

# Leveraging ON/OFF Traffic Model – Literature Example

Backup Material for the Presentation on Support for eCPRI in 802.1CM

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# Examples from Related Literature

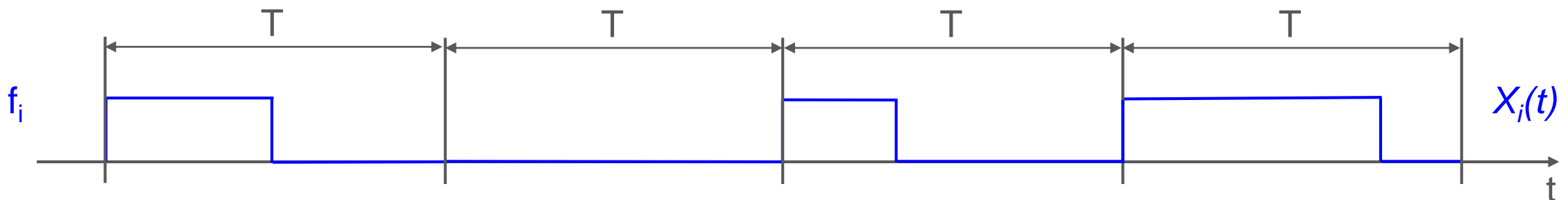
- [1] B. József, Z. Heszberger, M. Martinecz. "A family of performance bounds for QoS measures in packet-based networks." *NETWORKING 2004. Networking Technologies, Services, and Protocols; Performance of Computer and Communication Networks; Mobile and Wireless Communications* (2004): 1108-1119.
- [2] Yin, N., & Hluchyj, M. G. „Analysis of the leaky bucket algorithm for on-off data sources”, *Journal of High Speed Networks*, 2(1), 81-98, 1993.
- [3] Akar, N., & Arikan, E. „Markov modulated periodic arrival process offered to an ATM multiplexer”, *Performance evaluation*, 22(2), 175-190, 1995.
- [4] Malomsoky, S., Rácz, S., & Nádas, S., „Connection admission control in UMTS radio access networks”, *Computer Communications*, 26(17), 2011-2023, 2003.
- [5] Nádas, S., Rácz, S., Malomsoky, S., & Molnár, S., „Connection admission control in the UTRAN transport network”, *Telecommunication Systems*, 28(1), 9-29, 2005.

# The Paper Discussed in the Following

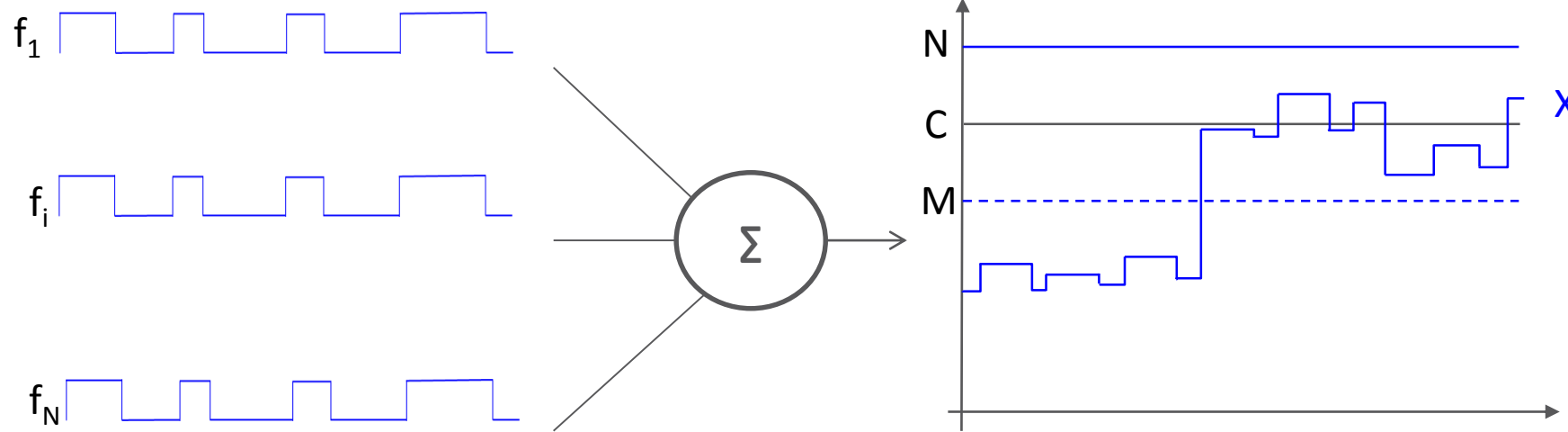
[1] B. József, Z. Heszberger, M. Martinecz. "A family of performance bounds for QoS measures in packet-based networks." *NETWORKING 2004. Networking Technologies, Services, and Protocols; Performance of Computer and Communication Networks; Mobile and Wireless Communications* (2004): 1108-1119.

# Notations

- ›  $N$  = total number of flows
- ›  $f_i$  = a generic stationary, ergodic flow
- ›  $X_i(t)$  = traffic generation rate of  $f_i$  at moment  $t$ 
  - It is a stochastic variable
- ›  $\max\{ X_i(t) \}$  = peak rate of  $f_i$  – must be known
- ›  $m_i = E(X_i(t))$  = average rate – must be known
- ›  $M = \sum_{i=1}^N m_i$  – must be known (even if  $m_i$  is not known)
- ›  $C$  = service rate (link capacity)
- › The flows can be observed periodically ( $T$  period)



# Aggregated Flows



› If the flows are independent, then

$$Pr(X > C) \leq \left(\frac{M}{C}\right)^C \left(\frac{N-M}{N-C}\right)^{N-C} \leq \epsilon \quad [1]$$

› Typical use case:

- Tolerable loss rate  $\epsilon$ , is given
- What  $C$  is needed?

# Numerical Examples

Max FLR =  $10^{-7}$  for IQ data in P802.1CM D0.7

Max FLR  $> \epsilon + \Pr\{\text{BER}\}$ ,  $P(\Pr\{\text{BER}\} = 6.6 \cdot 10^{-8} [4])$

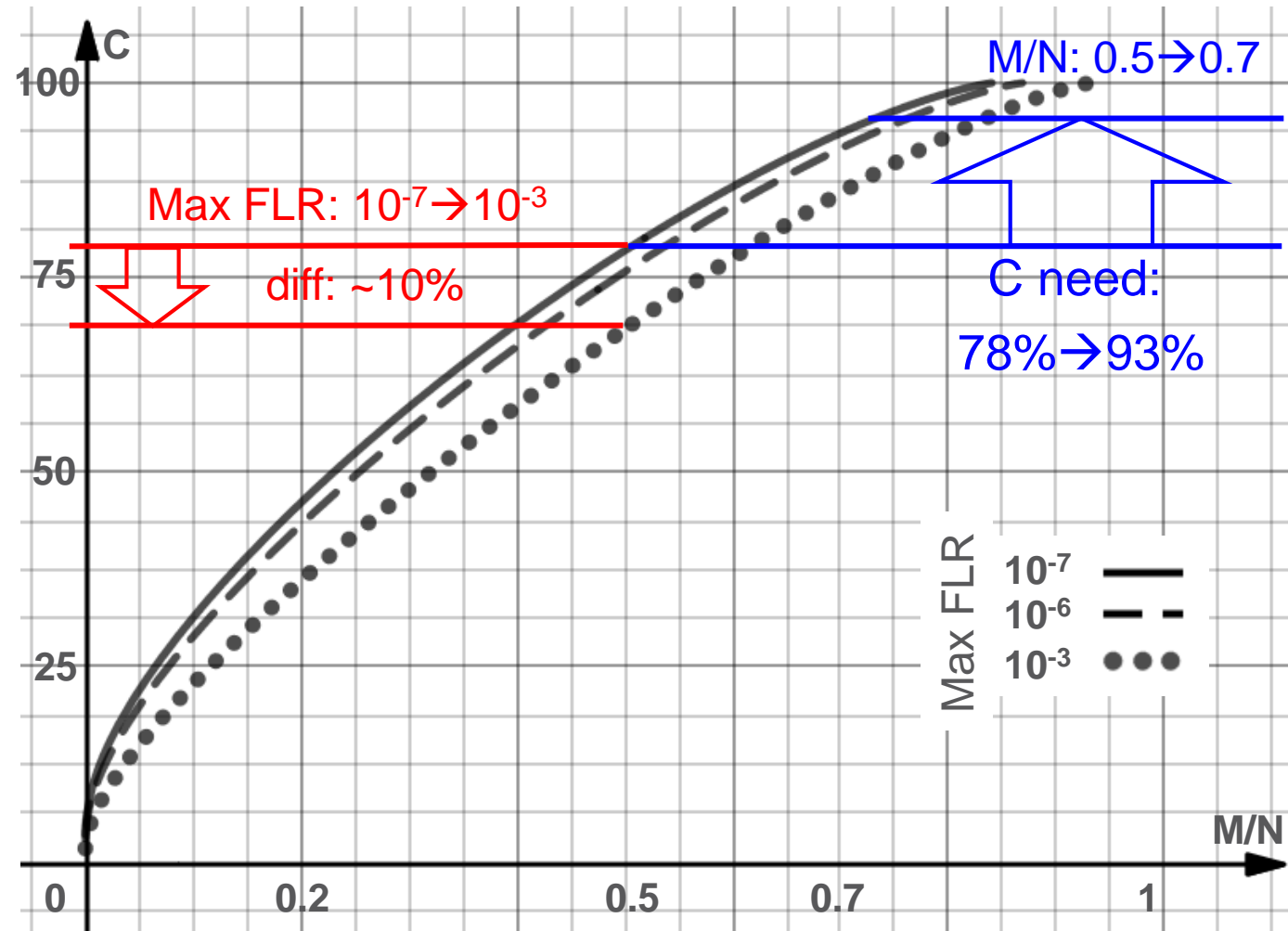
$$\Pr(X > C) \leq \left(\frac{M}{C}\right)^C \left(\frac{N-M}{N-C}\right)^{N-C} \leq \epsilon$$

$$C = f(M/N)$$

- X axis: normalized average arrival rate
- Y axis: C = capacity need for the aggregate
  - C corresponds to CIR
- Plotted for N = 100 flows
  - N corresponds to CIR + EIR
- Relaxing the loss requirement
  - Max FLR =  $10^{-7}$  vs.  $10^{-6}$  vs.  $10^{-3}$
- Capacity need difference if  $\frac{M}{N} = 0.5$ 
  - Max FLR =  $10^{-7}$  vs. Max FLR =  $10^{-3} \rightarrow \sim 10\%$  difference

**Worst-case values:**

$\frac{M}{N}$	$P(\text{each flow is sending data})$
0.25	$6.22 \cdot 10^{-61}$
0.50	$7.89 \cdot 10^{-21}$
0.75	$3.21 \cdot 10^{-13}$



Graph generated with: <https://www.desmos.com/calculator>

# Acknowledgements

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