P802.1CQ/D0.6 Preview v04

Roger Marks (EthAirNet Associates) P802.1CQ Editor 2021-06-11

thanks to Antonio de la Oliva and Lily Lv for review and constructive comments

Summary

- P802.1CQ/D0.5 was reviewed in TG Ballot.
- Comment resolution was completed in November.
 - Address Blocks were introduced for address claiming
- In March, Editor presented "Block Address Registration and Claiming (BARC)"
 - cq-marks-BARC-0321-v00.pdf
 - address blocks used for registrar-managed addresses as well
 - Address Registration and Claiming (ARC)
 - address blocks, and also claiming address ranges using MAAP
 - presented also to IEEE 1722 Working Group
- Main issues raised in March concerned VLAN operation
- This contribution previews P802.1CQ/D0.6
 - refinements and details since March presentation
 - discussion on improved VLAN support
 - v01 presented to TSN at May 802.1 Interim; this version (v03) adds detail

BARC assigns MAC Addresses in Address Blocks

- 1) Address Blocks (ABs) are sets of local addresses.
- 2) An AB includes equal-sized unicast and multicast address subblocks.
- 3) No BARC address falls within more than one AB.
- 4) An Address Block Designation (ABD) is a CABA or a RABI.
- 5) Claimable AB Address (CABA) is claimable by a Claimant without using a Registrar.
 - identifies Claimable Address Blocks (CABs) holding Claimable Addresses (CAs)
 - CABA is a multicast MAC address, not in any AB, and used as a DA.
- 6) RABI
 - identifies a Registrable Address Block (RAB) holding Registrable Addresses (RAs)
 - Registrable Address Block Indicators (RABIs): held in inventory of a Registrar
 - may be assigned to Claimants
 - may be claimed by Registrants
- 7) A large set of Temporary Unicast Addresses (TUAs) is specified
 - useful for initial discovery by Claimant lacking a unicast address

MAC Address Categorization

determinable via inspection:	Expanded name	I/G	indicates, by inspection
CA [ICA=unicast GCA=multicast]	claimable address, in claimable address block (CAB)	U,M	CABA, CAB Size, CAB (including all other CAs in CAB)
CABA	CAB Address	М	CAB Size, CAB
RA [IRA=unicast GRA=multicast]	registrable address, in registrable address block (RAB)	U,M	BABI Size [Basic ABI Size]
TUA	temporary unicast address	U	note: ~6.9E10 to choose among

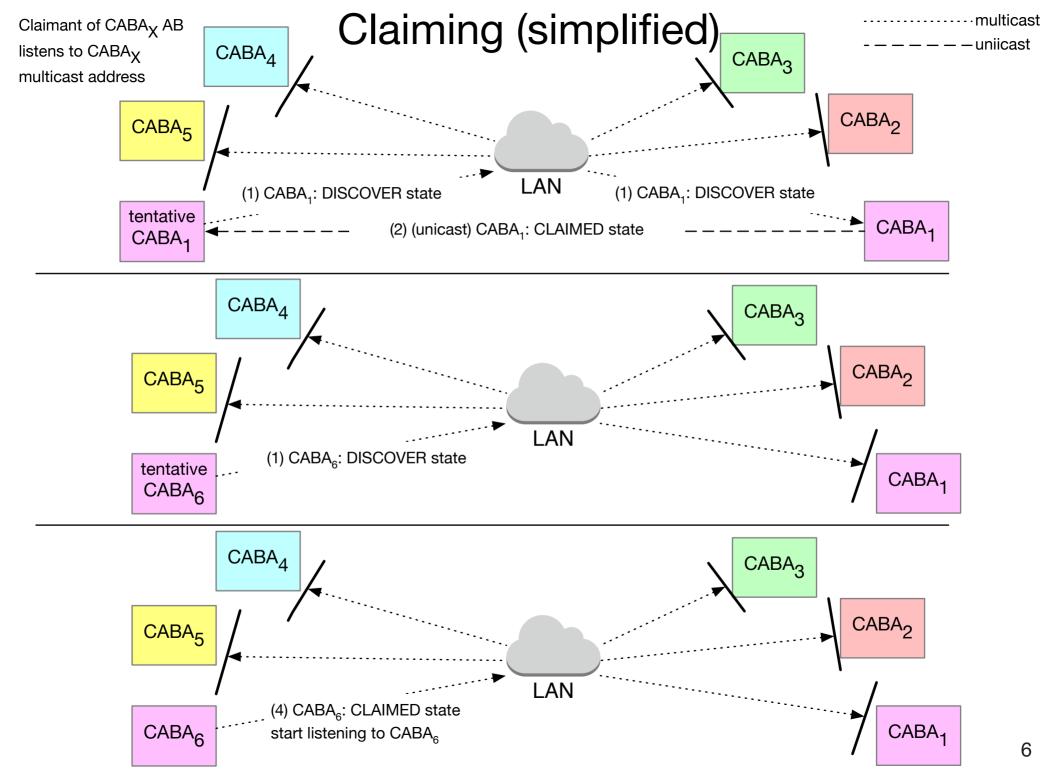
Address Block Designation (ABD) Categorization

ABD	Address Block Designation (CABA or RABI)	not an AB address (including all RAs in AB)		
RABI	RAB Identifier		RAB Size, BABI Size, RAB, s MABI Size [Multiple ABI Size]	
CABA	CAB address	М	CAB Size, CAB	

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BARC MAC Address Structure

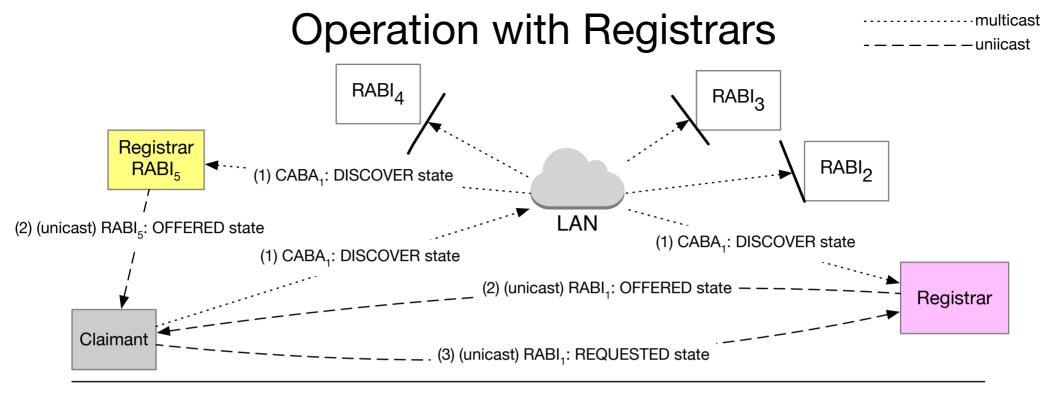
N11	r	i	j	k	for registrable addresses, r=1; for claimable addresses, r=0						
N10	1	1	1	m	m is the usual multicast (I/G) bit; 111 is local "SAI" range per IEEE Std 802c						
N9					0000 for claimable addresses						
N8					. odd	roop blo				¢	
N7	 • address block includes subblocks of - 16^{jk} claimable addresses, or 										
N6	 16^{<i>jk</i>} registrable addresses (or aggregated into larger blocks) for claimable addresses, <i>i</i> distinguishes 										
N5	- Claimable Addresses (CAs) from										
N6		 CABAs identifiers that are also used as addresses 									
N5	see Appendix for details										
N2	r							i	jk	т	
	-					СА	0	1	CAB	I/G	
N1						CABA	0	0	Size	1	
NO						TUA	0	0	0	0	
12 nibbles per 48-bit address						RA	1	RABI Option	BABI Size	I/G	

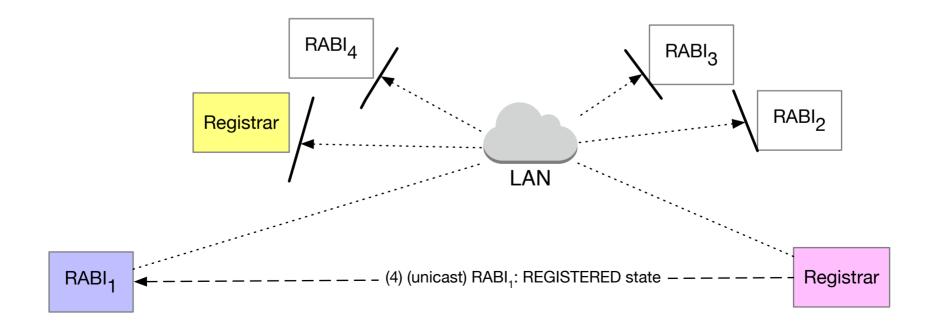


Registrar

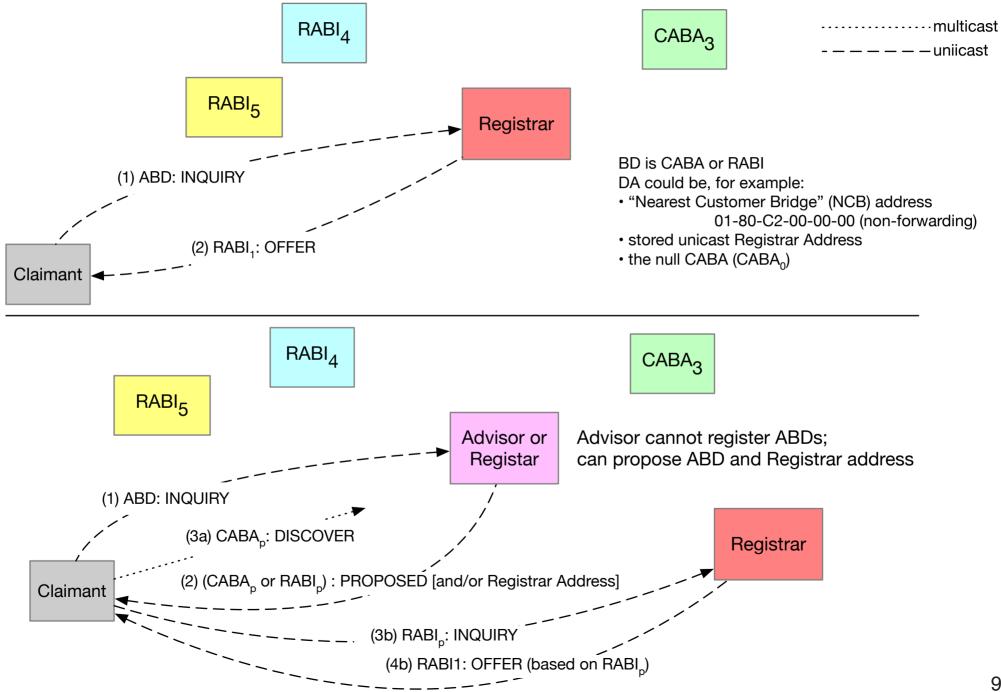
- Claimant need not be aware of Registrar when initiating a claim.
- Registrar maintains an inventory of RABIs.
 - a protocol specifies how Registrars acquire RABIs.
 - set of RABs is disjoint from the set of CABs
 - AB is either claimable (CAB) or registrable (RAB); not both
- Registrar listens for all messages to a CABA.
 - r=0, i=0, m=1, i.e. DA begins 00**-1111
 - [MMRP NumberOfValues field is 13 bits]
- Registrar can respond to a DISCOVER with an offer of a RABI in its inventory.
 - The offer can also defend the DISCOVER's CABA.
 - Registrar confirms registration of request for offered RABI.
- Pre-claim Inquiry lets Claimant reach Registrar or Advisor.
 - Client can learn of Registrars and received Claim proposals.

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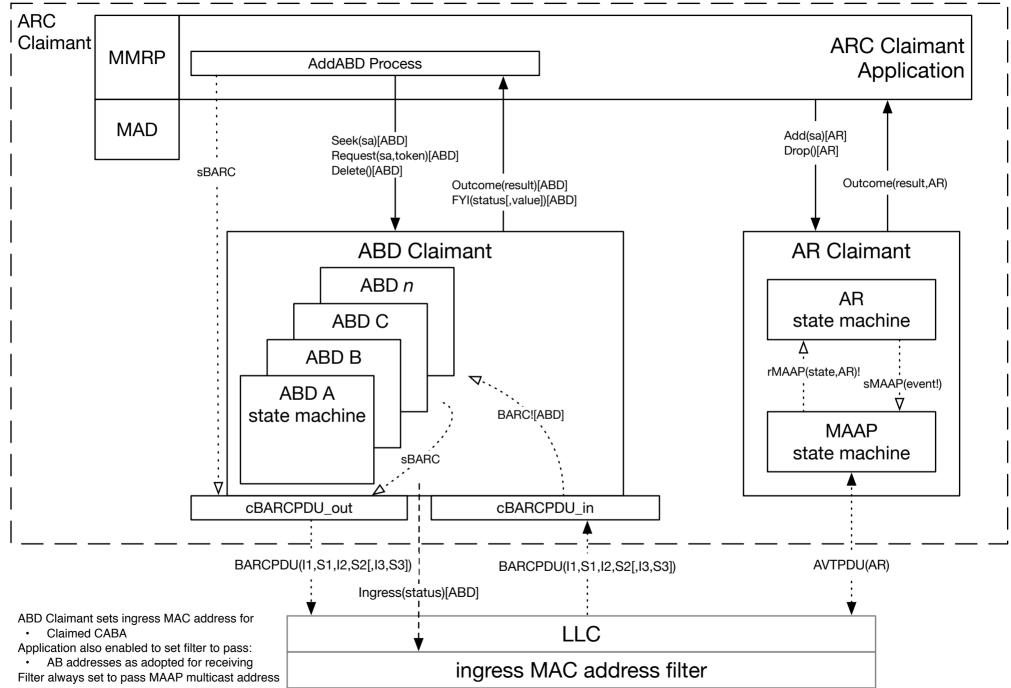
Inquiry to (anticipated) Registrar or Advisor



BARC Design

- A BARC architecture follows, with details including state machines.
 - additional details in Appendix
- BARC (Block Address Registration and Claiming) is put into the broader context of Address Registration and Claiming (ARC), which supports both:
 - address blocks (ABs), identified by Address Block Identifiers (ABIs)
 - address ranges (ARs), excluding addresses specified by BARC
- ARC is the general protocol
 - BARC handles ABI Registration and CABA Claiming
 - existing MAAP handles AR Claiming

ARC Architecture – ARC Claimant

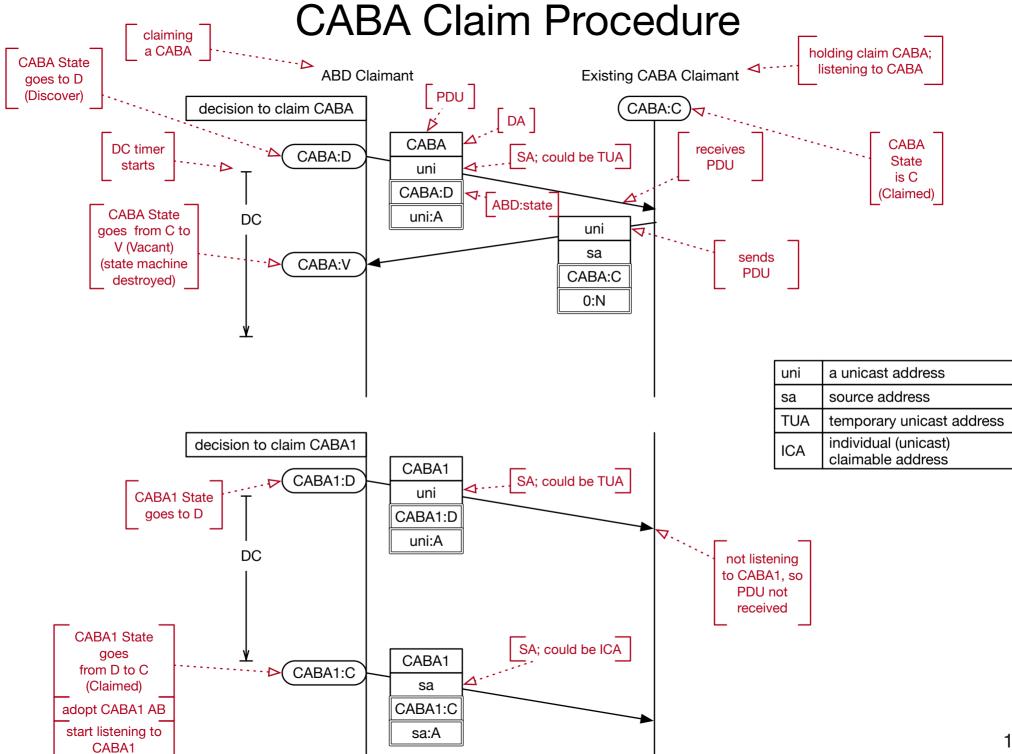


BARCPDU Summary

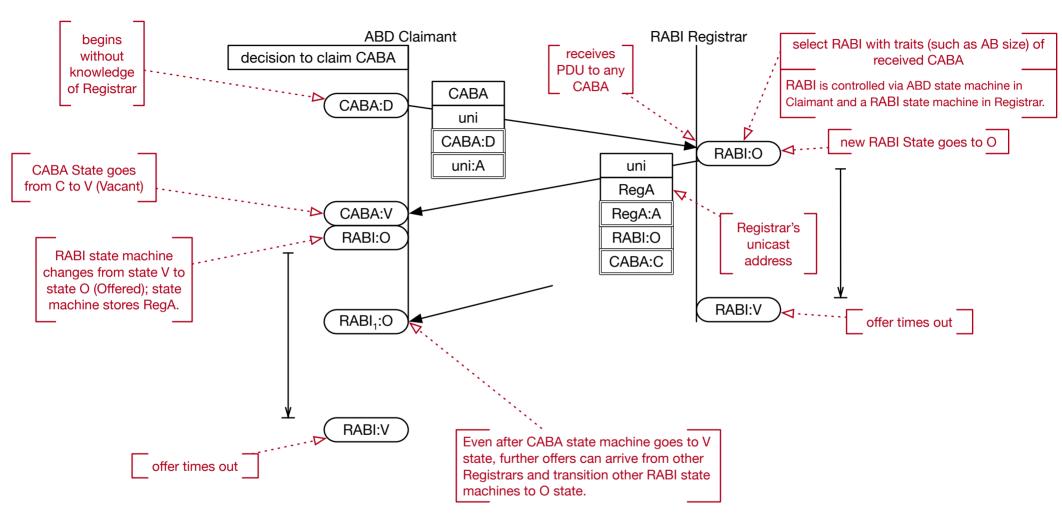
field name	purpose	content
DA	dest addresss	DA
SA	source address	
E	Ethertype	[tbd; could be 22F0 (MAAP Ethertype)]
t	subtype	[tbd, per 1722 WG; see IEEE 1722 Table 6]
S1	State 1	D (Discover), C (Claimed), V (Vacant), R (Registered), I (Inquiry), P (Proposal), A (address), RD, RC, RV, RX, N(null)
11	Identifier 2	48-bit address or ABI
S2	State 2	O (Offered), A (address), N (null)
12	Identifier 2	48-bit address or ABI
S3	State 3	A (address), T (token)
13	Identifier 3	48-bit address or token

AVTPDU Summary

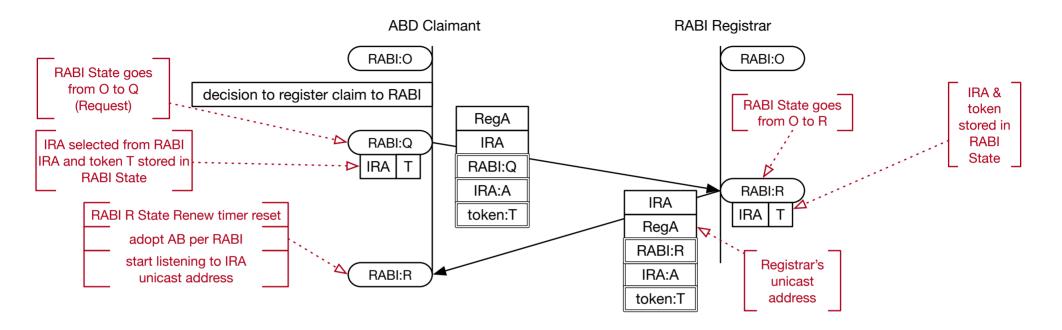
field name purpose		content		
DA	dest addresss	91:E0:F0:00:FF:00 for MAAP multicast		
E	Ethertype	22F0 (MAAP Ethertype)		
t	subtype	FE per IEEE 1722 Table 6		



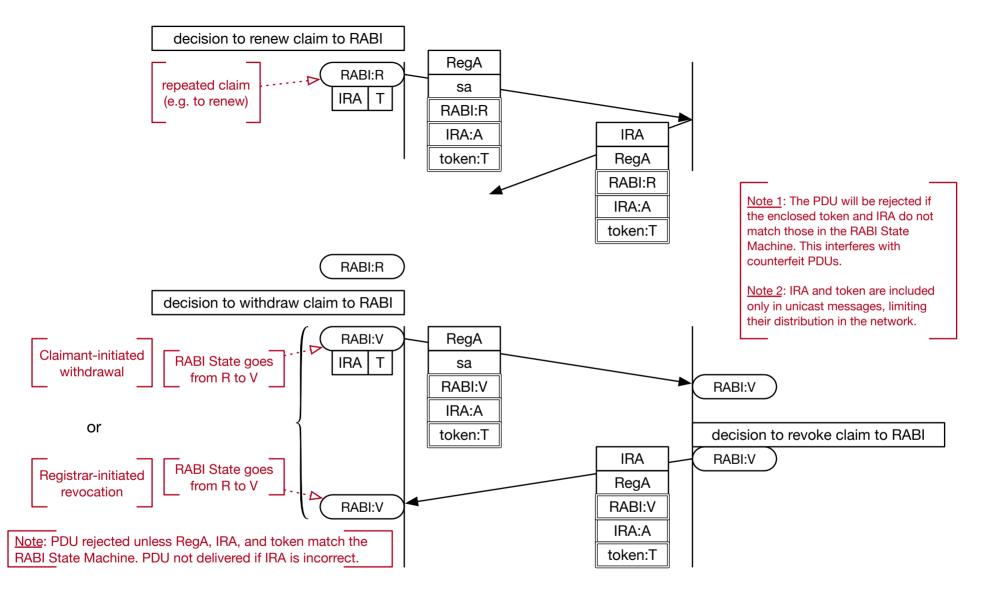
ABD Claimant/Registrar Procedure



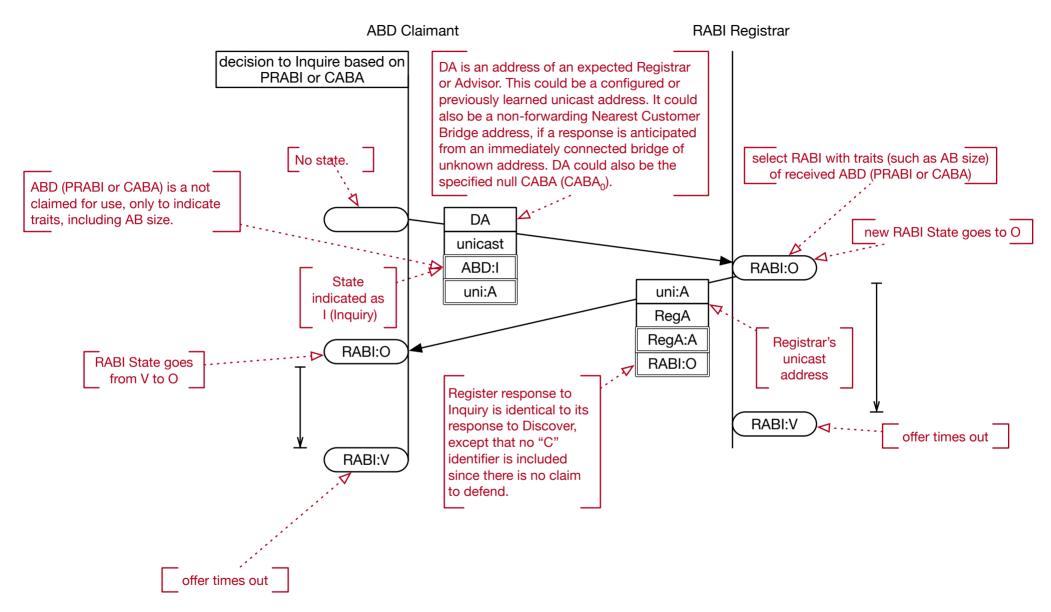
RABI Registration



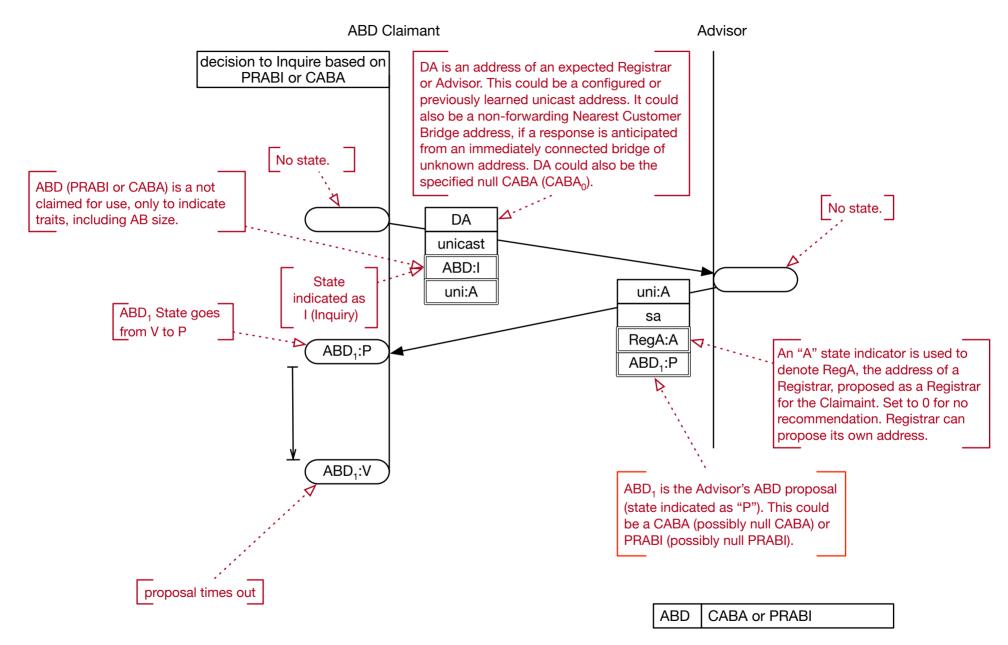
Renewal and Withdrawal of a Registration



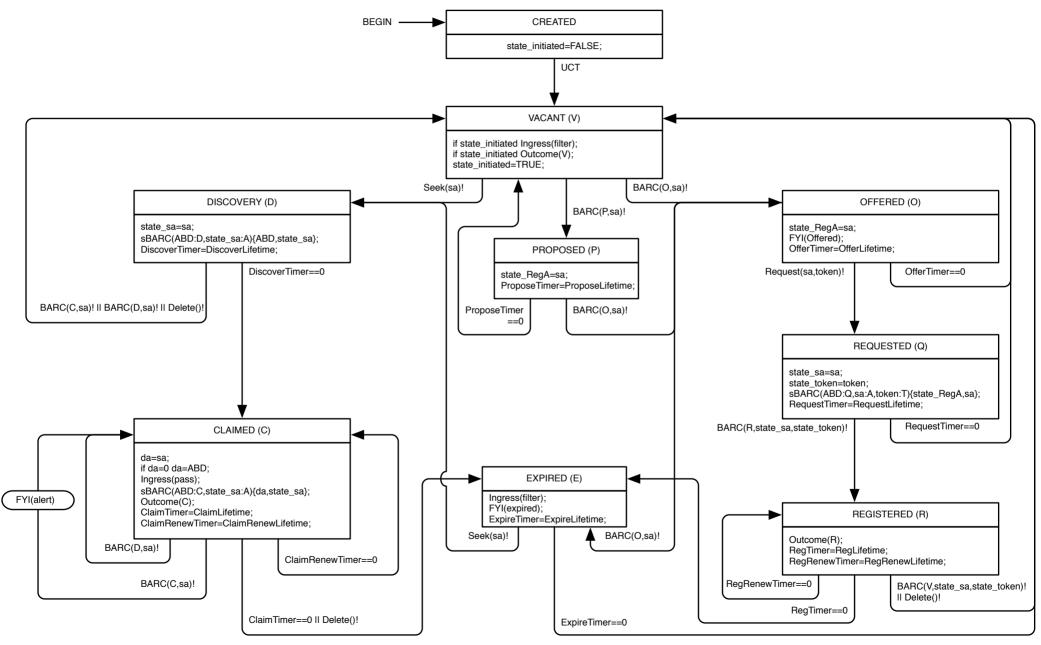
Inquiry followed by Registrar Offer



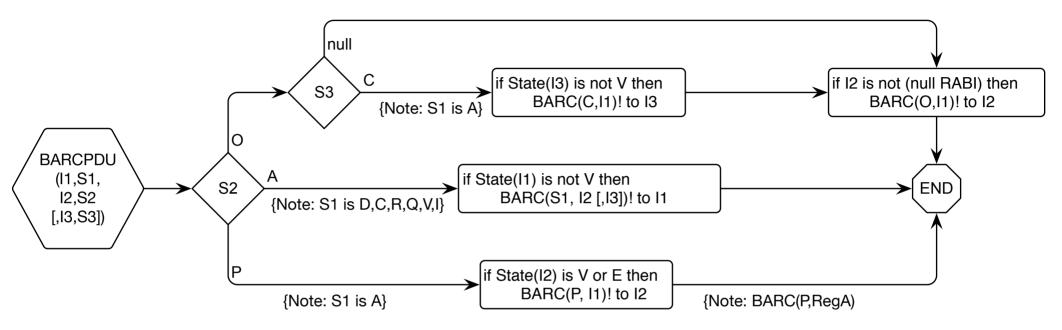
Inquiry followed by Advisor Proposal



ABD Claimant: ABD State Machine

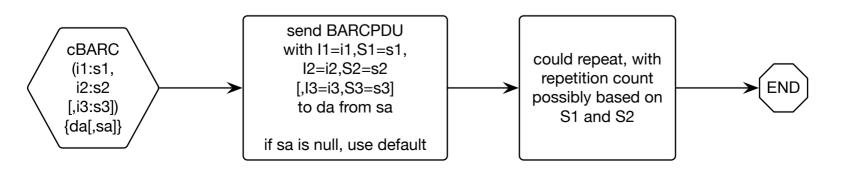


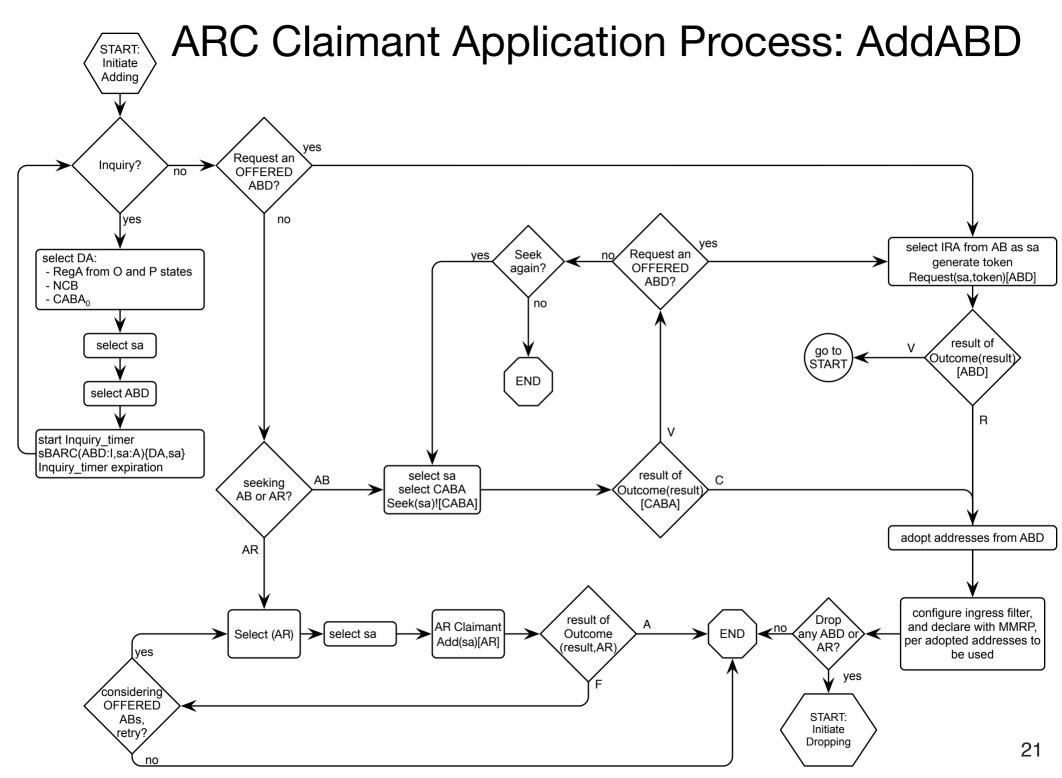
ARC Claimant: cBARCPDU_in



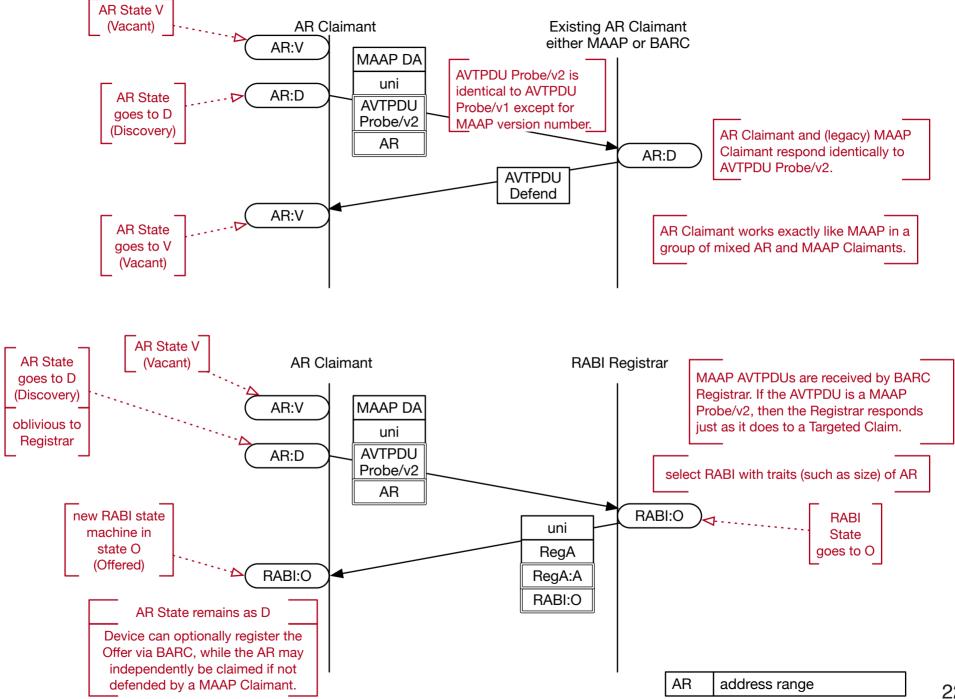
{Note: square-bracketed parameters are sometimes absent.}

ARC Claimant: cBARCPDU_out

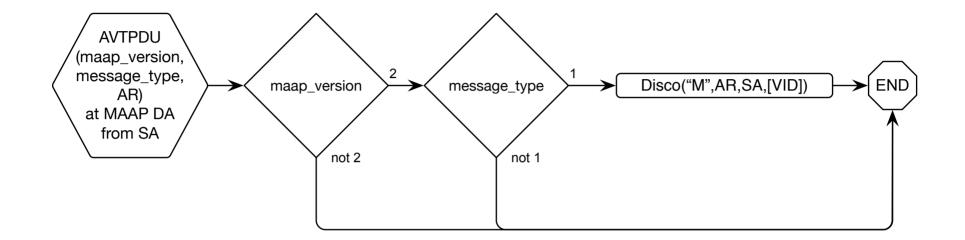




AR Claimant Procedure



BARC Registrar: AVTPDU Processor



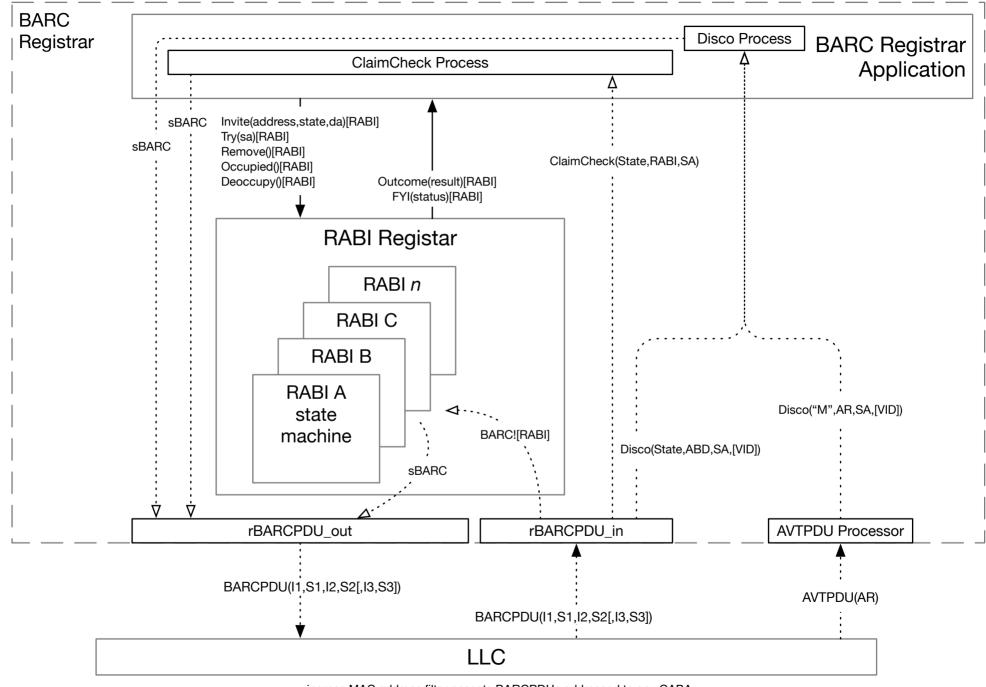
AR State Transition Table

	State					
Event	VACANT (V)	DISCOVERY (D)	ACQUIRED (A)			
Add(sa)!	sMAAP(Begin(AR,sa)!) DISCOVERY					
rMAAP(AR:Defend)!		Outcome(A,AR) ACQUIRED				
rMAAP(AR:Initial)!		Outcome(F,AR) VACANT	Outcome(X)[AR] VACANT			

rMAAP(AR:State!) invokes an event at the state machine when the MAAP state changes to State

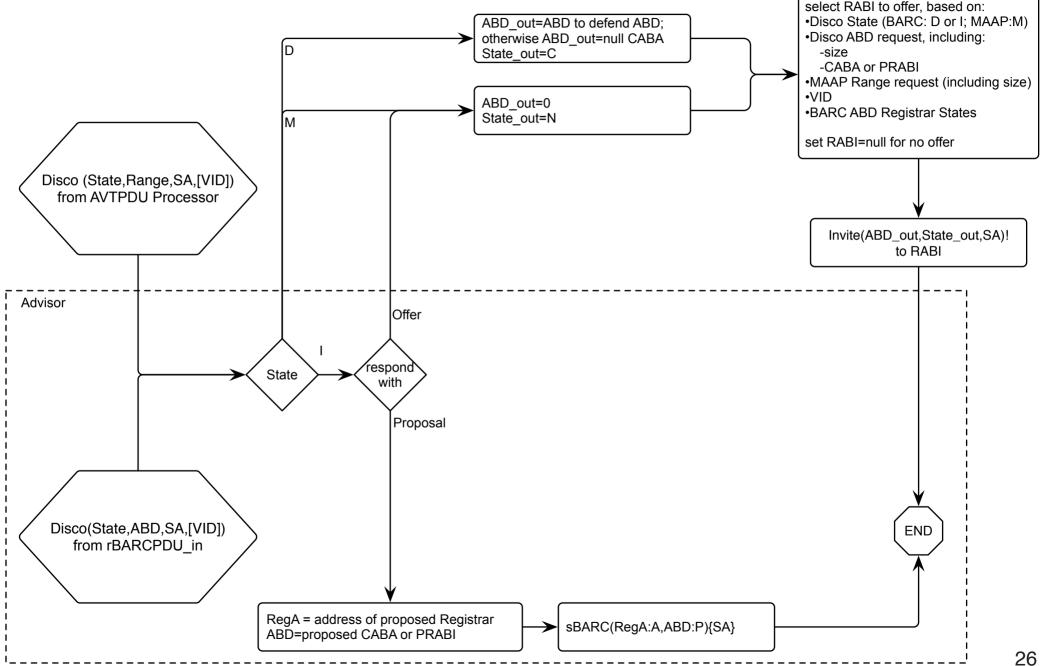
sMAAP(Action!) invokes Action! event at MAAP state machine

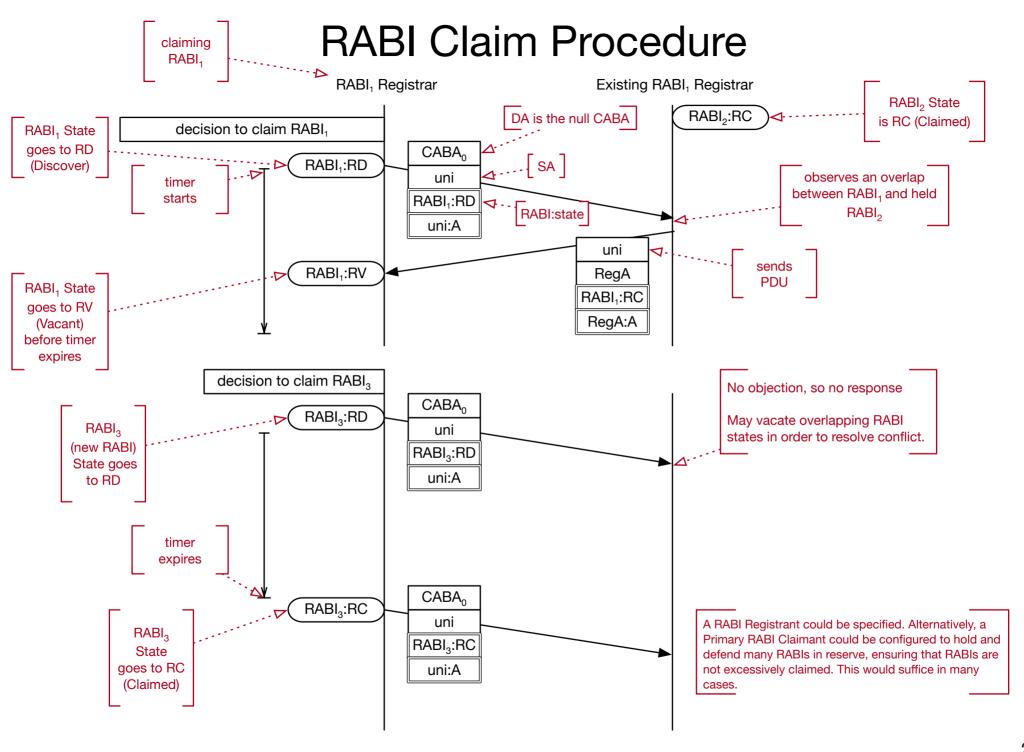
BARC Architecture – Registrar



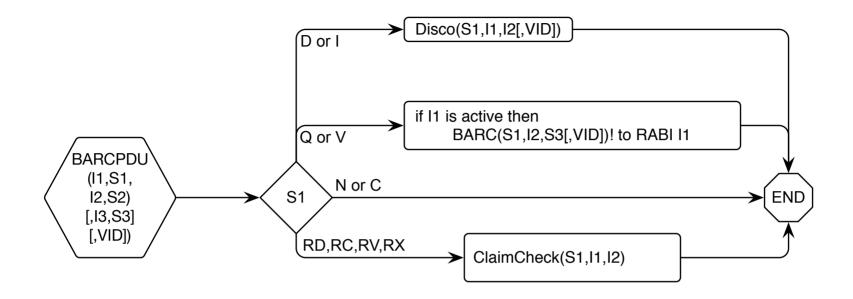
ingress MAC address filter accepts BARCPDUs addressed to any CABA and MAAP multicast address for AVTPDUs

BARC Registrar Application: Disco Process

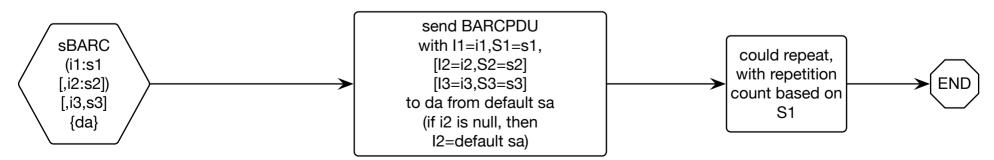




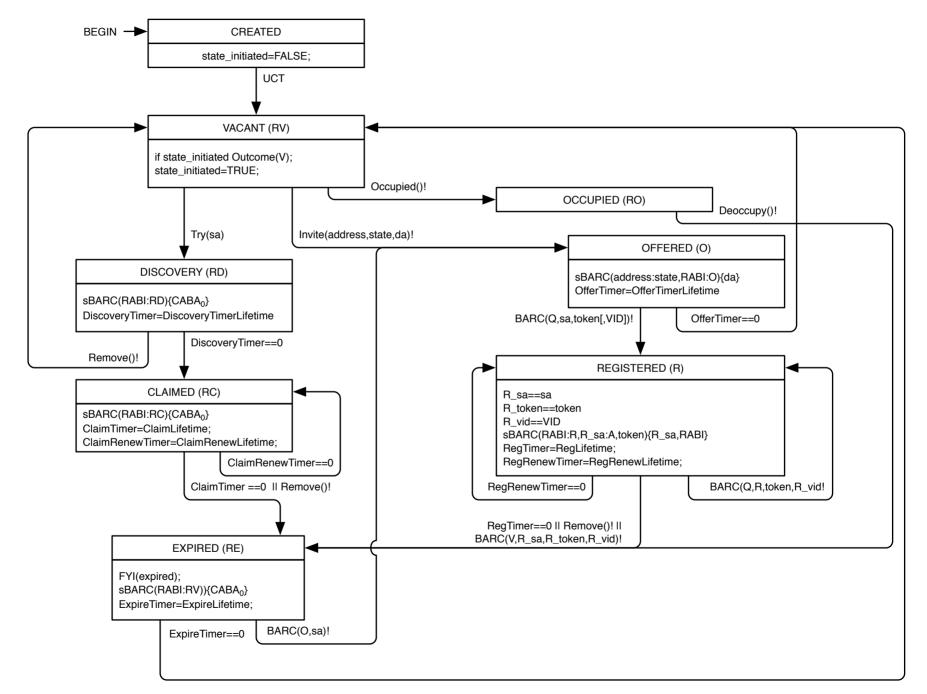
BARC Registrar: rBARCPDU_in



BARC Registrar: rBARCPDU_out



Registrar: RABI State Machine



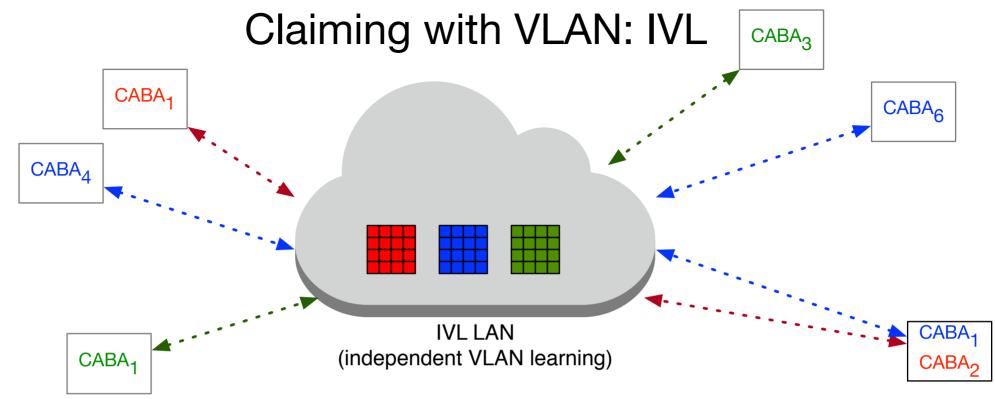
VLANs

• All address assignments are specific to the VLAN (or VLANs) in which messaging is communicated and under which the assignment was completed.

• Usage of any address may need to be limited to the VLAN or VLANs under which it was obtained.

• Due to the possibility that the same unicast address may be assigned in different VLANs, Independent VLAN Learning (IVL) may be needed in bridges, per IEEE Std 802.1Q Annex F (F.1.2).

- This requirement could be relaxed in some cases
- e.g. when assigned unicast addresses are declared via MMRP (instead of learning)
- This issue is common to claiming protocols generally.
- Some approaches follow.



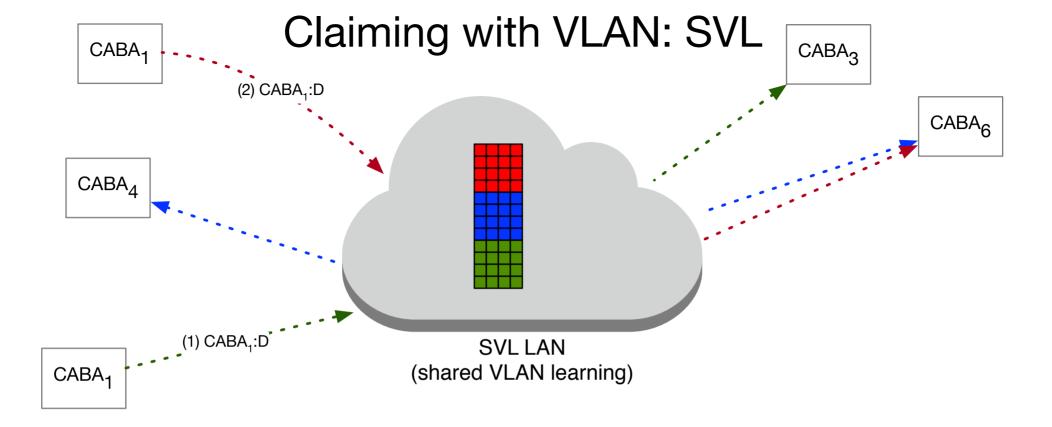
IEEE Std 802-2014 says "Local MAC addresses need to be unique on a LAN or bridged LAN unless the bridges support VLANs with independent learning."

With IVL, each VLAN has an independent forwarding table. -but IVL is not always possible

BARC claiming on each VLAN is independent

a duplicate address may occur in more than one VLAN; that is not harmful if managed carefully

A claimant with multiple VLANs needs to claim in each VLAN. Claimed address is usable only in claimed VLAN: Claimant needs to bind address to VLAN For convenience, Claimant may claim the same address in each of its VLANs -Still, requires multiple claim messages and multiple forwarding table entries. -Device needing many VLANs should consider an EUI

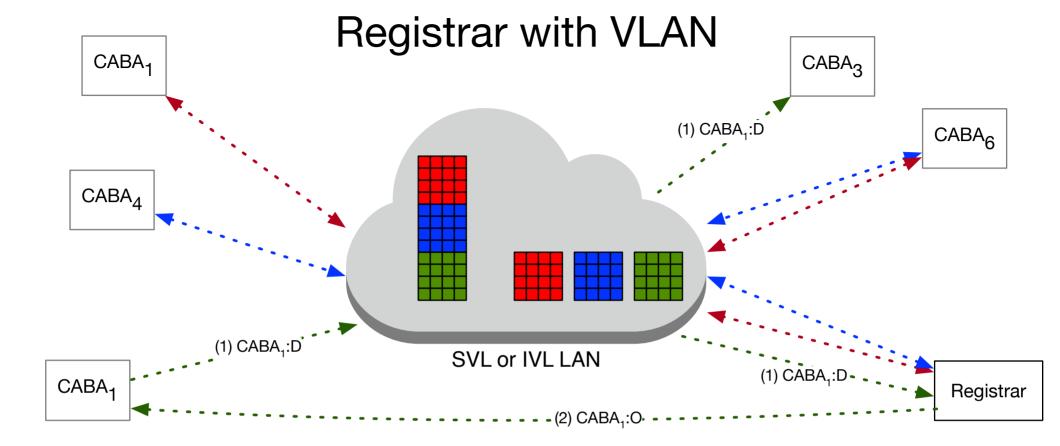


With SVL, VLANs share a forwarding table.

BARC claiming on each VLAN is independent

an address could become a duplicate, existing in more than one VLAN

forwarding table is limited to one entry per address, so duplication is catastrophic.



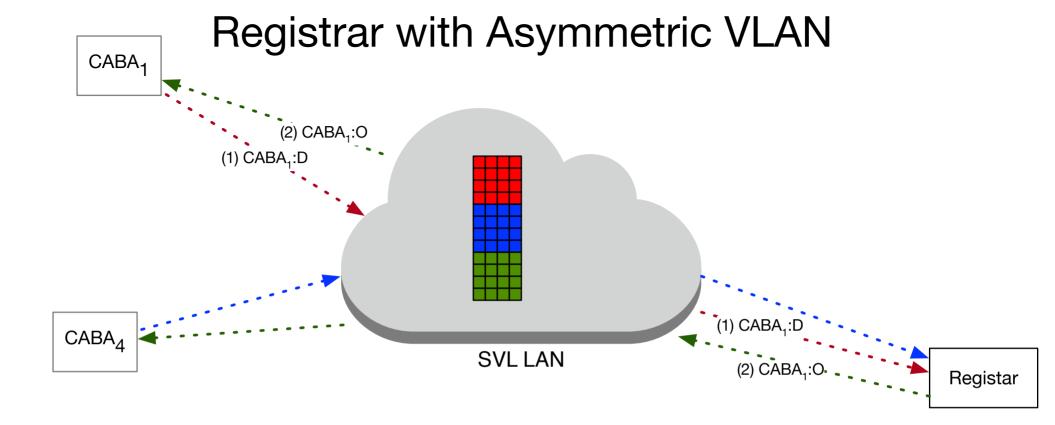
Network is configured with Registrar on all active VLANs on which BARC is used.

BARC claim from any VLAN is delivered to Registrar.

-Offer delivered on Claimant's VLAN

Registrar ensures that registered address is unique across all (or perhaps only some) of its VLANs. -SVL or IVL will work

Registrar needs to remember over which VLANs the address was assigned. -should be retained in State Machine



SVL is used for Asymmetric VLAN (IEEE Std 802.1Q Annex F.1.3)

Registrar can assign address to be unique across all VLANs available to the Registrar.

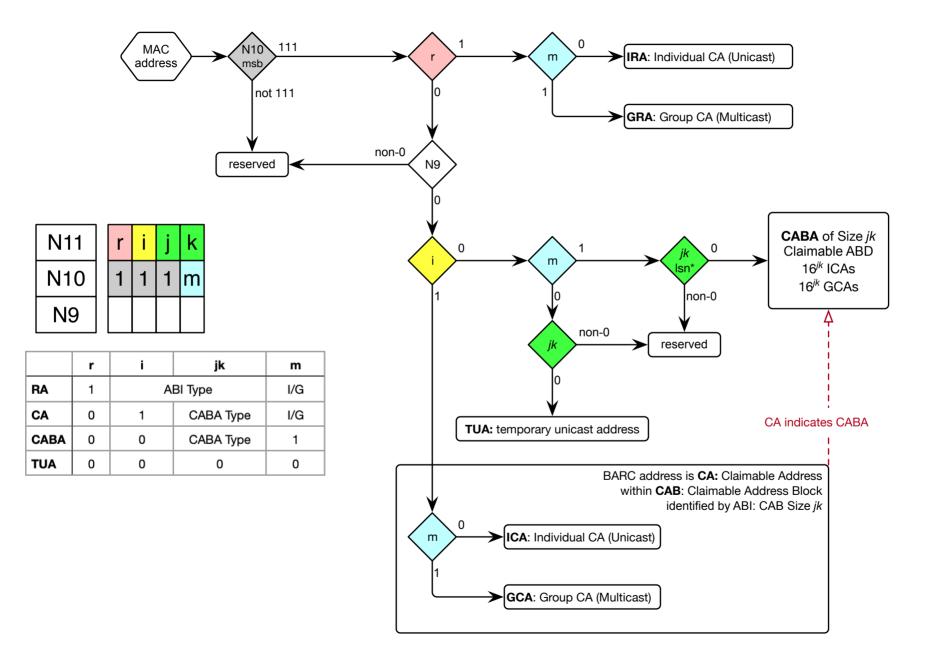
Summary

- Claimants operate with or without Registrars.
- Multiple registrars are supported, holding claims of disjoint RABIs.
- The block discretization provides:
 - a vast set of addresses to a LAN
 - a large set of temporary unicast addresses
 - operational efficiency and simplicity
 - both unicast and multicast addresses to Claimant
 - unicast and multicast subblocks share the same range, except for the m bit
 - could be exploited
 - devices needing both unicast and multicast addresses need make only one claim
- Could integrate with MMRP to limit propagation and eliminate learning of unicast AB content.
 - MMRP needs to efficiently handle address ranges
 - BARP could be specified as alternative MRP application
 - (e.g. would understand an ABD)

Appendix 1

• additional details on BARC addresses and identifiers

BARC Address Parsing



CABA and CA, CAB Size 0-3

CAB Si	ze C=0	CAB Siz	ze C=1	CAB Si	ze C=2	CAB Si	ze C=3
CABA	CAB	CABA	CAB	CABA	CAB	CABA	CAB
0000	0 1 0 0	0 0 0 1	0 1 0 1	0 0 1 0	0 1 1 0	0 0 1 1	0 1 1 1
1 1 1 1	1 1 1 *	1 1 1 1	1 1 1 *	1 1 1 1	1 1 1 *	1 1 1 1	1 1 1 *
0000	0 0 0 0	0 0 0 0	0000	0 0 0 0	0 0 0 0	0 0 0 0	0000
X8							
X7							
X6							
X5							
X4							
X3							
X2	X2	X2	X2	X2	X2	0	*
X1	X1	X1	X1	0	*	0	*
X0	X0	0	*	0	*	0	*

2 contiguous subblocks per CABA (one unicast, one multicast)

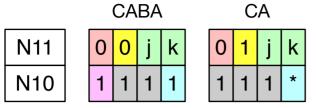
6.9E10 Size 0 CABAs
1 CA/subblock

• 4.3E9 Size 1 CABAs • 16 CAs/subblock • 2.7E8 Size 2 CABAs

256 CA/subblock

1.7E7 Size 3 CABAs
4096 CAs/subblock

CABA/CAB Math



 CABA indicates the CAB CA indicates the CABA

•C=ik indicates the CAB Size

•each CAB subblock includes 16^c contiguous addresses

&	bitwise AND
I	bitwise OR
~	bitwise NOT
/	divide

 $Cmask(C) = \sim (0 \times 410000000000 + 0 \times 10^{**}C - 1)$ [CABA mask, per Size; used to create CABA from CA1 $CABA(CA) = CA\&Cmask(C(CA)) + 0 \times 010000000000$

Example: [Note: Underlining on the middle four nibbles is shown only as a reading aid.]

- CA = 0x6F0123456789 = 0110-1111-0000-0001-0010-0011-0100-0101-0110-0111-1000-1001
- $C(CA) = 0 \times 200000000000 / 0 \times 10000000000 = 2$
- Cmask $(0 \times 2) = \sim (0 \times 410000000000 + 0 \times 0100 1) = \sim (0 \times 4100000000FF) = 0 \times BEFFFFFF00$
- CABA(CA) = CA&0×BEFFFFFF00 + 01000000000 = 0×2F0123456700

A CA is within CABA, if and only if CABA(CA) = CABA,

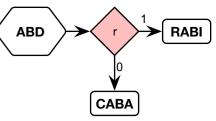
•this requires identical CAB Size

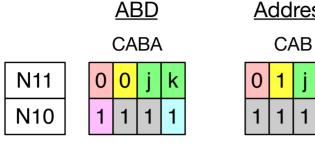
The CAB of CABA, is the set of all CAs that satisfy CABA(CA) = CABA,

Lowest ICA in CABA: ICAmin = (CABA-0x01000000000) | 0x400000000000 (example: 0x6E0123456700)Highest ICA in CABA: ICAmax = ICAmin + 0x10**C(CABA) - 1 (example: 0x6E01234567FF)

Lowest GCA in CABA: GCAmin = CABA $| 0 \times 410000000000$ (example: $0 \times 6F0123456700$) Highest GCA in CABA: GCAmax = GCAmin + $0 \times 10^{**}$ C(CABA) - 1 (example: $0 \times 6F01234567$ FF)

ABD (CABA/RABI) Format









•CAB indicates the CABA

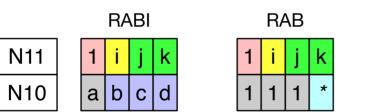
•ik indicates the CABA size C

•the C least significant nibbles of the CABA are 0

•each CAB subblock includes 16^c contiguous addresses)

each CAB includes a unicast subblock and a multicast subblock

•no CA within a CAB is within any other CAB (that is, a CAB with a different CABA



•RABI is an ABD (indicating RAB) and never used as a MAC address

•RAB does not fully indicate the RABI

•i indicates RABI Option (Options 0 and 1 provide independent RABIs/RABs) •*ik* indicates the BABI Size B

- "BABI" for "Basic Address Block Indicator"

•abcd indicates MABI Size M: those bits are not in the RAB addresses

- "MABI" for "Multiple Address Block Indicator"

•RAB Size R=B+M

•the R least significant nibbles of the RABI are 0

•each RAB subblock includes 16^R contiguous addresses

•each RAB includes a unicast subblock and a multicast subblock

RABI Aggregation

•RAB Size R=B+M

• for 48-bit addresses, set a=0; then $R = ik + bcd \le 3 + 7 = 10$, matching the 10 available nibbles N0-N9

• could use ijk as the BABI Size, and/or the full abcd as the MABI Size; e.g., for 64-bit addresses

•A RABI may aggregate other RABIs.

•A RABI of RAB Size R and MABI Size M can be disaggregated into:

- •16 RABIs of RAB Size *R*-1 (MABI Size *M*-1), or
- •16² RABIs of RAB Size *R*–2 (MABI Size *M*–1), or
- •16ⁿ RABIs of RAB Size *R*–n (MABI Size *M*–n), or
- ... 16^M RABIs of RAB Size B (MABI Size 0), or

•A RABI of RAB Size B (MABI Size 0) cannot be disaggregated. •An RA appears in one and only RABI of each M.

RABI and RA, MABI Size 0, BABI Size 0-3

BABI	Size 0	BABI Size 1		BABI Size 2			BABI Size 3		
RABI	RAB	RABI	RAB	RABI	RAB		RABI	RAB	
1 i 0 0	1 i 0 0	1 i 0 1	1 i 0 1	1 i 1 0	1 i 1 0		1 i 1 1	1 i 1 1	
0000	1 1 1 *	0000	1 1 1 *	0000	1 1 1 *		0000	1 1 1 *	
X9	X9	X9	X9	X9	X9		X9	X9	
X8	X8	X8	X8	X8	X8		X8	X8	
X7	X7	X7	X7	X7	X7		X7	X7	
X6	X6	X6	X6	X6	X6		X6	X6	
X5	X5	X5	X5	X5	X5		X5	X5	
X4	X4	X4	X4	X4	X4		X4	X4	
X3	X3	X3	X3	X3	X3		X3	X3	
X2	X2	X2	X2	X2	X2		0	*	
X1	X1	X1	X1	0	*		0	*	
X0	X0	0	*	0	*		0	*	

2 contiguous subblocks per RABI (one unicast, one multicast)

1.1E12 Size 0 RABIs
1 RA/subblock

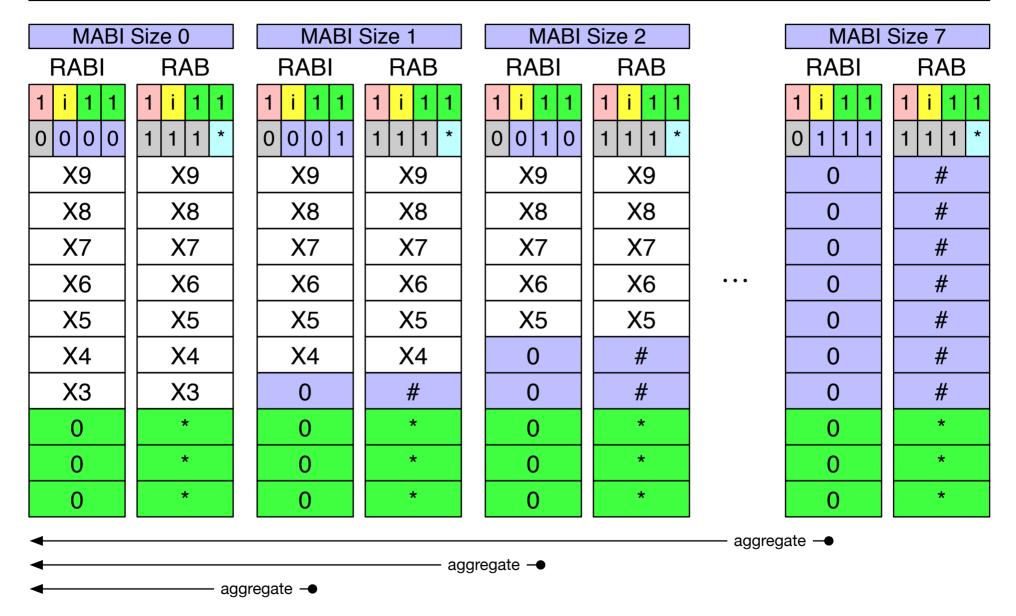
6.9E10 Size 1 RABIs
16 RAs/subblock

• 4.3E9 Size 2 RABIs • 256 RA/subblock

2.7E8 Size 3 RABIs
4096 RAs/subblock

Aggregation Example: BABI Size 3, various MABI Sizes

BABI Size 3



Example:

Hierarchical RABI Addressing with common Registrar

common RAB prefix

_										
		MA	BI	ç	Siz	e 2	2			
	RA	B	l			R/	٩B	,		
1	i	1	1		1	i	1	1		
0	0	1	0		1	1	1	*		
	Х	9				Х	9			
	Х	8			X8					
	X7				X7					
	X6				X6					
	X5					X5				
	0				#					
	0					#				
0					*					
	()			*					
	()			*					

Held by Registrar

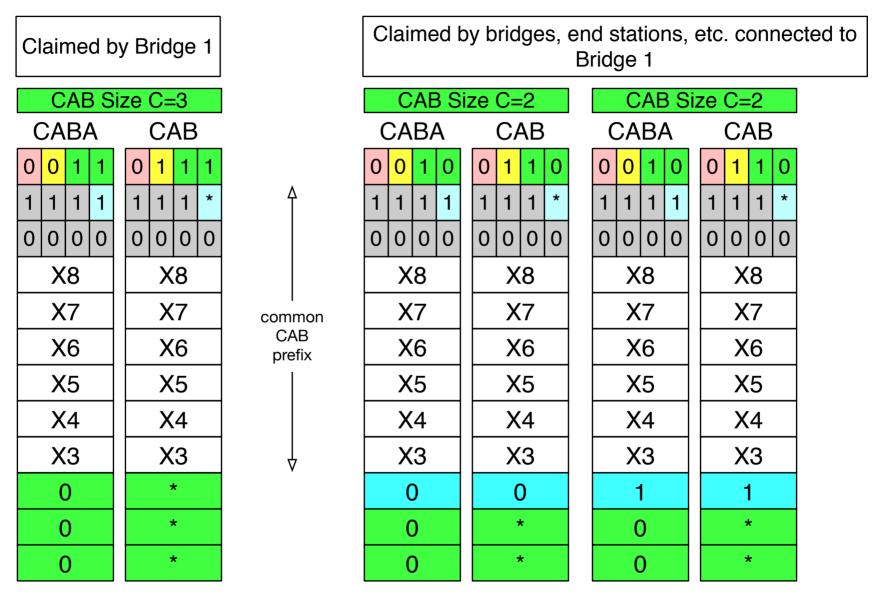
MABI Size 1 RABI RAB 1 i 1 i 1
$\begin{array}{c c c c c c c c c c c c c c c c c c c $
$\begin{array}{c c c c c c c c c c c c c c c c c c c $
X9 X9 X8 X8 X7 X7 X6 X6 X5 X5
X8 X8 X7 X7 X6 X6 X5 X5
X7 X7 X6 X6 X5 X5
X6 X6 X5 X5
X5 X5
8-15 8-15
0 #
0 *
0 *
0 *

Assigned by Registrar to bridges, end stations, etc. connected to Bridge 1

	-							
MABI	Size 1	MABI Size 1						
RABI	RAB	RABI	RAB					
1 i 1 1	1 i 1 1	1 i 1 1	1 i 1 1					
0001	1 1 1 *	0001	1 1 1 *					
X9	X9	X9	X9					
X8	X8	X8	X8					
X7	X7	X7	X7					
X6	X6	X6	X6					
X5	X5	X5	X5					
0	0	1	1					
0	#	0	#					
0	*	0	*					
0	*	0	*					
0	*	0	*					

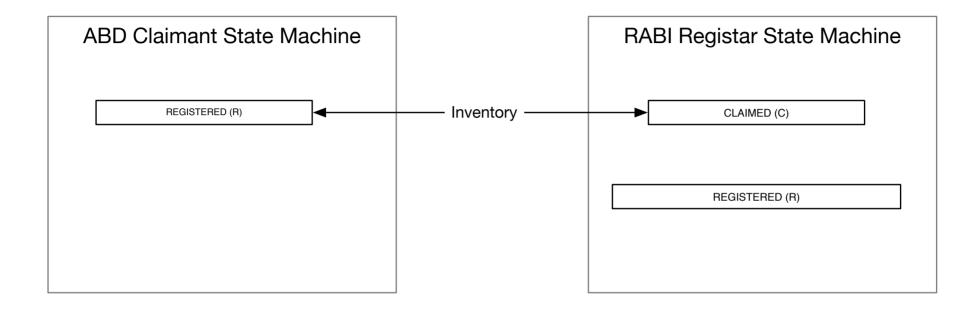
Bridge serves as Advisor to connected station; proposes PRABI, and Registrar Address, in response to an Inquiry.

Example: Hierarchical CABA Addressing



Bridge serves as Advisor to connected stations; proposes CABA in response to an Inquiry.

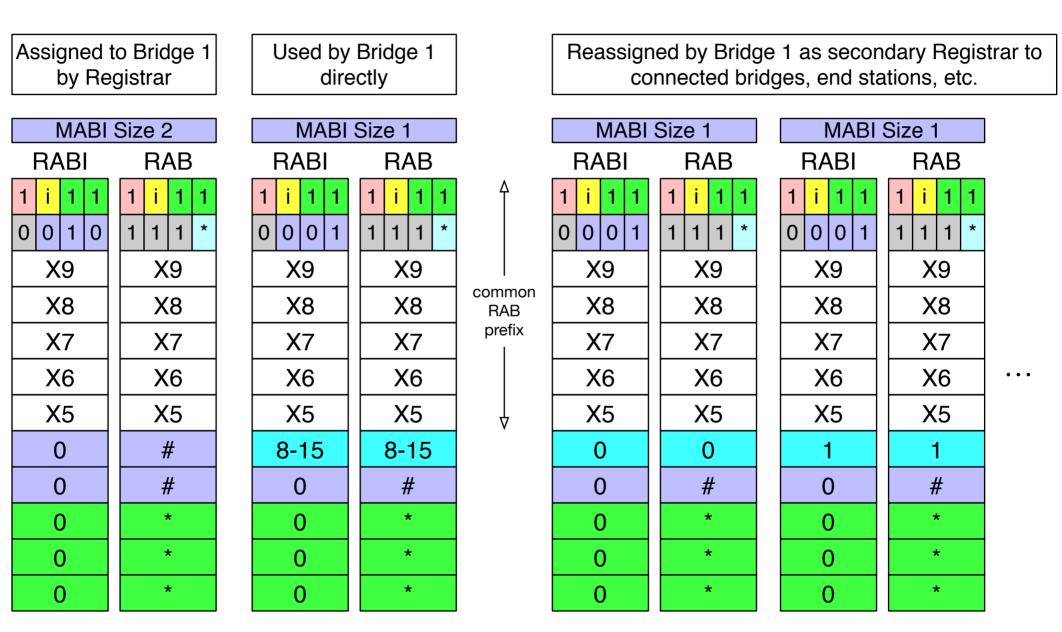
Claimant as Registrar

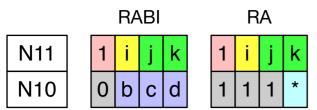


RABI registered as a Claimant could be disaggregated and reassigned by a Registrar function managed jointly with the Claimant.

A RABI in the "REGISTERED" state of the ABD Claimant State Machine could be considered to be in the Inventory of the RABI Registrar State Machine (along with Claimed RABIs) and could be disaggregated, Offered and Registered by that Registrar.

Example: Hierarchical RABI Disaggregation with Tiered Registrars





RABI Math 101

ik indicates the BABI Size B *bcd* indicates the MABI Size M

&	bitwise AND
۲	bitwise NOT
/	divide
==	binary equality

- i(X) = (X&0×40000000000) / 0×40000000000 [RABI Option, when X is RABI or RA]
- B(X) = (X&0x30000000000) / 0x10000000000 [BABI Size, when X is RABI or RA]
- M(X) = (X&0×07000000000) / 0×010000000000 [MABI Size when X is a RABI]
- R(X) = B(X) + M(X) [RAB Size, when X is a RABI]

 $Rmask(N) = \sim (0 \times 0F 000000000 + 0 \times 10^{**}N - 1)$ [mask, used below]

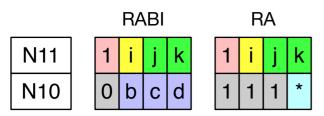
An RA is within RABIx if and only if RA&Rmask(R(RABIx)) = RABIx&Rmask(R(RABIx))

this requires identical RABI Option and BABI Size

The RAB of RABIx is the set of all RAs that satisfy RA&Rmask(R(RABIx)) = RABIx&Rmask(R(RABIx))

Example:

- $\bullet \mathsf{RABI} = 0 \times \mathsf{F20123400000} = 1111 0010 0000 0001 0010 0011 0100 00000 00000 0000 0000 -$
- M(RABI) = 0x020000000000 / 0x01000000000 = 2
- R(RABI) = 3+2 = 5
- Rmask(5) = $\sim (0 \times 0F000000000 + 0 \times 10^{**}5 1) = \sim (0 \times 0F00000FFFFF) = 0 \times F0FFFF000000$
- RABI&Rmask(5) = 0xF00123400000
- $\bullet \mathsf{RA} = 0 \times \mathsf{FE0123456789} = 1111 1110 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 0000 0001 00000 0000 0000 0000 0000 00$
- RA&Rmask(5) = 0×F00123400000 = RABI&Rmask(5)
- so RA is within RABI



RABI Math 102

ik indicates the BABI Size B*bcd* indicates the MABI Size M

&	bitwise AND
2	bitwise NOT
/	divide
==	binary equality

RABIcheck(RABI1,RABI2) determines whether RABI RABI1 overlaps RABI2.

RABIcheck(R1,R2) = R1&Rmask(R(R2))

RABI1 shares RAs with RABI2 if and only if:

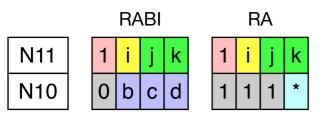
RABIcheck(RABI1, RABI2) = RABIcheck(RABI2, RABI1)

•Note: This can be ruled out by inspection if the two N11 nibbles differ

Example:

- B(RABI2) = 0x300000000000 / 0x1000000000 = 3
- M(RABI2) = 0×070000000000 / 0×10000000000 = 5
- R(RABI2) = 3+5 = 8
- Rmask(8) = 0×F0FF00000000
- RABI1 = $0 \times F20123400000$
- RABI1&Rmask(8) = $0 \times F0010000000$
- RABI2&Rmask(5) = $0 \times F0010000000$
- so RABI1 and RABI2 have RAs in common

```
Example: 0xFF01234FFFFF)
```



RABI Math 103

ik indicates the BABI Size B *bcd* indicates the MABI Size M

&	bitwise AND
۲	bitwise NOT
/	divide
==	binary equality

An RA exists in one and only one RABI (called $RABI_W$) of each MABI Size *W*. Given the MABI Size *W*, what is $RABI_W$?

A RABI exists in one and only one aggregated RABI called RABI) of each larger MABI Size W. Given the MABI Size W, what is RABI_W?

 $RABI_{W}(X, W) = X\&Rmask(W+B(X)) + W^*0 \times 010000000000$ [X is an RA, or a RABI with M(X)<W]

Example:

- RABI1 = 0xF20123400000 B(RABI1)=3; M(RABI1)=2
- *W* = 5; *W* + B = 8
- RABI1&Rmask(8) = 0×F20123400000&0×F0FF00000000 = 0×F00100000000
- RABIm(RABI1,5) = 0×F00100000000 + 0×050000000000 = 0×F50100000000

Example:

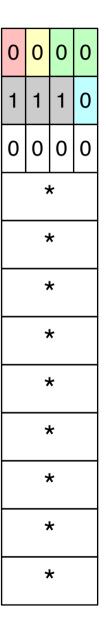
- RA = 0×FE0123456789 [B(RA)=3]
- *W* = 5; *W* + B = 8
- RA&Rmask(8) = 0xFE0123456789&0xF0FF00000000 = 0xF00100000000
- RABIm(RA,8) = 0×F00100000000 + 0×05000000000 = 0×F50100000000

Example:

- RA = 0xFE0123456789 [B(RA)=3]
- W = 2; W + B = 5
- RA&Rmask(5) = 0×FE0123456789&0×F0FFFF00000 = 0×F00123400000
- RABIm(RA,5) = 0×F00123400000 + 0×02000000000 = 0×F20123400000

Temporary Unicast Address (TUA)

- For temporary use
- device without a source address selects a random temporary unicast address for initial discovery <u>only</u>
 - protocol then assigns at least one persistent unicast address
- simultaneous duplicate temporary addresses may lead to message loss in some circumstances
 - network learns route to source as initial message crosses the network
 - before response is returned, another initial message with duplicate source address crosses the path and rewrites the route
 - unlikely to be disastrous
 - loss of initial message will be corrected eventually
- · nevertheless, need to consider the likelihood of duplication
- Temporary address range includes 9 full nibbles of 16 values each (0-F)
 - $-16^9 = 68,719,476,736$ (= *N*) temporary addresses in the pool
 - chance of no duplicates with k randomly selected addresses is approximated exp(-k*(k-1)/(2*N))
 - with k=1000 devices <u>simultaneously</u> using a temporary address, chance of no duplicates is ~0.99988
 - can add *jk* bits to the pool if more entropy is needed



null CABA (CABA₀)

CABA₀

• Null CABA (CABA₀) is not an assignable address.

- Registrar listens to $CABA_0$.
- No Claimant listens to $CABA_{0}$.
- Can be used as the DA of BARC Inquiry; e.g., when Registrar address is unknown.

	0	0	0	0					
	1	1	1	1					
		()						
		()						
F	0 0 0 0 0 0 0 0 0 0 0								
	0								
	0								
	0								
	0								
	0								
		()						
		()						

Proposed RABI (PRABI) and null RABI/PRABI

 Usable only as the content of a **BARC Inquiry or BARC Proposal** message.

 Indicates a set of RABIs characterized by RABI Option *i*, BABI Size B *jk*, and MABI Size M bcd.

• If the PRABI has the form of a RABI (with 0 in the B+M least significant nibbles), then the PRABI set is only that RABI.

 Non-zero values in the lower B+M nibbles can signify that some bits of the higher nibbles are "don't care."

 For example, the lower B+B nibbles of the PRABI could form a bitmask of the higher nibbles.

PRABI					proposed RABIs					
1	i	1	1		1 i 1 1					don't
0	0	1	0		1 1 1 *					care:
	X9					X	9			F
	Х	8				Х	8			0
	X7					X 7				F
	X6					X6				0
	X5				X5					0
F					0					
0					0					
F			0							
	()				0				
0						()			

• As a PRABL in a **BARC Inquiry or** Proposal message. indicates a set of **RABIs characterized** by RABI Option *i*, BABI Size B *jk*, and MABI Size M bcd, without expressing a preference for the RABI values of the 0 nibbles.

• As a RABI in a BARC Offer message. indicates to Claimant that no RABI is offered.

Proposed RABIs are those that satisfy: RABI | Pmask = (PRABI&~(0x10**(R(PRABI)) - 1)) | Pmask where $Pmask(PRABI) = 10^{**}(R(PRABI))^*PRABI&(0x10^{**}(Rcap) - 1)$ and Rcap=min(R(PRABI), 10-R(PRABI))

Appendix 2

additional procedural details

BARC Address Propagation with MMRP

The ARC Claimant Application AddABD Process includes "declare with MMRP". This entails declaring, to MMRP (when available), MMRP attributes, using an MMRPDU per IEEE Std 802.1Q § 10.12.1.6:

- The multicast address represented by the CABA
 - FirstValue field = CABA/NumberOfValues=1
- The unicast address subblock indicated by the ABD (CABA or RABI)
 - FirstValue field = first address in unicast subblock
 - NumberOfValues = 16^{Size} per ABD Size
 - limited to Size up to 3: 16³=4096, and MRP provides 13 bits of NumberOfValues (2¹³=8192)

The ARC Claimant Application AddABD Process includes ("select CABA"); this selection should consider any local MMRP registration database to avoid selecting a registered CABA.

Unicast MMRP declaration can be useful because:

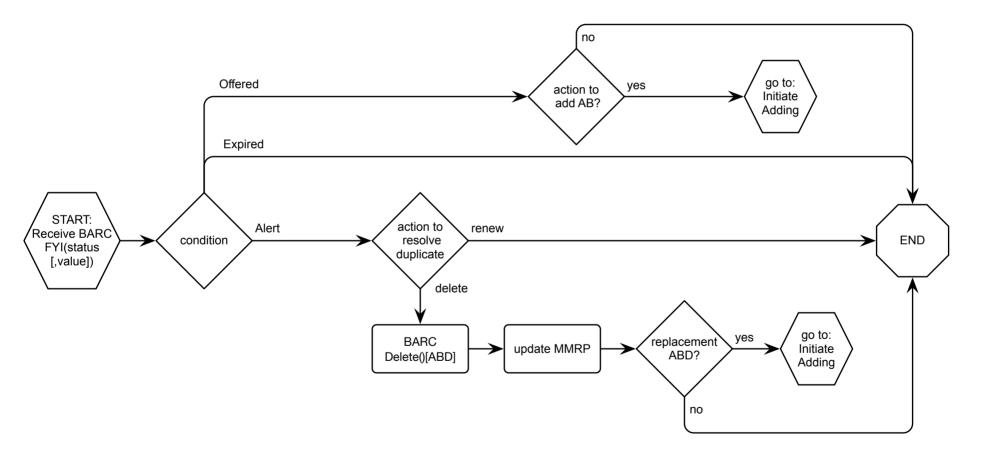
- (1) A one-step declaration covers a contiguous range of self-assigned unicast addresses.
- (2) Eliminates flooding for all the unicast addresses in the assignment.
- (3) Eliminates the need for learning of each unicast address when used.
- (4) Precludes erroneous re-learning of an address when a false duplicate is used elsewhere in the network.
 - Could be a way to control duplication.
 - Security issues to study.

BARC could alternatively specify "BARP," a new MRP application. This could entail the following changes:

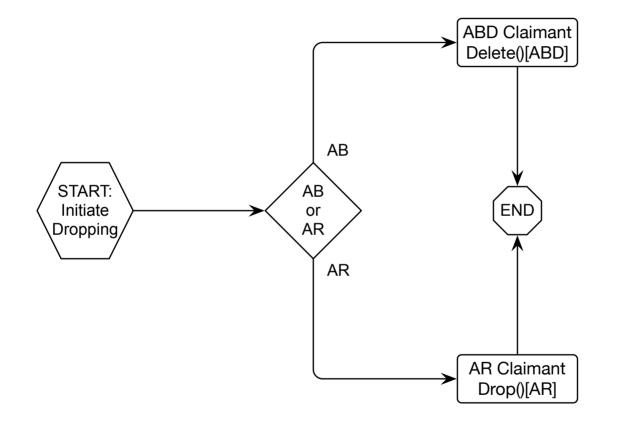
(a) the BARP application would be enabled to Join and Leave with the ABD as the declared attribute
(b) the BARP application would be specified to understand the semantics of the ABD and extract from it the indicated ABD multicast address and the indicated unicast address set, then use it to populate the FDB
(c) In the BARC Claimant ABD State Machine, the CABA claim ["sBARC(CABA:C)]" might not be needed, since the a BARP declaration could convey the claim to the CABA as well as the declaration of interest in receiving at the CABA multicast address

BARP might be better suited to specification within IEEE Std 802.1Q instead of 802.1CQ.

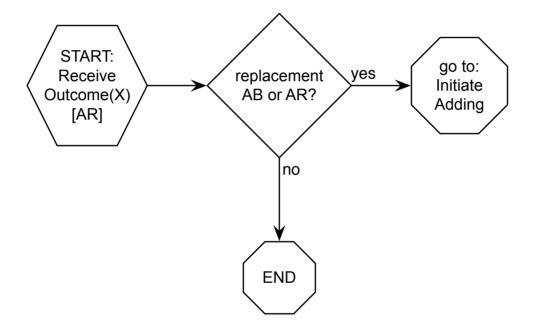
ARC Claimant Application Process: BARC Management



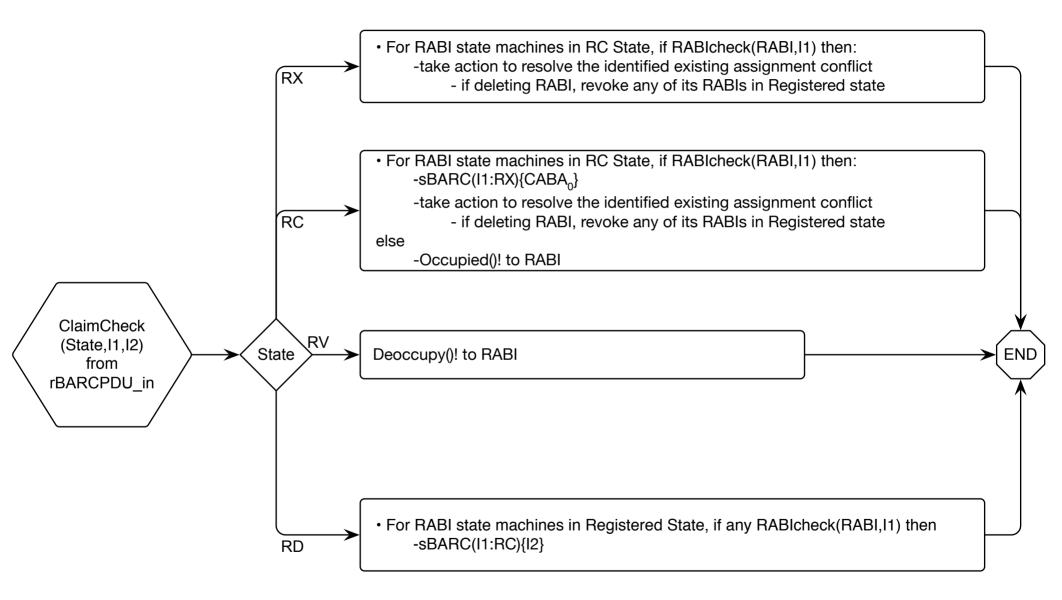
ARC Claimant Application Process: Drop Claim



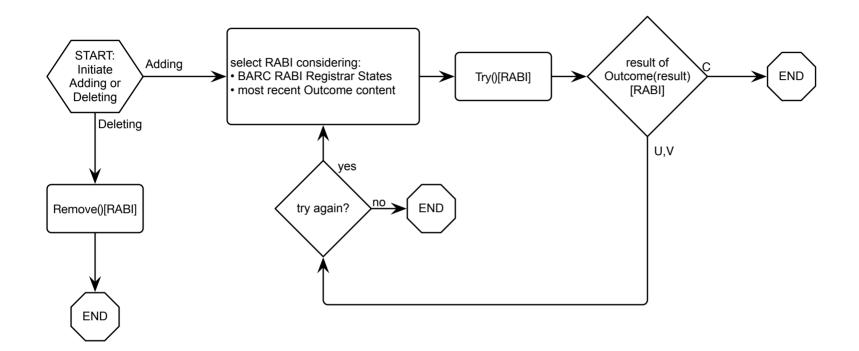
ARC Claimant Application Process: MAAP Management



BARC Registrar Application: ClaimCheck Process



BARC Registrar Application: RABI Claiming



RABI Claimant Application Process: BARC Management

