2 and 3
 - ▶ T_r may repetitively expire in TPMR2, triggering retransmission of "add"
 - ▶ T_w will expire in TPMR4, triggering rightwards "blip" and leftwards "confirm"
 - ▶ Confirm cancels T_r in TPMR2



Loss of Connectivity

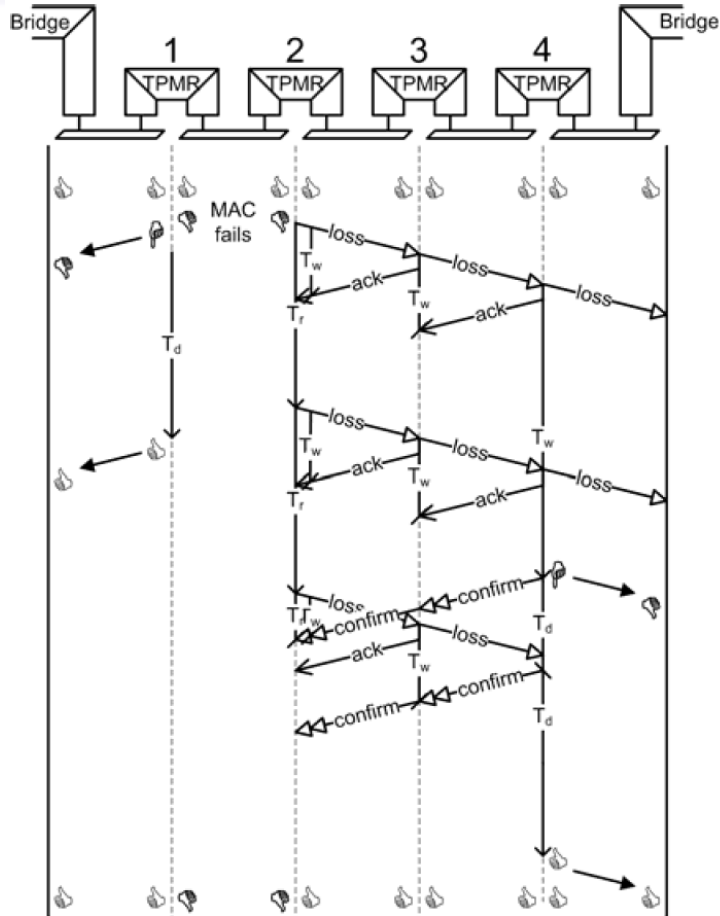


Figure 23-6—Connectivity failure

Key



OperUp



OperDown but locally
MAC_Enabled



MAC_Enabled FALSE




MAC_Operational, but
reporting OperDown
to clients

Timers

- ▶ T_d = blip time (longish)
- ▶ T_r = retransmit time
- ▶ T_w = ACK wait time



Loss of Connectivity (general case)

- ▶ In the more general case, TPMRs don't know whether their link partners speak LSPP and must try it with timeouts
- ▶ To illustrate this, consider TPMR2 and its rightward neighbours
- ▶ Link fails between TPMR1 and TPMR2: MAC_Operational FALSE
- ▶ Both TPMRs enter 
- ▶ As TPMRs 2-4 don't know if their neighbours speak LSPP:
 - ▶ TPMR2 sends "loss" to its right neighbour and starts T_r and T_w
 - ▶ TPMRs 3 and 4 forward the "loss", ack leftward, and start T_w
 - ▶ Received acks cancel T_w in TPMRs 2 and 3
 - ▶ T_r may repetitively expire in TPMR2, triggering retransmission of "loss"
 - ▶ T_w will expire in TPMR4, triggering rightwards "blip" and leftwards "confirm"
 - ▶ Confirm cancels T_r in TPMR2



Questions and Answers

- ▶ How does LAN state propagation protocol interact with CFM?
 - ▶ This question is particularly relevant if RSTP is not being used
 - ▶ LSPP tells you “something changed”; CFM on top confirms connectivity
 - ▶ CFM could be used underneath to control MAC_Operational
- ▶ At initialisation of the system, will the TPMR learn whether its neighbours speak LSPP?
 - ▶ When TPMR comes up, its MACs will come up
 - ▶ That should trigger “add” messages going out
 - ▶ If there are responses, then there’s an LSPP neighbour
 - ▶ If not, there isn’t
 - ▶ So as long as initialisation is carefully handled, the TPMR will know its neighbours’ LSPP capability
 - ▶ This means response to link loss can be immediate



Questions and Answers

- ▶ Can one assume that same link partner is present after blipping the link?
 - ▶ Even if the link partner changes, the protocol will still work
 - ▶ Once you learn your link partner doesn't speak the protocol, you won't learn they do if they are changed.



Use with telco-style protection

- ▶ LSPP is designed for RSTP-based protection systems
- ▶ It uses single events to signal a change
 - ▶ “loss” or “add” messages to partners speaking LSPP
 - ▶ “blipping” the link to partners who don’t speak LSPP
- ▶ These events can trigger protection switching
 - ▶ “blipping” will automatically cause switching
 - ▶ Receipt of LSPP “loss” messages could do so
- ▶ However nothing in LSPP marks a link as down
 - ▶ Other events might trigger a switch back to this bad link
 - ▶ Examples include a break on the other link or revertive switching back
 - ▶ Needs an additional mechanism to verify recovery (e.g. slow CFM)
 - ▶ Mark the link as bad until you know it’s good
- ▶ LSPP doesn’t have the high traffic load of fast CFM



Thank You

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