

## 9. MAC ~~fairness~~fairness

**Editors' Notes:** To be removed prior to final publication.

References:  
None.

Definitions:

Abbreviations:

Revision History:

Draft 0.1, February 2002

Draft 0.2, April 2002

Draft 0.3, June 2002

Initial draft document for RPR WG review.

Revised according to WG comments for TF review.

Revised according to WG comments for TF review.

### 9.1 Overview

**Editors' Notes:** To be removed prior to final publication.

Missing things in the overview section include:

- Functional block diagram and contents of clause
- Notational conventions

Detailed state diagrams are needed that describe:  
the events that cause fairness messages to be sent;  
the actions when a fairness message is received;  
congestion and congestion subsided, etc.

The PICS needs to be added.

Need to make sure that reference model clause is in sync with this one.

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the events that cause fairness control messages to be sent;  
the actions when a fairness control message is received;  
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#### 9.1.1 Scope

~~This clause defines the fairness protocol for RPR MACs. The fairness protocol ensures a fair distribution of available bandwidth across all stations on a ring even during times when the ring is heavily utilized or congested. This clause specifies the MAC fairness protocol and how the MAC uses the protocol to enforce fairness among stations on the ring.~~

This clause defines the fairness algorithm for RPR MACs how the MAC uses the algorithm to enforce fairness among stations on the ring. The RPR fairness algorithm (RPR-FA) is a local fairness mechanism that enforces fairness among all the stations on a ring, even during times when the ring is heavily utilized or congested. Only fairness eligible traffic (the class-B traffic in excess of the provisioned amount, and all class-C traffic) is subject to the fairness algorithm.

### 9.1.2 Goals and objectives

**Editors' Notes:** To be removed prior to final publication.

*We could not agree on a definition for stability, and hence the corresponding objective has been removed. If we find an acceptable way to describe stability, then we can add it later. For now, note that contributions are invited on a suitable definition (and corresponding objective) for stability. One possible definition is:*

**Stability**—The protocol should not oscillate regardless of the type of traffic presented to the network.

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**Stability**—The protocol should not oscillate regardless of the type of traffic presented to the network.

The fairness protocol has the following objectives:

- a) **Source-based weighted fairness**—On any given ~~segment-link~~ on the ringlet, the available bandwidth is allocated to each ~~node-station~~ in proportion to its relative weight. For example, if every ~~node-station~~ has an equal weight, then the available bandwidth on the ~~segment-link~~ should be shared equally by all ~~nodesstations~~. On the other hand, if one ~~node-station~~ has a higher weight, the bandwidth allocated to that ~~node-station~~ should be in proportion to the ~~nodestation~~'s weight divided by the sum of the weights of all the ~~other~~ stations.
- b) **Reclamation of unused committed bandwidth**—The fairness protocol should be able to reclaim unused ~~bandwidth-bandwidth, including~~ that ~~which is provisioned but not reserved~~ ~~provisioned~~.
- c) **Support for ~~non~~-VDQ-single choke and VDQ-multi choke capable clients**—The fairness protocol should be able to support clients whether ~~they utilize a single choke point or not they perform multiple choke points (such as those clients performing~~ virtual destination queueing (VDQ)). ~~VDQ is a solution to the head-of-line (HOL) blocking problem.~~
- d) **Fast response time**—Because data traffic tends to be bursty, in order to ensure maximum ring bandwidth utilization and to ensure that the protocol is responsive to instantaneous changes in traffic load, it must have a fast response time.
- e) **High bandwidth utilization on the ring**—The protocol should be able to achieve very high levels of bandwidth utilization even under heavy loads approaching 100% of the ring capacity.
- f) **Scalability**—The protocol should be scalable and should be able to function predictably for all ringlet speeds and ring diameters allowed by this standard.

### 9.1.3 Relationship to other clauses

~~TBD. Will be done by referring to the MAC reference model. Also need to refer to the appropriate primitives in the MAC service interface definitions.~~

The RPR-FA is implemented within a control entity called the Fairness Control Unit (FCU) located in the MAC Control Sublayer, as described in Clause 5.

The FCU uses the following variables defined in Clause 6: add rate, add rate congested, fw rate, fw rate congested.

The FCU provides the following variables for use in Clause 6: allowed rate, allowed rate congested, TTL to congestion. The usage of these variables by Clause 6 is described in 9.10 .

## 9.2 Acronyms

This clause contains the following acronyms:

<u>FA</u>	<u>fairness algorithm</u>
<u>FCM</u>	<u>fairness control message</u>
<u>FCU</u>	<u>fairness control unit</u>
<u>FE</u>	<u>fairness eligible</u>
<u>LR</u>	<u>line rate</u>
<u>MC</u>	<u>multi choke</u>
<u>RTT</u>	<u>round trip time</u>
<u>SC</u>	<u>single choke</u>

## 9.3 Variables and terminology used

**Editors' Notes:** *To be removed prior to final publication.*

*This section has been added by the editor. We need to have a list of definitions of all the variables used by the fairness algorithm which is completely missing from the present text.*

This clause contains the following definitions and formula variables:

**9.3.1 active stations:** Calculated. The number of stations that were active (as measured by having at least 1 transited fairness eligible frame) on the outgoing link of a ringlet during the last aging interval.

**9.3.2 add rate:** Calculated. A running byte count of all FE marked packets transmitted to the ringlet from the add queues of the local station during the previous AGECOEF aging intervals (including the current interval), and smoothed over the previous AGECOEF aging intervals. This rate is enforced by the Sd shaper to be less than or equal to the allowed rate (+ 1 MTU - 1 byte).

**9.3.3 add rate congested:** Calculated. A running byte count of all FE marked packets transmitted to the ringlet from the add queues of local station during the previous AGECOEF aging intervals (including the current interval), and smoothed over the previous AGECOEF aging intervals, that have a destination past the congestion point. This rate is enforced by the Sc shaper to be less than or equal to the allowed rate congested (+ 1 MTU - 1 byte).

**9.3.4 add rate congested ok:** Calculated. A boolean value indicating if the current add rate congested is within acceptable limits. This value is an input to the sendC signal defined in Clause 6.

**9.3.5 add rate ok:** Calculated. A boolean value indicating if the current add rate is within acceptable limits. This value is an input to the sendC signal defined in Clause 6.

**9.3.6 advertised fair rate:** Calculated. The normalized rate (in bytes) advertised in an FCM giving the fair number of bytes that can be sent during the next aging interval.

**9.3.7 advertisement interval:** Configured. The interval at which FCMs are sent. This time is less than the aging interval. The default value for this constant is a rate that uses .125% of the link bandwidth (e.g. 40 usec for a LINK RATE of 2.5 Gbps).

**9.3.8 AGECOEF:** Configured. The coefficient used for aging the running add rate, add rate congested, fw rate, and fw rate congested. The default value for this constant is 4. Values other than factors of 2 are not required to be supported.

**9.3.9 aging interval:** Configured. The interval at which aging and low pass filtering functions are performed. The value to use for each supported link speed is shown in Table 9.1.

**Table 9.1—Rate Coefficient**

Link Speed	aging_interval
OC-12 and higher rates	100 usec
OC-3	400 usec

**9.3.10 aggressive:** Configured. Aggressive at maximizing link utilization. Local station fairness (during transitions). Trades off stability for utilization (during transitions).

**9.3.11 allowed rate:** Calculated. The rate at which the local station is allowed to transmit FE marked packets to the ringlet, specified as the number of bytes per AGECOEF aging intervals. During any aging interval, the MAC may send up to allowed rate minus add rate FE bytes. This value is an input to the Sd rate shaper/limiter defined in Clause 6.

**9.3.12 allowed rate congested:** Calculated. The rate at which the local station is allowed to transmit FE marked packets to the ringlet that have a destination past the congestion point, specified as the number of bytes per AGECOEF aging intervals. During any aging interval, the MAC may send up to allowed rate congested minus add rate congested FE bytes beyond the congestion point. This value is the input to the Sc rate shaper/limiter defined in Clause 6.

**9.3.13 congested:** Calculated: The indication of whether the local station considers itself to be congested in trying to add traffic.

**9.3.14 congestion point:** Calculated. The most congested station on the ringlet reporting a congestion condition (as indicated by the SA and TTL in the most recent SC-FCM).

**9.3.15 conservative:** Configured. Conservative at maximizing link utilization. Congested span fairness (during transitions). Trades off utilization for stability (during transitions).

**9.3.16 fairness eligible:** Definition. A quality of a packet that indicates (as specified in 8) if it is subject to the fairness algorithm. Packets that are subject to fairness are the class-B packets added in excess of the station's provisioned amount, and all class-C traffic. See also 6.

**9.3.17 FULL RATE:** Constant. A special value for control rate, indicating no congestion. The value for this constant is an all "1"s value.

**9.3.18 full threshold:** Configured. The STQ threshold that indicates that the STQ of a dual-queue MAC is almost full. Same as NEEDY in Clause 6. The default value for full threshold is STQ - 1 MTU.

**9.3.19 fw rate:** Calculated. A running byte count of all FE marked packets transmitted to the ringlet from the transit queues of the local station during the previous AGECOEF aging intervals (including the current interval), and smoothed over the previous AGECOEF aging intervals.

**9.3.20 fw rate congested:** Calculated. A running byte count of all FE marked packets transmitted to the ringlet from the transit queues of the local station during the previous AGECOEF aging intervals (including the current interval), and smoothed over the previous AGECOEF aging intervals, that have a destination past the congestion point.

**9.3.21 high threshold:** Constant. The threshold that indicates that a MAC is experiencing some congestion. When the STQ depth or link utilization exceed this value, the outgoing link is considered congested.

**9.3.22 LINK RATE:** Configured. Conceptually, the actual maximum transmission rate for the outbound link. For fairness formulas, it is interpreted as the maximum number of bytes that can be transmitted in AGECOEF aging intervals. This rate is less than AGECOEF \* aging interval worth of bytes.

**9.3.23 local fair rate:** Calculated. The rate at which the local station may transmit FE marked packets, regardless of downstream congestion, specified as the number of bytes per AGECOEF aging intervals.

**9.3.24 local station:** Definition. The station in which the local MAC resides.

**9.3.25 low threshold:** Constant. The threshold that indicates that a MAC may soon start experiencing congestion. When the STQ depth or link utilization exceed this value, the outgoing link is considered imminent to being congested.

**9.3.26 lp add rate:** Calculated. A low pass filtered value of add rate.

**9.3.27 lp add rate congested:** Calculated. A low pass filtered value of add rate congested.

**9.3.28 lp fw rate:** Calculated. A low pass filtered value of fw rate.

**9.3.29 lp fw rate congested:** Calculated. A low pass filtered value of fw rate congested.

**9.3.30 lp nr xmit rate:** Calculated. A low pass filtered value of nr xmit rate.

**9.3.31 LPCOEF:** Configured. The coefficient used for low pass filtering the add rate and fw rate to result in lp add rate and lp fw rate. The default value for this constant is 64. Values other than factors of 2 are not required to be supported.

**9.3.32 MAX ALLOWED RATE:** Configured. The maximum value for allowed rate. The default value for this constant is LINK RATE.

**9.3.33 multi choke:** Definition. The ability to work with multi choke points (of different values) at the same time, e.g. as in a VDO implementation. The primary intended use of the multi choke fairness control message (MC-FCM).

**9.3.34 norm local fair rate:** Calculated. The (NORMCOEF) normalized value of local fair rate. It is the value for advertised fair rate in SC-FCMs when the local station is more congested than downstream stations; and always the value for advertised fair rate in MC-FCMs.

**9.3.35 norm lp fw rate congested:** Calculated. The (NORMCOEF) normalized value of lp fw rate congested.

**9.3.36 NORMCOEF:** Configured. The coefficient used for normalizing advertised fair rates to a common rate. The value for this constant is the product of AGECOEF, RATECOEF, and WEIGHT.

**9.3.37 nr\_xmit\_rate:** Calculated. A running byte count of all non reserved (non subclass-A0 marked) packets transmitted to the ringlet (from all add queues and transit queues) of the local station during the previous AGECOEF aging intervals (including the current interval), and smoothed over the previous AGECOEF aging intervals.

**9.3.38 num\_stations:** Calculated. The number of stations that are known to be on the ringlet. Reported via rprIfNodesOnRing.

**9.3.39 RAMPCOEF:** Configured. The coefficient used for ramping up or down the allowed rate congested or local fair rate. The default value for this constant is 64. Values other than factors of 2 are not required to be supported.

**9.3.40 RATECOEF:** Configured. The coefficient used for normalizing advertised fair rates to a common link rate. The RATECOEF is used to ensure that the control value does not overflow 16-bits. The value to use for each supported link speed is shown in Table 9.2.

**Table 9.2—Rate Coefficient**

Link Speed	Rate Coefficient
Up to and including 2.5 Gbps	1
10 Gbps	4
40 Gbps	16

**9.3.41 rcvd\_fair\_rate:** Calculated. The normalized rate (in bytes) received in an FCM, giving the fair number of bytes that can be sent during the next aging interval.

**9.3.42 reserved\_rate:** Calculated. The total class A traffic that is reserved (i.e. the total of all sub class A0 traffic) on the outgoing link for all stations, including the local station, specified as the number of bytes per AGECOEF aging intervals.

**9.3.43 single choke:** Definition. The ability to work with only one choke point at one time. The primary intended use of the single choke fairness control message (SC-FCM).

**9.3.44 add\_rateTTL to congestion:** Calculated. The ~~actual amount~~ ~~number~~ of ~~fairness-eligible traffic~~ ~~sourced by hops to the node for the ringlet~~ ~~congestion point~~.

- a) ~~add\_rate\_congestion:~~ The amount of fairness-eligible traffic sourced by the node for the ringlet, and destined beyond the congestion point.
- b) ~~advertisement interval:~~ The time interval at which Type A fairness messages are sent.
- c) ~~allow\_rate\_congestion:~~ The maximum amount of fairness-eligible traffic that the node is allowed to source for destinations beyond the congestion point.
- d) ~~allowed rate:~~ The rate at which a node is allowed to transmit fairness-eligible traffic on a given ringlet.
- e) ~~available\_bandwidth:~~ This is the (link speed—reserved bandwidth).
- f) ~~decay\_interval:~~ The decay interval is a timer that is used to perform fairness-specific tasks every time it expires. It is defined in microseconds and defaults to 100-usec. Every time the decay interval timer expires, the MAC fairness control unit performs the following tasks:

- 1) ~~Set allow\_rate\_congestion with the weighted value in the received fairness message if the value is not null, otherwise ramp it up.~~
- 2) ~~Set the fill rate for the token bucket for fairness eligible traffic with the value corresponding to the allowed\_rate\_congestion (See 6.X).~~
- 3) ~~Determine if the node is congested or not.~~
- 4) ~~Low pass filtering operations are performed for add\_rate, forward\_rate, add\_rate\_congestion, and forward\_rate\_congestion.~~
- 5) ~~add\_rate, add\_rate\_congestion, forward\_rate, and forward\_rate\_congestion are decremented (decayed).~~
- g) fair\_rate: The fair rate is a normalized byte count measured over a decay interval.
- h) fairness\_eligible\_traffic: Traffic whose access to the ring is controlled based on the fairness algorithm. It consists of EIR Class B and Class C traffic only.
- i) forward\_rate: The amount of fairness eligible traffic that transits through the station.
- j) forward\_rate\_congestion: The amount of fairness eligible traffic that transits through the station and that goes beyond the point of congestion.
- k) line\_rate: The speed of the link expressed in bytes per decay interval.
- l) max\_rate: The maximum possible transmit rate of fairness eligible traffic for the node.

**Editors' Notes:** To be removed prior to final publication.

More detail is probably needed for the definition of ramp up.

- m) ~~ramp\_up: Increasing the allowed rate at a station decay interval.~~
- n) ~~reserved\_rate: The reserved bandwidth at a node. This includes the total outgoing Class A0 traffic from the node. It is the sum of the transit Class A0 traffic and the Class A0 traffic sourced by the node.~~
- o) ~~revd\_advertised\_rate: The control value received in a Type A message.~~
- p) ~~STQ\_HI\_THRESH: A high threshold for the secondary transit queue in dual queue MAC.~~
- q) ~~STQ\_LO\_THRESH: A low threshold for the secondary transit queue in a dual queue MAC.~~

**Editors' Notes:** To be removed prior to final publication.

Clarify which rates have a corresponding low pass filtered value that is maintained. Also clarify that the advertised rates are the low pass filtered ones. The low pass filtering options should be defined in the fairness clause. Contributions are requested for this.

**9.3.45 unreserved\_rate:** Calculated. The rate available for the outbound link that is not reserved, specified as the number of bytes per AGEcoef aging intervals.

**9.3.46 WEIGHT:** Configured. The local weight by which allowed rate, allowed rate congested, and local fair\_rate are calculated. The default value for this constant is 1.



## 9.4 MAC fairness operation

**Editors' Notes:** To be removed prior to final publication.

The following summary describes the generation of fairness messages. Make sure that the document is consistent with the following description.

1) When a node is congested.

How to generate Type As?

If the node is more congested than the node from which it received the Type A message, then the local node advertises its fair rate (with its SA). Otherwise, the downstream node's value is passed without modifying the SA.

How to generate Type Bs?

Local node advertises its fair rate using its own SA and with a TTL of 255.

2) When a node is not congested.

How to generate Type As?

If the following condition is true:

(forward\_rate\_congestion > allow\_rate\_congestion)

the downstream node's value is passed without modifying the SA. Otherwise, a NULL rate is sent in the Type A's.

How to generate Type Bs?

Local node advertises its fair rate using its own SA and with a TTL of 255. It may use a NULL value for the fair rate.

Clarify that fairness messages are sent only at fairness intervals.

A figure to be added describing operation of Type A's and Type B's.

~~The RPR-FA is a local fairness mechanism that enforces fairness among the stations on the ring. It controls the amount of fairness-eligible traffic (i.e. EIR Class B and Class C traffic) that the MAC is allowed to insert or each ringlet. RPR-FA is implemented within a control entity called the Fairness Control Unit (FCU). The MAC utilizes the topology information including the ringlet selection preference passed by the MAC client to perform fairness and policing functions.~~

The fairness algorithm implemented within the FCU consists of the following functions:

- a) Determining when the congestion threshold is crossed and when the congestion has subsided;
- b) Determining the fair rate for advertisement;
- c) Determining the station's allowed rate;
- d) Sourcing and consuming fairness control messages;
- e) Communicating the allowed rate to the data path shapers for controlling access to the medium;
- f) Providing the information contained in MC-FCMs to the client.

Each station is assigned a weight, which allows the user to allocate more ring bandwidth to certain stations as compared with other stations. This is referred to as the *weighted fairness* property of RPR-FA.

In RPR-FA, a ~~node-station~~ advertises a fair rate (~~i.e. add\_rate~~) to upstream ~~nodes-stations~~ via the ringlet opposite ~~ring of the ringlet upon which the algorithm is running~~. The fair rate ~~counter~~ is run through a low pass filter ~~function and divided function, with the result known as the local fair rate~~. The local fair rate is normalized by a weighting function (i.e. the local station weight WEIGHT), the aging coefficient (AGECOEF), and the rate coefficient (RATECOEF) to yield the norm local fair rate and the advertised fair rate. The low-pass filter stabilizes the feedback, ~~and the division by weight-WEIGHT~~ normalizes the transmitted value to a weight of 1.0, the division by AGECOEF normalizes the transmitted value to one aging interval, and the division by RATECOEF normalizes the transmitted value to a link speed of 2.5 Gbps. When the upstream stations receive an advertised fair rate, they will adjust their transmit rates for fairness eligible traffic so as not to exceed the advertised value (adjusted by their respective weights).



Propagation of the advertised value to other ~~nodes-stations~~ on the ring is done using fairness control messages. The format of the fairness control messages is described in 9.12. There are 2 types of fairness control messages—~~Type A~~ Single Choke and ~~Type B~~ Multi Choke.

Single Choke messages are propagated hop-by-hop around the opposite ringlet and are processed by the FCU. They are sent every advertisement interval.

A wrapped ring shall be treated as a folded ring from the point of view of SC-FCMs. A station with a wrapped attachment point receiving a SC-FCM shall change the ringlet id and wrap the SC-FCM. SC-FCMs are stripped only when received with SA equal to the local MAC address and with the correct ringlet id.

Single Choke fairness control messages contain the SA of the most congested station. If a station experiences congestion, it will include a non-FULL RATE value for the fair rate in its Single Choke fairness control messages. A station that receives a Single Choke message with a SA of its MAC address shall treat the message as having a control value of FULL RATE regardless of the actual control value.

A station that in the congested state shall advertise the minimum of its local fair rate and the last received fair rate (known as the rcvd fair rate) in its Single Choke message.

~~Type A~~ messages are propagated hop-by-hop around the opposite ringlet and are processed by the FCU. They are sent every advertisement interval. They contain the SA of the most congested node. If a node experiences congestion, it will include a non-NULL value for the fair rate in its ~~Type A~~ fairness messages. A node that receives an advertised value in a ~~Type A~~ message, and station that is also in the congested state, shall advertise the minimum of its normalized low pass filtered fair rate and the received fair rate in its ~~Type A~~ message. A node that is not congested and that receives a ~~Type A~~ Single Choke message containing a non-NULL rcvd\_advertised\_rate FULL RATE rcvd fair rate shall either propagate the rcvd\_advertised\_rate rcvd fair rate to its upstream neighbor (leaving the SA the same), or it will send a value of NULL FULL RATE and set the SA to its own. Which of these is sent is determined by the following. If the low pass filtered forward\_rate fw\_rate congested is less than the ~~allow\_rate\_congestion~~ allowed rate congested divided by the local station weight, then a NULL FULL RATE value is propagated to the upstream neighbor instead of the rcvd\_advertised\_rate rcvd fair rate. Otherwise, there cannot be an upstream ~~node-station~~ that is the cause of congestion. Thus there is no need to propagate the ~~Type A~~ Single Choke fairness control message indicating congestion upstream of this ~~node-station~~.

The value of the rate that is allowed ~~rate~~ for the fairness eligible traffic (allowed rate congested) that can be added by the ~~node-station~~ is derived from the received advertised rate in the ~~Type A~~ Single Choke messages and the local station ~~weight~~ normalization factor (NORMCOEF). If the number of hops to the destination of the a data packet from the MAC client is less than the number of hops to the ~~node-station~~ that generated the ~~Type A~~ Single Choke message, then the MAC client can take advantage of the available bandwidth on the ringlet; otherwise, if the destination of the a data packet would allow the packet to go beyond the ~~node-station~~ that generated the fairness control message, the client will at least receive its fair share of the bandwidth from the most congested span that it contends for. ~~The fair rate is equal to the allowed rate adjusted by the local station weight.~~

~~Type B~~ Multi Choke messages are broadcast on the ringlet opposite to the ringlet upon which the fairness algorithm is running and contain the SA of the ~~node-station~~ that originated the message. They are sent every 10 ~~advertisement intervals~~ advertisement intervals. If a ~~node-station~~ experiences congestion, it will include a non-NULL FULL RATE value for the fair rate in its ~~Type B~~ Multi Choke fairness control messages; otherwise it will include a FULL RATE value for the fair rate. ~~Type B~~ Multi Choke messages are not required to

be processed by the FCU, but the information is passed to the MAC client by the FCU and is used by the MAC client as described below.

**Editors' Notes:** To be removed prior to final publication.

*What do Type B messages contain if the node is not congested? According to the description in the previous editors' note (which was agreed to at the last comment resolution session) the rate may be a "fair rate" or it may be NULL. Since this seems to be a contentious issue, I have left it open.*

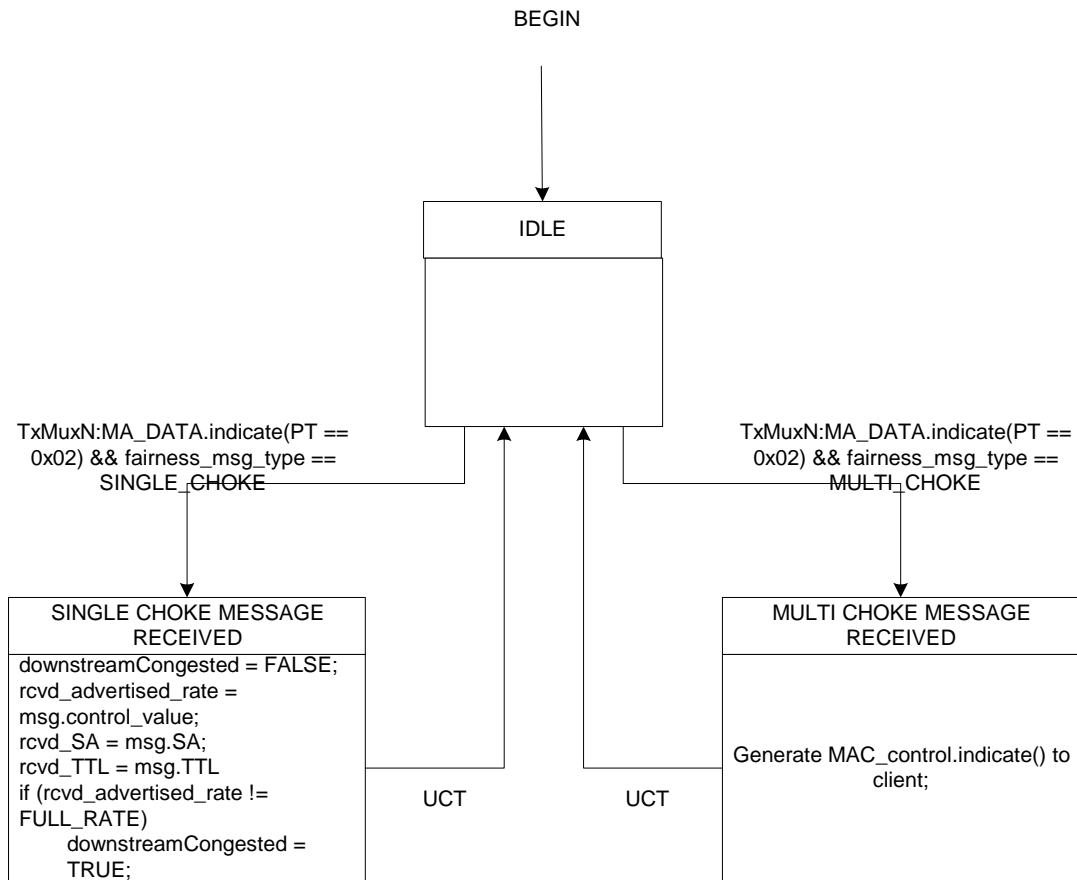
NOTE—The client can also take advantage of Virtual Destination Queueing (VDQ) by utilizing the multi-choke concept of RPR-FA. VDQ combined with RPR-FA can increase ring utilization. The multi-choke concept deals with the case where a ~~node-station~~ wants to send traffic to a destination that is closer than a congested link. As an example, consider the case where ~~node-station~~ 1 wants to send traffic to ~~node-station~~ 2, and the link between ~~nodes-stations~~ 2 and 3 is congested. RPR-FA will allow ~~node-station~~ 1 to send as much traffic as it wants to ~~node-station~~ 2, and will only limit traffic to ~~nodes-stations~~ beyond the congested link to the fair rate. In a multi-choke implementation of the RPR-FA, each client will track advertised fair rates for congested ~~nodes-stations~~. A ~~node-station~~ is allowed to send unlimited traffic to any ~~node-station~~ between itself and the first congested ~~node-station~~ (choke point). It can send traffic to ~~nodes-stations~~ between the first and second choke point based on the first choke point's advertised fair rate. In general, a ~~node-station~~ can send traffic to a particular destination if it has satisfied the fair rate conditions for all choke points between itself and the destination. The maximum possible number of choke points is equal to the number of ~~nodes-stations~~ on the ring.

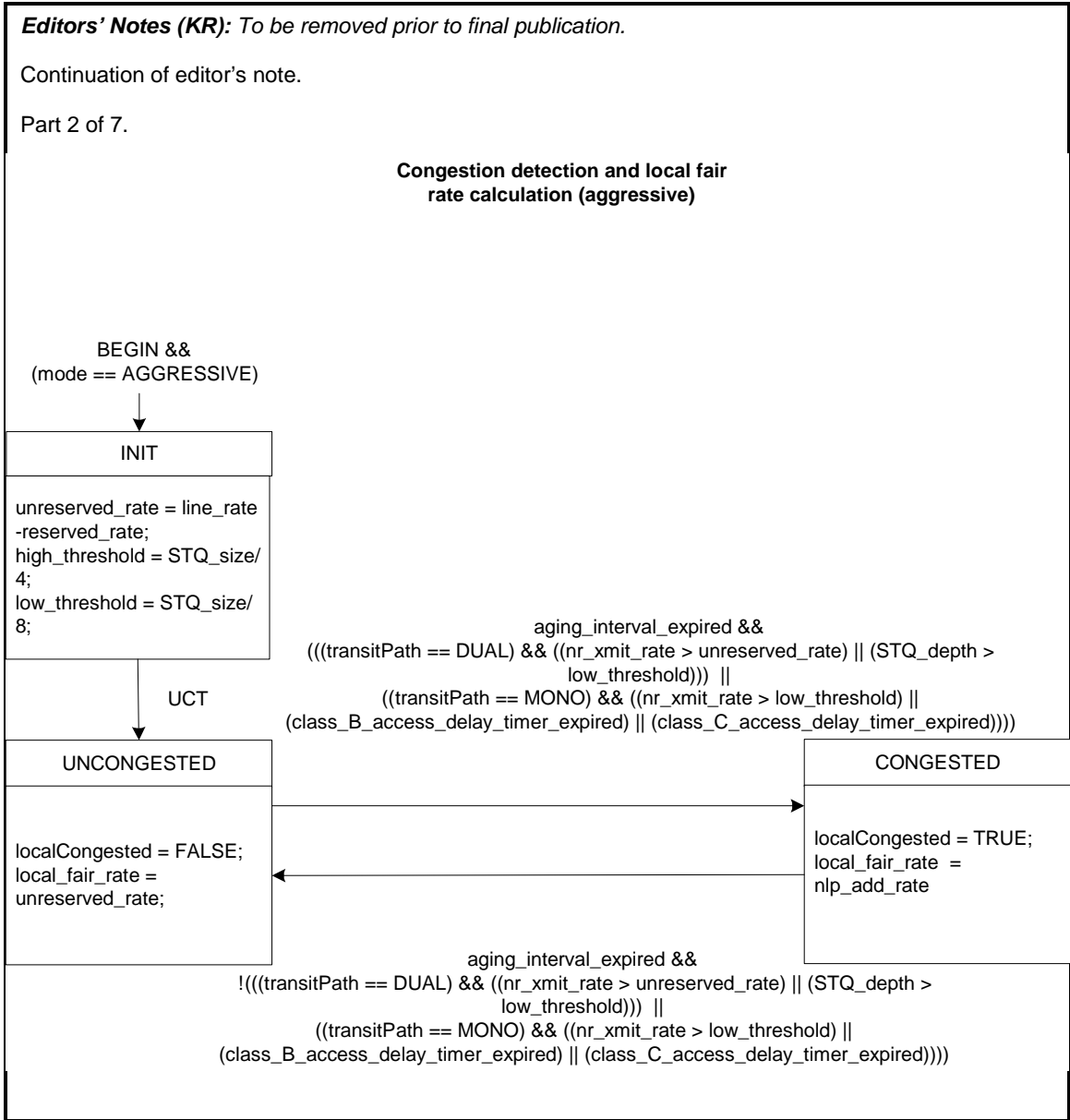
**Editors' Notes (KR):** To be removed prior to final publication.

The following state diagrams are intended to provide additional information to help in the understanding and completion of the state tables that follow them. These need to be reconciled with the tables and the text.

Part 1 of 7.

### Receive Fairness Message





98

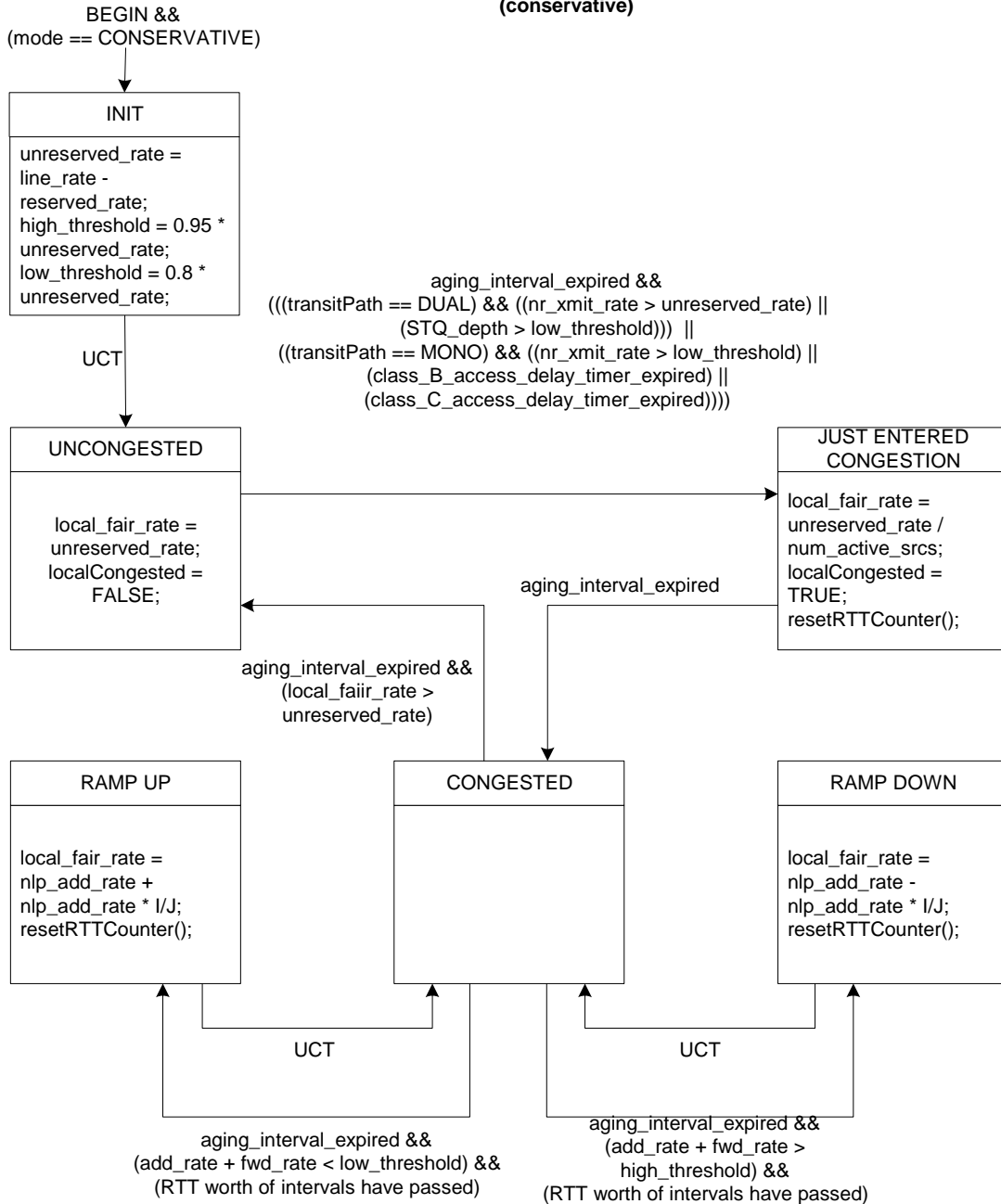
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**Editors' Notes (KR):** To be removed prior to final publication.

Continuation of editor's note.

Part 3 of 7.

**Congestion detection and Local  
fair rate calculation  
(conservative)**

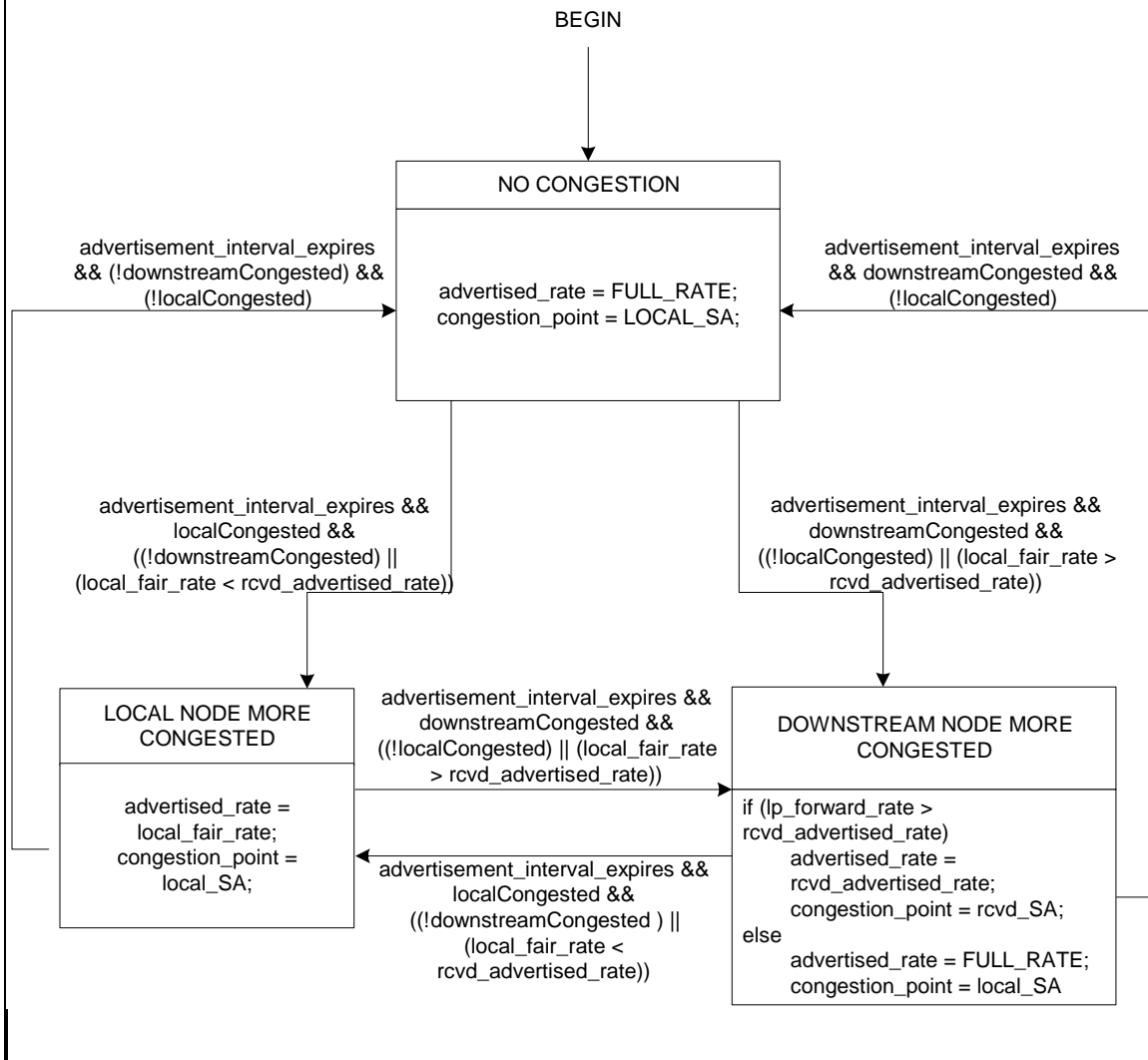


**Editors' Notes (KR):** To be removed prior to final publication.

Continuation of editor's note.

Part 4 of 7.

### Generation of single choke message

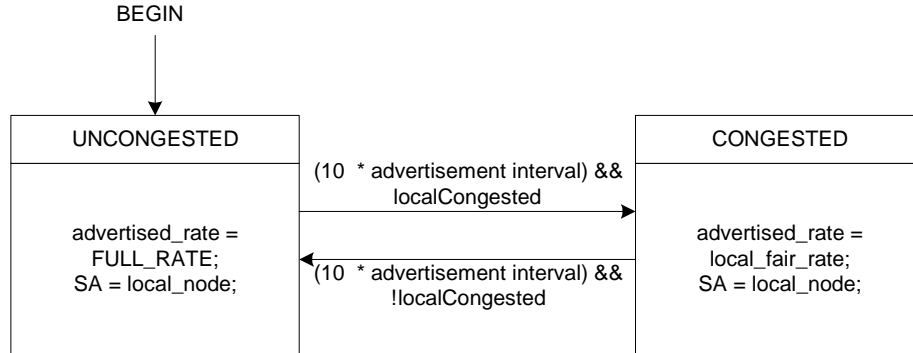


**Editors' Notes (KR):** To be removed prior to final publication.

Continuation of editor's note.

Part 5 of 7.

#### Generation of multi choke message

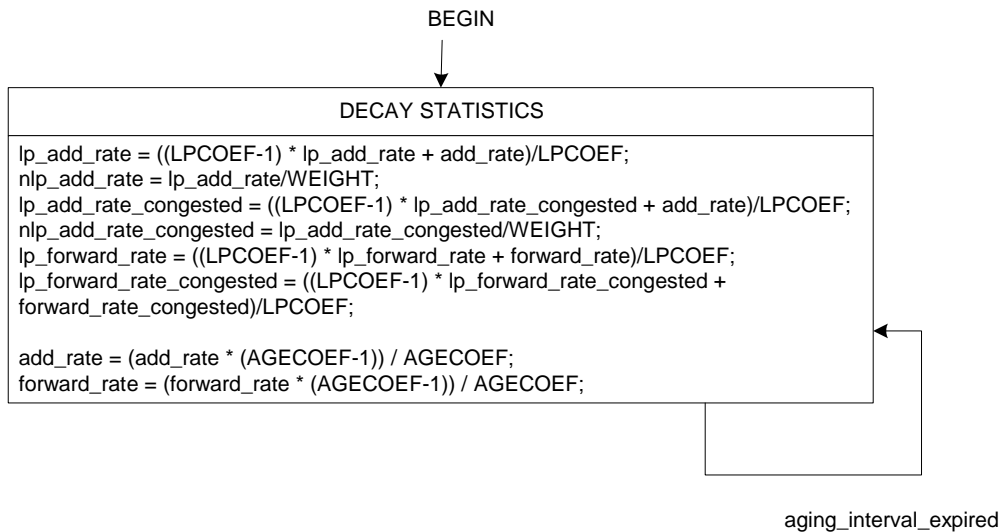


**Editors' Notes (KR):** To be removed prior to final publication.

Continuation of editor's note.

Part 6 of 7.

#### Statistics Decay

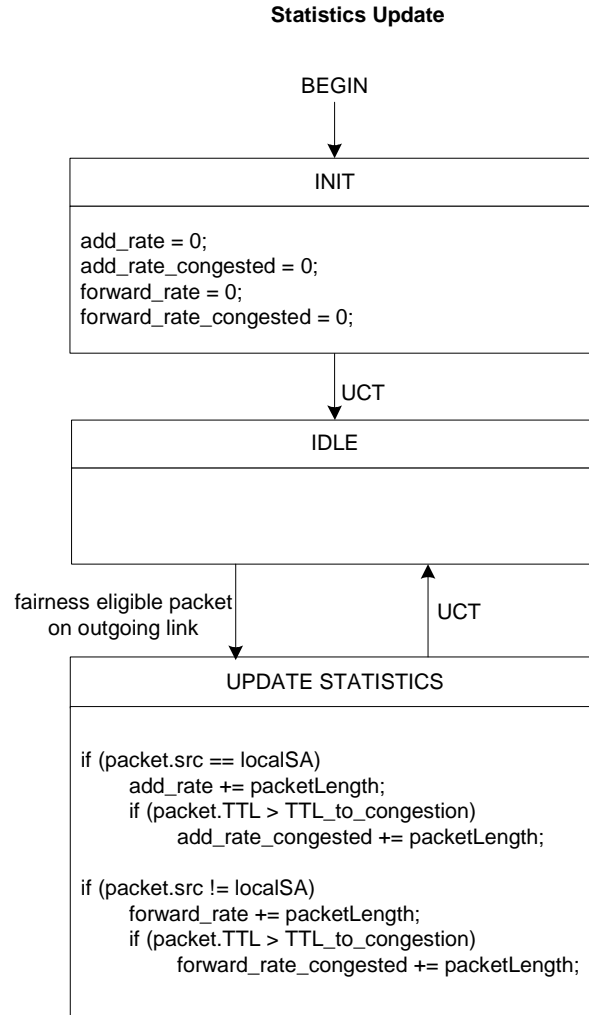




**Editors' Notes (KR):** To be removed prior to final publication.

Continuation of editor's note.

Part 7 of 7.



9.4.1 FCM Processing

9.4.1.1 FCM Receive

For every aging interval, all variables used to determine congestion condition, congestion location, and rate control are recomputed. Those values that depend upon values received in SC-FCMs use the most recently received SC-FCM. The flow chart and state machine for processing received FCMs are as follows:

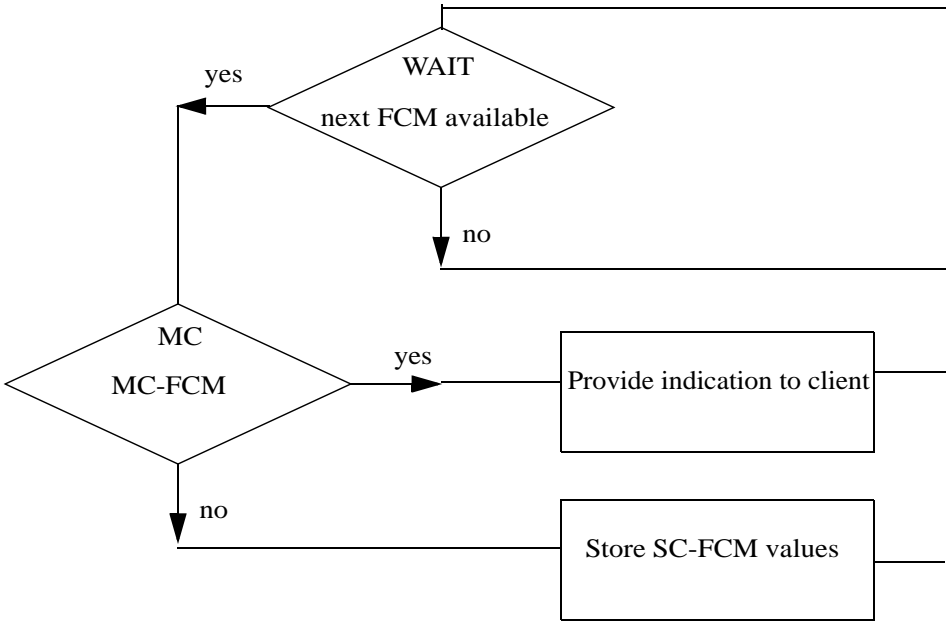


Figure 9—1— Fairness Control Message Reception

Table 9.3—FCM Reception States

Last state		Row	Next state	
state	condition		action	state
MC	MC-FCM	9.0.1	MA_CONTROL.indication(msg);	WAIT
	—	9.0.2	rcvd_fair_rate = msg.control_value; rcvd_SA = msg.SA; rcvd_TTL = msg.TTL;	
WAIT	next FCM available	9.0.3	—	MC
	—	9.0.4	—	WAIT

**Row 9.0.1:** The FCM is a MC-FCM. Make its relevant values available to the client.

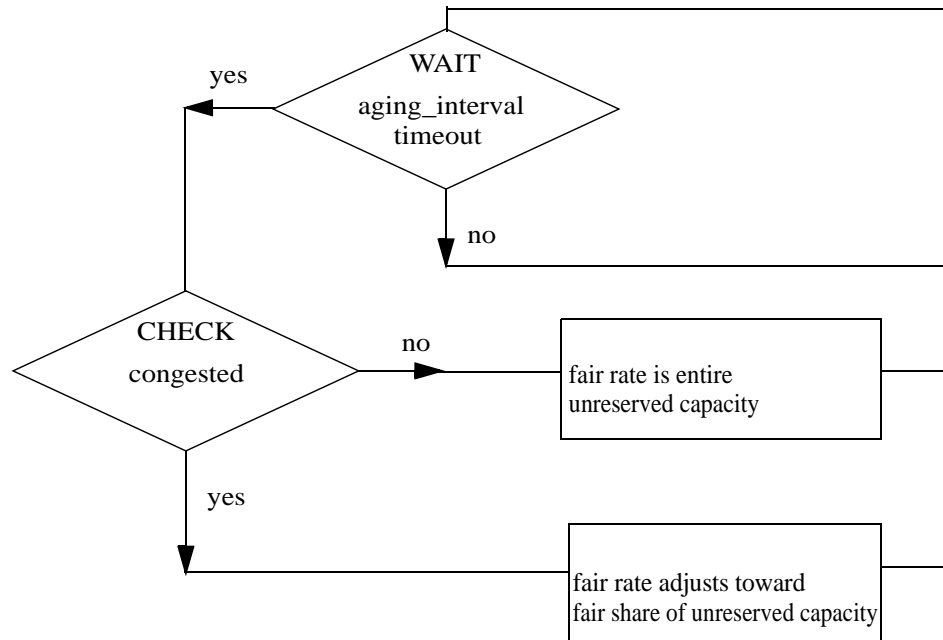
**Row 9.0.2:** The FCM is a SC-FCM. Store its relevant values for use after the next aging interval expires.

Row 9.0.3: When a new FCM is received, start processing it.

Row 9.0.4: Wait for another FCM.

#### 9.4.1.2 Local fair rate calculation

Using the most recently received SC-FCM, the local fair rate is recomputed every aging\_interval.



**Figure 9—2—Local Fair Rate Calculation**

**Table 9.4—Local Fair Rate Calculation States**

Last state		Row	Next state	
state	condition		action	state
CHECK	!congested	9.0.1	local_fair_rate = unreserved_rate;	WAIT
	aggressive	9.0.2	local_fair_rate = lp_add_rate;	
	just_entered_congestion	9.0.3	set local_fair_rate to weight adjusted equal share of unreserved_rate	
	FE transmit rate below low_threshold	9.0.4	ramp up local_fair_rate	
	FE transmit rate above high_threshold	9.0.5	ramp down local_fair_rate	
	—	9.0.6	—	
WAIT	aging_interval timeout	9.0.7	recalculate the world;	CHECK
	—	9.0.8	—	WAIT

**Row 9.0.1:** Outgoing link is not congested. Local station can try to use entire (unreserved) capacity.

**Row 9.0.2:** Aggressively try to use current add rate.

**Row 9.0.3:** When first becoming congested, try to use an equal share of the unreserved capacity, adjusted for station weight.

**Row 9.0.4:** FE transmit rate is too low. Ramp up fair rate towards unreserved capacity.

**Row 9.0.5:** FE transmit rate is too high. Ramp down fair rate towards 0.

**Row 9.0.6:** FE transmit rate is not too low and not too high. Don't change amount of outgoing link currently trying to use.

**Row 9.0.7:** When the next aging\_interval expires, recalculate all aging\_interval dependent variables, except for local\_fair\_rate and norm\_local\_fair\_rate.

**Editors' Notes (JL):** *To be removed prior to final publication.*

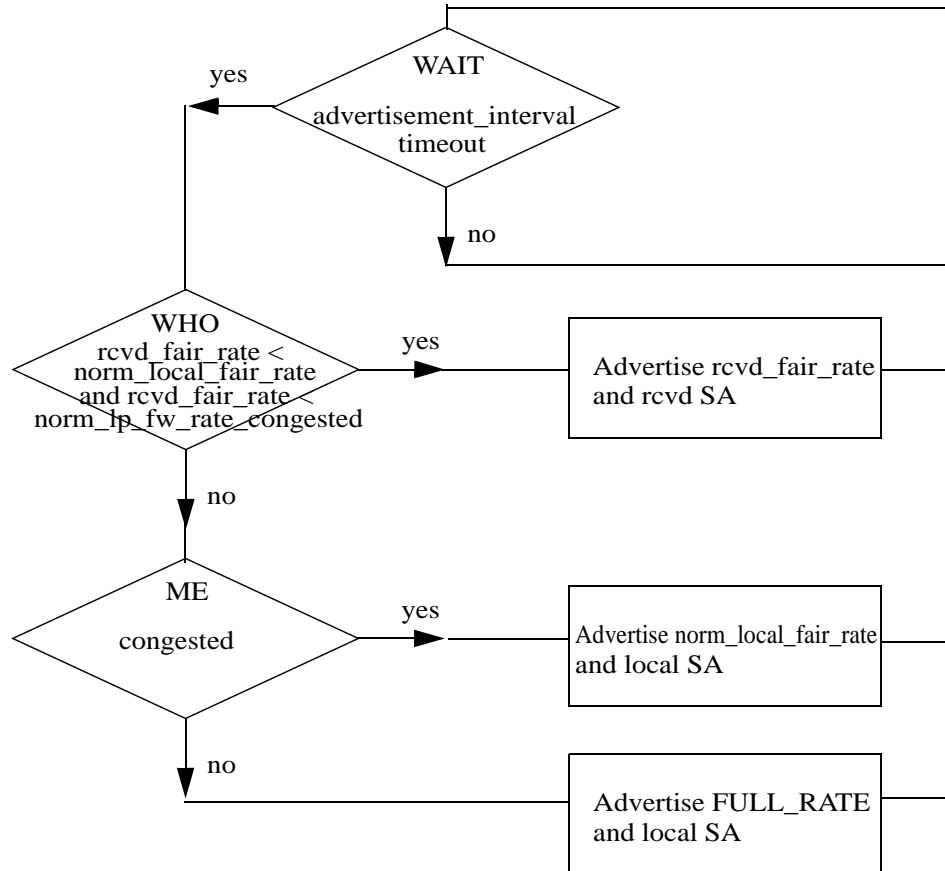
Should 9.8.3 be referenced, or should the calculations/formulas be stated in this section? Also, is it helpful to describe the calculation of local\_fair\_rate as state machine instead of just a formula?

**Row 9.0.8:** Wait for the aging\_interval to expire.

## 9.4.2 FCM Transmit

### 9.4.2.1 SC-FCM Transmit

For every advertisement interval, the advertised fair rate and congested SA are recomputed (as described in 9.13.4 ) and advertised in a SC-FCM.



**Figure 9—3—Single Choke Fairness Control Message Advertisement**

**Table 9.5—SC-FCM Advertisement States**

Last state		Row	Next state	
state	condition		action	state
WHO	rcvd_fair_rate < norm_local_fair_rate && rcvd_fair_rate < norm_lp_fw_rate_congested	9.0.1	advertised_fair_rate = rcvd_fair_rate; congestion_point = rcvd_SA;	WAIT
	—	9.0.2	—	ME
ME	congested	9.0.3	advertised_fair_rate = norm_local_fair_rate; congestion_point = local_SA;	WAIT
	—	9.0.4	advertised_fair_rate = FULL_RATE; congestion_point = local_SA;	
WAIT	advertisement_interval timeout	9.0.5	—	WHO
	—	9.0.6	—	WAIT

**Row 9.0.1:** [A downstream station is more congested than the local station.](#)

**Row 9.0.2:** [No downstream station is more congested than the local station](#)

**Row 9.0.3:** [The local station is congested.](#)

**Row 9.0.4:** [The local station \(and therefore all stations\) is not congested.](#)

**Row 9.0.5:** [When the next advertisement interval expires, transmit another SC-FCM.](#)

**Row 9.0.6:** [Wait for the advertisement interval to expire.](#)

9.4.2.2 MC-FCM Transmit

For every 10 advertisement intervals, the local fair rate is advertised in a MC-FCM.

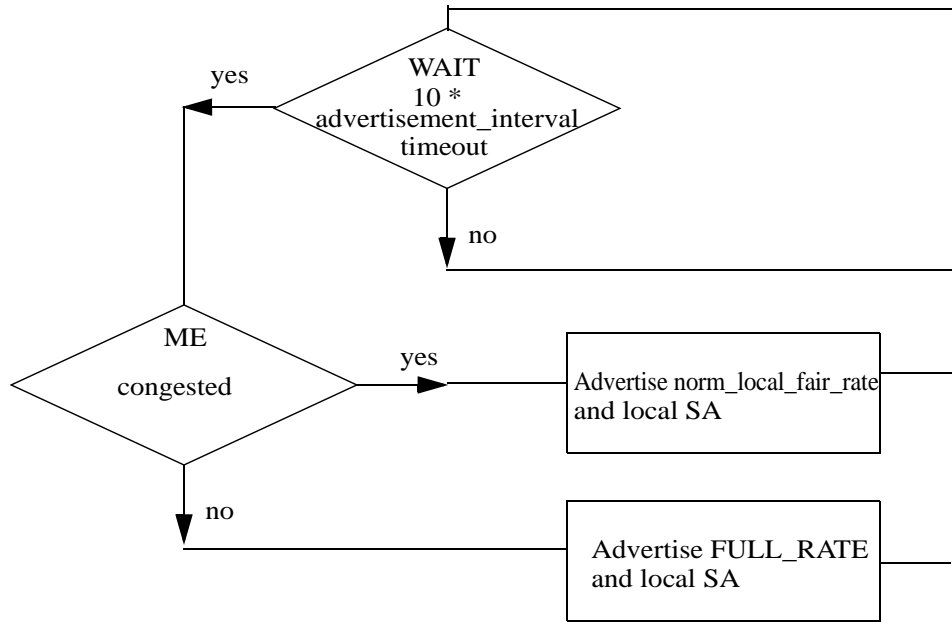


Figure 9—4—Multi Choke Fairness Control Message Advertisement

Table 9.6—MC-FCM Advertisement States

Last state		Row	Next state	
state	condition		action	state
ME	congested	9.0.1	advertised_fair_rate = norm_local_fair_rate; congestion_point = local_SA;	WAIT
	—	9.0.2	advertised_fair_rate = FULL_RATE; congestion_point = local_SA;	
WAIT	10 * advertisement_interval timeout	9.0.3	—	ME
	—	9.0.4	—	WAIT

Row 9.0.1: The local station is congested.

Row 9.0.2: The local station is not congested.

Row 9.0.3: When 10 advertisement intervals expire, transmit another MC-FCM.

Row 9.0.4: Wait for 10 advertisement intervals to expire.



## 9.5 Congestion detection

**Editors' Notes:** To be removed prior to final publication.

In the text below, we need to have a better definition for "configured thresholds" to detect coming out of congestion.

Congestion is declared by using the following ~~criteria depending on the number of transit queues used by the MAC~~ criteria. The formula for determining congestion is specified in 9.13.3.4.

~~If the MAC uses a mono transit queue design, congestion~~ Congestion is detected when any of the following conditions is true:

**Editors' Notes:** To be removed prior to final publication.

How does this handle CIR class B traffic? It is not accounted for in either reserved rate not in the add and forward rates. If is included in the reserved rate, then it prevents unused CIR Class B bandwidth from being reclaimed.

Also, we should probably define HI and LO link utilization thresholds and queue thresholds that we use for hysteresis for congestion detection and coming out of congestion. Contributions are invited for this.

- ~~a) The rate of outgoing non-reserved traffic passes a configured congestion threshold on the output link; i.e. if the sum of Class A1, Class B, and Class C traffic on the outgoing link is more than (line\_rate - reserved\_rate). The outgoing link rate is measured using a byte counter which is passed through a low pass filter before the comparison. Further, the node will be considered to be congested until its outgoing link utilization falls below a configured threshold.~~
- a) The rate of outgoing non-reserved traffic (lp\_nr\_xmit\_rate) is more than the unreserved rate (i.e. LINK\_RATE - reserved\_rate), or the low threshold, if specified as a rate.
- b) The access delay timer for Class B add packets expires.
- c) The ~~access delay timer for Class B add packets or the~~ access delay timer for Class C add packets expires.

~~If the MAC uses a dual transit queue design, congestion is detected when any of the following conditions is true:~~

- ~~a) The rate of outgoing non-reserved traffic passes a configured congestion threshold on the output link; i.e. if the sum of Class A1, Class B, and Class C traffic on the outgoing link is more than (line\_rate - reserved\_rate). The outgoing link rate is measured using a byte counter which is passed through a low pass filter before the comparison. Further, the node will be considered to be congested until its outgoing link utilization falls below a configured threshold.~~
- b) ~~The depth of the STQ reaches a congestion threshold.~~
- c) The depth of the STQ exceeds the low threshold congestion depth threshold, or the rate of non-A0 traffic exceeds the low threshold congestion rate threshold.

Conditions b) and c) are optional for dual-queue MACs. Condition c) is not applicable for mono-queue MACs.

The above metrics are maintained independently for each ringlet. Any of them could indicate the onset of congestion for the ringlet. The absence of all that are maintained indicates the lack of congestion.

**Editors' Notes (JL):** To be removed prior to final publication.

Adding (optional) hysteresis to each of a - d need to be added.

The access delay timers are maintained by the MAC for each ringlet. There are separate access delay timers for Class B and Class C traffic. The timer corresponding to the class of packet is started at the time the packet becomes the head-of-line packet in the MAC. The timer is reset when the MAC successfully transmits that packet. If the MAC is unable to transmit the packet before the timer expires, the MAC enters a congested state. The access delay timers will have a range of 1-10 msec, with a default value of 5 msec.

**Editors' Notes:** To be removed prior to final publication.

*Is there any justification for the above numbers? It should probably depend at least on the link speed and the RTT.*

*For Class A traffic there is no access delay timer. The access delay timer expiring for Class A traffic means there is something wrong with respect to provisioning. Perhaps OAM&P should look into providing an alarm for this condition. However, this issue doesn't concern fairness.*

~~For Class B and Class C traffic, the access delay timer will have a range of 1-10 msec with a default value of 5 msec.~~

**Editors' Notes:** To be removed prior to final publication.

*Is there any justification for the above numbers? It should probably depend at least on the link speed and the RTT.*

*For Class A traffic there is no access delay timer. The access delay timer expiring for Class A traffic means there is something wrong with respect to provisioning. Perhaps OAM&P should look into providing an alarm for this condition. However, this issue doesn't concern fairness.*

## ~~9.6 Interoperability between mono and dual queue MACs~~

**Editors' Notes:** To be removed prior to final publication.

*There is no new information in this subclause. Suggest deleting it. Unless we want to say something about shaping the STQ in a dual queue transit path design to enable interoperability between mono and dual transit queue MACs.*

~~If a ring consists of mixed RPR MACs (mono and dual transit queue nodes), the fairness scheme will need to interact without disadvantaging any other node on the ring. The actions required for the interaction between mono and dual transit queue nodes is the same as the interaction between similar nodes. Upon receiving a Type A fairness message the RPR MAC shall reduce its allowed rate to the fair rate received in the fairness message multiplied by the local weight, and then forward the message upstream with the minimum of its own advertised rate and the received fair rate.~~

~~The RPR MAC should rate shape its transmit traffic using the dynamic traffic shaping algorithm described in XX—Dynamic traffic shaping.~~

## 9.7 ~~Threshold settings for dual transit queue RPR MACs~~

**Editors' Notes:** To be removed prior to final publication.

Since this clause deals with the onset of congestion, the actual threshold settings probably belong in the clause on congestion detection. The rationale, etc. should probably be moved to Annex I: Implementation Guidelines.

~~Two registers, STQ\_HI\_THRESH and STQ\_LO\_THRESH, are used to specify congestion thresholds within a large STQ. When the occupancy of the STQ crosses STQ\_LO\_THRESH, the node is considered to be in a congested state, and no more fairness eligible traffic is accepted from the MAC client for transmission on to the ringlet. If the occupancy of the STQ crosses STQ\_HI\_THRESH, it is almost full and the traffic from the STQ will be prioritized above all traffic from the MAC client until the occupancy of the STQ drops below STQ\_HI\_THRESH.~~

When a station determines that it is a congestion point, it will send a fairness control message containing a non-FULL\_RATE fair rate, based on a normalized value of its low pass filtered add rate.

### 9.7.1 Threshold settings

Three thresholds, full threshold, high threshold, and low threshold, are used to specify congestion thresholds. For dual-queue MACs, they are based on occupancy of the STQ. For mono-queue MACs, they are based on the rate at which traffic is flowing through the MAC. The formulas for calculating these thresholds are provided in 9.13.1. When the low threshold is crossed, the station is considered to be approaching a congested state, and advertises a reduced fair rate. When the high threshold is crossed, the station is considered to be in a congested state, and no more fairness eligible traffic is accepted from the MAC client for transmission on to the ringlet. For dual-queue MACs, if the occupancy of the STQ crosses full threshold, it is almost full, and the traffic from the STQ will be prioritized above all other traffic until the occupancy of the STQ drops below full threshold (by any amount).

**Editors' Notes (JL):** To be removed prior to final publication.

The remainder of this section deals with rationale, and should probably be moved to Annex I: Implementation Guidelines.

The setting of ~~STQ\_LO\_THRESH~~high threshold depends on the following factors:

- a) If set too low, it will trigger congestion reports too frequently which in turn will adversely affect the link utilization.
- b) If set higher, the difference between the ~~STQ\_LO\_THRESH~~high threshold and ~~STQ\_HI\_THRESH~~full threshold is reduced. Because this reduces the time between congestion onset and the STQ filling up, it has the following effect:
  - 1) For a given level of reclaimable Class A bandwidth (i.e. Class A1 bandwidth), the link distances must be shorter.
  - 2) For a given link distance, a smaller level of reclaimable Class A traffic can be supported. (See Clause 6.x where this calculation is provided.)

The more severe consequences of a) vs. b) are the reason that a threshold of less than 50% is chosen. This recommendation applies only when a STQ of many MTUs is used to support high amounts of reclaimable Class A bandwidth.

**Editors' Notes:** To be removed prior to final publication.

As requested at the May meeting, Vasanth Karigattam will have a contribution for an analytical method to compute the STQ size.

The primary transit queue needs to hold at least 2 MTUs.

**Editors' Notes:** To be removed prior to final publication.

As requested at the May meeting, Necdet Uzun and David James will have a justification for this threshold.

~~STQ\_LO\_THRESH should be set to about 25% of the total buffer available. STQ\_HI\_THRESH should be set to at least 1 MTU less than the total buffer size.~~

NOTE—~~high threshold should be set to about 25% of the total buffer available. full threshold should be set to at least 1 MTU less than the total buffer size.~~ The recommendation for the setting of the threshold values is done based on delivering the best possible end-to-end delay for the low priority traffic without penalizing the Class A traffic. Lower values will result in higher end-to-end delays for low priority data packets. If either Class A, B, or C traffic is extremely bursty, then a lower threshold value should be considered. A ~~STQ\_HI\_THRESH~~ full threshold of at least 1 MTU less than the total buffer size is needed to ensure that the ~~node~~ station has sufficient buffer space for a packet that may arrive on the transit path. If the Class A traffic has a bursty ~~nature~~ nature, a more conservative (i.e. lower) value of ~~STQ\_LO\_THRESH~~ high threshold is recommended in order to avoid overflow of the STQ.

## 9.8 Traffic policing function

~~RPR-FA utilizes the allow\_rate\_congestion and TTL\_to\_congestion registers along with the add\_rate\_congestion counter to police the fairness eligible traffic that is added by the node to the ringlet. The allow\_rate\_congestion and TTL\_to\_congestion registers store the values of the fair rate and the distance to the choke point of the most recently received Type A fairness packet, respectively. The add\_rate\_congestion counter accumulates the number of bytes of fairness eligible traffic sent beyond the congestion point on the ringlet (i.e., only packets whose destination is further than TTL\_to\_congestion contribute to add\_rate\_congestion). If the add\_rate\_congestion exceeds allow\_rate\_congestion, backpressure signals are generated to the client.~~

## 9.9 Dynamic traffic shaping

**Editors' Notes:** To be removed prior to final publication.

This section should probably be combined with the previous one, into a subsection called Access Control.

~~RPR-FA limits the ring access rate for fairness eligible traffic destined beyond the congestion point to within allow\_rate\_congestion. Due to the dynamic nature of RPR-FA, allow\_rate\_congestion can vary widely from one decay interval to the next. When access traffic contains transmission bursts, its rate can vary significantly and is only limited by allow\_rate\_congestion. As a result, traffic on the ring can be bursty. To achieve better jitter performance, the ring access rate should be dynamically shaped to the allow\_rate\_congestion, which allows ring access rate conform to a fair rate as governed by RPR-FA. This reduces extra bursty transmission.~~

## 9.10 Interaction with data path

RPR-FA uses the add\_rate, add\_rate\_congested, fw\_rate, and fw\_rate\_congested values provided by the data path, as specified in Clause 6.

RPR-FA provides the add\_rate\_ok, add\_rate\_congested\_ok, allowed\_rate, allowed\_rate\_congested, and TTL\_to\_congestion values to the data path rate control function to police the fairness eligible traffic that is added by the station to the ringlet. The add\_rate\_ok and add\_rate\_congested\_ok variables are used to enable or disable the sendC signal and to determine if the TTL\_to\_congestion value should be applied. The TTL\_to\_congestion variable stores the value of the distance to the present or more recent choke point. The allowed\_rate\_congested variable stores the value of the fair rate for fairness eligible traffic passing through the congestion point, as given by the most recently received Single Choke fairness control message. The allowed\_rate variable stores the value of the rate for all fairness eligible traffic. The data path maintains rate shapers to limit the add\_rate (Sd) and add\_rate\_congested (Sc) values below the provided allowed\_rate and allowed\_rate\_congested values, respectively, and to shape their output at or below the provided rates to smooth bursty transmissions.

**Editors' Notes (JL):** To be removed prior to final publication.

This following paragraph belongs in clause 6.

The RPR-FA dynamic shaper consists of a token bucket. The token bucket has a default maximum depth of MTU bytes. The tokens are generated at a rate equal to the ~~allow\_rate\_congestion~~ allowed\_rate\_congested. For every byte that is transmitted, the number of tokens is decremented by one. When a packet is waiting for access to the ring at the head of the queue, it will be granted ring access only if its rate conforms to the fair rate as computed by RPR-FA and there is at least one byte in the token bucket. It is therefore possible for the number of tokens to become negative after subtracting the number of bytes in the transmitted packet. The maximum number of tokens that can accumulate in the token bucket is equal to the size of an MTU.

## 9.11 ~~RPR ring access operation~~

**Editors' Notes:** To be removed prior to final publication.

This whole section is kind of redundant with the clause on "RPR fairness overview". These should probably be merged.

~~The RPR-FA governs access to the ring. The RPR-FA only applies to fairness eligible traffic. Class A and "within CIR" Class B traffic does not follow RPR-FA rules and may be transmitted at any time as long as traffic shaping/policing allow it, STQ is not almost full and PTQ does not have a packet.~~

**Editors' Notes:** To be removed prior to final publication.

Once we define all the variables, we don't need the following paragraph.

~~The RPR-FA requires few counters/registers which control the traffic forwarded and sourced on the RPR ring. Those counters are add\_rate (tracks the amount of Class C and EIR Class B traffic sourced on the ring), add\_rate\_congestion (tracks the amount of Class C and EIR Class B traffic sourced on the ring and destined beyond the congestion point) and forward\_rate (amount of Class C and EIR Class B traffic transited through the station from the STQ), allow\_rate\_congestion (the current maximum Class C+EIR Class B transmit rate~~

~~for that node beyond the congestion point) and max\_rate (the current maximum Class C+ EIR Class B transmit rate for that node).~~

**Editors' Notes:** To be removed prior to final publication.

*This paragraph belongs to the section on access control.*

~~The traffic policer shall not allow add\_rate\_congestion and add\_rate to exceed allow\_rate\_congestion and max\_rate, respectively. It is accomplished by generating "not SendC" signal to the client when add\_rate\_congestion and add\_rate exceed allow\_rate\_congestion and max\_rate, respectively. (See 6.X for details on the SendC signal.)~~

**Editors' Notes:** To be removed prior to final publication.

*The next few sections should probably be merged with the "MAC fairness operation section".*

~~When a node receives a Type A fairness message that contains a NULL value, it shall increase its allow\_rate\_congestion every decay interval. The maximum value for allow\_rate\_congestion is max\_rate. The max\_rate is a per node parameter that limits the maximum amount of fairness eligible traffic that a node can send.~~

~~When a node sees congestion it starts to advertise a normalized add\_rate value to upstream nodes. The normalized fair rate (nlp\_add\_rate) is obtained by passing add\_rate through a low pass filter and then dividing by its weight. In this way, the fair rates passed on the links are always normalized to a weight of 1.0. Congestion detection is described in 9.5~~

~~When a node determines that it is a congestion point, it will send a fairness message containing a non-NULL fair rate. A node that receives a fairness message with a non-NULL fair rate (revd\_advertised\_rate) will set its allow\_rate\_congestion and TTL\_to\_congestion to the revd\_advertised\_rate value multiplied by its weight and 256-received TTL value, respectively. This allows a node with a weight of N to utilize N times as much bandwidth as a node with a weight of 1.0. If the source of the revd\_advertised\_rate is the same node that received it then the revd\_advertised\_rate should be disregarded. When comparing the revd\_advertised\_rate source address the ring identifier of the fairness packet must match the receiver's ring identifier in order to qualify as a valid compare. The exception is if the receive node is in the wrap state in which case the fairness packet's ring identifier is ignored.~~

**Editors' Notes:** To be removed prior to final publication.

*Per a comment on D0.2, the reference to convergence time of 100 msec for rings of several hundred miles was removed. The same comment also requested a formula for the convergence time, for which contributions are solicited.*

## 9.12 Fairness control messages

### 9.12.1 Generation of fairness control messages

~~Type A-Single Choke~~ and ~~Type B-Multi Choke~~ fairness control messages shall be sent periodically to propagate fair rate information to upstream stations. ~~Type A-Single Choke~~ messages are sent every ~~advertisement interval~~advertisement interval. The recommended ~~advertisement interval~~advertisement interval is between the transmission delay for a single MTU and one-half of the ~~decay interval~~aging interval. ~~Type B-Multi Choke~~ messages are sent every 10 ~~advertisement intervals~~advertisement intervals.

### 9.12.2 Receipt of fairness messages

**Editors' Notes:** To be removed prior to final publication.

Since this is covered by the MAC fairness operation clause, we can probably remove this (or move the appropriate text here).

### 9.12.3 Impact of lost fairness control messages

The ~~allow\_rate\_congestion~~ allowed\_rate\_congested maintains its current value until the next ~~Type-A-Single Choke~~ fairness control message is received. Therefore, the loss of a ~~Type-A-Single Choke~~ fairness control message would result in one of the following. If the lost message contained a ~~NULL value~~ new value of FULL RATE, it would prevent the ~~node station~~ from ramping up its allowed rate for the ringlet. If the value contained in the lost message was a new non-NULL FULL RATE value, the MAC will not be able to react by setting its allowed rate to the received advertised rate, which could have been higher or lower than the current allowed rate. If the lost message did not contain a new value from the previous advertisement, there will be no effect on the fairness algorithm.

Refer to Clause 10 for protection implications of not receiving consecutive Single Choke fairness control messages.

The ~~Type-B-Multi Choke~~ messages are not processed by the MAC, and therefore loss of ~~Type-B-Multi Choke~~ messages doesn't impact behavior of the MAC.

~~Refer to Clause XX for protection implications of not receiving consecutive Type A fairness messages.~~

### 9.12.4 Validation of fairness control messages

**Editors' Notes (JL):** To be removed prior to final publication.

This following paragraph probably belongs in clause 6 and should be generalized for all control messages.

If the source of the rcvd fair rate is the same station that received it then the rcvd fair rate will be disregarded. When comparing the rcvd fair rate source address, the ring identifier of the fairness control message must match the receiver's ring identifier in order to qualify as a valid compare. The exception is if the receive station is in the wrap state, in which case the fairness control message's ring identifier is ignored.

### 9.12.5 Packet format

The fairness control message format is as shown in Figure 9.6. The fairness protocol determines the number of hops to the station that originated the fairness control message by ~~(256—subtracting the~~ message) message from 256. Therefore all fairness control messages shall be sourced with a TTL of 255. And all single choke fairness control messages shall reset the TTL to 255 upon changing the value of the SA to the local SA.



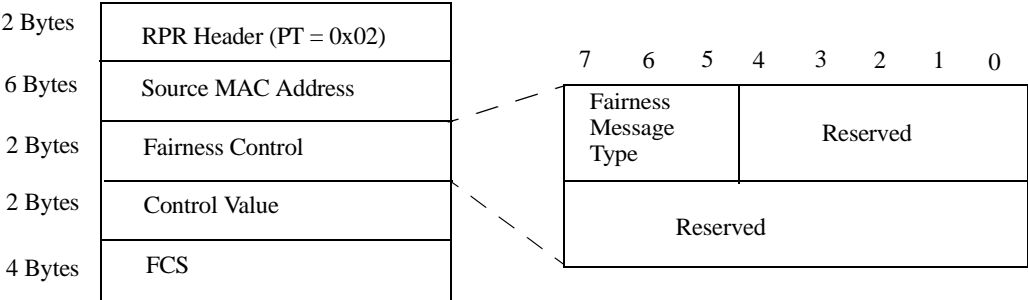


Figure 9.5—Fairness Packet Format

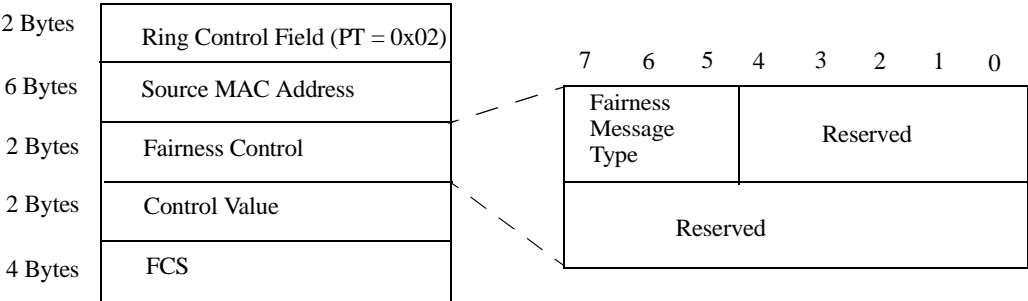


Figure 9.6—fairness control message Format

9.12.5.1 Fairness Control Message Type (3 bits)

This field specifies the type of fairness control message. Table 9.7 shows the values of the Fairness Control Message Type. Type-A-Single Choke fairness control messages are used to implement RPR-FA and provide information about the most congested span on the ringlet. Type-B-Multi Choke fairness control messages are needed to support a multi-choke implementation.

9.12.5.2 Reserved field (~~42~~13 bits)

These bits are ignored on receipt and set to zero on transmit for both Type-A-Single Choke and Type-B-Multi Choke fairness control messages.

9.12.5.3 Control value (16 bits)

This field carries the fair rate encoded as a 16-bit quantity. A fair rate of all ones indicates a value of ~~NULL~~ FULL RATE which corresponds to advertising a fair rate equal to the line rate. The fair rate, when not ~~NULL~~ FULL RATE, indicates the number of bytes that ~~the a station was able to may~~ add to the ringlet during the ~~last decay interval divided next aging interval~~. The reported fair rate is a rate that has been normalized by the normalization coefficient (NORMCOEF), to normalize the number of aging intervals incorporated into the rate (AGECOEF), the local station weight. To ensure that this value does not overflow

**Table 9.7—Fairness [Control](#) Message Type values**

Value (binary)	Type of fairness <a href="#">control</a> message	How is it used
000	<del>Type-A</del> <a href="#">Single Choke</a>	Generated by the FCU every <del>advertisement-interval</del> <a href="#">advertisement interval</a> . SA is the MAC address of the most congested <del>node-station</del> <a href="#">node</a> in the fairness domain
001	<del>Type-B</del> <a href="#">Multi Choke</a>	Generated by the FCU every 10 <del>advertisement-intervals</del> <a href="#">advertisement intervals</a> . The SA is the SA of the MAC, and the messages are broadcast. The received fair rate is passed to the client.
010 to 111	Reserved	For future use.

~~16-bits within a decay-interval~~ [weight \(WEIGHT\)](#), ~~a normalization factor is used depending on~~ [and](#) the ~~link speed as shown in Table 9.2~~ [rate coefficient \(RATECOEF\)](#).

**Table 9.8—Normalizaiton Factors**

Link Speed	Normalization factor
Up to and including OC-48	1
OC-192	4
OC-768	16

### [9.13 Fairness algorithm formulas](#)

**Editors' Notes (JL):** *To be removed prior to final publication.*

A large amount of the following section probably belongs in clause I. Some of the following will probably remain as equations inlined in the text to augment descriptions that need a more precise definition.

[This state machines in this clause use the following formulas:](#)

#### [9.13.1 Initialization](#)

[At the beginning of the state machine, initialize:](#)

```

— allowed_rate
— full_threshold
— high_threshold
— low_threshold
— unreserved_rate

allowed_rate = MAX_ALLOWED_RATE;

unreserved_rate = LINK_RATE - reserved_rate;

if (dual_queue_MAC)
  full_threshold = STQ_SIZE - MTU_SIZE; // can subtract more than MTU_SIZE

if (dual_queue_MAC && aggressive)
  high_threshold = STQ_SIZE / 4;
else
  high_threshold = .95 * unreserved_rate;

if (dual_queue_MAC && aggressive)
  low_threshold = STQ_SIZE / 8;
else
  low_threshold = .8 * unreserved_rate;

```

### 9.13.2 Per byte

**Editors' Notes (JL):** *To be removed prior to final publication.*

This section (per byte operations) should be moved to Clause 6.

For each byte transmitted, update:

```

— add_rate
— add_rate_congested
— fw_rate
— fw_rate_congested
— nr_xmit_rate

if (FE byte added by local station)
  add_rate = add_rate + 1;
if (FE byte added by local station && sent beyond congestion_point)
  add_rate_congested = add_rate_congested + 1;
if (FE byte forwarded by local station)
  fw_rate = fw_rate + 1;
if (FE byte forwarded by local station && sent beyond congestion_point)
  fw_rate_congested = fw_rate_congested + 1;
if (non-A0 byte transmitted by local station)
  nr_xmit_rate = nr_xmit_rate + 1;

```

### 9.13.3 Per aging interval

For each aging interval, based upon the last received FCM, update:

- active stations
- add rate
- add rate ok
- add rate congested
- add rate congested ok
- fw rate
- fw rate congested
- nr xmit rate
- congested
- allowed rate
- allowed rate congested
- congestion point
- TTL to congestion
- local fair rate
- norm local fair rate
- lp add rate
- lp add rate congested
- lp fw rate
- lp fw rate congested
- lp nr xmit rate
- norm lp fw rate congested

#### 9.13.3.1 active stations

```

for (active_stations = 0; station = 0; station < num_stations; station++)
    if (sent_FE_during_aging_interval(station))
        active_stations = active_stations + 1;

```

#### 9.13.3.2 lp add rate, lp add rate congested, lp fw rate, lp fw rate congested, lp nr xmit rate, norm lp fw rate congested

The following low pass filterings must be done before the first pass agings specified in 9.13.3.3 .

```

lp_add_rate = ((LPCOE-1) * lp_add_rate + add_rate) / LPCOE;
lp_add_rate_congested = ((LPCOE-1) * lp_add_rate_congested +
    add_rate_congested) / LPCOE;
lp_fw_rate = ((LPCOE-1) * lp_fw_rate + fw_rate) / LPCOE;
lp_fw_rate_congested = ((LPCOE-1) * lp_fw_rate_congested +
    fw_rate_congested) / LPCOE;
lp_nr_xmit_rate = ((LPCOE-1) * lp_nr_xmit_rate + nr_xmit_rate) / LPCOE;
norm_lp_fw_rate_congested = lp_fw_rate_congested / NORMCOEF;

// alternatively, each can be simplified as shown for lp_add_rate
// lp_add_rate = lp_add_rate -
//     lp_add_rate >> (log2(LPCOE)) + add_rate >> (log2(LPCOE));

```

#### 9.13.3.3 add rate, add rate congested, fw rate, fw rate congested, nr xmit rate

The following first pass agings must be done after the low pass filterings specified in 9.13.3.2 .

```

add_rate = (add_rate * (AGECOEF - 1)) / AGECOEF;
add_rate_congested = (add_rate_congested * (AGECOEF - 1)) / AGECOEF;
fw_rate = (fw_rate * (AGECOEF - 1)) / AGECOEF;
fw_rate_congested = (fw_rate_congested * (AGECOEF-1)) / AGECOEF;
nr_xmit_rate = (nr_xmit_rate * (AGECOEF - 1)) / AGECOEF;

```

#### 9.13.3.4 congested

**Editors' Notes (JL):** To be removed prior to final publication.

This does not currently take into account any hysteresis. This needs to be added. The hysteresis should allow delta to be set to a range of values including 0 (no hysteresis).

```

if (dual_queue_MAC)
{
    if (lp_nr_xmit_rate > unreserved_rate)
    || (STQ_depth > low_threshold))
        congested = TRUE;
    else if (aggressive) // clearing of congested for conservative mode
        congested = FALSE; // is handled in calculation for local_fair_rate
}
else // mono_queue_MAC
{
    if (lp_nr_xmit_rate > unreserved_rate)
    || (lp_nr_xmit_rate > low_threshold)
    || (access_delay_timer for add class B expired)
    || (access_delay_timer for add class C expired))
        congested = TRUE;
    else if (aggressive) // clearing of congested for conservative mode
        congested = FALSE; // is handled in calculation for local_fair_rate
}

```

#### 9.13.3.5 allowed rate

```

if (aggressive)
    allowed_rate = MAX_ALLOWED_RATE;
else // conservative
{
    if (congested) // local station is congested
        allowed_rate = local_fair_rate;
    else // local station is not congested
        allowed_rate = allowed_rate+(MAX_ALLOWED_RATE-allowed_rate)/RAMPCOEFF;
}

```

#### 9.13.3.6 allowed rate congested, congestion point, TTL to congestion

```

if (rcvd_fair_rate != FULL_RATE)
{
    allowed_rate_congested = rcvd_fair_rate * NORMCOEF;
    congestion_point = rcvd_SA;
    TTL_to_congestion = 256 - rcvd_TTL;
}
else // no congestion downstream, ramp up
{
    allowed_rate_congested = allowed_rate_congested +
        (MAX_ALLOWED_RATE - allowed_rate_congested) / RAMPCOEFF;
    congestion_point = local_SA;
    TTL_to_congestion = TTL_to_congestion;
}

```

```

// TTL_to_congestion should be left at the last congestion point so that
// the allowed_rate is not all of a sudden presented to the former
// congestion point. One could add the following statement:
// if (allowed_rate_congested == allowed_rate)
//     TTL_to_congestion = 255;
// But it is unnecessary since it has no effect upon the rate at which
// the station may transmit.
}

```

### 9.13.3.7 local fair rate, norm local fair rate

```

if (aggressive)
{
    if (congested)
        local_fair_rate = lp_add_rate;
    else // not congested
        local_fair_rate = unreserved_rate;
}
else // conservative
{
    if (congested)
    {
        if (just_entered_congested) // first time when congested
            local_fair_rate = (unreserved_rate / active_stations) * WEIGHT;
            // optionally, instead of active_stations, which uses a weight of 1
            // for each of the other stations, one could use sum(ActiveWeights)
        else // was congested last time
        {
            if ((add_rate + fw_rate < low_threshold)
                && (RTT_worth_of_aging_intervals_passed_since_last_update))
            {
                local_fair_rate = min(unreserved_rate,
                    local_fair_rate +
                    (unreserved_rate - local_fair_rate) / RAMPCOEFF;
                if (local_fair_rate >= unreserved_rate)
                    congested = FALSE;
            }
            else if ((add_rate + fw_rate > high_threshold)
                && (RTT_worth_of_aging_intervals_passed_since_last_update))
                local_fair_rate = local_fair_rate - local_fair_rate / RAMPCOEFF;
            else
                local_fair_rate = local_fair_rate; // no change
        }
    }
    else // not congested
        local_fair_rate = unreserved_rate;
}

norm_local_fair_rate = local_fair_rate / NORMCOEFF;

```

### 9.13.3.8 add rate ok, add rate congested ok

```

add_rate_ok =
    (add_rate < allowed_rate) \\ current allowance is not exceeded
    && (lp_nr_xmit_rate < unreserved_rate) \\ space for reserved traffic
    && ((STQ_depth = 0) \\ upstream stations are not in need
    || ((fw_rate > add_rate) \\ upstream stations are not starved
        && (STQ_depth < high_threshold))); \\ node is not fully congested

```

NOTE—An optional optimization is to change (fw\_rate > add\_rate) to (fw\_rate > add\_rate/WEIGHT). This is not needed under almost all circumstances. The one exception is if the sum of weights of upstream stations is less than the weight of the local station. If this optimization is chosen, the implementation may require that WEIGHT values be factors of 2.

```
add_rate_congested_ok = add_rate_ok  
&& (add_rate_congestion < allowed_rate_congestion);
```

#### **9.13.3.9 sendC**

**Editors' Notes (JL):** *To be removed prior to final publication.*

This section (sendC) should be moved to Clause 6.

```
if (!add_rate_ok)  
    sendC = 0;  
else if (!add_rate_congested_ok)  
    sendC = TTL_to_congestion;  
else // see comment in 9.13.3.6 on TTL_to_congestion  
    sendC = TTL_to_congestion; // calculation with no downstream congestion
```

#### **9.13.4 Per advertisement interval**

For each advertisement interval, update:

— advertised fair rate

```
if ((rcvd_fair_rate < norm_local_fair_rate) // downstream is more congested,  
&& (rcvd_fair_rate < norm_lp_fw_rate_congested)) // upstream is contributing  
    advertised_fair_rate = rcvd_fair_rate; // to downstream congestion  
else if (congested) // local station is congested (and more than downstream)  
    advertised_fair_rate = norm_local_fair_rate;  
else // no downstream or local congestion  
    advertised_fair_rate = FULL_RATE;
```