



Institute of Computer Science Chair of Communication Networks Prof. Dr. Tobias Hoßfeld



Latency Model and Example Reservation Flow in RAP IEEE 802.1 Meeting, January 2023

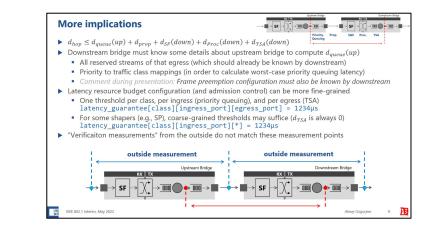
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Overview for this Presentation

- Recap: last presentation (measurement points for latency models)
 - cf. <u>dd-grigorjew-measurement-points-0522-v02.pdf</u>
 - What are the measurement points (and resulting delay segments)?
 - What are the reasons for this change?
 - Some implications
- More implications
 - Visualization of delay segments
 - What happens at the Listener?
 - Suggestion: merge two configurations per delay segment
 - What happens with different Shapers?
 - Suggestion: communicate the behavior at the Priority Transmission Selection Queue
- Example reservation flow
 - Very simple scenario (2 switches, 2 streams)
 - Clarify general procedure
 - Clarify the implications of delay segments





if we have time SW 1 SW 2

Recap: Last Presentation

MEASUREMENT POINTS FOR LATENCY MODELS



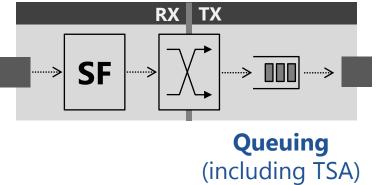


Extended delay model, including transmission selection algorithm

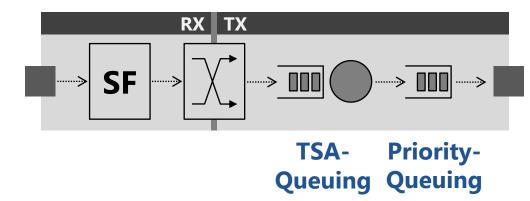
- Split "queuing" latency of formal latency models into...
 - Transmission Selection Algorithm (TSA)
 - Priority-Queuing, where only the eligible frames interfere

Previous model:

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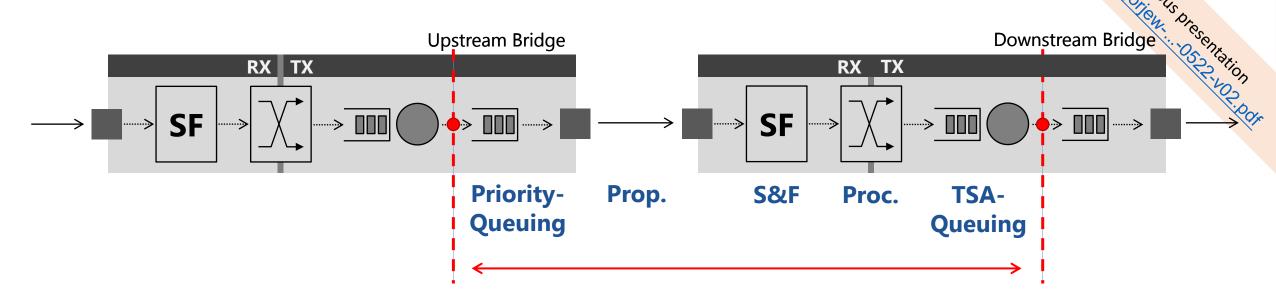
Extended model:



- Add measurement point during queuing when frame becomes eligible for transmission
 - SP: Immediately after enqueuing
 - CBSA: When credits >= 0, the head of the queue becomes eligible for transmission
 - ATS: When the defined eligibility time for that frame is reached (cf. Qcr)
 - CQF: When queues swap roles (receive \rightarrow send), all frames in the send queue become eligible

previous presentation

Suggestion: Use ATS measurement points for all shapers in RAP



- Suggestion: Use the ATS measurement points for all TSAs & latency models in RAP
- Per-hop latency is given by...
 - Queuing after eligibility time was reached (upstream)
 - Propagation

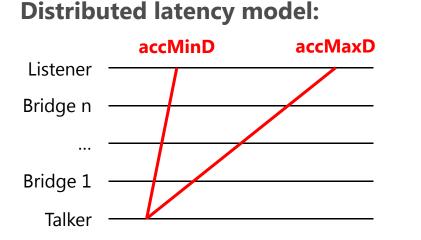
WÜ

- Store-and-Forward (downstream)
- Processing (downstream)
- Queuing until eligibility time is reached (downstream)
- Comment during presentation: PHY can often introduce a delay after priority queuing. The simple suggestion is to account for it as part of the upstream processing delay, even if it technically occurs after the measurement point.

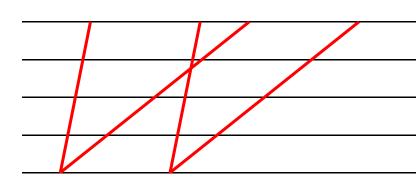
// queuing for priority transmission selection

// queuing for transmission selection algorithm

Why is shaper-to-shaper latency beneficial?



CQF (edge to edge measurement):



CQF (shaper to shaper):

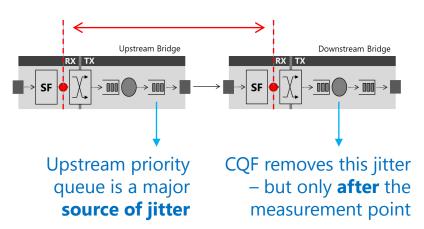
dd-grigorjew-strict-priority-latency-0320-v02.pdf

Generally:

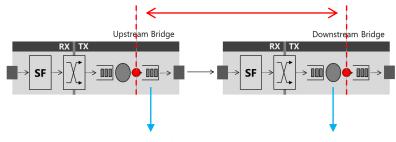
WÜ

- Minimum delay and maximum delay accumulated per hop
- Accumulating bursts are calculated based on (accMaxD - accMinD)
- A lower latency variance is better for downstream delay computation

Fully-received to fully-received:



Shaper to shaper:



All sources of jitter can be removed; Well-defined traffic pattern, as intended by the TSA, is measured directly after the TSA.

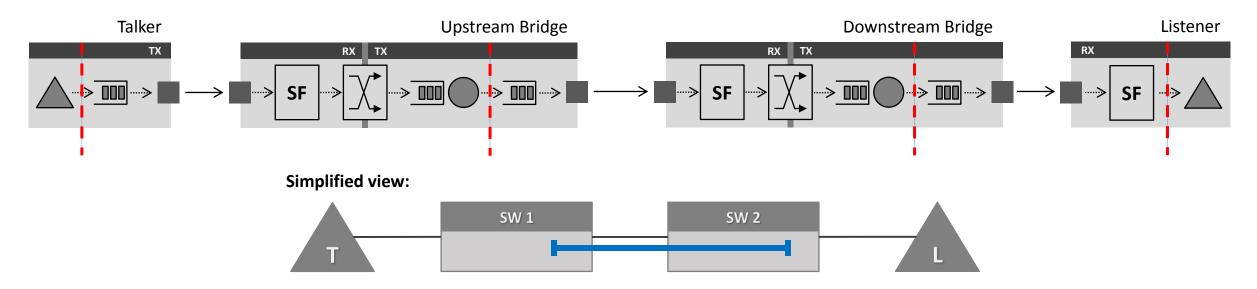
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Visualization and new Suggestions

MORE IMPLICATIONS

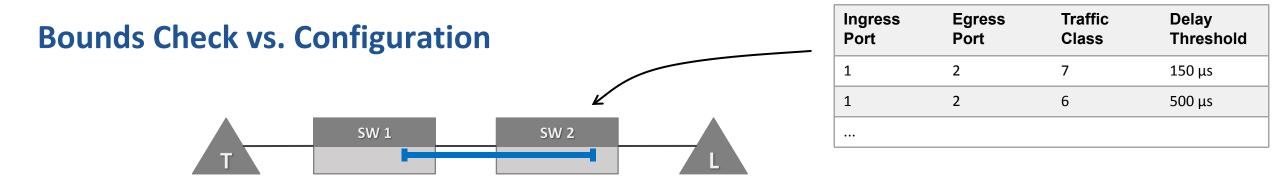


Full End-to-end Path with Delay Segments



One delay segment includes

- TX of upstream bridge (SW 1)
- RX of downstream bridge (SW 2)
- TX of downstream bridge (SW 2)
- Downstream bridge (SW 2) performs the bounds check during reservation
- But where does the configuration (delay threshold) come from? SW 1 or SW 2?
- General problem: on any path with n bridges (2 bridges), we have n+1 delay segments (3 delay segments)



- Initial suggestion: SW 2 performs bounds check and contains the delay threshold config
 - But: we don't really like the fact that SW 1 has no say, although it is involved in the delay segment

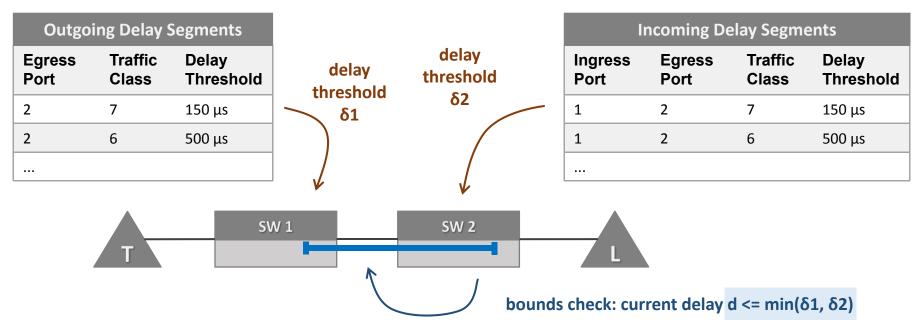
In addition: what happens at the Listener?



It can perform bounds checks

- But we don't really want to configure that aspect in our end devices
- (Config sources can be: default configuration, profile, CLI, Network Management System)

Implications of Having two Delay Threshold Tables

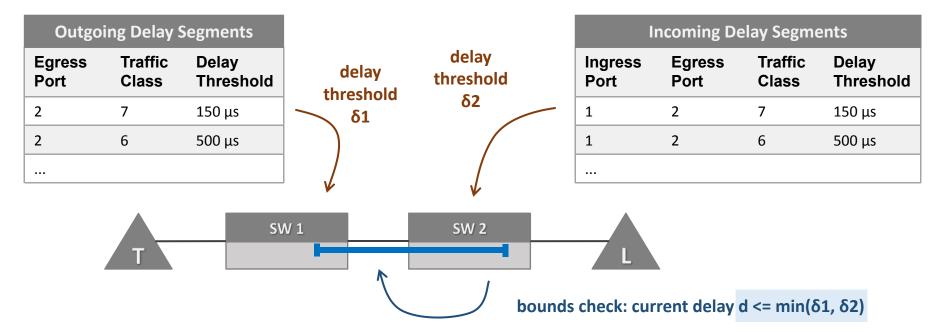


► The listener no longer needs a delay threshold configuration

- It can simply use $\delta 1$ of upstream (SW 2 in that case)
- It can still specify its own δ2 where necessary (e.g., routers are listeners from layer 2 RAP point of view)
- When optimizing a network's configuration (e.g., via NMS), simply use the same value for $\delta 1$ and $\delta 2$
 - It is the same delay segment after all
- Upstream (SW 1) could specify "don't care" in order to prevent unnecessary resource constraints
 - Technically, it still **needs** a valid outgoing $\delta 1$ config in case an end device connects to that port



Suggestion: Both Devices Suggest a Delay Threshold



Suggestion: split threshold configuration for each delay segment into two configs

- Upstream bridge (SW 1) has one config for each egress port and traffic class
- Downstream bridge (SW 2) has one config **for each in ingress/egress port pair** and traffic class
- Each bridge will have two delay config tables: one for outgoing delay segments, one for incoming segments
- Upstream (SW 1) communicates the outgoing $\delta 1$ with the downstream neighbor (SW 2)
 - Downstream aggregates both configurations and selects the minimum of both for bounds checking



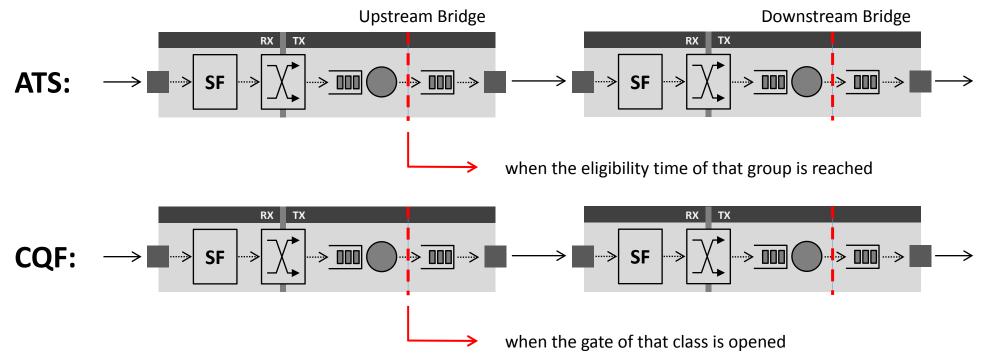
A Closer Look at Delay Segments with Different Shapers

- Recap old presentation: delay segments begin when the frame "becomes eligible for transmission"
- More specifically, we want delay segments to be tied to the events that change the shaper's state
- ► This ensures that the shaper has the intended effect on the latency model
- ► This is simple for ATS and CQF:

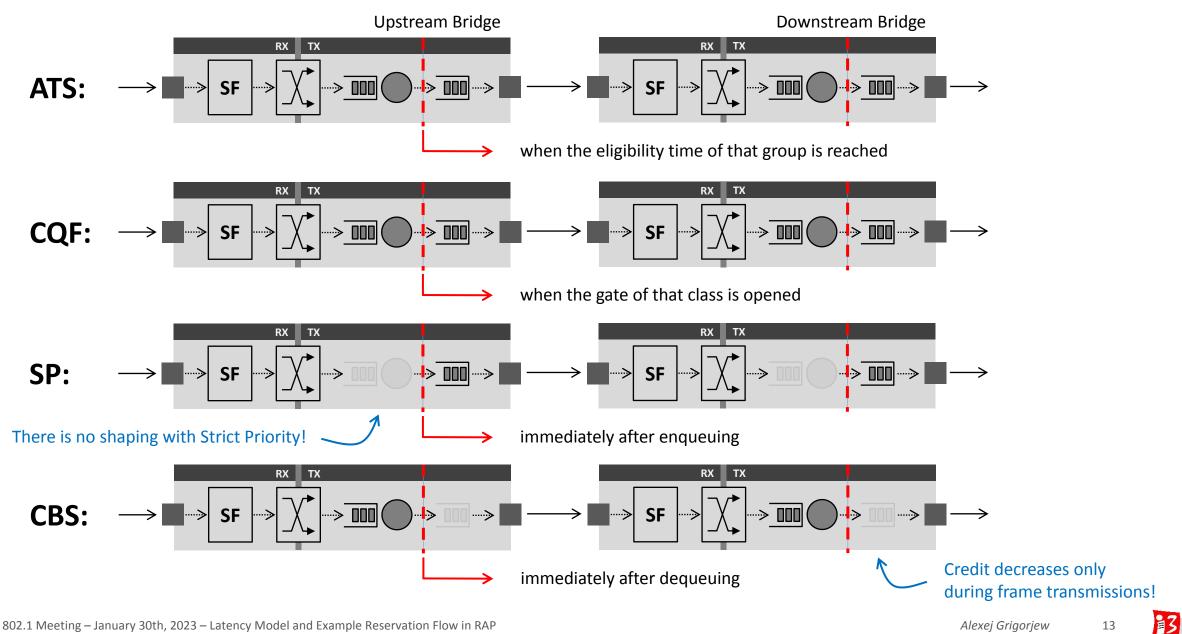
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CQF (shaper to shaper):



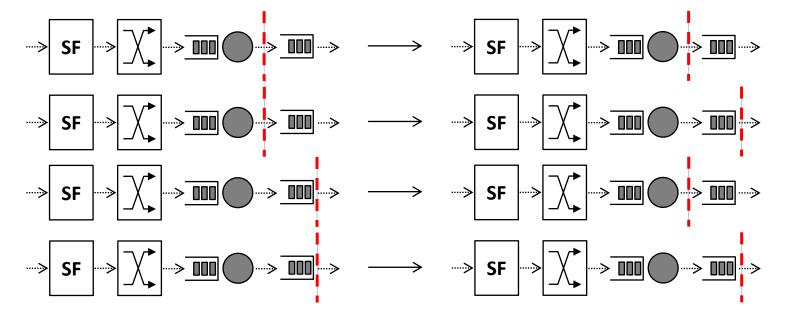
Not All Shapers Use All Delay Segments



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Suggestion: Communicate Whether the Last Queue is Part of the Delay Segment

- Suggestion: Instead of fixed, shaper-specific behavior, introduce a variable that indicates whether the (pure) priority transmission selection is part of the next delay segment
- For heterogeneous networks, this creates four scenarios:



Each device can now specify which delay segment the last transmission queue belongs to

- ATS and CQF do not include it
- CBS does always include it

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SP can now decide! (this can help with the vast zoo of end devices)

This includes the bounds check and the accMaxLatency field



TAs, LAs, Bounds Checks, Example Values

EXAMPLE RESERVATION PROCESS

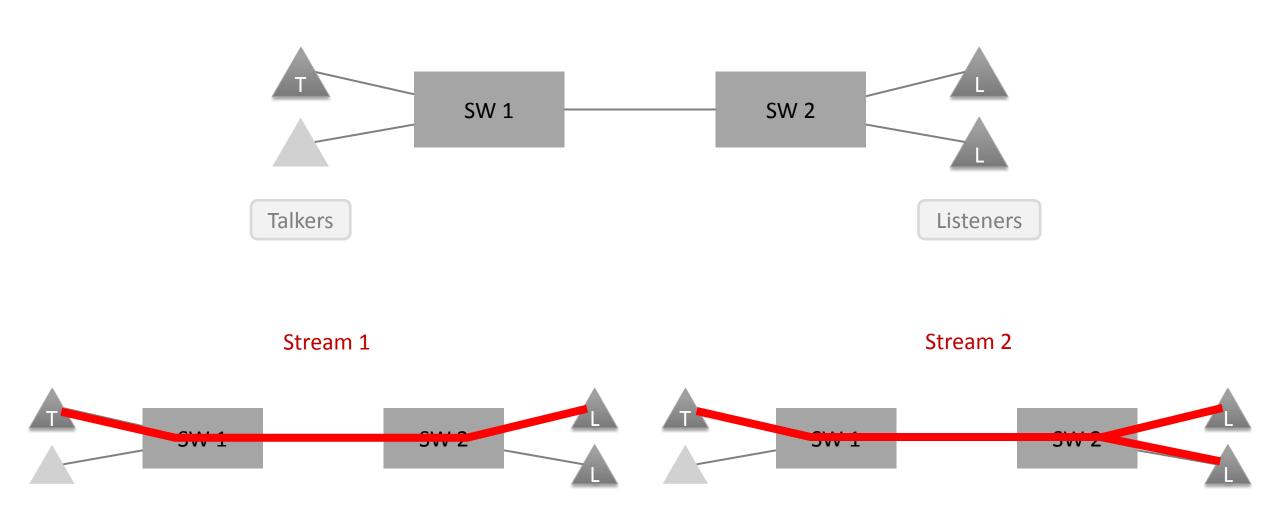


Disclaimer

- Just a simple example!
- Many things are simplified
- Some things are only suggestions
- Some things are subject to change in the standard
- See this as a means for easy introduction
- Please do not cling to the details

E3

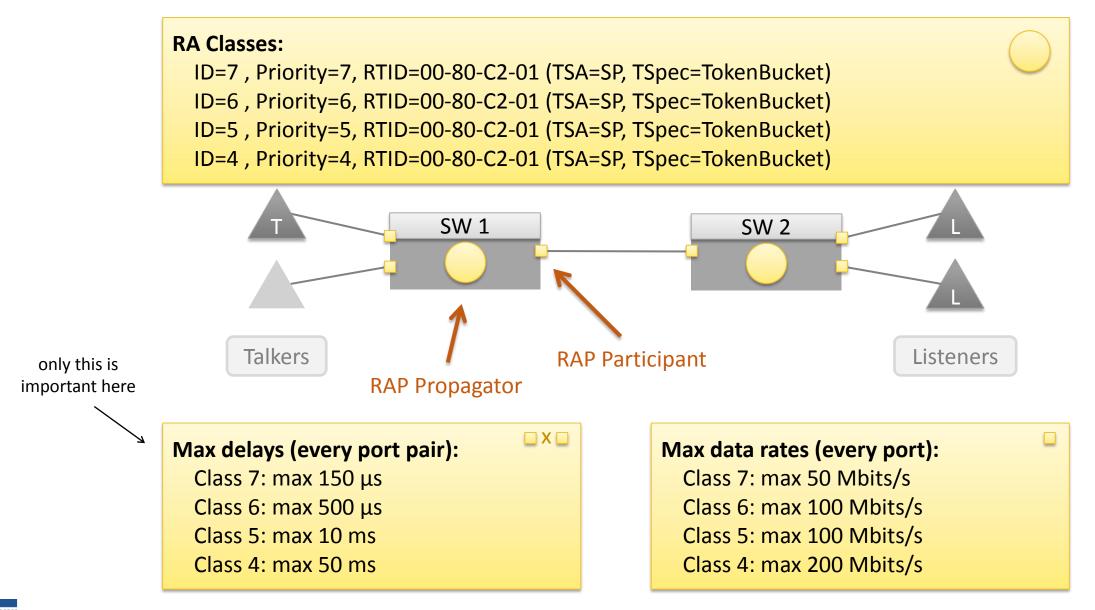
Example Topology Overview



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Configuration

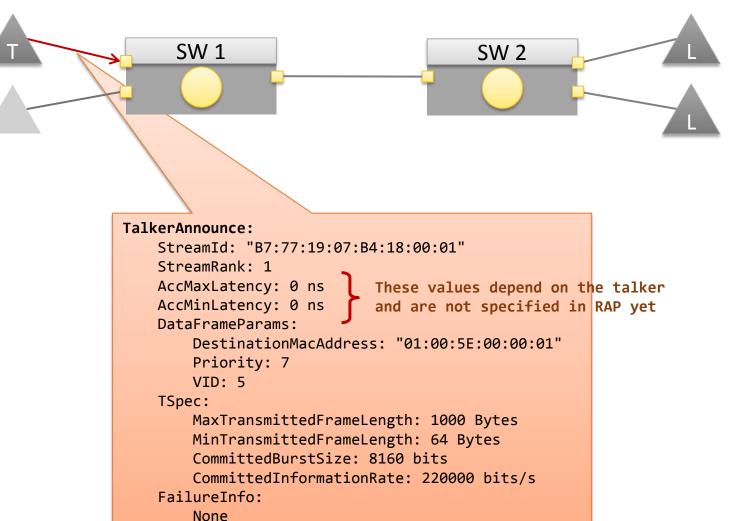
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IEEE 802.1 Meeting – January 30th, 2023 – Latency Model and Example Reservation Flow in RAP



Stream 1

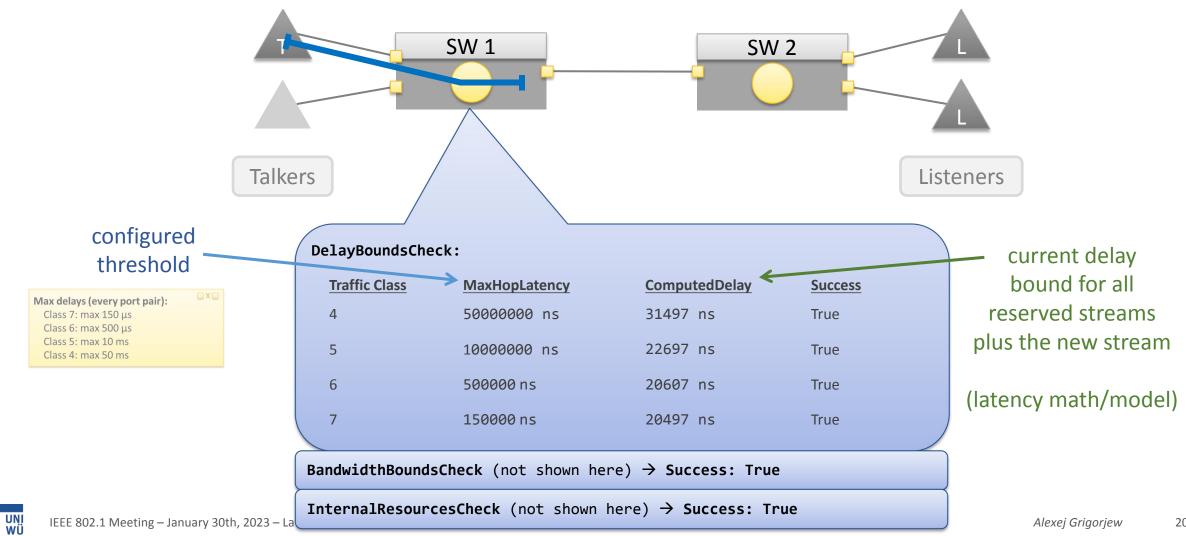


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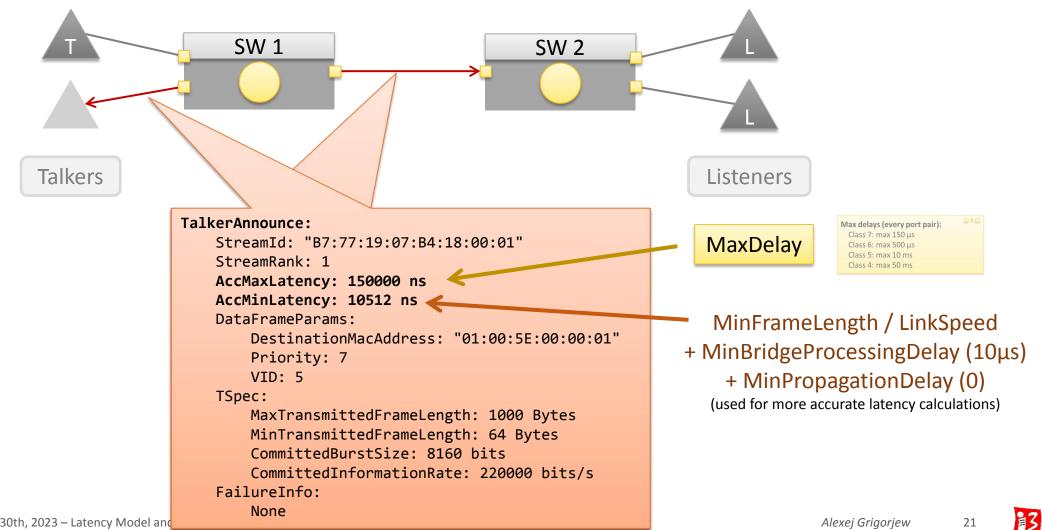
i3

Bounds Check on SW 1 (EgressPort: SW 2)



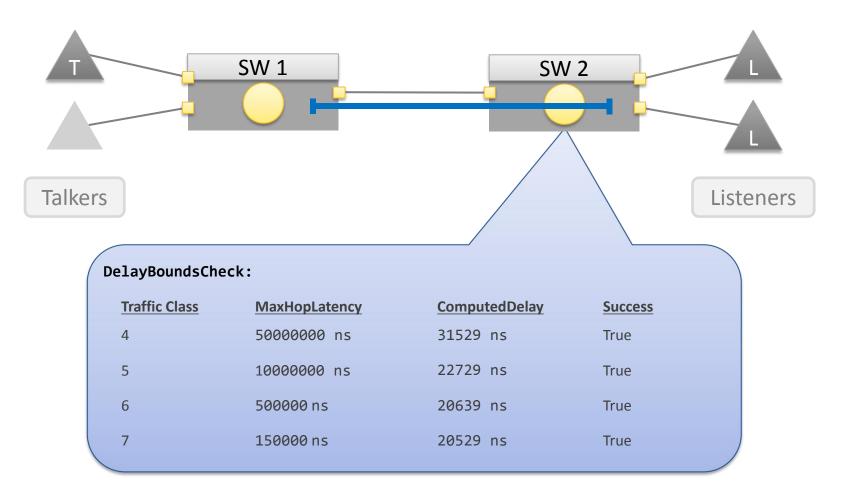
E3

Adjusted TA is Propagated to other Ports



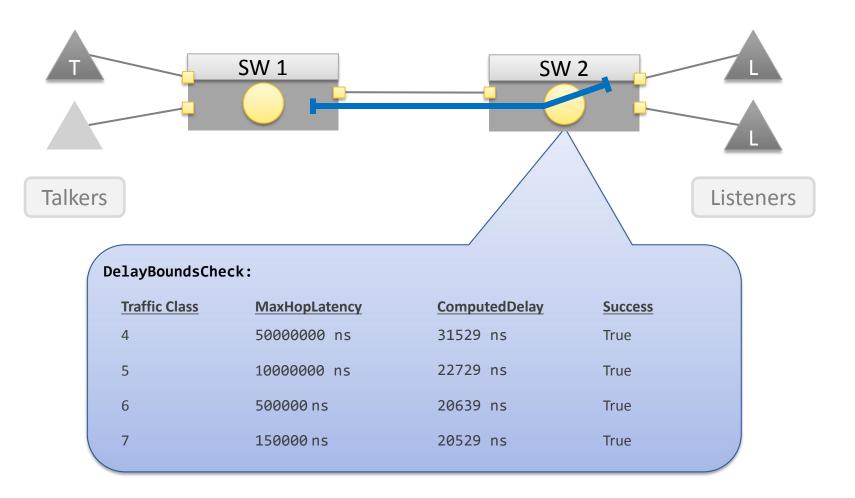
UN

Bounds Check on SW 2 (EgressPort: Listener 2)



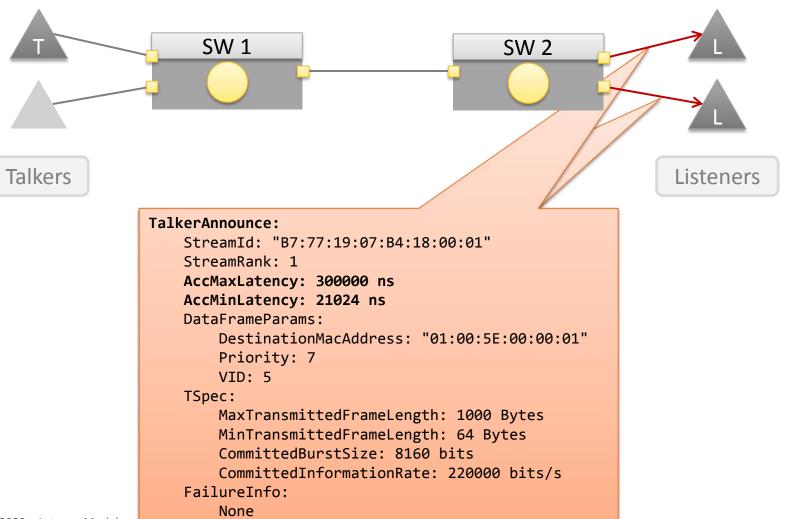
UNI WÜ 3

Bounds Check on SW 2 (EgressPort: Listener 1)



UNI WÜ 3

Adjusted TA is Propagated to other Ports

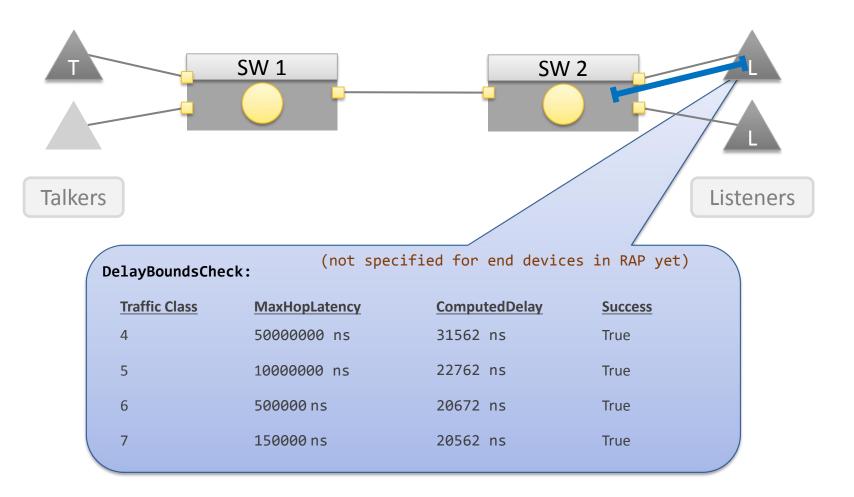


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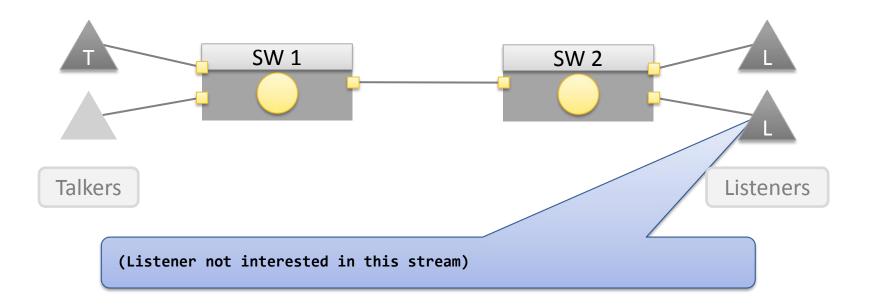
i3

Bounds Check on Listener 1



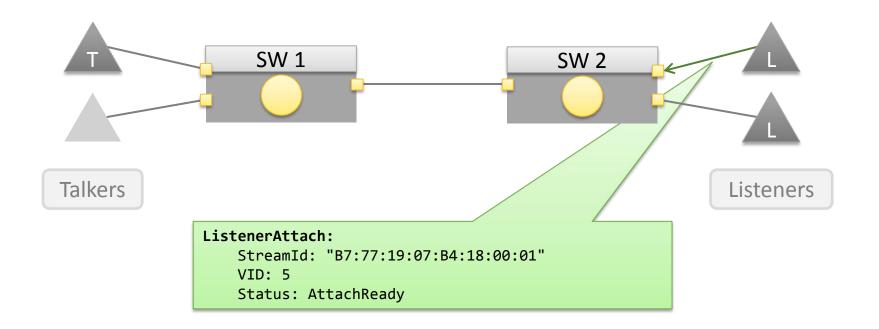


Listener 2 is not Attaching



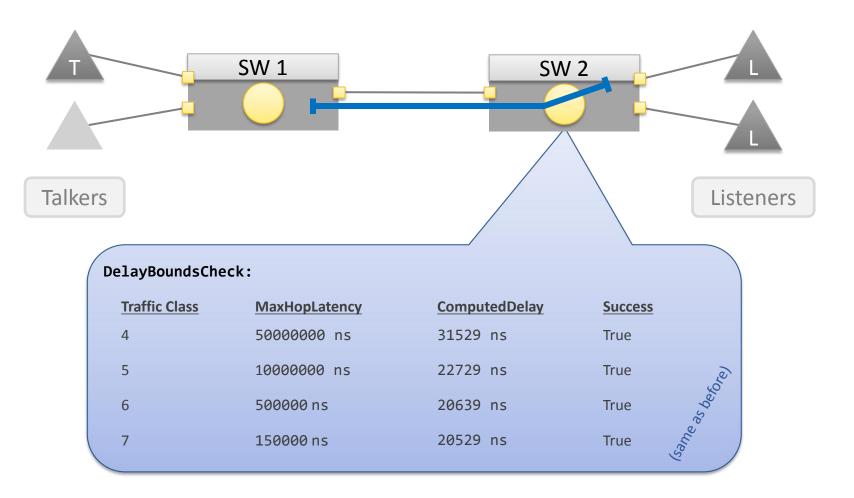


Listener 1 sends LA



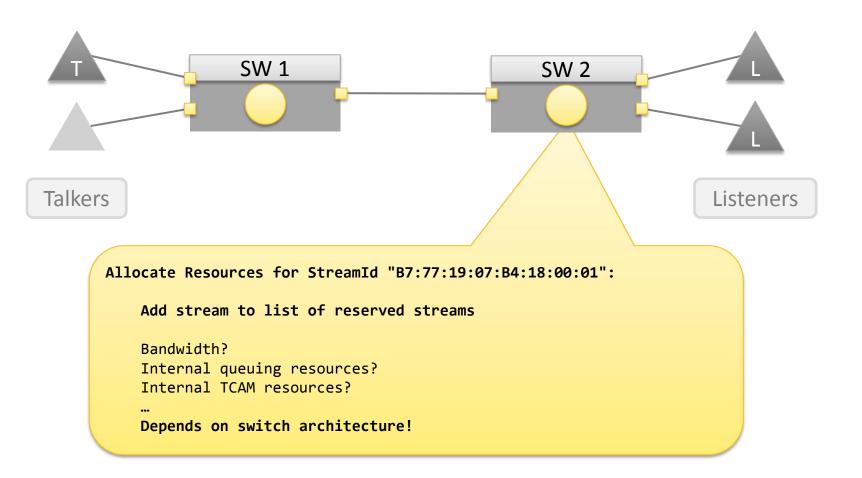


Bounds Check on SW 2 (EgressPort: Listener 1)



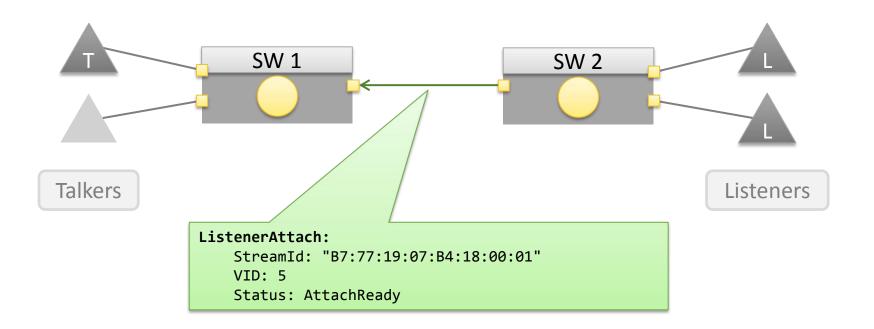


Reservation on SW 2 Successful



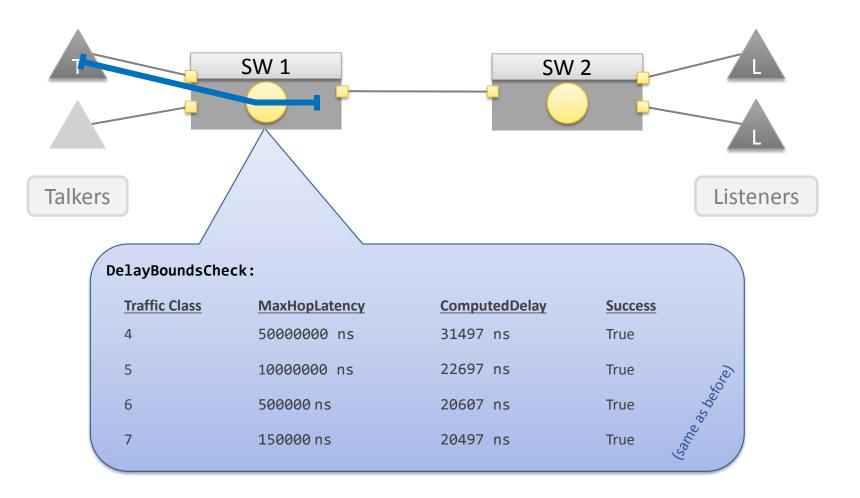
UNI WÜ 13

SW 2 Forwards the LA to SW 1



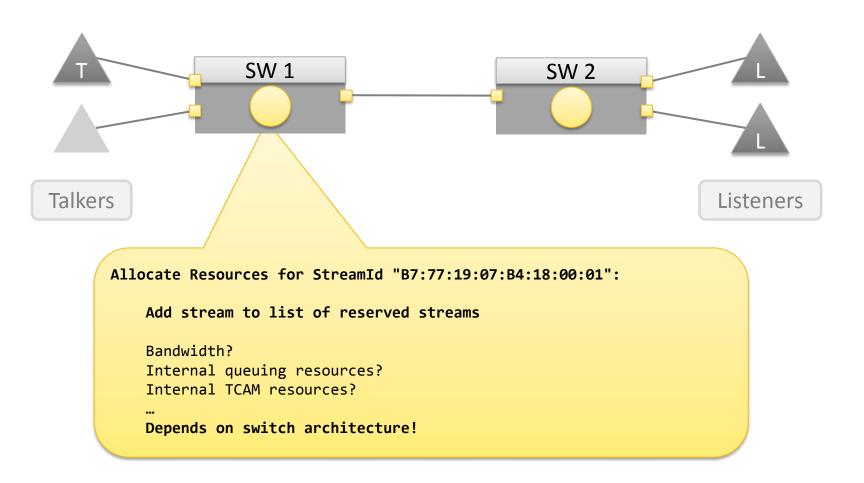


Bounds Check on SW 1 (EgressPort: SW 2)



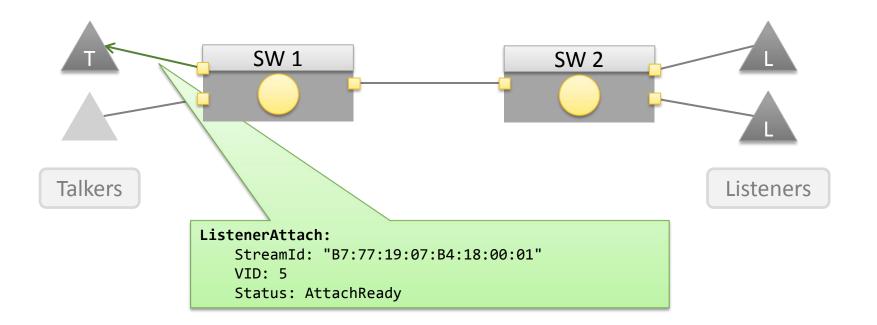


Reservation on SW 1 Successful



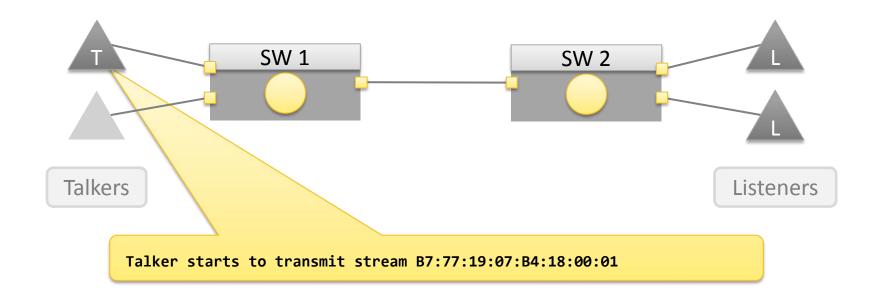
UNI WÜ 3

SW 1 Forwards the LA to the Talker



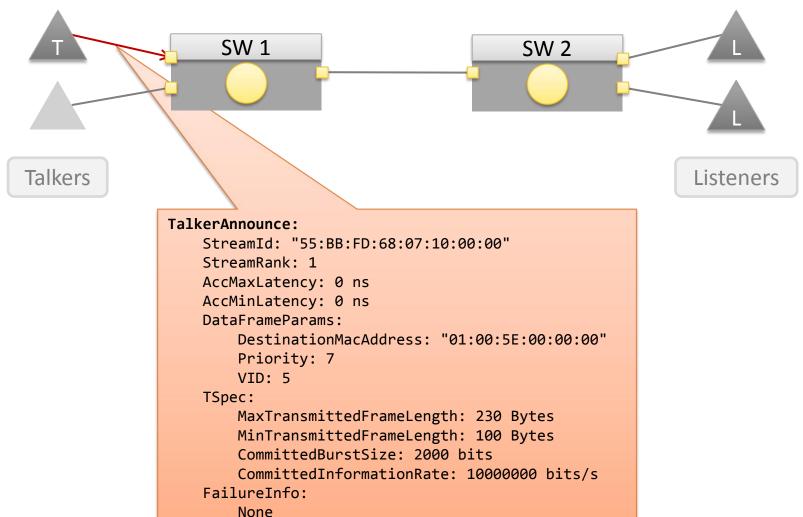


Talker Received Successful LA and Starts to Transmit



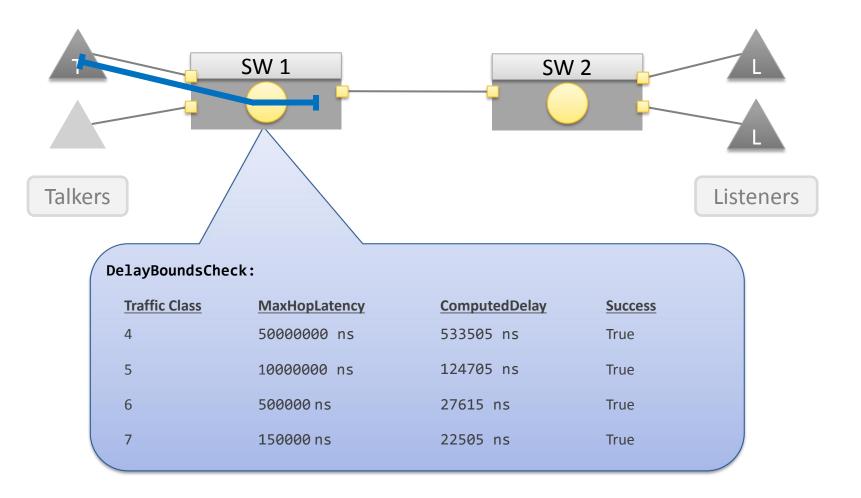


Stream 2



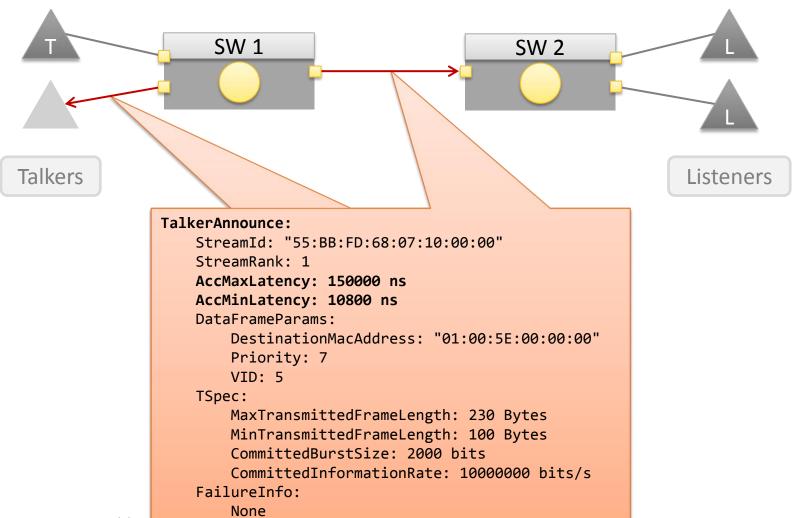
UNI WÜ 13

Bounds Check on SW 1 (EgressPort: SW 2)



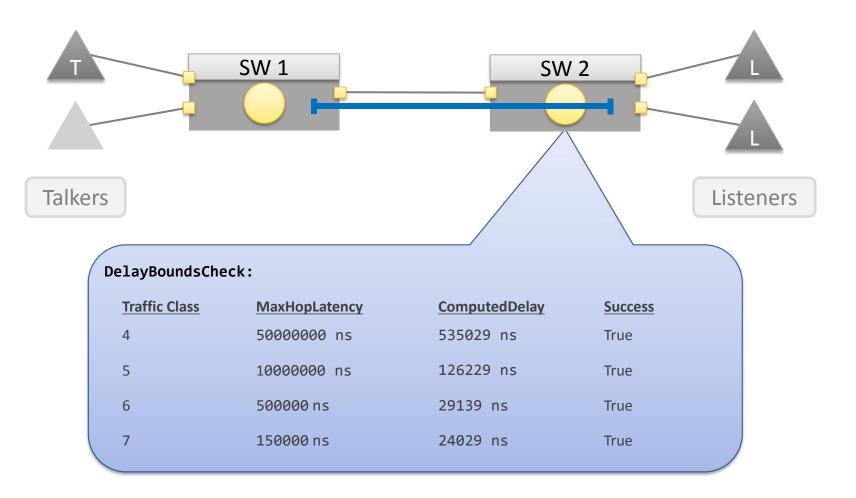


Adjusted TA is Propagated to other Ports



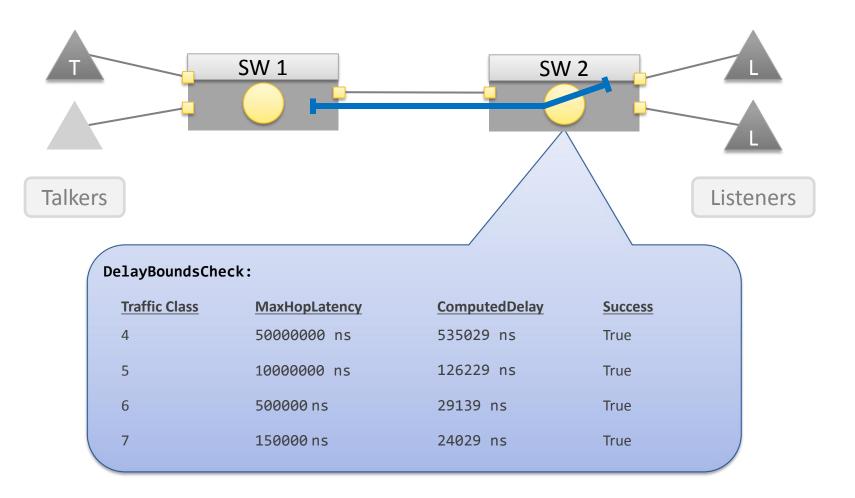
UNI WÜ **i**3

Bounds Check on SW 2 (EgressPort: Listener 2)



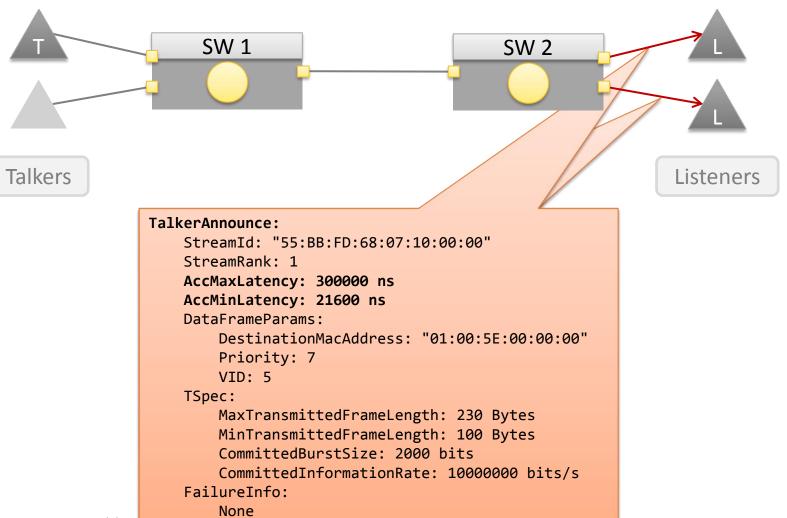


Bounds Check on SW 2 (EgressPort: Listener 1)





Adjusted TA is Propagated to other Ports

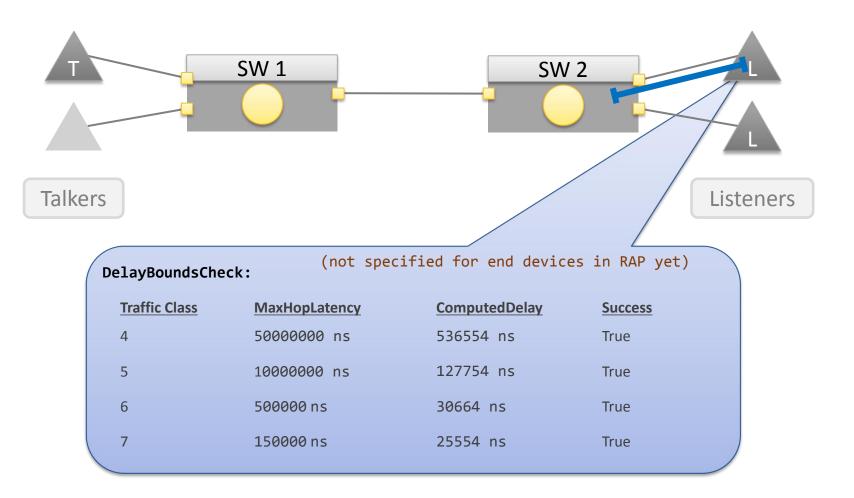


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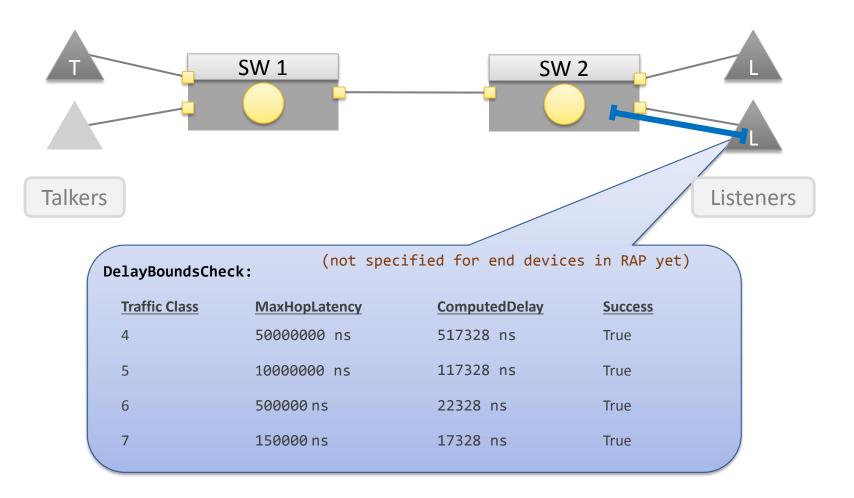
i3

Bounds Check on Listener 1



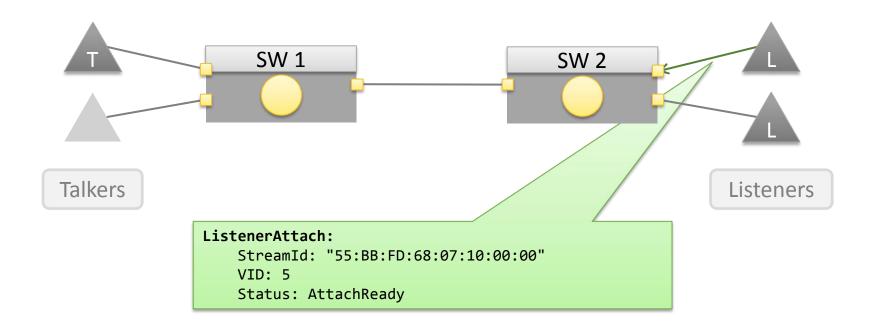
UNI WÜ

Bounds Check on Listener 2



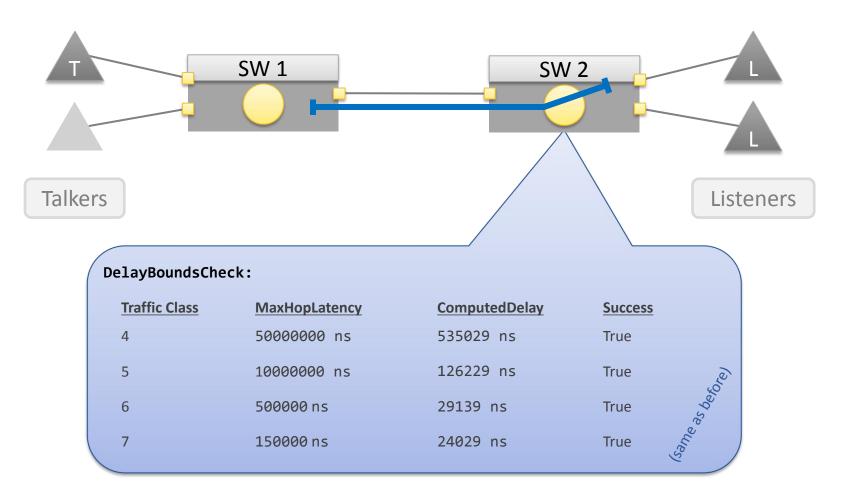
UNI WÜ

Listener 1 sends LA



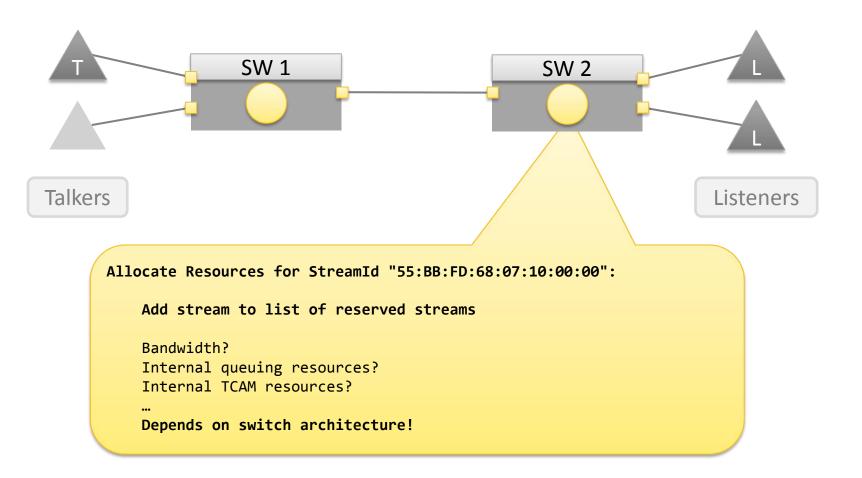


Bounds Check on SW 2 (EgressPort: Listener 1)



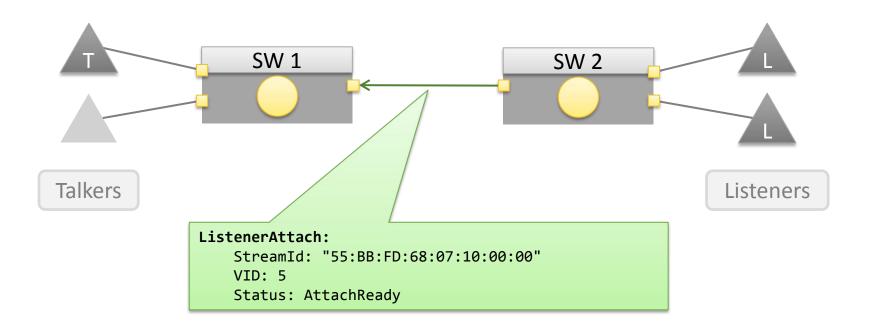
UNI WÜ

Reservation on SW 2 (Port 1) Successful



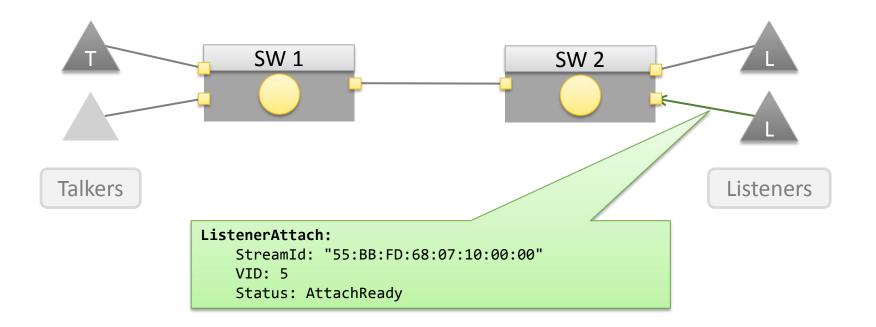
UNI WÜ **B**3

SW 2 Forwards the LA to SW 1



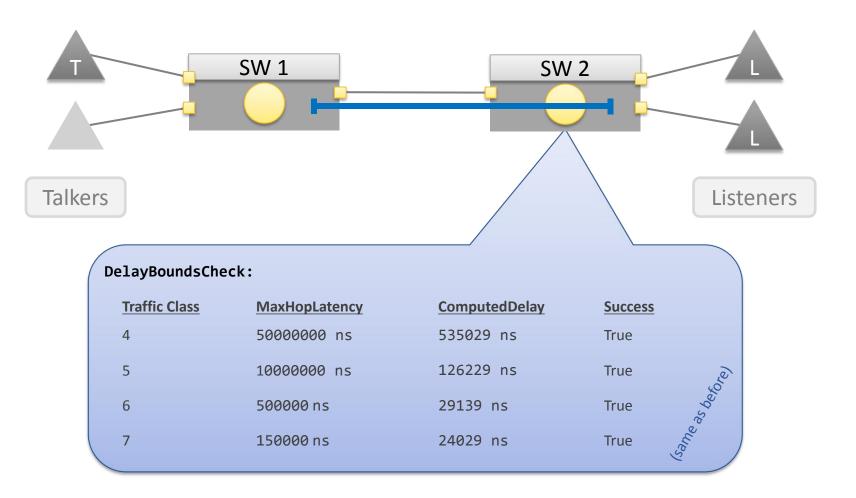


Listener 2 sends LA



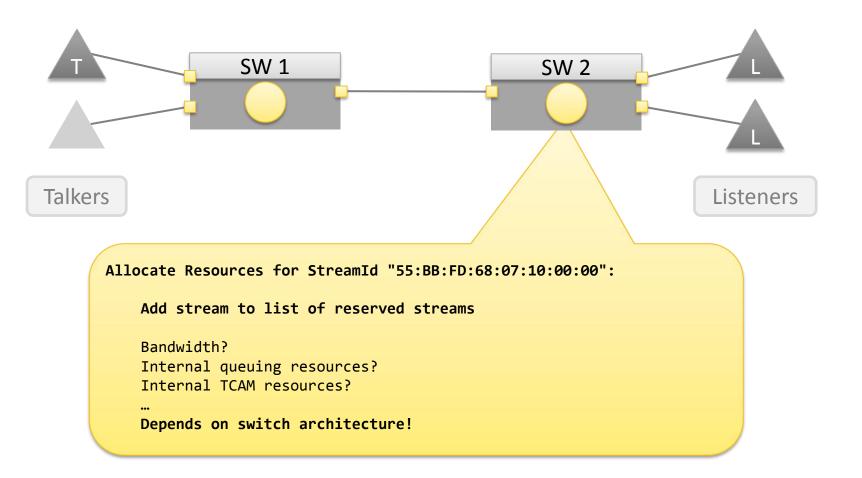


Bounds Check on SW 2 (EgressPort: Listener 2)



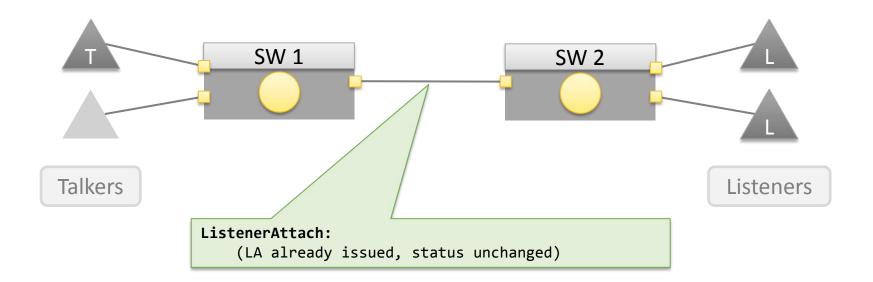


Reservation on SW 2 (Port 2) Successful



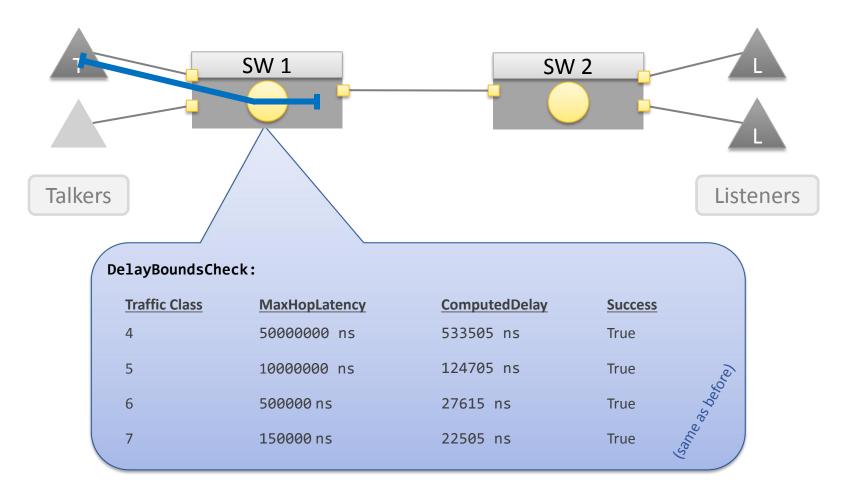
UNI WÜ **B**3

SW 2 does Nothing, Existing LA for this Stream Unchanged



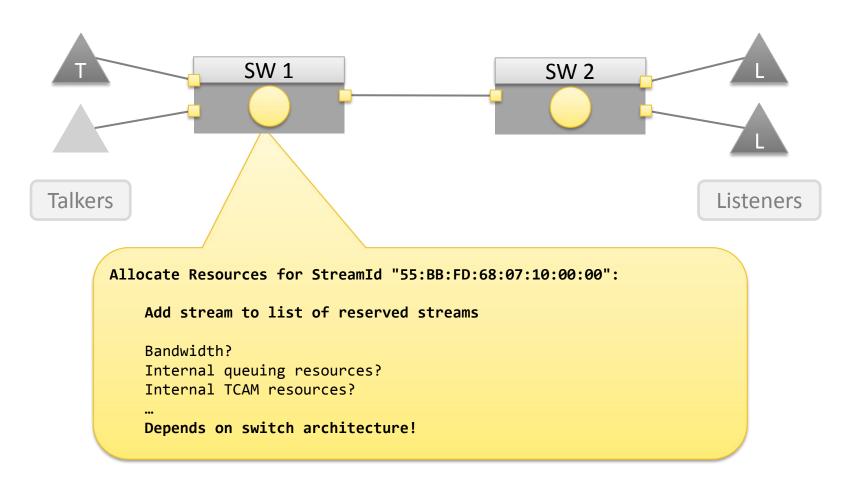


Bounds Check on SW 1 (EgressPort: SW 2)



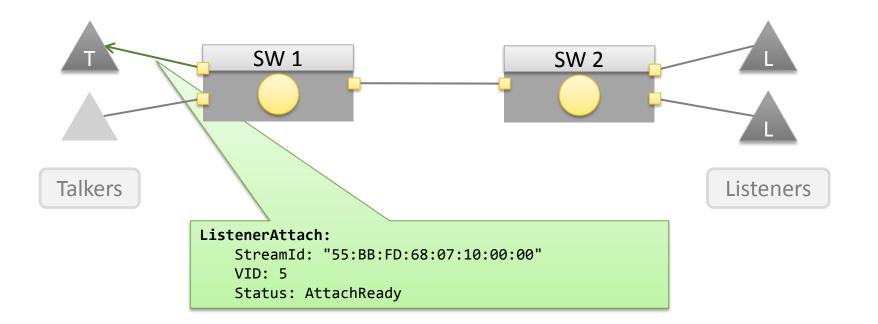
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Reservation on SW 1 Successful



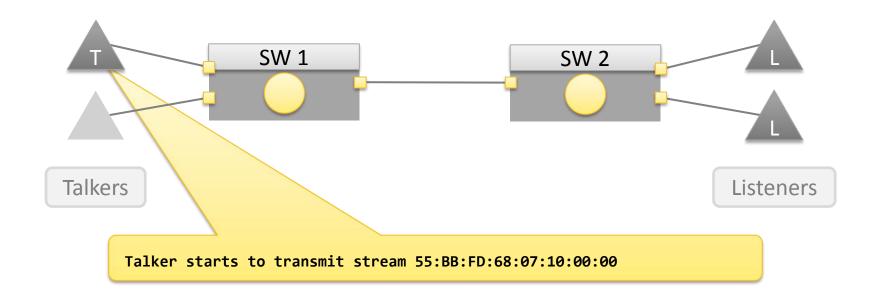
UNI WÜ

SW 1 Forwards the LA to the Talker





Talker Received Successful LA and Starts to Transmit





THANK YOU!

Questions, comments, suggestions?



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