

Using Strict Priority to meet latency targets

a method for cyclic traffic

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Inspired by

- Strict Priority with **proper** Traffic Engineering (TE) can be TSN - <https://www.ieee802.org/1/files/public/docs2025/detnet-tsn-farkas-tsn-update-0725-v00.pdf>

- > IEEE Std 802.1CM™-2018 Clause 8.1.3 gives a guideline for worst-case delay calculation for the highest priority traffic based on SP.
- > And Annex B gives an example (Route-based calculation).

Table B-1—Bridge delays for Profile A

	Bridge 12	Bridge 13	Bridge 14	Bridge 15	Total
$t_{MaxBridge}$	9.9344 μ s	9.9344 μ s	7.4672 μ s	7.4672 μ s	34.8032 μ s

- If the end-to-end latency budget is 100 μ s for the high priority data, then the total propagation delay of these links can be 65.1968 μ s, i.e., the distance between can be approximately 13 km.

- The latency guarantee of a specific TSN solution can be ‘modularized’ with preconditions and reservations.

- > IEEE Std 802.1BA™-2011 Clause 6.5 specifies the way to meet latency targets for SR classes A and B.
- > There is a calculation process behind, but users can focus solely on the outcome as in Table 6-2.

Table 6-2—Latency targets for SR classes A and B

SR class	Max end-to-end latency
A	2 ms
B	50 ms

- NOTE 1—The choice of latency targets shown in Table 6-2 reflects the requirements of some typical deployment scenarios, and should not be taken as hard-and-fast limits on the end-to-end latency in an AV network. They do, however, form a useful basis for achieving “plug-and-play” interoperability.
- NOTE 2—The 2 ms figure for SR Class A can be met for 7 hops of 100 Mb/s Ethernet if the maximum frame size on the LAN is 1522 octets.

- A significant portion of the most critical flows exhibit periodic characteristics in industrial automation, automotive, aerospace, etc. Can we find similar ‘modularized’ easy-to-compute solutions based on Strict Priority?

- > E.g., in 60802 v2 for ccB.

High bandwidth brings feasibility

- An example:
 - > 21 end-stations, 20 pairs of 1ms control with a 200B packet.
 - > There are background flows with Max. 1500B packet.
 - > End-to-end Max. 5 hop with 5us fixed delay (includes all other elements) per hop, 1G bandwidth.
 - > End-to-end latency requirement: 3ms.
- Calculate the pessimistic worst-case delay:
 - > Imagine a packet, goes the longest way, and on every hop has to wait a 1500B lower priority packet, and is ranked last among the 20 packets. When everything can go wrong does go wrong, the e2e delay goes to 123.4μs, far less than 3ms.

$$5 * [5 + (1500 + 200) * 8/1000] + \frac{200 * 8 * 19}{1000} = 123.4$$

propagation delay neglected

all other considered as a fixed delay

queuing delay

frame transmission / S&F delay

self-queuing delay

Benefits (from the perspective of 60802 v1)

- Still using centralized configuration, but can be much easier.
- No harm if the calculated result doesn't meet the requirement.
 - > Use other methods with less pessimism to calculate to see if they work.
 - > Choose ccA.

Potential use cases

- Within a machine / production cell where the numbers of end-stations and networking hops aren't too high.
 - > Even for control loops with very high frequencies, it may still be feasible.
 - > Besides, we still have time synchronization in ccB.
- Within a factory where the important control traffics come across have a relatively low frequency.
 - > <https://www.ieee802.org/1/files/public/docs2025/60802-Steindl-Proell-IA-Controller-ConfDomain-Cloud-0725-v01.pdf>
 - > It may be feasible when there is no strong demand for placing the control that requires very high frequency and precision (such as motion control) on the edge controller (container).
- Others.

'It' in this page refers to the preliminary and vague concept of Strict Priority based simplified Traffic Engineering.

Discuss

- Possible outcome: Maybe in 60802 v2, or a new Annex in .1Q. It's too early to tell.
- The author wants to see if there is any comments or common interests on this matter.