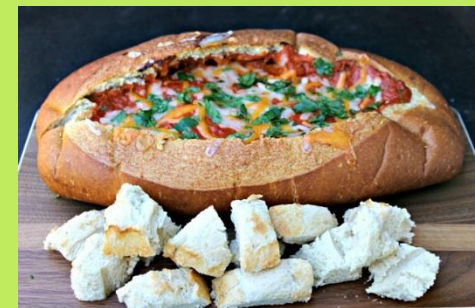


STATELESS RESOURCE SHARING AND ATS

PUTTING TOGETHER THE BEST OF BREADS

IEEE P802.1Qcr, 2016-07-27, San Diego, CA
Szilveszter Nádas, Balázs Varga (A), János Farkas



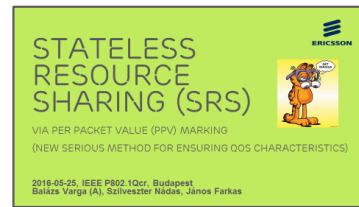
INTRODUCTION

SRS OVERVIEW AND PRESENTATION GOAL



› Stateless Resource Sharing (SRS) via per packet value (PPV) marking

- Basic concept was presented in Budapest
- [Link-to-contribution](http://www.ieee802.org/1/files/public/docs2016/cr-varga-srs-ppv-0526-v02.pdf)



› Goal of this presentation to show

- Combination of SRS and ATS
- Achievable advantages by such a combination

SUMMARY

UNIQUE CHARACTERISTICS

SRS via PPV at-a-glance

- › SRS essentials:
 - Share of available BW between flows is encoded in the packets (i.e. shade of a packet)
 - Shade based dropping in queues does not need traffic situation specific pre-configuration at the congestion point(s)
 - BW share is automatically controlled by the shades of packets participating in congestion situation
 - Accuracy of BW share depends on the number of shades used (Note: predefined ratio exactly ensured if drop level = border between shades)
 - Explicit feedback is possible for systems that need congestion notification (i.e. what level of congestion they will face during transport is known from the congestion shade)

Stateless Resource Sharing (SRS) | 2016-05-25 | Page 8

SRS PROVIDES ZERO CONGESTION LOSS AND DETERMINISTIC LATENCY



- › SRS can achieve zero congestion loss and deterministic latency

- Congestion results in packet drop:
 - › Darker packets kick-off brighter ones
- Congestion level:
 - › Can be defined by a „congestion shade”

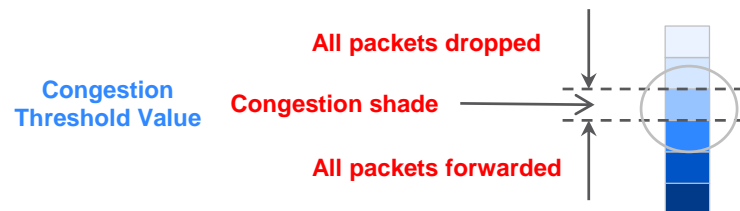
- › **Note:** All traffic having darker shades than the ”Congestion Threshold Value” will experience zero loss and deterministic delay

- › SRS shortcomings

- Congestion Threshold Value is not set in advance, it is rather the result of actual load and bottleneck capacity
- SRS (itself) does not protect against bad behaving nodes



- › **BUT:** Congestion Threshold Value can be dimensioned by resource allocation and worst case delay calculation (by e.g. ATS)



SRS ADD-ONS

DETERMINISTIC DELAY FOR LOSSLESS & LOSSY TRAFFIC



- › It is possible to extend the ATS scenario with traffic that has the same delay guarantee as “guaranteed” TSN traffic, but some loss is allowed, i.e., it is lossy
- › What does it mean in practice?
 - Allow reservation of less resources for ATS flows which have a loss tolerant component but needs in-order-delivery for all packets of a data-flow
 - Additional loss tolerant flows that require the same deterministic delay can be served easily
 - Other scenarios may also exist ...
- › The positive effect
 - It can highly increase link utilization when some flows do not use their reservation all the time
 - Under the prerequisite that (some) lossy traffic has delay guarantee too
- › This works hop-by-hop, not just for flows following the same path

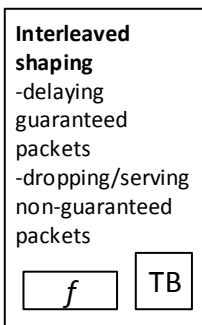
CHANGED COMPONENTS FOR COMBINING SRS AND ATS



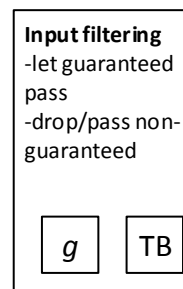
- › Need to be able to identify additional traffic (at all hops)
- › Slightly larger bucket sizes are needed to allow excess traffic



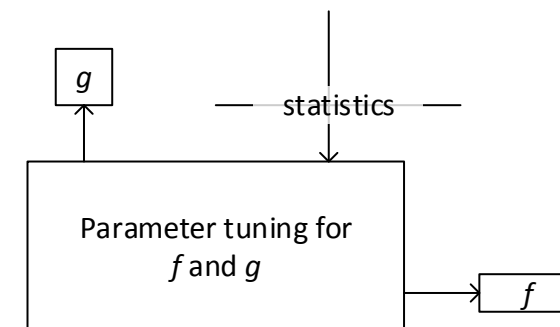
- › Interleaved shaping has to be able to
 - drop excess packets
 - read packet values, and based on that, influence whether a packet is dropped or not



- › Input filtering for excess packets is needed in order to avoid flooding the queues
 - Slightly larger buffers are needed

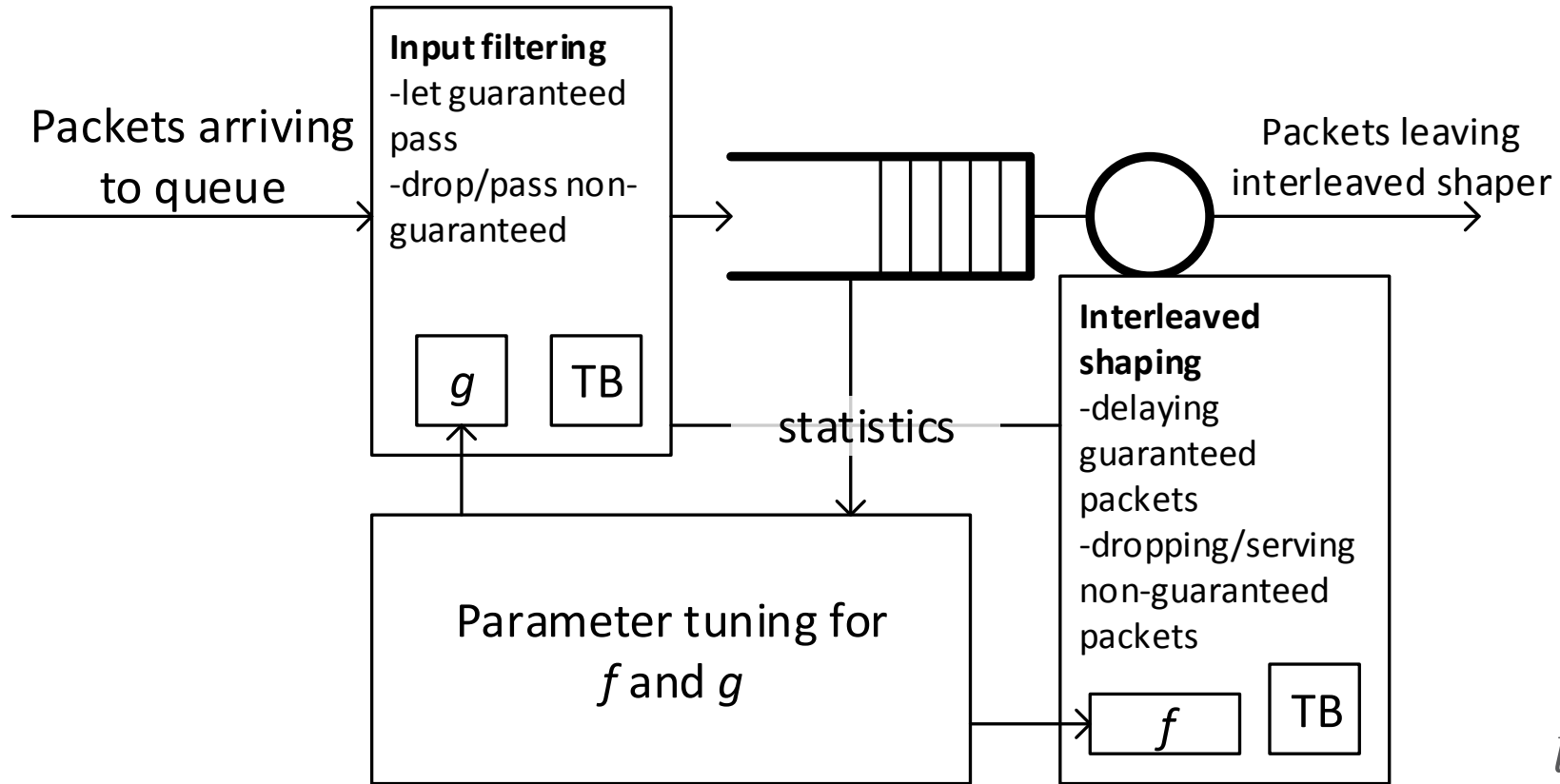


- › Statistics from all boxes are needed to tune packet value aware dropping

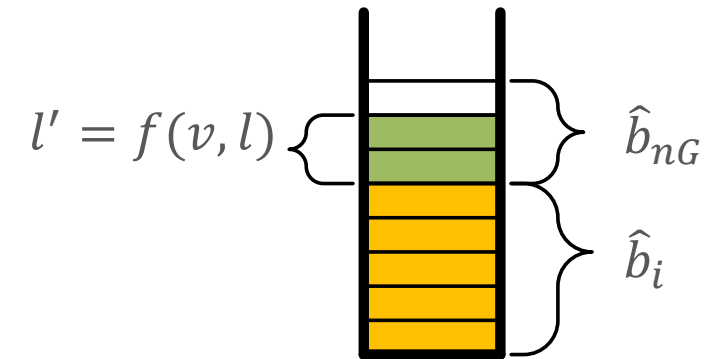


ARCHITECTURE

PUTTING IT ALL TOGETHER



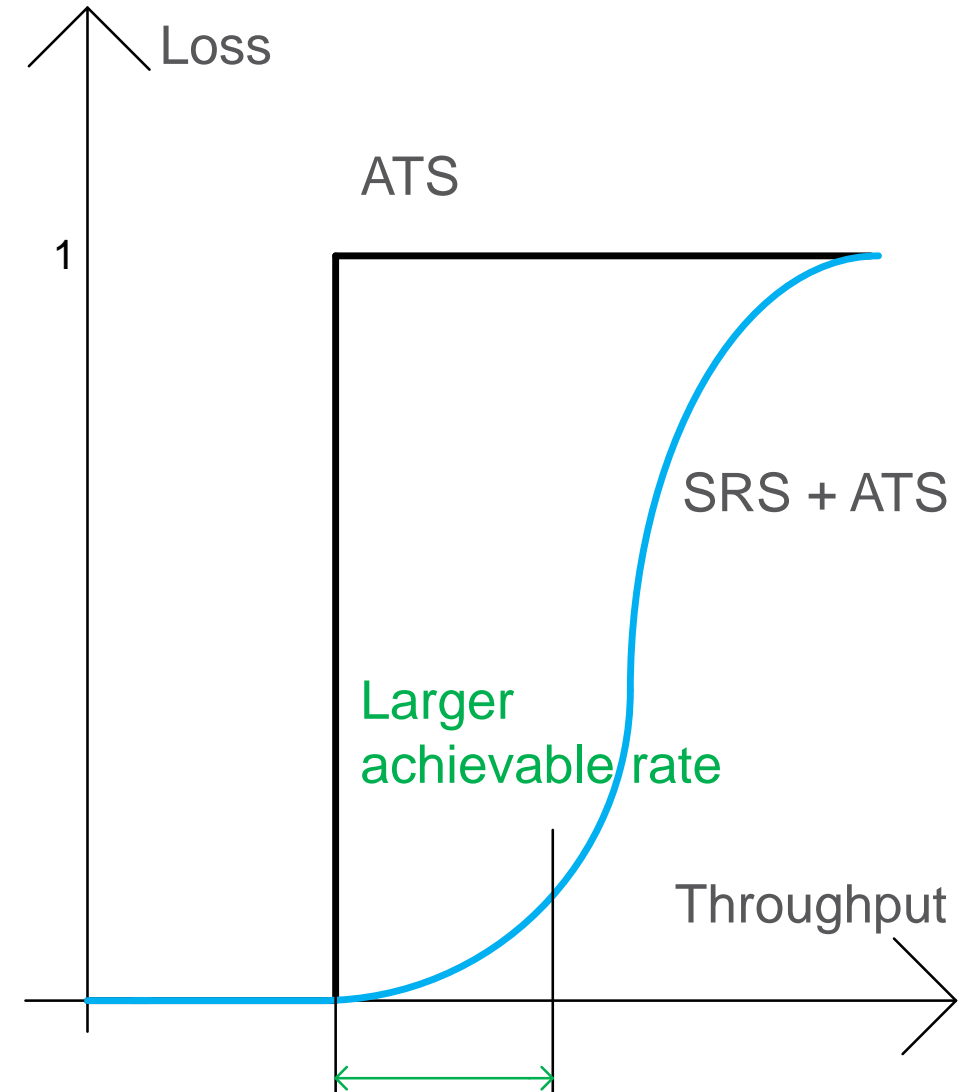
Non-guaranteed eligible of at least $\hat{b}_i + l'$ tokens in bucket



OUTCOME OF COMBINATION LOSS VS. THROUGHPUT



- › A slightly larger bucket size might be needed for each flow
- › The possibility of forwarding non-guaranteed packets results in a larger **achievable flow rate**
- › The packet value determines
 - Whether or not a packet is guaranteed
 - Whether a non-guaranteed packet is dropped or forwarded (note: there can be more important and less important non-guaranteed packets)
- › The size of the larger bucket is chosen based on a compromise between (1) delay, (2) allowed total guaranteed rate, and (3) allowing excess traffic
 - In many cases, slightly increased delay is still within the E2E delay budget



SUMMARY

FURTHER WORK



- › This is an initial proposal to evolve ATS further
- › Delay model and math to be discussed and verified in detail
- › Creating a list of use cases



ERICSSON