

# Proper Layering for the TSN Layer 2 Data Plane

Norman Finn Version 2

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# Moving ahead TSN L2/L3 work

- At the AVnu face-to-face meeting in December, Norm Finn made presentations on the subject of ensuring that the IEEE 802.1 Time-Sensitive Networking (TSN) efforts are compatible with the established body of work from Internet Engineering Task Force (IETF) on the Internet Protocol (IP).
- A four-step plan for advancing the TSN work was presented.

# Moving ahead TSN L2/L3 work

- A. Pick at least one data plane model for joint L2/L3 stream identification.
- B. Pick at least one data plane model for joint L2/L3 duplicate packet deletion.
- C. Pick at least one data plane model for the dual-homed end station.
- D. Pick exactly one (hopefully) suite of protocols to fit the data plane choices made.

# This presentation

- This is <u>tsn-nfinn-L2-Data-Plane-0114-v02</u>. It offers an improved model for layering the AVB/ TSN concepts, which leads to a set of choices for answering the data plane questions A, B, and C. Contents of this presentation:
  - What's broken about AVB/TSN?
  - What's the Fix?
  - Alternative TSN Encapsulations

Summary

• This presentation does not make choices, and so cannot advance to Question D, protocols.

# What is Broken?



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#### What's broken?

- This author perceives a **layering problem** in the current AVB/TSN standards suite that is the source of a number of difficulties:
  - Integration/reconciliation with the way the IP layer uses Ethernet, especially hosts that use VLANs for separate addressing or broadcast domains.
  - Meeting the need for carrying unicast streams and the need for multiple fixed paths through the network, in a manner that does not confuse MAC address learning or waste VLAN ID space.

## What's broken

- Without detracting from its achievements, especially plug-and-play, AVB remains a solution for a few very specific communities, where the vendors and users, together:
  - Have complete control over the application.
    Have complete control over the host stack.
    Have complete control over the network.
- Fortunately, it's not hard to fix.
- And once fixed, we have less work to do.

# Complete control over the application

- How do existing applications address each other?
  - >Universal Resource Locator (HTTP).
  - >Unicast and multicast IPv4 and IPv6 address.
  - ≻Individual and Group MAC address.
  - Locally-administered MAC address.
  - >VLAN, often one VLAN = one IP subnet.
- How must AVB applications address each other?

➢ Group MAC address on a particular VLAN.

# Complete control over the host stack

- How do existing higher layers obtain MAC addresses?
  - Domain Name Servers are used to map host addresses to IP (or other!) addresses.
  - >IPv4 stacks issue ARP requests/replies.
  - >IPv6 stacks use Neighbor Discovery.
  - Fixed mapping of IPv4 multicast to MAC addresses.
- How must AVB stacks obtain MAC addresses?

>Override the stack to use a 1722 address.

#### Complete control over the network

- How are **existing** networks designed?
  - Bridges provide "who" addresses, and learn geography, thus encouraging broadcast plugand-play protocols like Bonjour and ARP.
  - Network administrators break up the network into bridged domains, separated by routers, to prevent those broadcasts from overwhelming the network.
- How must AVB networks be designed?
   Everything must be one flat bridge domain.
   10,000 stations, 5000 streams in a stadium?!?!

## :You can make it work

- Demonstrably, there is enough value in AVB for communities of developers to create applications, host stacks, and networks that operate to the AVB specification.
- But, we can make TSN applicable to users who already have their networks defined.

# **Driving assumption**

- The goal of the TSN TG should be to write standards for new Quality of Service (QoS) classes for high reliability and low latency, that offer incremental benefit to any network, whether L2, L3, or mixed, that follows established general-purpose operational models.
- To the extent that TSN standards require variations from those models, their adoption will be hindered.

- What is the industry standard term for standards like, "Stream Reservation"?
- Connection-oriented services (or circuits) over connectionless networks.
  - When you have per-flow state in every node along the path, and rebuild the per-flow state when the path changes, that's a "connection" or a "circuit".

- How does the industry set up a circuit from A to B over a connectionless network?
- A gives the network address of B to a protocol that sets up a connection, and that protocol returns to A (and B) an encapsulation, with a connection ID, to use for the connection's packets.

- How does the packet get from source to destination?
- A encapsulates the packet in a form that:
  The network will get to B; and
  Can be recognized by B, and by participating network nodes along the way, as a packet belonging to this particular connection.

- Why was Nigel Bragg saying that we need one VLAN per source per path through the network (worst case) for nailed-up paths?
- Because we didn't realize that we are tunneling circuits through the network, and thought we had to have the connectionless network addresses on the outside of the frame.

- What's wrong with the AVB frame format?
- Nothing!

- Then, what's the fix?
- Read on.

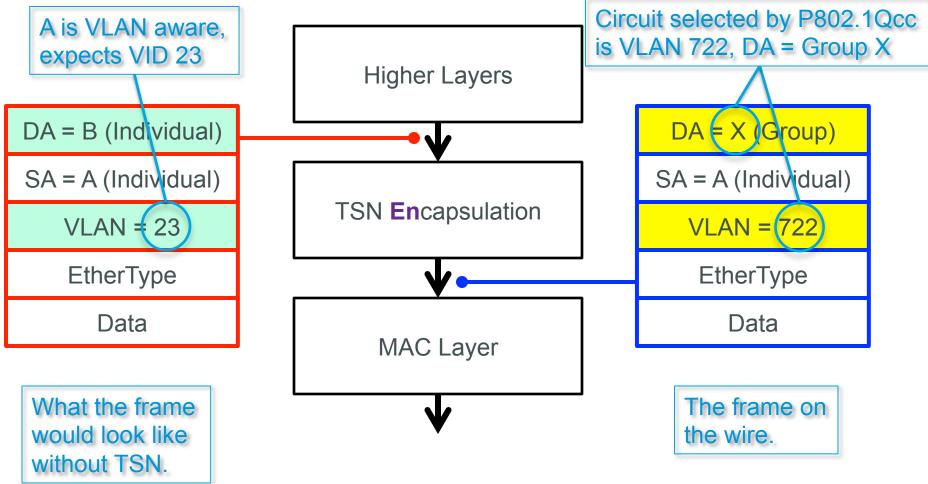
# What is the Fix?



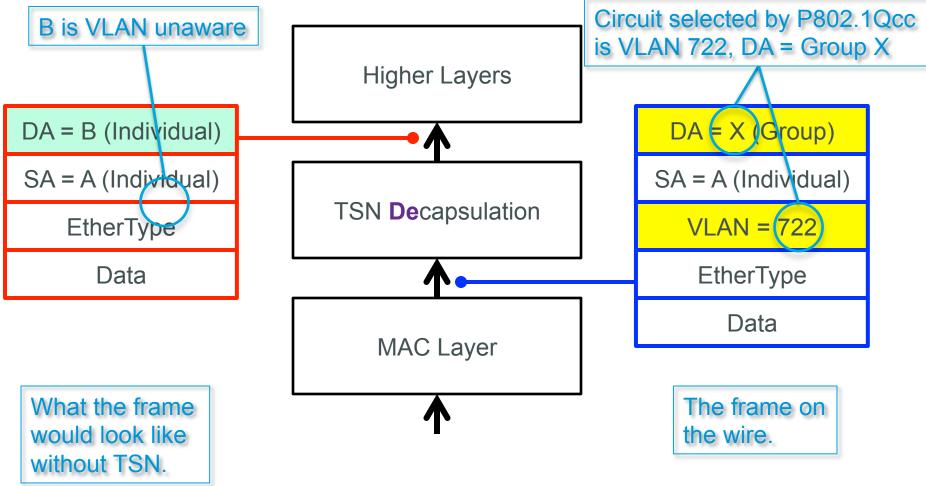
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# Correct stack in Talker A



# Correct stack in Listener B

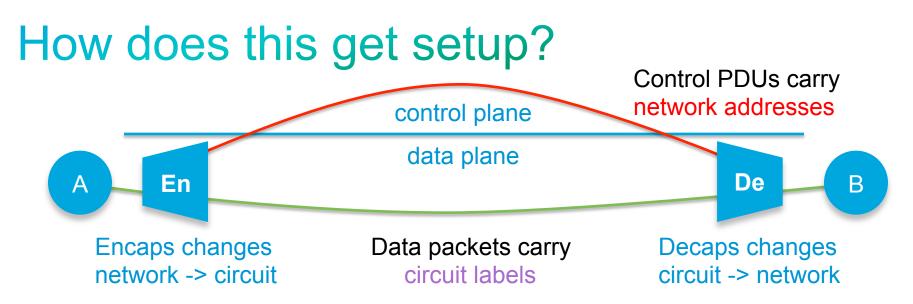


# How does this get setup?

- A's control functions use P802.1Qcc to say, "I want a connection to the function reached by a connectionless packet (an Ethernet frame) with VLAN 23 and Destination MAC address **B**."
- P802.1Qcc picks the encapsulation {VID 722, Group Address X}.
- Along the path from Bridge to Bridge, each bridge makes whatever modifications to the end-point addresses inside the P802.1Qcc PDU that would be made to the network packet.
   For example, VID translations or tag removal.

# How does this get setup?

- The edge Bridge adjacent to **B**, in particular, removes the VLAN tag from the P802.1Qcc request, and tells **B** about the encapsulation.
  - B's TSN Encaps/Decaps function is VLANaware, even if the rest of B's protocol stack is not!



 The information required by the TSN Encapsulation and Decapsulation functions came from the control plane.

# The end result?

- No problem with IP, including unicast streams.
  - > The transformations are transparent to the IP stack.
  - > The IP stack works just like it always has with ARP, etc.
- No problem with any other protocol that knows about or requires particular MAC addresses or VLANs
  - > The {VLAN, address} is restored as it comes up the receiving stack.
- No problem with fixed paths.
  - Every stream has a multicast address. One or two VLANs can support all fixed-path streams, even with VLAN-aware host stacks.
- No problem with MAC address learning.
  - > Because fixed paths are on a VID that doesn't do learning.
- No problem with backwards compatibility.
  - Existing AVB stations see the same encapsulation as always.
  - We'll craft P802.1Qcc to allow the Talker to supply the tunnel address.

# Why didn't we see this, before?

- The circuit encapsulation was so trivial, we didn't realize that we were doing circuits with labels.
- The circuit encapsulation is so trivial, you can implement an Ethernet-only application that uses the circuit label as a network address.

#### What do we do?

- 1. Revise 802.1Q to have the TSN Encaps/ Decaps layer, working as described.
- 2. Fix P802.1Qcc to:
  - Take the connectionless target {VID, MAC address} pair as an input.
  - Modify that pair, as needed, as P802.1Qcc proceeds through the network.
  - Give the circuit {VID, MAC address} to the endpoints, rather that taking it as an input (except as necessary for legacy reasons).

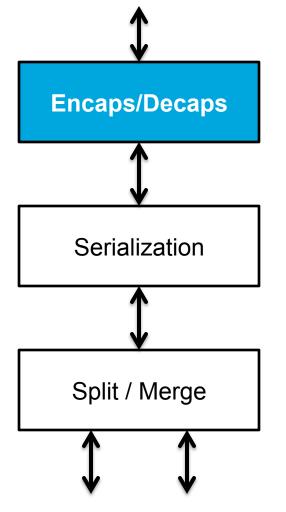
# Do I have to change my ASICS?

- If what you're doing now works, don't change it. We're not changing the bits on the wire.
- How or whether you do the full reconstitution of the frame inside your host stack is an implementation matter.
- If you want to provide a totally transparent service to the upper layers, you have the option of implementing the TSN Encaps/Decaps.
- Or, of implementing (or using your existing implementation of) several other protocols.

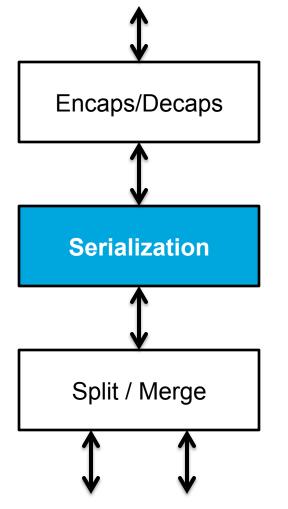
# Serialization and Seamless Redundancy



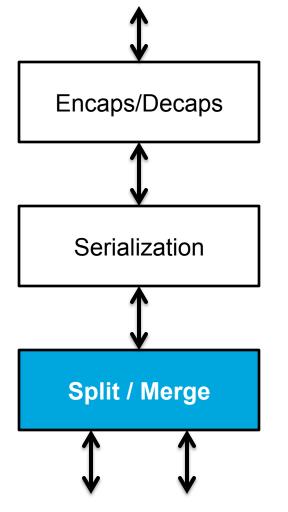
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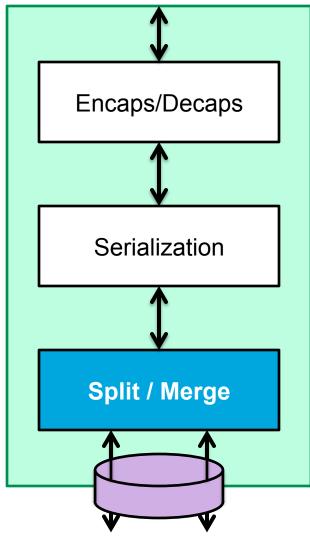
- Individual flows must be identified, and packets encapsulated, for:
   Fixed paths;
  - Per-flow resources;
  - Seamless redundancy.



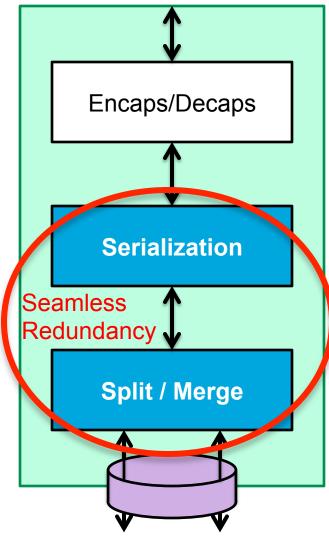
- If stream requires either seamless redundancy or long-range out-of-order delivery protection, then its packets must be:
  - Serialized on transmit.
  - Reordered on receive.



- If stream requires seamless redundancy, then its packets must be:
   Serialized.
  - ≻Split on transmit.
  - ≻Merged on receive.

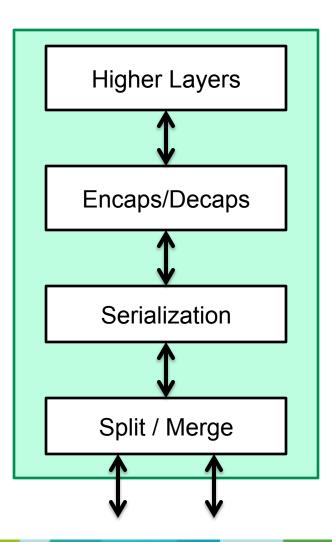


- Note that the Split Merge function does not require two physical ports.
- It may replicate / merge the streams by using different explicit in-band markers, leaving it to normal networking to make the physical split.



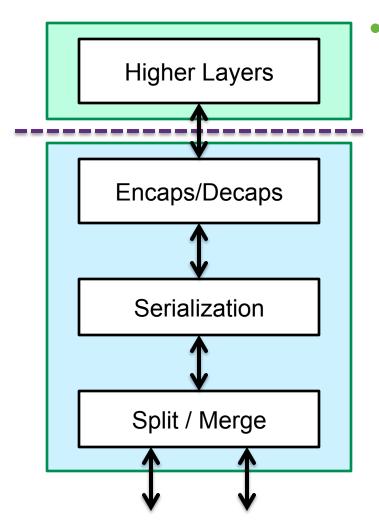
- One may want Serialization without Split/Merge.
- Split/Merge and Serialization often are combined into a single seamless redundancy function.
- All three can be combined, for that matter.

# Single system



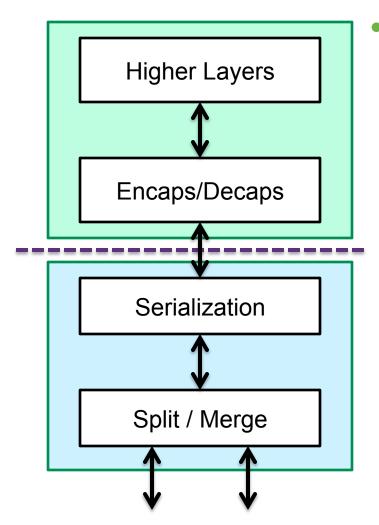
- Within a single system, Flow Identification can be out-of-band or in-band.
   Out-of-band methods include:
  - ≻Socket ID.
  - Separate service instances.
  - Multiple Serialization functions.

# Single system



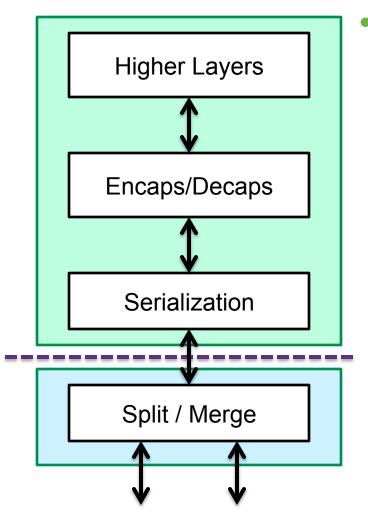
- For multiple systems, flow Identification must be in-band:
  - ≻Some form of tag.
  - One or more layers of explicit addresses (e.g. VLAN ID or IP 5-tuple).
  - A flow ID buried in an application.

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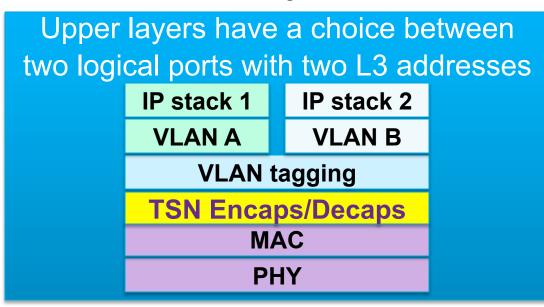
# Single system



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  - ➢A flow ID buried in an application.

#### Single-port multiple VLAN host

 Common model for a multi-VLAN host with a single physical port (router or multi-VLAN server). TSN Encaps/Decaps works, but has no seamless redundancy.

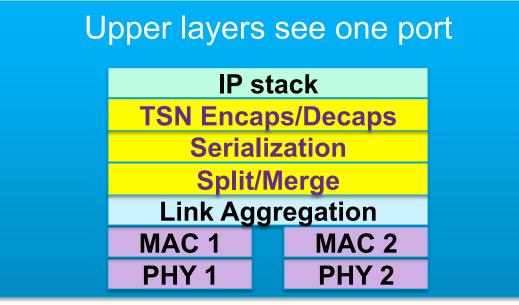


#### Single-port multiple VLAN host

- Even if the VLAN Tagging layer is not present (i.e., the host is VLAN-unaware), the TSN Encaps/Decaps function is VLAN-aware.
- No serialization or split/merge functions were shown in the diagram. They could be present. Often, however, their functions would be proxied by the adjacent bridge, in which case the TSN MAC address provides a great circuit ID.

# Dual-port non-relay host (Link Aggregation)

 The DRNI model works transparently with Seamless Redundancy. Link Aggregation layer splits regular traffic normally, and splits SR traffic using the circuit label.

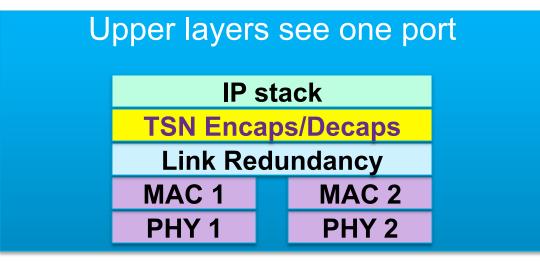


#### **DRNI** Host

- The DRNI host cannot be part of a ring it can only be a dual-ported end station. (That's a use case, not a problem.)
- The non-TSN applications work just fine, without replication, because DRNI works.

#### Dual-port host (HSR or PRP)

 IEC 62439-3 HSR/PRP supports dual-homed hosts along with TSN. The Link Redundancy layer provides Serialization and Split/Merge capabilities, within the limits of the HSR/PRP topology assumptions.

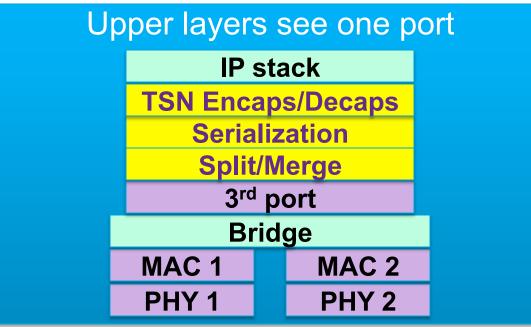


#### Dual-port host (HSR or PRP)

- As written, HSR supports a host that relays traffic from port to port, and PRP supports a host that does not.
- HSR requires a ring topology.
- PRP requires connections to separate networks.
- As we will see in the next section, both protocols can be easily adapted to work over a general purpose 802.1 network for TSN.

#### Dual-port relay host (bridged)

 The Bridge model works fine. The host stack creates differently-labeled circuits, and the bridge part directs them on different paths.



#### Dual-port relay host (bridged)

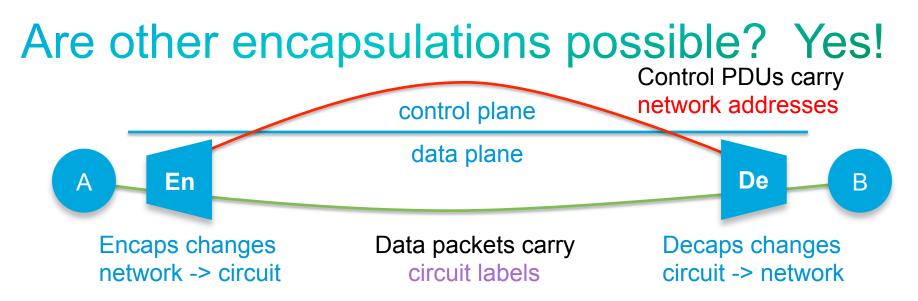
- This model should be of particular interest to TSN, as it is the most general, and of course, 802.1 defines bridges.
- It automatically supports rings or dual-homed stations.
- When combined with the <u>Simple 2-port Intermediate System</u> concept, it makes a powerful ring/chain node implementation.

# Alternative TSN Encapsulations



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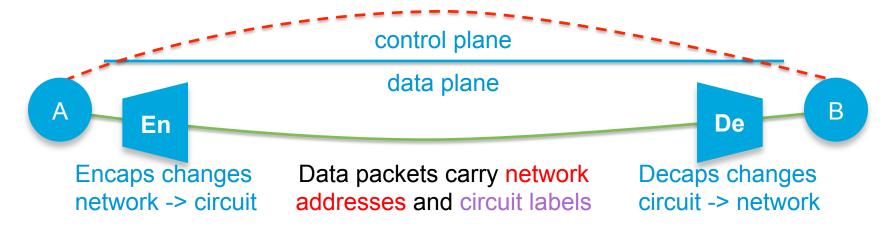
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 To use the current AVB format, while making TSN services transparent to the user, the Encapsulation function destroys information, and the Decapsulation function restores it.

#### • Let's call this **out-of-band tunneling**.

#### Are other encapsulations possible? Yes!



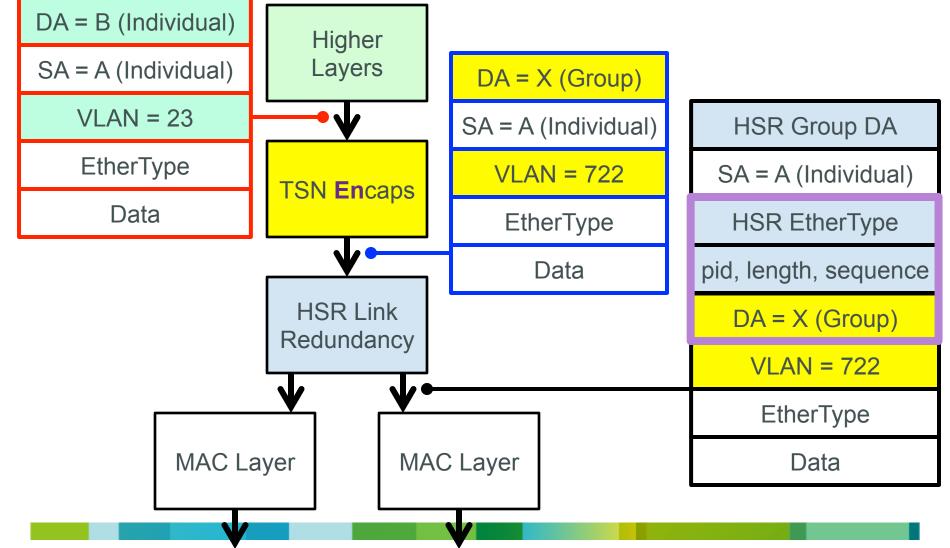
- We could do in-band tunneling, and encapsulate the B's address and VLAN in the data frame, itself.
- There are existing, applicable protocols that work both ways.

# 1. HSR seamless redundancy (1)

Fixed Group DA
user SA
optional VLAN Tag
HSR EtherType
pid, length, sequence
saved user DA
opt. user VLAN Tag
user data

- The IEC 62439-3 High-Speed Seamless Redundancy (HSR) encapsulation provides in-band tunneling.
- Almost. On the good side:
  - HSR encapsulates the original destination MAC and VLAN, rather than using the data plane.
  - HSR includes a sequence number for seamless redundancy. (That is the name of the protocol!)

# HSR seamless redundancy (1)



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#### 2. HSR seamless redundancy (2)

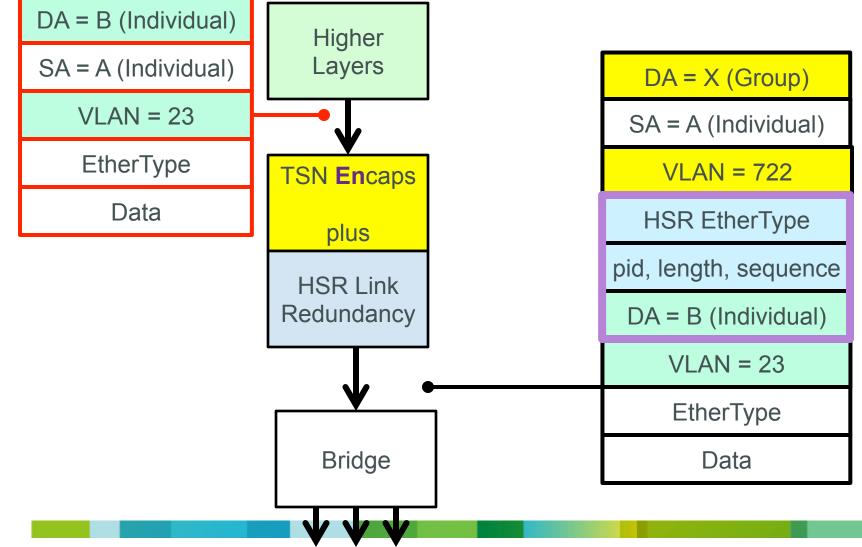
- Oops! HSR only works if the network is an HSR ring, because the TSN Circuit address has been buried behind the fixed HSR DA.
- To make HSR work over 802.1 networks:
   > One change to the PDU format:

 $_{\odot}$  We substitute the TSN circuit DA for the HSR fixed DA.

>One change to the use of that format:

 The Link Redundancy layer does not forward packets – that is left to a bridge function, below it.

# HSR seamless redundancy (2)

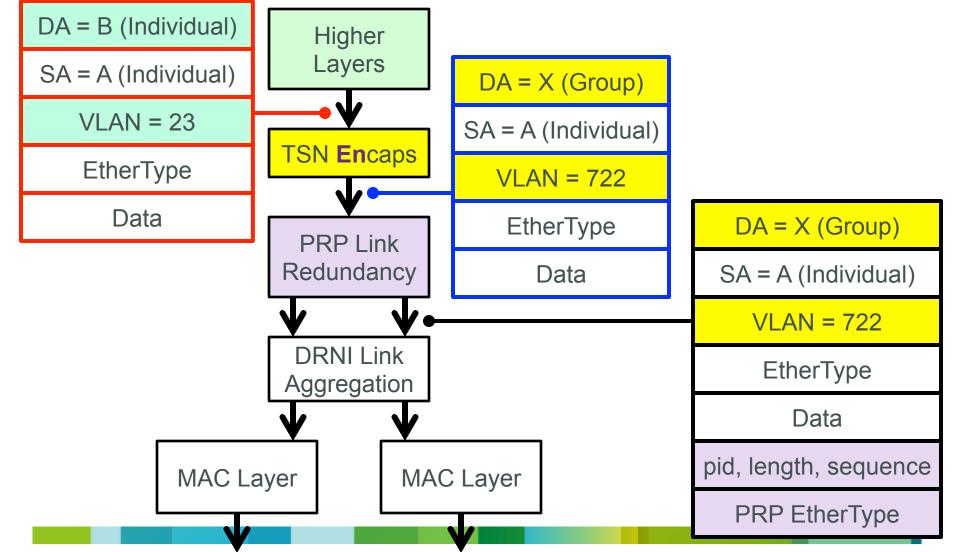


#### 3. PRP seamless redundancy + TSN

user SA
user SA
opt. user VLAN Tag
user data
pid, length, sequence
PRP EtherType

- The IEC 62439-3 Parallel Redundancy Redundancy (PRP) encapsulation will work with the TSN encapsulation.
- There are serious issues with an 802.3 frame trailer. Let's leave that for another discussion.
- TSN supplies the circuit ID, and PRP supplies the split/merge capability.

#### IEC 62439-3 PRP + TSN



#### IEC 62439-3 PRP + TSN

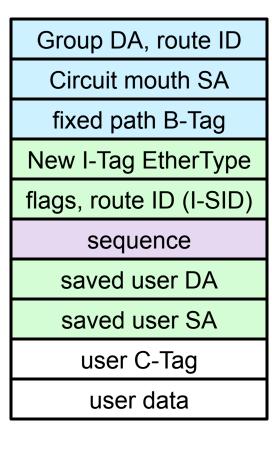
- PRP was intended to work only over separate networks; one port is connected to one network, and the other to another network. But, we can make it work over a single 802.1 network.
- It works, because the TSN circuits provide the logical equivalent of separate networks.
- A DRNI / Link Aggregation unit, as described earlier, would connect this stack to the physical ports.

#### 4. PBB-TE for TSN

Group DA, route ID
Splitter SA
fixed path B-Tag
I-Tag EtherType
flags, route ID (I-SID)
saved user DA
saved user SA
user C-Tag
user data

- MAC-in-MAC (802.1ah) solves the stream ID and fixed path problems using in-band tunneling.
  - It encapsulates the data now carried in the control plane.
  - But, it has no flow ID or sequence number for seamless redundancy.

#### **PBB-TE for seamless redundancy**

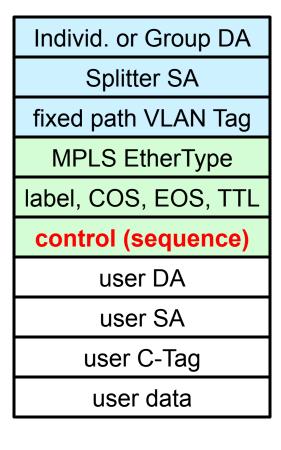


- We must either have an additional tag for the sequence number (not clear where), or a new I-Tag format that includes a sequence number (shown at left).
- Or, just use it for tunneling.

#### 5. MPLS Ethernet Pseudowire

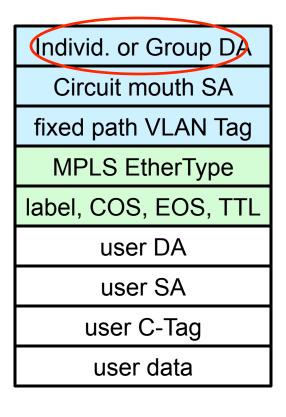
- An MPLS pseudowire provides almost exactly the same encapsulation structure as 802.1ah.
- Like the current AVB encapsulation, it provides out-of-band tunneling.

#### Pseudowires for seamless redundancy



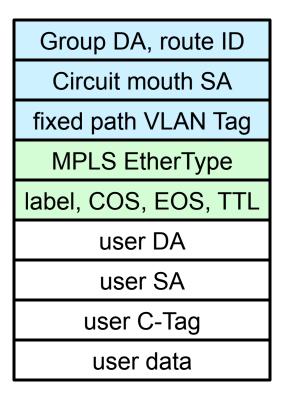
 Plus, Pseudowires have an optional control word that provides a sequence number for seamless redundancy.

#### **MPLS Ethernet Pseudowire**



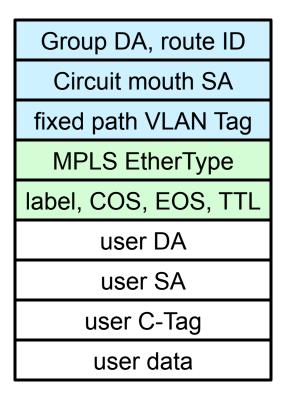
 There is one problem: MPLS uses the next-hop destination, or a fixed Group DA.

#### **MPLS Ethernet Pseudowire**



- There is a solution: MPLS uses a fixed Group DA OUI, with the outermost MPLS label value in the low-order bits of the Group DA.
- (Just like PBB-TE.)

#### **MPLS Ethernet Pseudowire**



- There another solution: We just talked about the TSN Encaps/Decaps layer changing the outer MAC address!
- Which solution is better is To Be Determined

#### 6. No protocol at all

- OpenFlow can recognize a stream based on common fields, such as the MAC addresses or IP 5-tuple.
- Any number of bridges have "Access Control Lists" (ACLs) that can inspect a frame and take special action.
- So, one can always leave the original frame intact, and use frame inspection to identify the circuit so that special actions can be taken.
- But, this layering model is still important, you're still doing circuits, and P802.1Qcc still has to change.

# Summary



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### Summary 1/2

- When you get the layering right, everything just works.
  - No need for deep packet inspection. No problems with existing IP stacks, too many VLAN IDs, MAC address learning, fixed paths, or backwards compatibility.
  - ≻We must <u>align P802.1Qcc with proper layering</u>.
  - We must define the TSN Encaps/Decaps function.

# Summary 2/2

- When you get the layering right, several existing protocols support TSN circuits.
  - Properly layered TSN, HSR(1), HSR(2), Pseudowires, or no protocol at all.

There are more.

 When you get the layering right, several existing protocols support seamless redundancy, already.

HSR (1), HSR (2), PRP + TSN, Pseudowires.There are more.

#### Thank you.

#