

# Comments for P802.1Qdq



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# Comments #1

- Current text

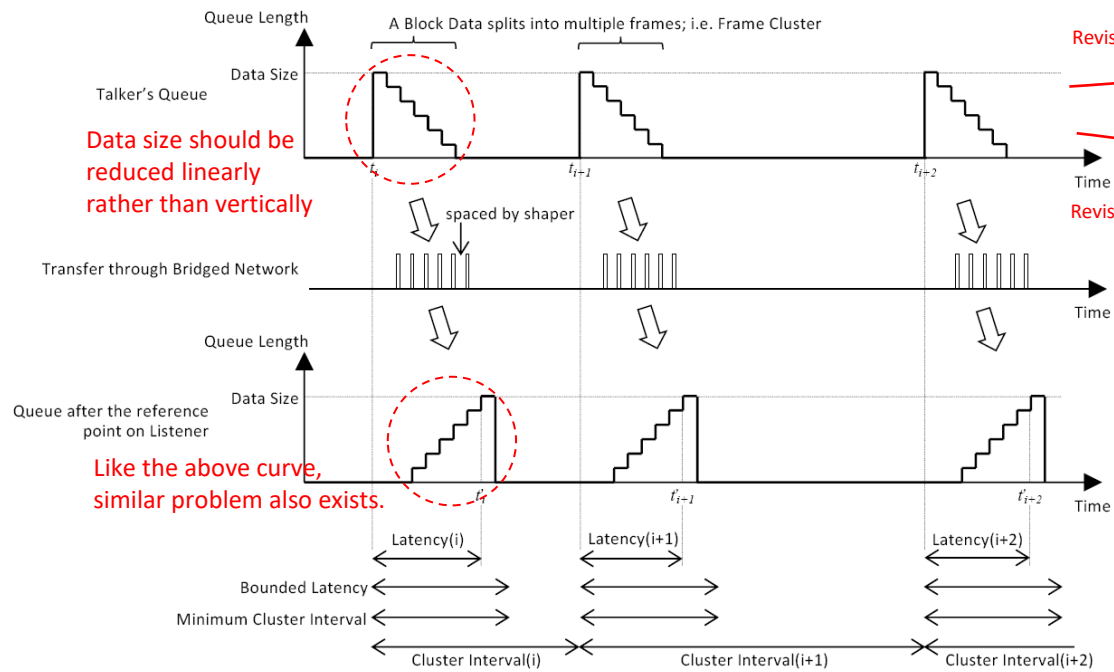
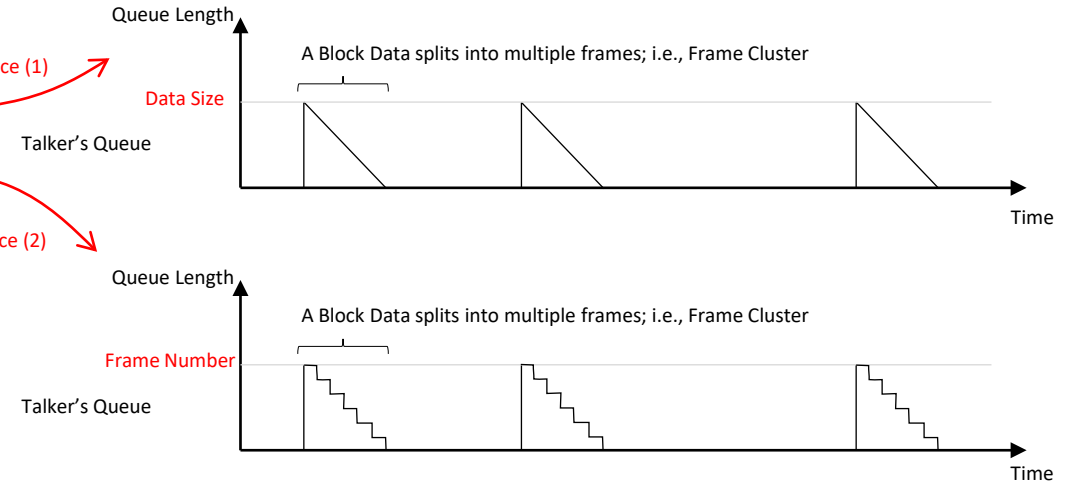


Figure Y-3— Traffic pattern in an application's point of view

- Suggested text



# Comments #2

- Current text

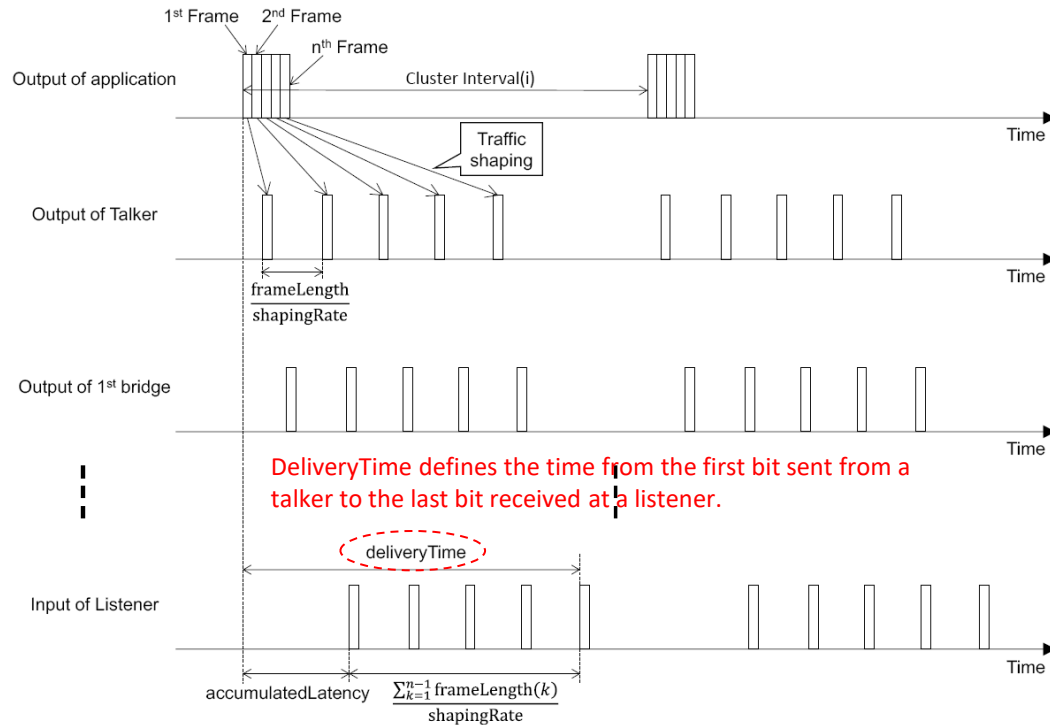


Figure Y-4— Frame propagation from Talker to Listener

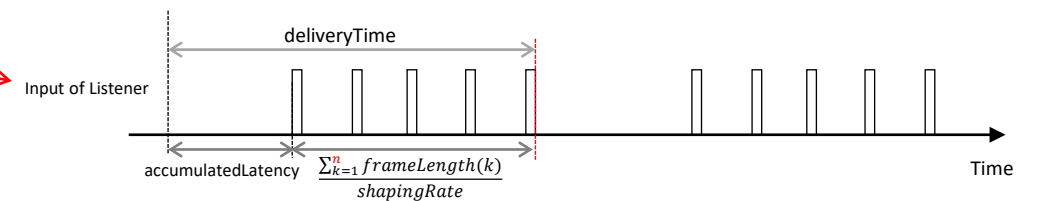
Accordingly, the formulas (Y-2) & (Y-3) should be modified

$$deliveryTime(i) = \frac{\sum_{k=1}^{n-1} frameLength(n)}{shapingRate} + accumulatedLatency \quad (Y-2)$$

$$requiredMinimumShapingRate = \frac{\sum_{k=1}^{n-1} frameLength(k)}{targetLatency} \quad (Y-3)$$

$$= \frac{dataSize - frameLength(n)}{targetLatency}$$

- Suggested text



$$deliveryTime(i) = \frac{\sum_{k=1}^n frameLength(k)}{shapingRate} + accumulatedLatency \quad (Y-2)$$

$$requiredMinimumShapingRate = \frac{\sum_{k=1}^n frameLength(k)}{targetLatency} = \frac{dataSize}{targetLatency} \quad (Y-3)$$

# Comments #3

- Current text

### Y.3.1 Accumulated Latency

Per-hop latency imposed by a bridged network against a single frame is discussed in Clause 35 and Annex L. Latency between Talker and Listener (hereinafter referred to as Network Latency) is derived from the sum of per-hop latency along the path between them. In bridged networks controlled by SRP, the value `portTcMaxLatency` and `AccumulatedLatency` can be obtained from the system, while it can be manually evaluated by the same way in networks without SRP.

The value `AccumulatedLatency` is used as one of inputs of the calculating procedure for shaper parameter settings defined in Y.4.

#### Weakness:

- ① The parameter 'portTcMaxLatency' is not used in the whole article, can be deleted?
- ② **Consistent issues:** the parameter 'AccumulatedLatency' should be `accumulatedLatency`, the same as in (Y-2). Other parameters, such as Data Size (`dataSize`), Bounded Latency (`boundedLatency`) should also follow a consistent style.
- ③ Add description. (see right)

- Suggested text

### Y.3.1 Accumulated Latency

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The value `accumulatedLatency` is used as one of inputs of the calculating procedure for shaper parameter settings 4 defined in Y.4.

As the calculation of `accumulatedLatency` (in Equation (V-6) in Annex V of IEEE Std 802.1Qcr, 2020) is highly related to other streams in the system, when the system load changes (especially new higher-class streams are added), `accumulatedLatency` should be recalculated. Correspondingly, the SRP may also need to be re-executed.

Added text

$$d_{max}(f) = \sum_{k=1}^n d_{BU,max}(k,f) + \sum_{k=1}^n d_{MD,max}(k) + \sum_{k=2}^n d_{AT,max}(k) + \sum_{k=2}^n d_{PR,max}(k), \quad (V-6)$$

$$d_{BU,max}(k,f) = \max_{h \in F_S(k,f)} \left\{ \frac{\sum_{g \in F_H(k,h) \cup F_S(k,h)} \begin{matrix} \text{Higher traffic class} & \text{Same traffic class} & \text{Lower traffic class} \\ b_{max}(k,g) - l_{min}(h) + l_{LP,max}(k,h) \end{matrix}}{R(k) - \sum_{g \in F_H(k,h)} r_{max}(k,g)} + \frac{l_{min}(h)}{R(k)} \right\}, \quad (V-2)$$

# Comments #4

- Current text

## Y.4 Shaper Parameter Settings

### Y.4.1 General Discussion of Shaping Rate

This standard defines several types of shapers. Any of those shapers makes intervals between frames, however its parameters vary according to the type of the shaper. Each shaper is discussed in the following subclauses.

In order to minimize over-provisioning of bandwidth reservation while ensuring the requirement for the delivery time is met, the bursty traffic should be shaped with the minimum shaping rate within the required bounded latency (required minimum shaping rate). Frame propagation within bounded latency while minimizing over-provision of bandwidth reservation is illustrated in Figure Y-5 and referred to as the target latency. From Figure Y-5, the target latency can be derived from bounded latency and accumulatedLatency. The required minimum shaping rate for traffic shaping is equal to:

$$\begin{aligned} \text{requiredMinimumShapingRate} &= \frac{\sum_{k=1}^{n-1} \text{frameLength}(k)}{\text{targetLatency}} \\ &= \frac{\text{dataSize} - \text{frameLength}(n)}{\text{targetLatency}} \end{aligned} \quad (Y - 3)$$

A formula can be added to expressed the calculation more clearly, such that formulas (Y-2) and (Y-3) can move forward more smoothly.

- Suggested text

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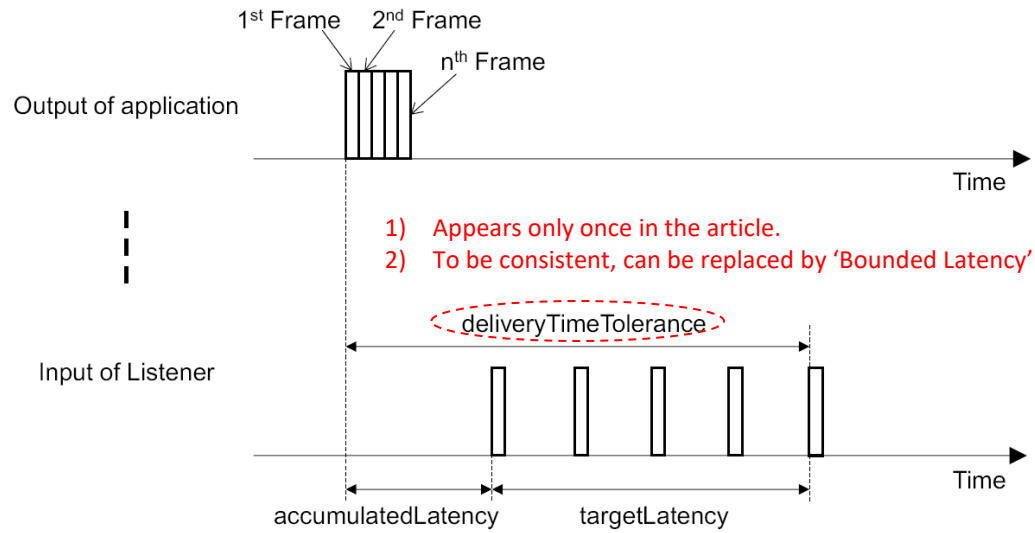
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$$\text{targetLatency} = \text{boundedLatency} - \text{accumulatedLatency}$$

$$\begin{aligned} \text{requiredMinimumShapingRate} &= \frac{\sum_{k=1}^{n-1} \text{frameLength}(k)}{\text{targetLatency}} \\ &= \frac{\text{dataSize} - \text{frameLength}(n)}{\text{targetLatency}} \end{aligned} \quad (Y - 3)$$

# Comments #5

- Current text



- 1) Appears only once in the article.
- 2) To be consistent, can be replaced by 'Bounded Latency'

Figure Y-5— Frame propagation within bounded latency while minimizing over-provision of bandwidth reservation

- Suggested text

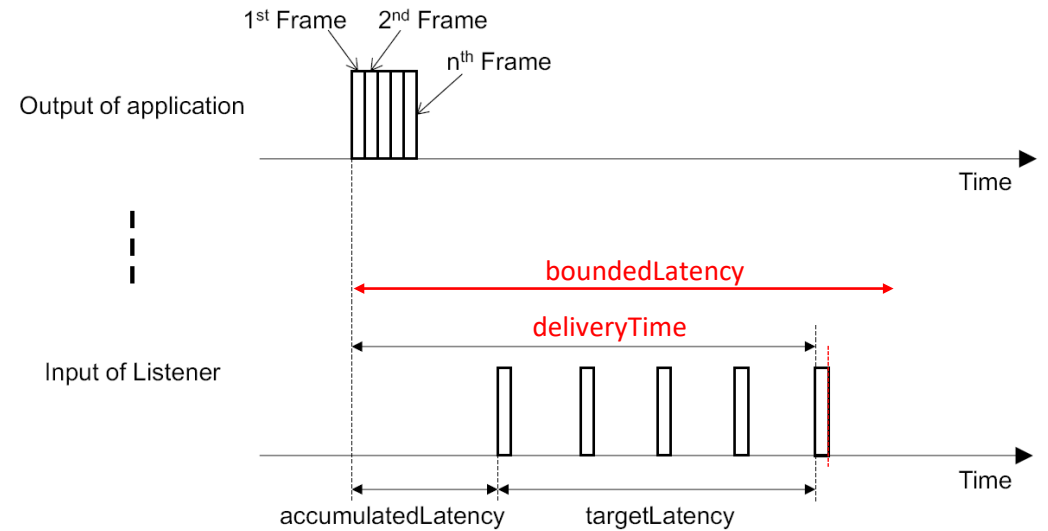


Figure Y-5— Frame propagation within bounded latency while minimizing over-provision of bandwidth reservation

# Comments #6

- Current text

## Y.4.3 Asynchronous Traffic Shaping

According to the definition of the ATS scheduler state machine in Clause 8.6.11 (IEEE Std 802.1Qcr-2020), CommittedBurstSize should be equal to or greater than frames sent by the Talker. In this case, it is recommended to be equal to the Maximum SDU Size. CommittedInformationRate is the data rate reserved for the stream and is recommended to be equal to the requiredMinimumShapingRate shown in Equation (Y-3). The approximation discussed in Clause Z.3 can also be applied. These lead to the following settings values:

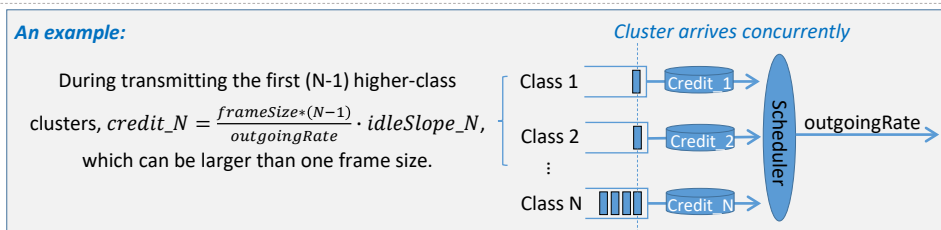
$$\text{CommittedBurstSize} = \text{Maximum SDU Size} \quad (Y-5)$$

$$\text{CommittedInformationRate} = \frac{\text{dataSize}}{\text{targetLatency}} \quad (Y-6)$$

Since the ATS scheduler state machine operation (8.6.11) assumes that the frame sizes that are processed are less than or equal to the associated CommittedBurstSize parameter (8.6.11.3.5), the CommittedBurstSize is set to be the maximum size of frame. That is equal to the Maximum SDU Size as shown in Equation (Y-5).

<< Contributor's Note: The ATS scheduler should also works for the case in which the CommittedBurstSize is greater than Maximum SDU Size. It does not affect other traffic for which the long-term averaged shaping rate. However, a small value of the CommittedBurstSize is desirable because the transient data rate, which is higher than the required minimum shaping rate, may be suppressed. This transient manner can be caused by the arrival of a new frame cluster at the shaper that has already accumulated large number of tokens causing some frames to be forwarded instantly. Such token-bucket state can occur when no frames arrive at the shaper for a period of time between clusters. >>

This part should be added and emphasized in the text.



- Suggested text

## Y.4.3 Asynchronous Traffic Shaping

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$$\text{CommittedBurstSize} = \text{Maximum SDU Size} \quad (Y-5)$$

$$\text{CommittedInformationRate} = \frac{\text{dataSize}}{\text{targetLatency}} \quad (Y-6)$$

Since the ATS scheduler state machine operation (8.6.11) assumes that the frame sizes that are processed are less than or equal to the associated CommittedBurstSize parameter (8.6.11.3.5), the CommittedBurstSize is set to be the maximum size of frame. That is equal to the Maximum SDU Size as shown in Equation (Y-5).

The ATS scheduler should also works for the case in which the CommittedBurstSize is greater than Maximum SDU Size. It does not affect other traffic for which the long-term averaged shaping rate. However, a small value of the CommittedBurstSize is desirable because the transient data rate, which is higher than the required minimum shaping rate, may be suppressed. This transient manner can be caused by the arrival of a new frame cluster at the shaper that has already accumulated large number of tokens causing some frames to be forwarded instantly. Such token-bucket state can occur when no frames arrive at the shaper for a period of time between clusters.

Similar issue also exists in the CBS scheduler. For example, when there are N classes of frame clusters arrive concurrently in the system, during the scheduling of the higher-class clusters, credits of the lowest-class cluster will be accumulated, which may results in consistent transmission of its frames. To alleviate this issue, we should also limit the credit upper-bound to the maximum SDU size.

Added text

# Conclusion and Proposal

## Conclusion:

- The project will play a great role for the parameter setting of shapers and make TSN easy to work.
- The current draft is good and the comments are trying to make it better/perfect.

## Proposal:

- It is proposed that these comments are taken into consideration or as input for the further drafting of the current draft.