

**Effects of Pause and QCN on TCP Sources:
Part 1
(Redo of simulation with
slightly different topology)**

Berk Atikoglu, Balaji Prabhakar

Analyze the Effects of Pause and QCN on TCP Sources: System Parameters

- **Congestion Management Schemes**

- TCP Only
- TCP + PAUSE
- TCP + QCN + PAUSE

- **Switch Parameters**

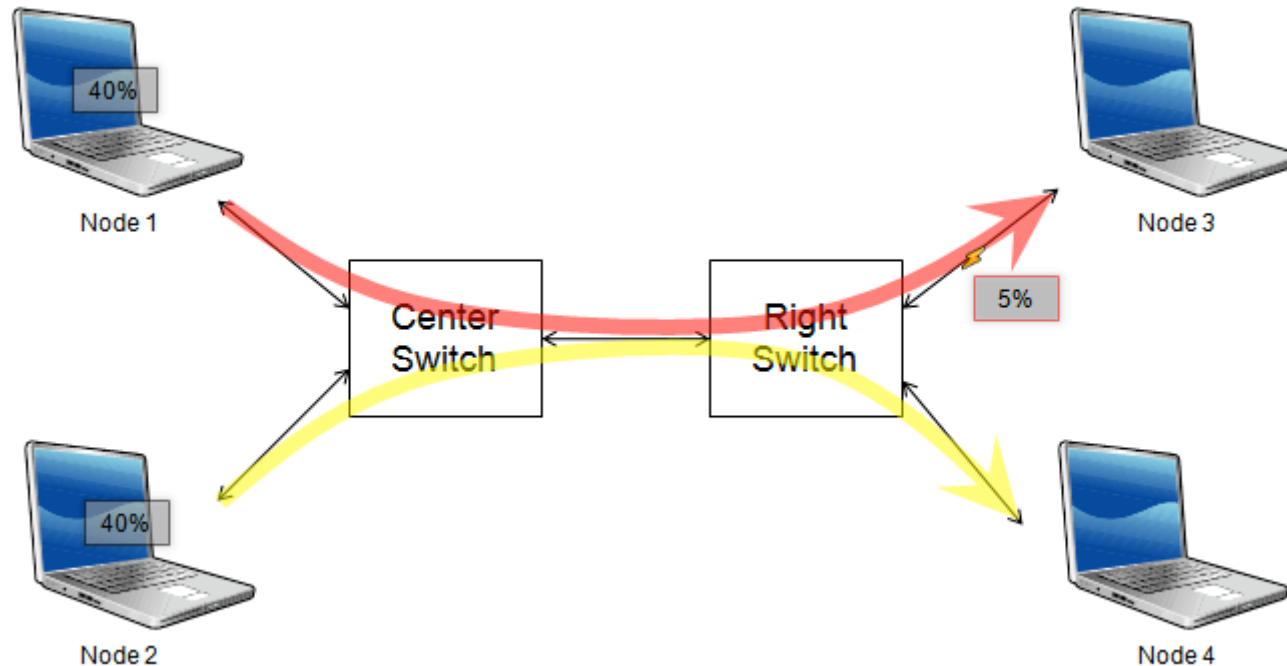
- PAUSE Disabled
 - Output queue limit of 150kbytes
- PAUSE Enabled
- No output queue limit
 - Applied on a per input basis based on watermarks
 - Watermark_hi = 130kbytes
 - Watermark_lo = 110kbytes

- **QCN Parameters**

- $W = 2.0$
- $Q_EQ = 26\text{kbytes}$
- $Gd = 1/128 = 0.0078125$
- Base marking: once every 150kbytes
- Jitter on marking: 30%
- Runit = 1Mb/s
- MIN_RATE = 10Mb/s
- BC_LIMIT = 150kbytes
- TIMER_PERIOD = 15ms
- R_AI = 5Mbps
- R_HAI = 50Mbps
- FAST_RECOVERY_TH = 5
- Quantized_Fb: 6 bits
- Jitter at RP: 30% (byte counter and timer)

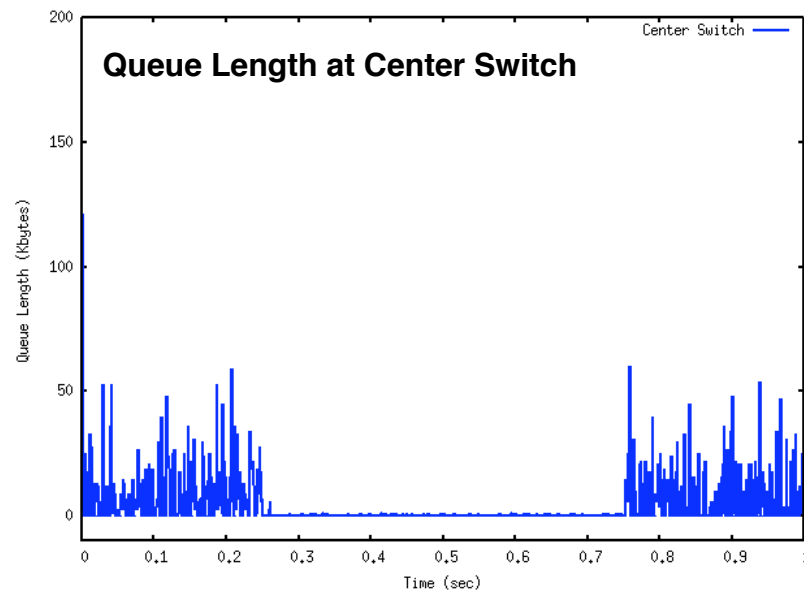
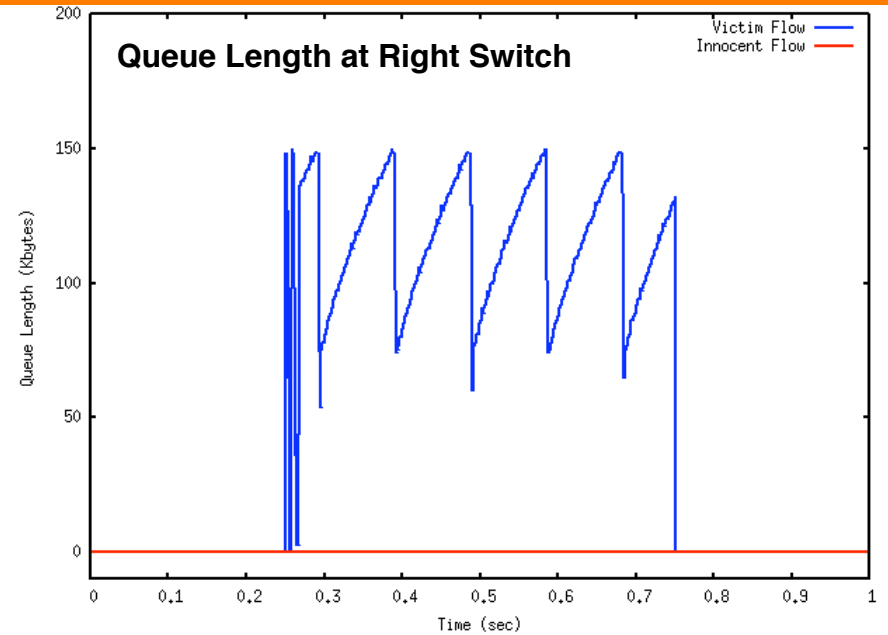
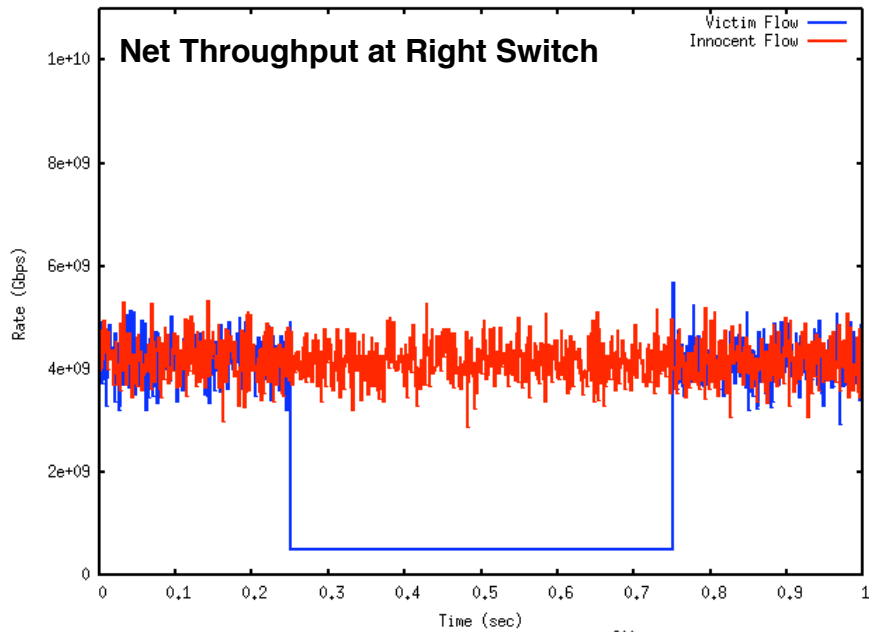
- **TCP Version → NewReno**

Topology and Workload

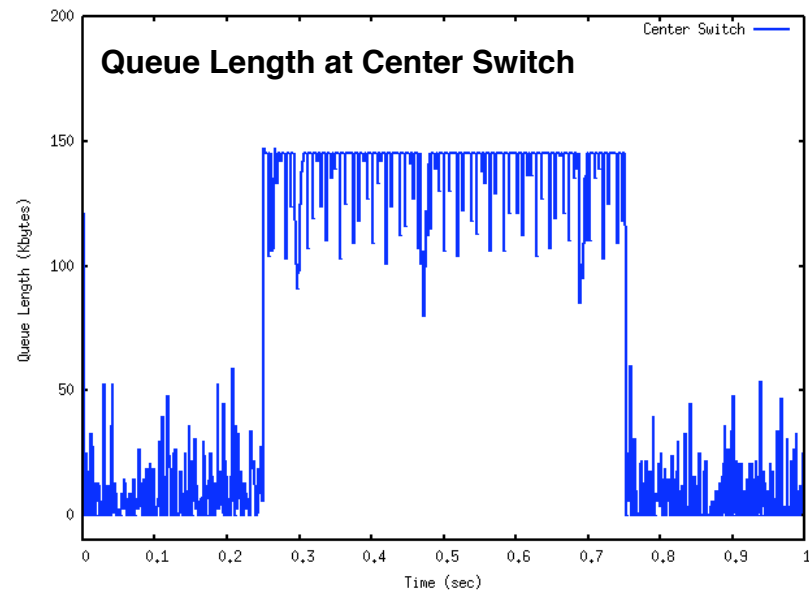
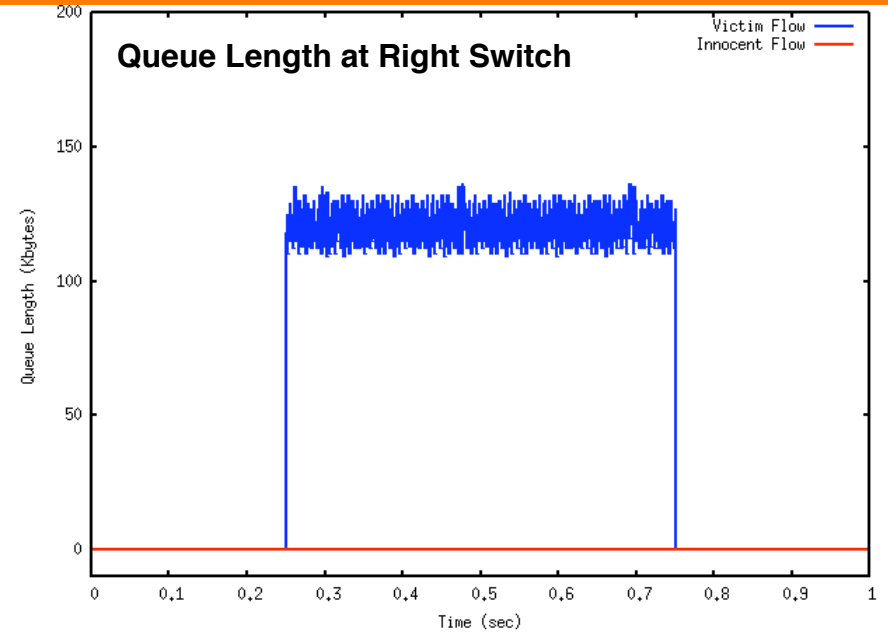
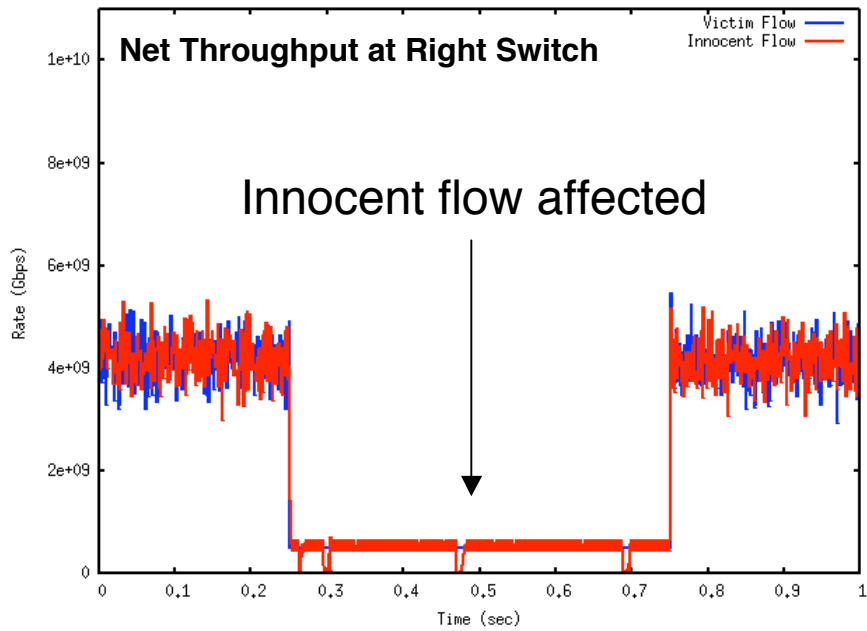


- Topology
 - Link Speed → 10Gbps for all links
 - Loop Latency → 18us
- Traffic Pattern
 - 9k byte transactions arriving with a Bernoulli distribution
 - Node 1 sends to Node 3 at 4Gbps (40%)
 - Node 2 sends data to Node 4 at 4Gbps (40%)
- Congestion Scenario
 - Node 3 temporarily reduces its service rate from 10Gbps to 500Mbps between [250-750ms]

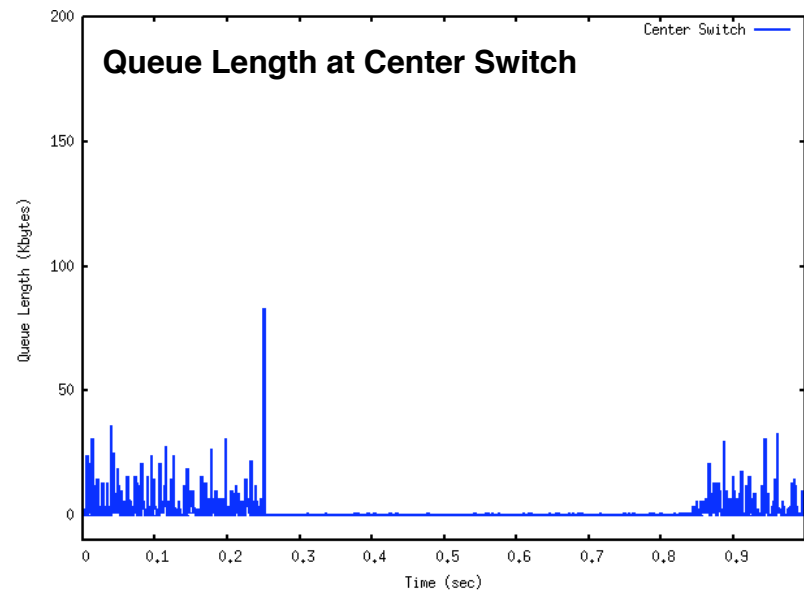
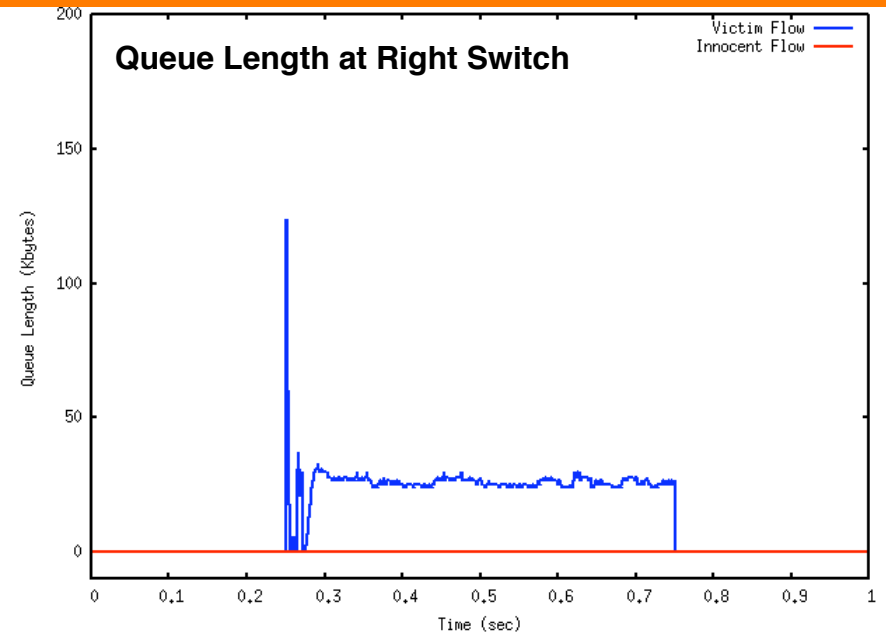
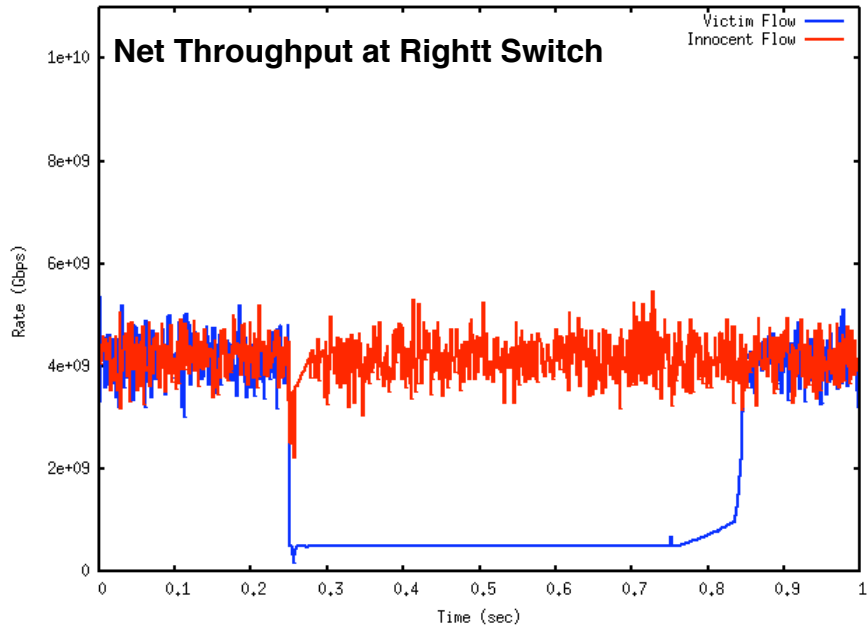
TCP Only



TCP and PAUSE



TCP, QCN and PAUSE



Interaction of QCN and TCP

Part 2

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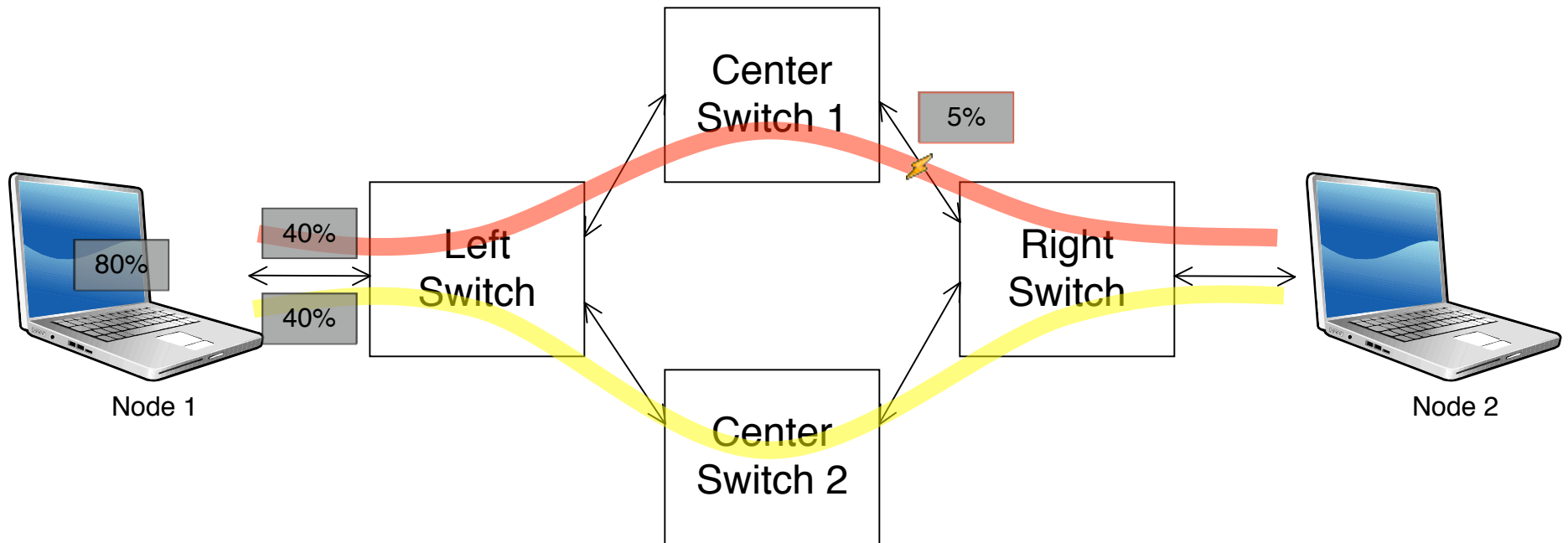
Overview

- In Part 1 we saw how QCN helps TCP discover the correct bottleneck bandwidth
 - In this case each TCP flow had its own RL
- In Part 2 (this presentation) we consider the case where RLs are shared among TCP flows
- Potential problem with the shared rate limiters
 - When there is a multipath scenario, a shared RL causes an innocent flow to lose throughput
- Proposed solution to the problem
 - A TCP flow can adjust its sending rate if it is aware of the congestion at Layer 2
 - This can be conveyed to it via a flowid
- Simulation results

QCN parameters

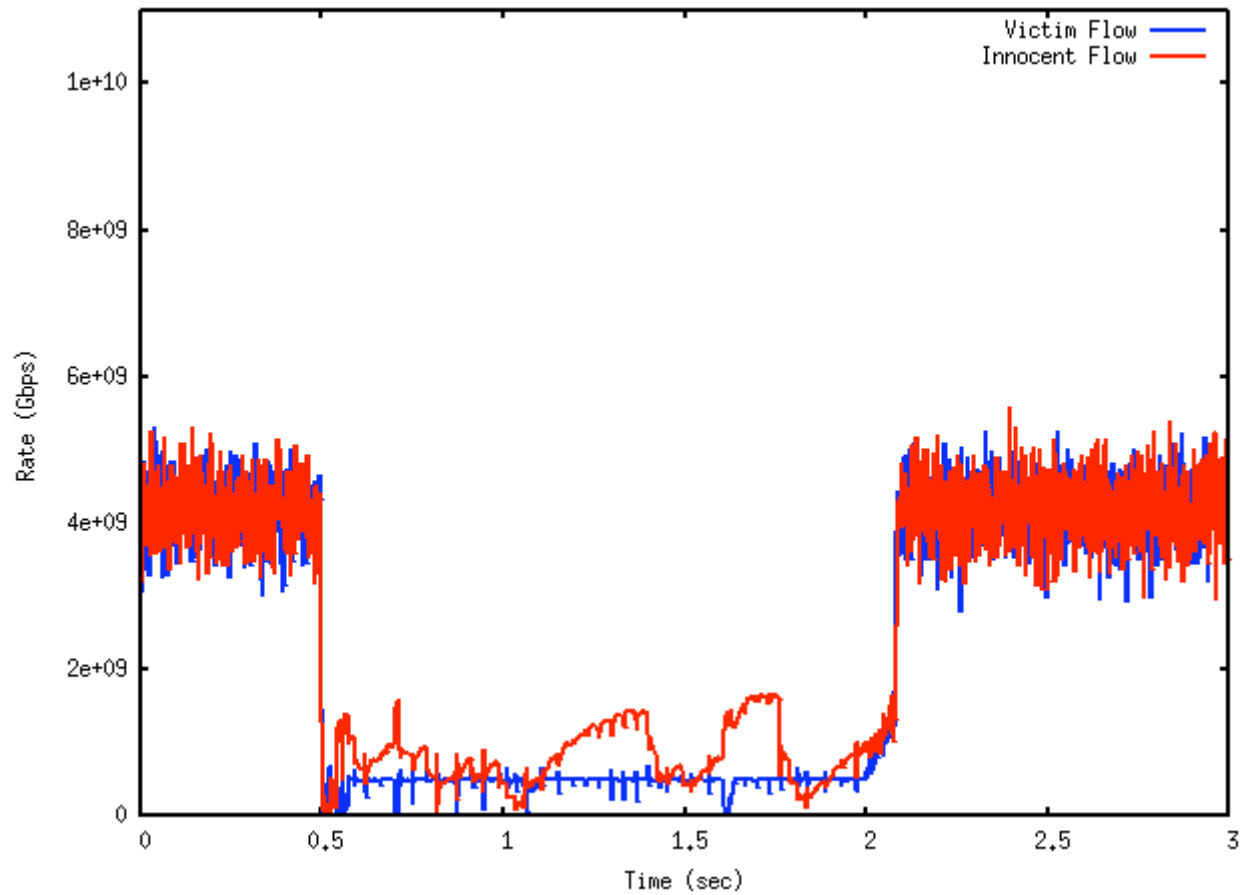
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- Base marking: once every 150kbytes
- Jitter on marking: 30%
- $MIN_RATE = 10$ Mb/s
- $BC_LIMIT = 150$ kbytes
- $TIMER_PERIOD = 15$ ms
- $R_AI = 5$ Mbps
- $R_HAI = 50$ Mbps
- $FAST_RECOVERY_TH = 5$
- Quantized_Fb: 6 bits
- Jitter at RP: 30% (byte counter and timer)

Topology and Workload



- Topology
 - Link Speed → 10Gbps for all links
 - Loop Latency → 200us
- Traffic Pattern
 - 9k byte transactions arriving with a Bernoulli distribution
 - Two TCP flows at Node 1 generate a total traffic of 8Gbps (80%)
 - Flow 1 sends data to Node 2 at 4Gbps (40%) using the path LS -> CS1 -> RS
 - Flow 2 sends data to Node 2 at 4Gbps (40%) using the path LS -> CS2 -> RS
- Congestion Scenario
 - CS1 temporarily reduces its service rate from 10Gbps to 500Mbps between [500ms-2s]

Throughput of the Flows (QCN and TCP)



Problem

- Ideally victim flow should get 500Mbps and innocent flow should get 4Gbps
- Shared RL causes innocent flow to also get around 500Mbps

Proposed solution

- If victim TCP flow was aware of the congestion, it could adjust its rate and innocent flow need not be affected
- RL is aware of the congestion
- RL can notify victim TCP flow
 - Drop a packet of the victim flow or,
 - Set the ECN bit of a packet of the victim flow
- In our scheme, we drop a packet of the victim flow

Simulation Results

