

# **Simulation Results** on **SRP Fairness Algorithm (SRP-fa)**

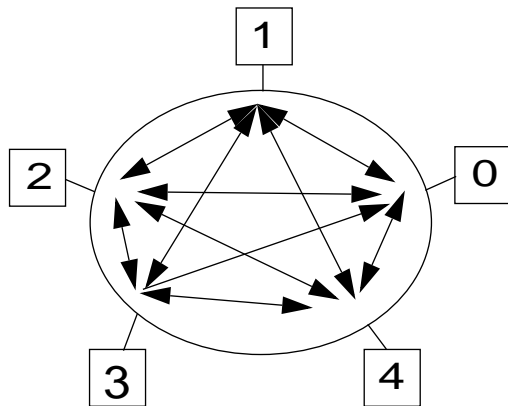
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March 8, 2000

# Overview

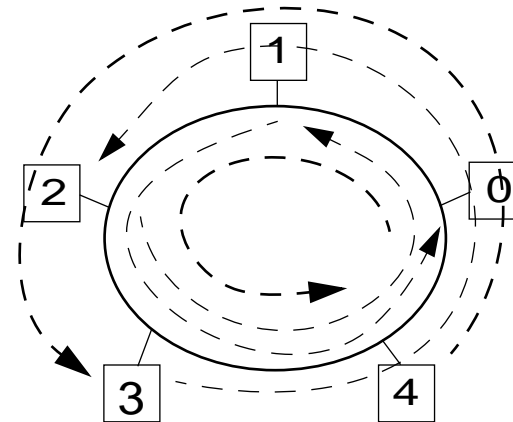
- Provide simulation results to demonstrate the behaviour of SRP fairness mechanism under various traffic pattern and load scenarios.
- Cases considered:
  - ❖ Homogeneous traffic
    - Mesh, Worst-case, Aggregation, Best-case
  - ❖ Heterogeneous traffic
    - Aggregation
  - ❖ Large number of nodes

# Traffic Pattern Scenarios



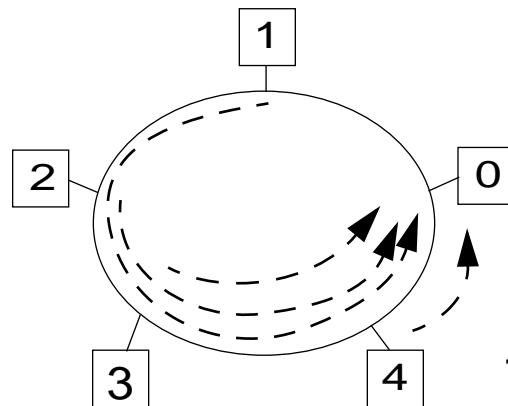
Mesh

200%



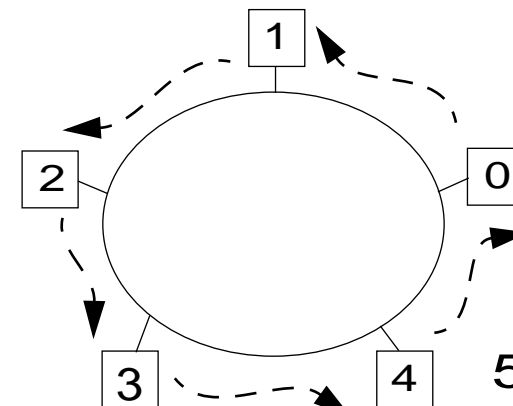
Worst-Case

120%



Aggregation

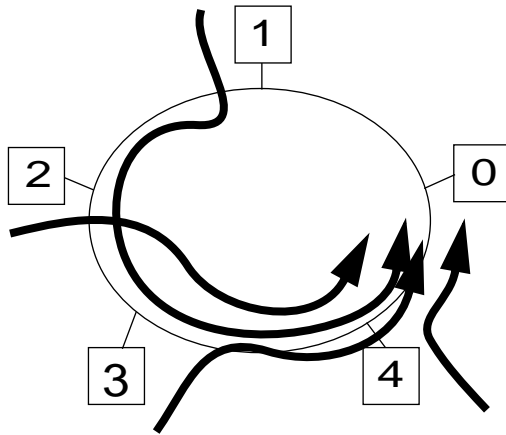
100%



Best-Case

500%

## Traffic Pattern Scenarios *(cont'd)*



Aggregation (heterogeneous input)

### ❖ case 1

1 -> 0:	5 %
2 -> 0:	25 %
3 -> 0:	50 %
4 -> 0:	75 %

(Total: 155%)

### ❖ case 2

1 -> 0:	5 %
2 -> 0:	15 %
3 -> 0:	30 %
4 -> 0:	45 %

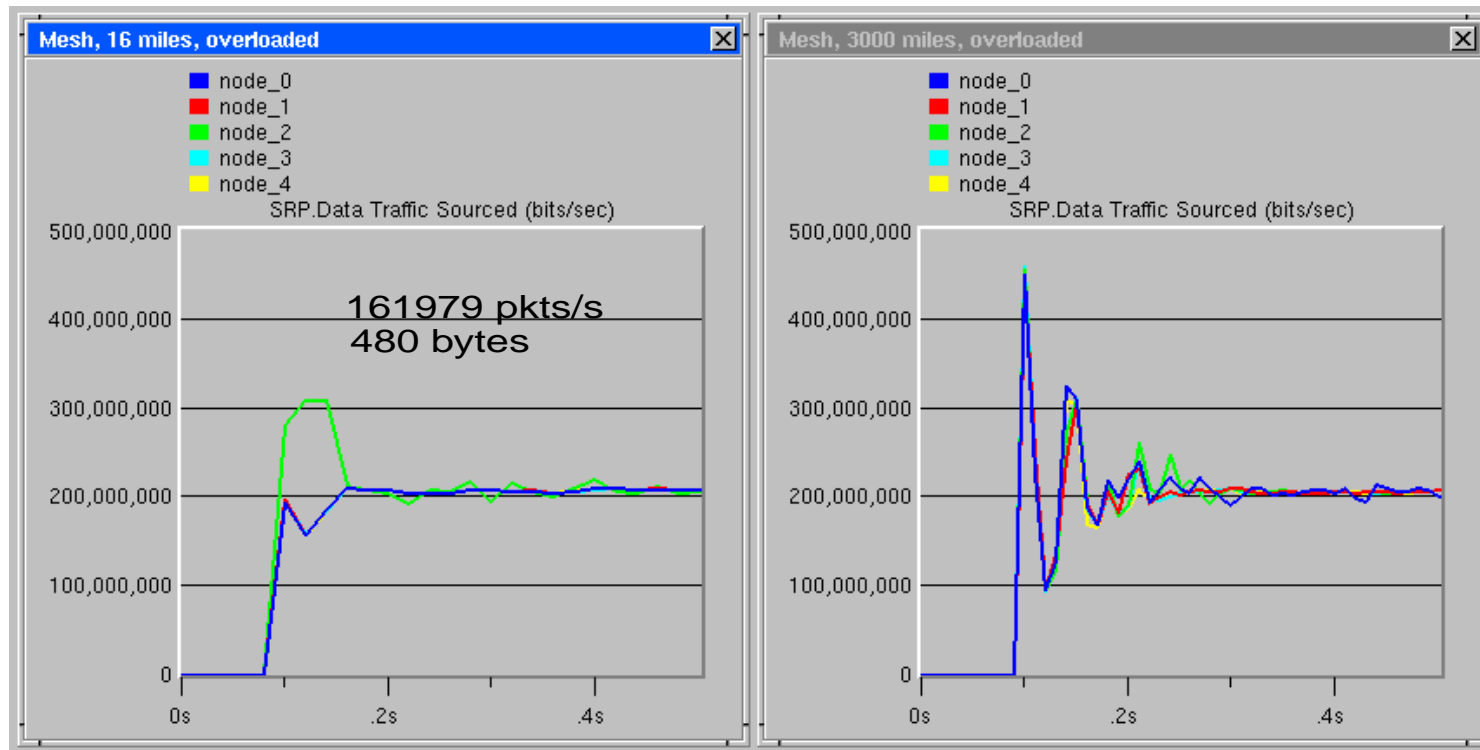
(Total: 95%)

# Simulation Parameters

\* 5 nodes with outer ring traffic only and equal-distanced

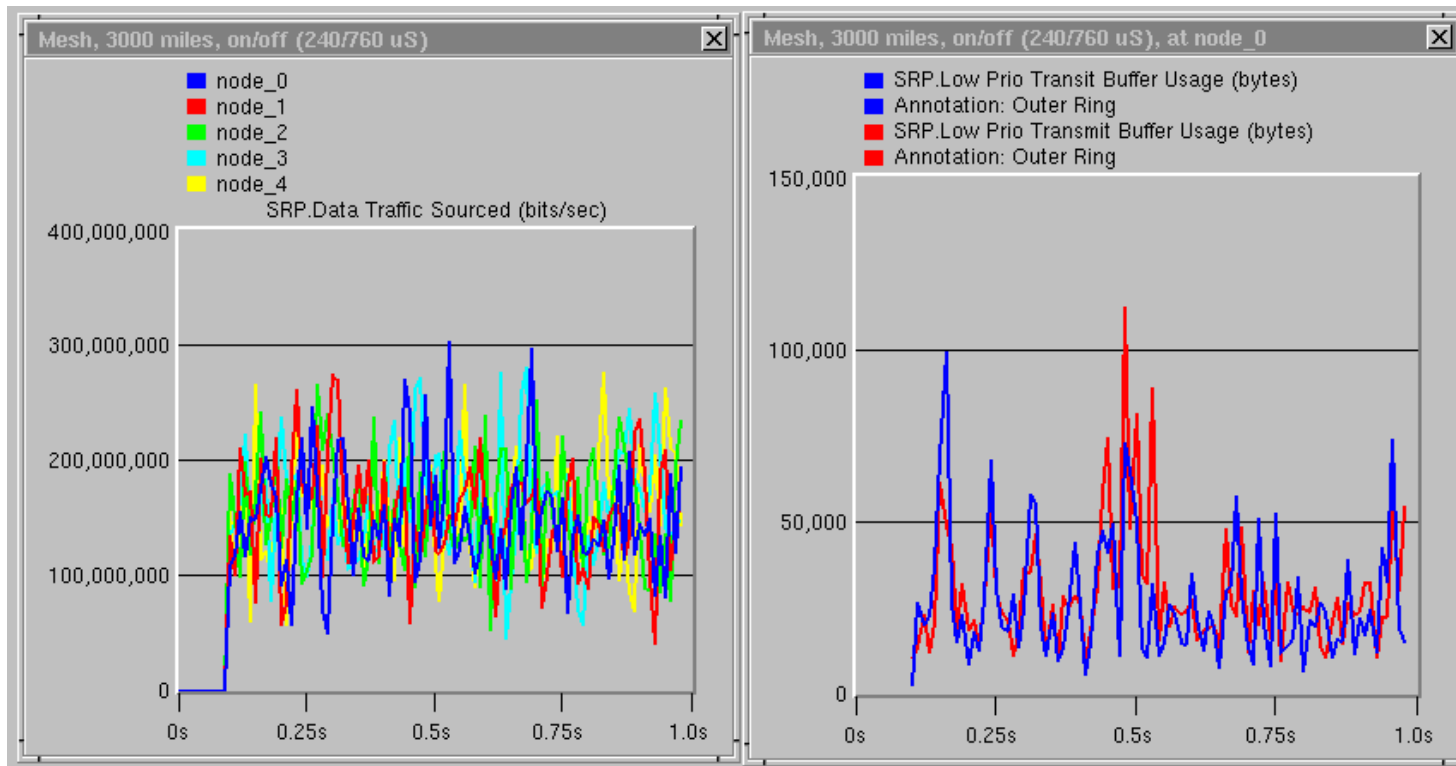
Ring circumference	:	16, 3000	(in miles)
Hi/Low priority ratio	:	varies	
Avg pkt rate	:	161979 pkts/sec	(exponential)
Avg pkt size	:	480 Bytes	(exponential)
Burst ON duration	:	varies	(exponential)
Burst OFF duration	:	varies	(exponential)
Tx high priority size	:	2.7 KB	
Tx low priority size	:	1 MB	
TB high priority size	:	2.7 KB	
TB low priority size	:	512 KB	(OC-12)
TB low threshold	:	256 KB	
TB high threshold	:	500 KB	(OC-12)
LP_FWD	:	64	(low pass filetering weight)
LP_MU	:	512	
LP_ALLOW	:	64	
AGECOEFF	:	4	(Aging Coefficient Value)
Decay Interval	:	32000	(in byte times, 100μS)
Usage Interval	:	Decay Interval + $\Delta$	

## Results: Mesh, Homogeneous, Overloaded, oc-12



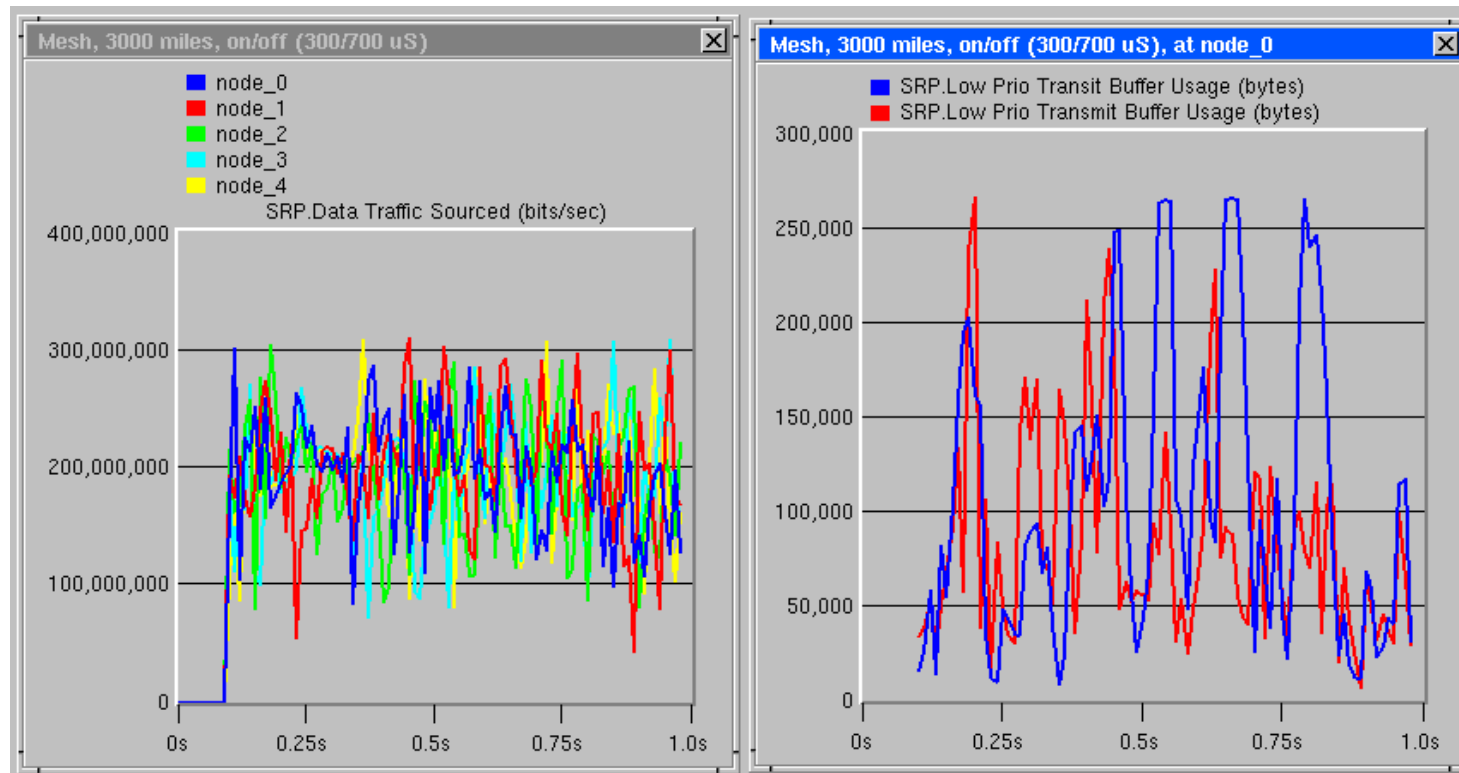
- Throughput achieves 160% of link capacity (spatial reuse)
- With small ring size, faster convergence (16 vs 3000 miles)

## Results: Mesh, on/off (240/760 uS), 3000 miles, OC-12



- $24\% \times 5 = 120\%$  of link capacity (spatial reuse)
- TB and Tx occupancy indicate underloaded condition

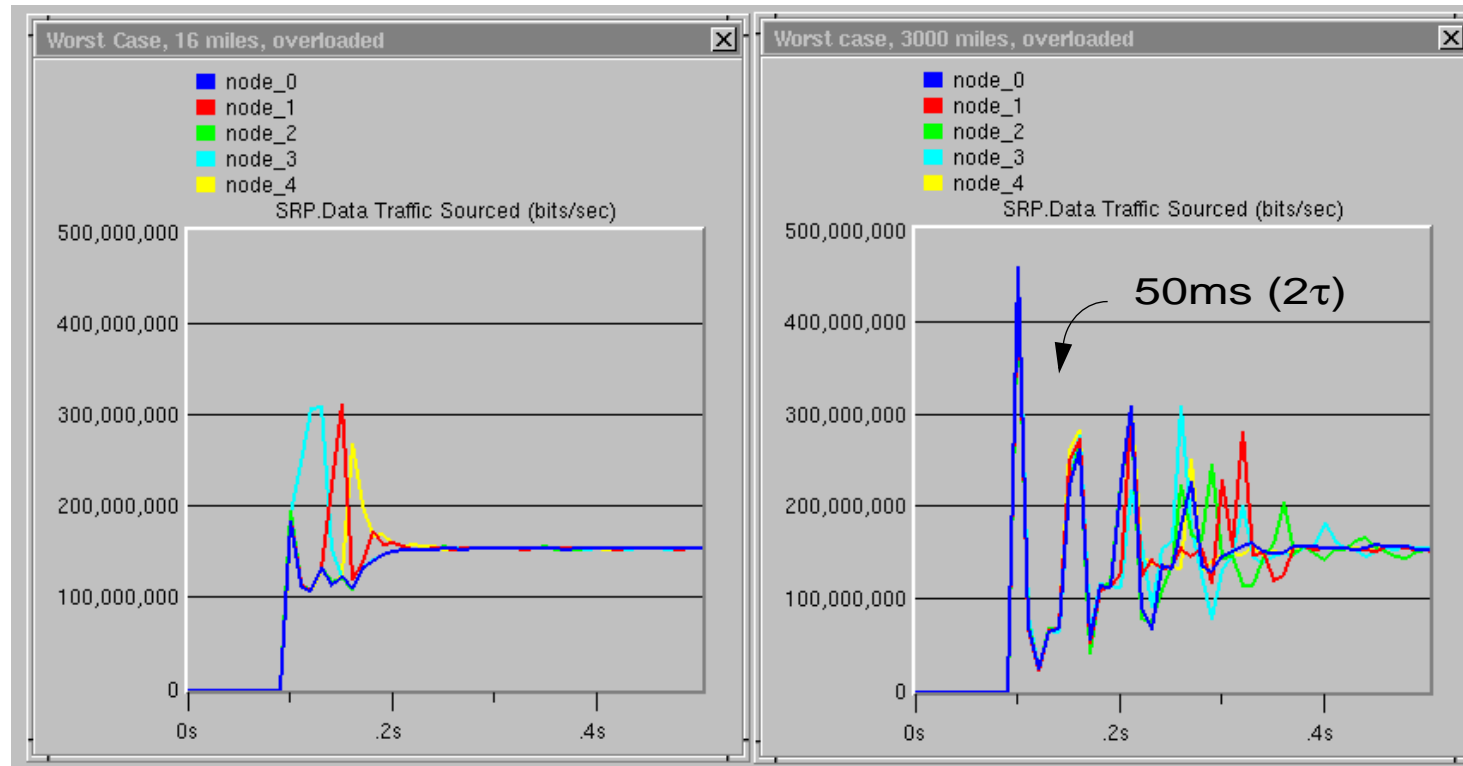
## Results: Mesh, on/off (300/700 uS), 3000 miles, OC-12



- $30\% \times 5 = 150\%$  of link capacity (spatial reuse)
- TB and Tx occupancy indicate heavier utilization

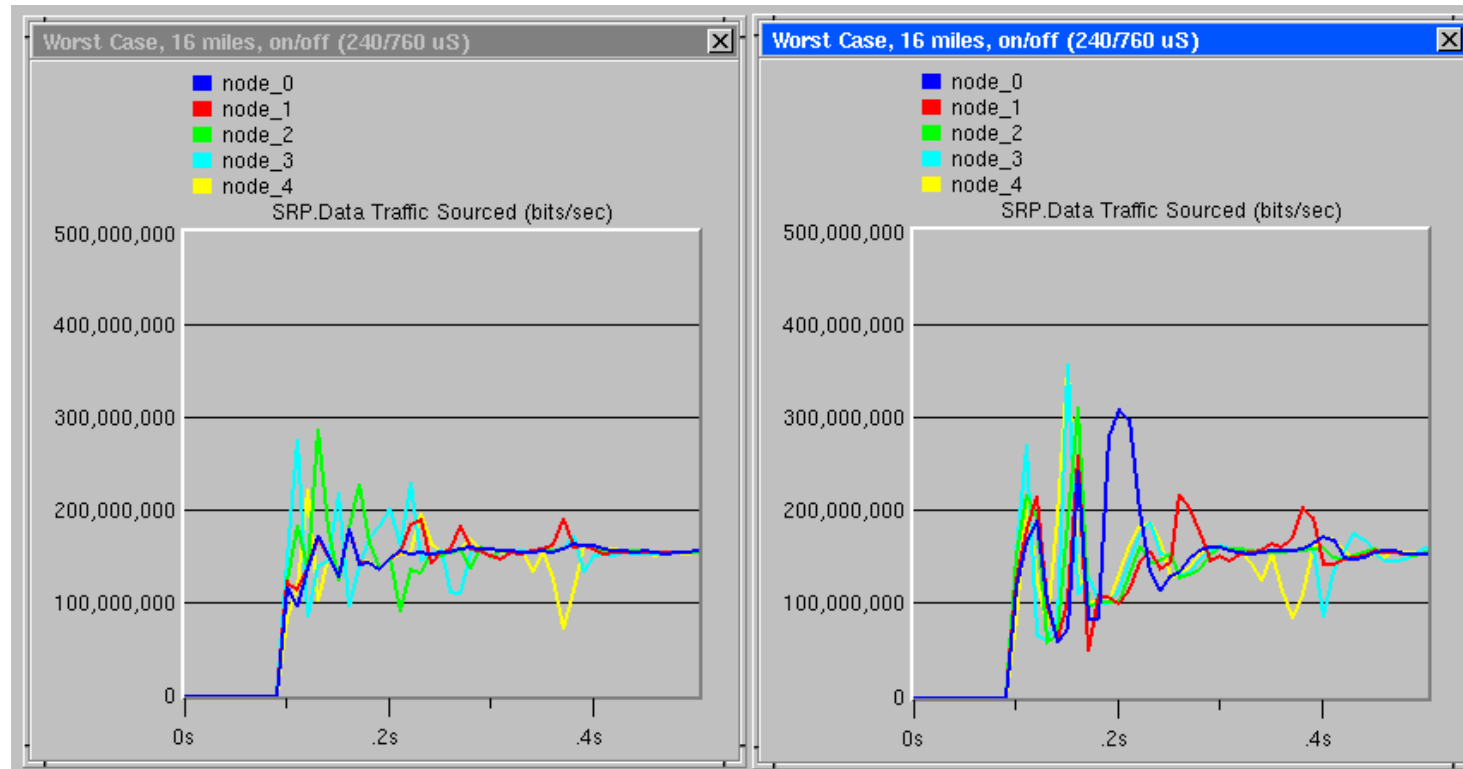


## Results: Worst Case, Overloaded, oc-12



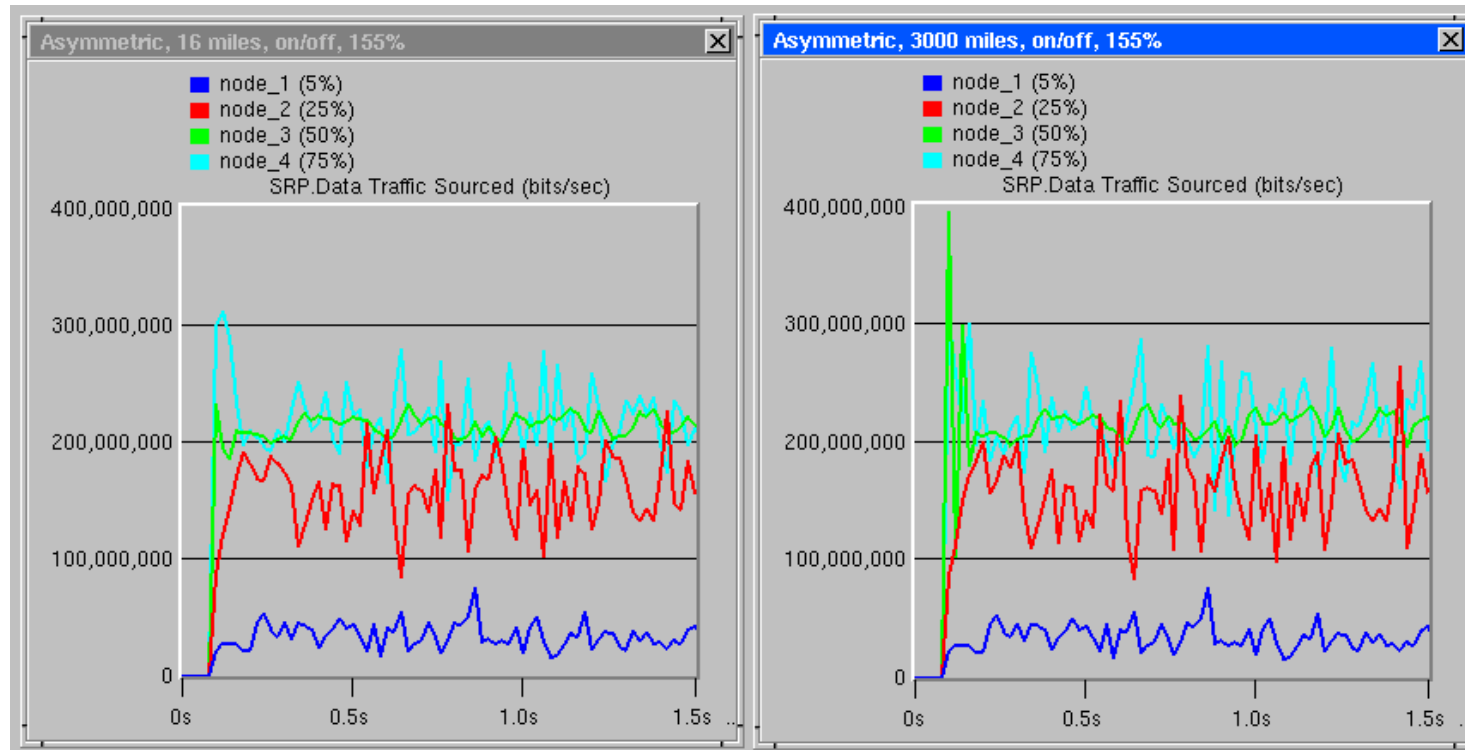
- Each node can source  $622/4=155\text{Mbps}$  (25% of link cap.)
- Convergence is worse when compared to mesh pattern

## Results: Worst Case, on/off (240/760 uS), oc-12



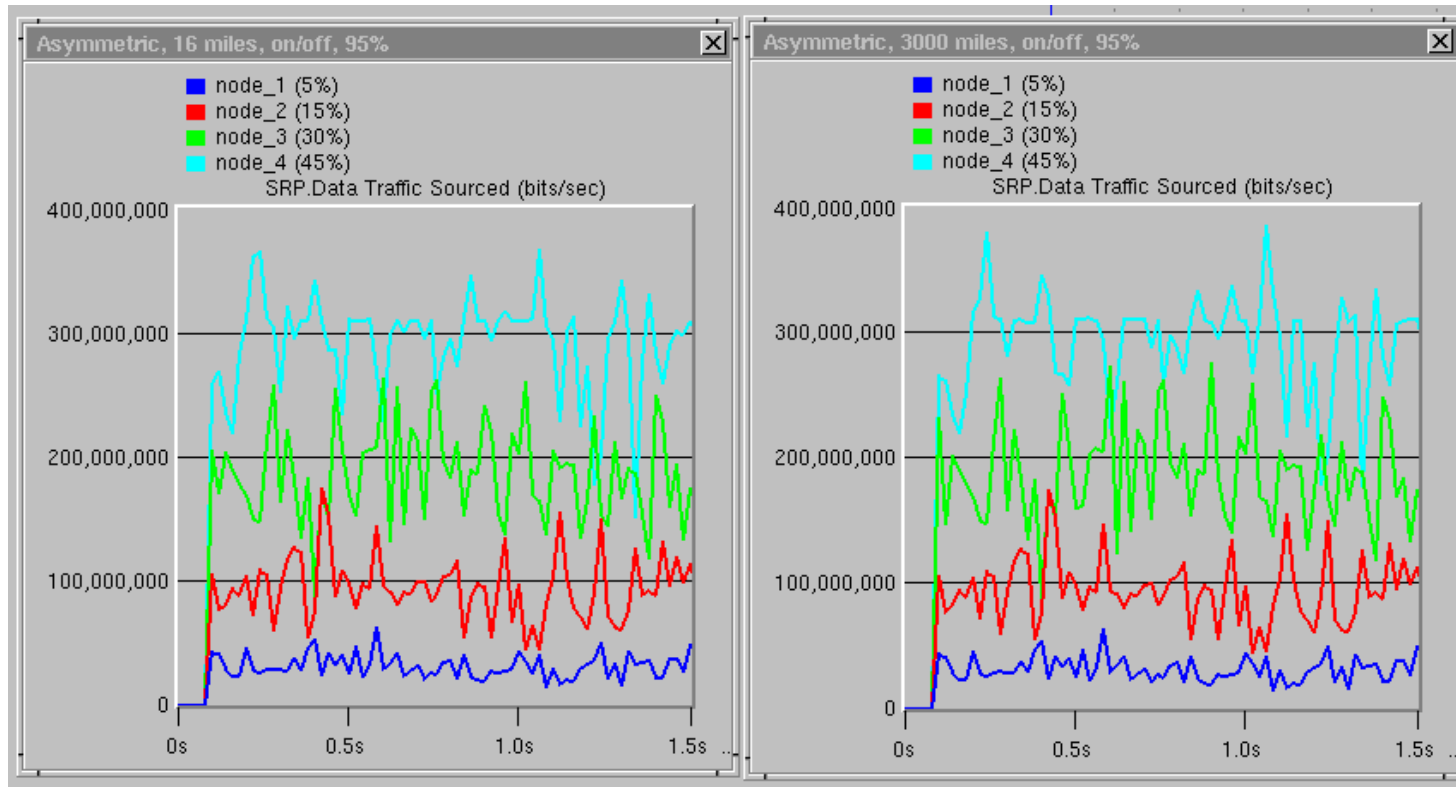
- Each node transmit 24% of link cap. on the average
- Heavy traffic condition with bursty input

## Results: Asymmetric (case 1), 155%, OC-12



- Under overload, 75% and 50% nodes share 30% each, while small input nodes are able to send all their inputs. SRP achieves total of 90% link utilization.

## Results: Asymmetric (case 2), 95%, OC-12



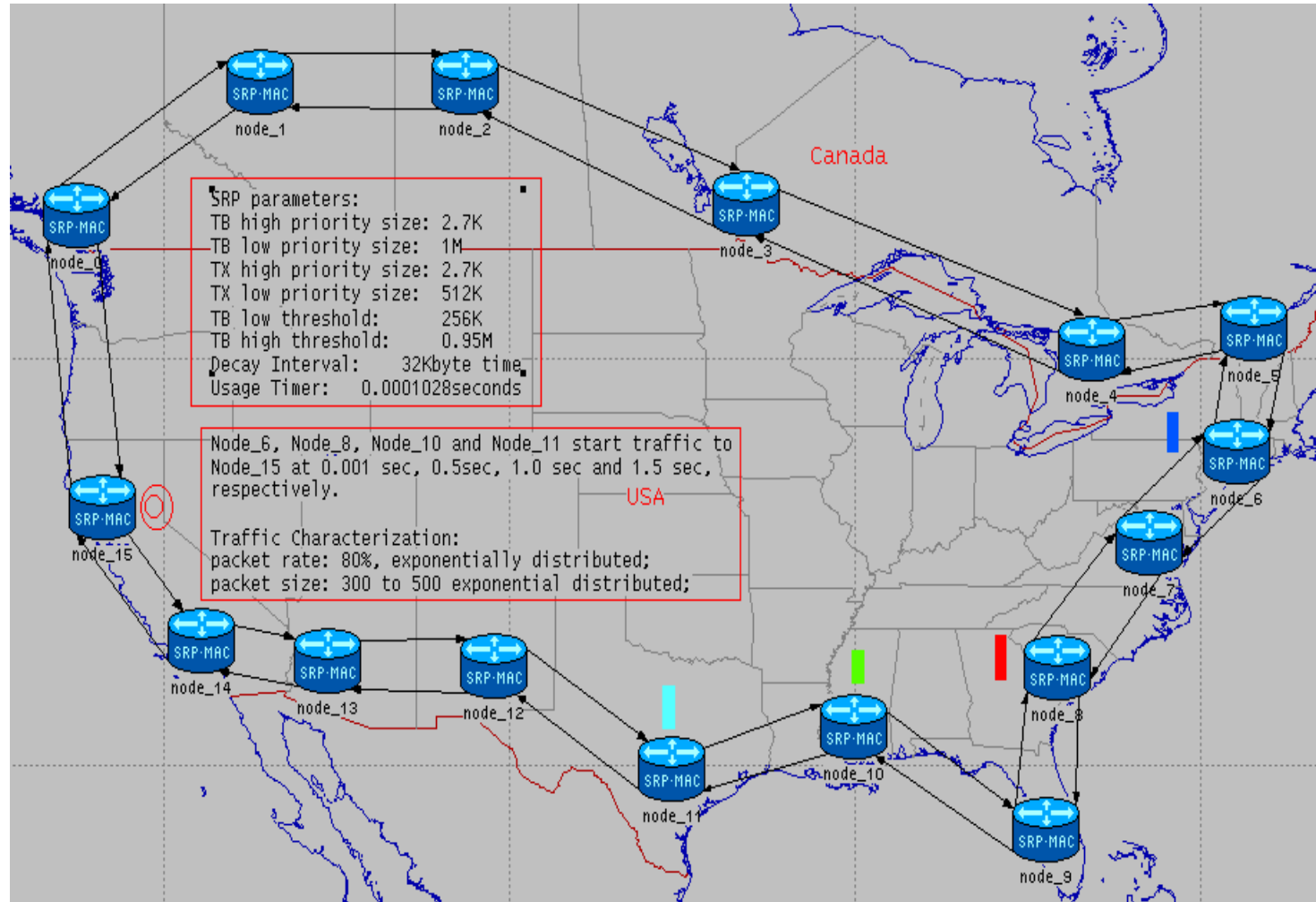
- Under 95% input load, all the nodes are able to send their full input traffic. Similar performance both with 16 and 3000 miles.

## Results: Asymmetric (case 2), 95%, OC-12

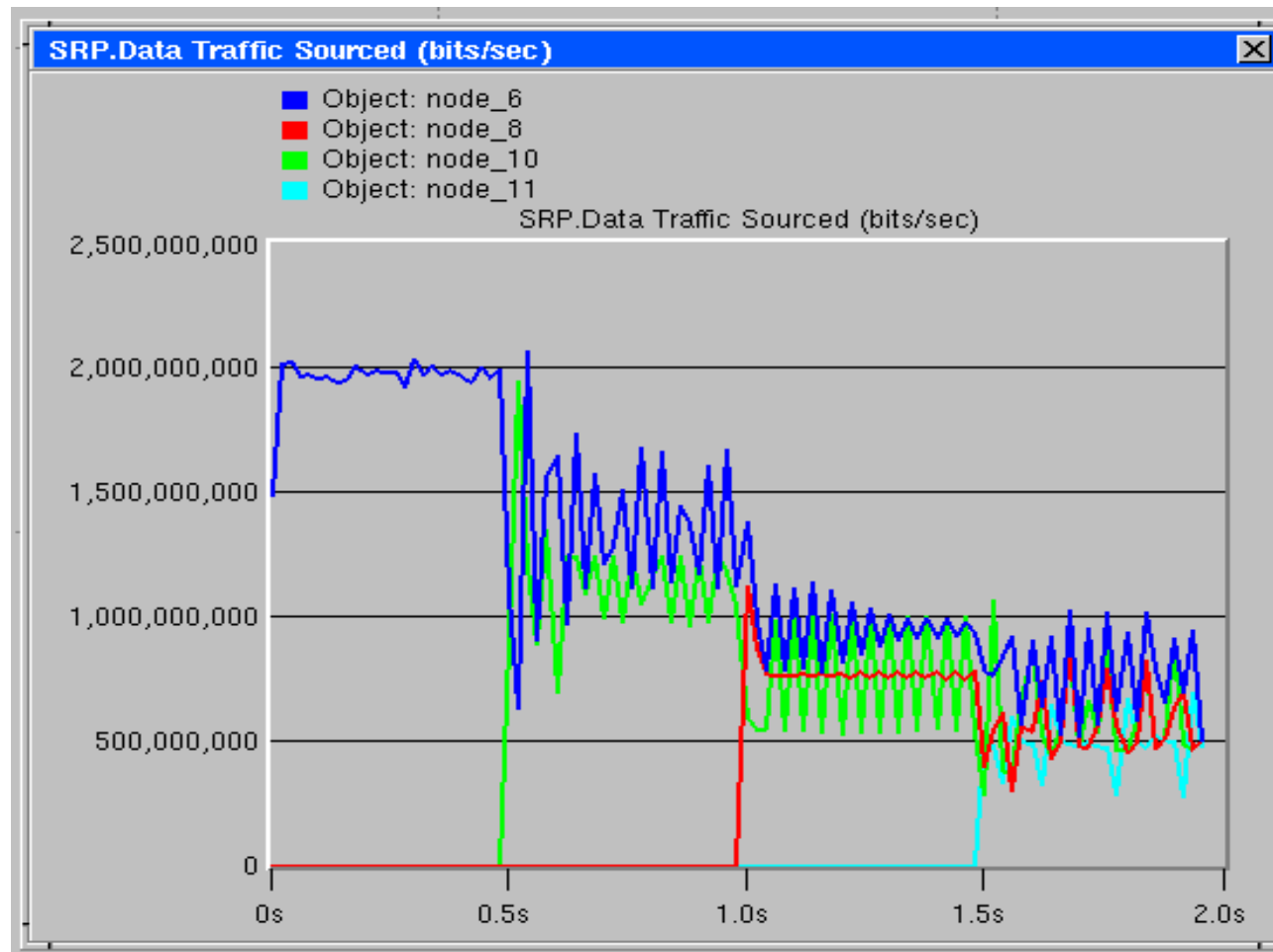


- TB Congestion only occurs at node 4 (aggregation point), and as a result, the allow usage at node 3 (upstream node) is often throttled.

## Results: WAN, Triggered Input, OC-48



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# Summary

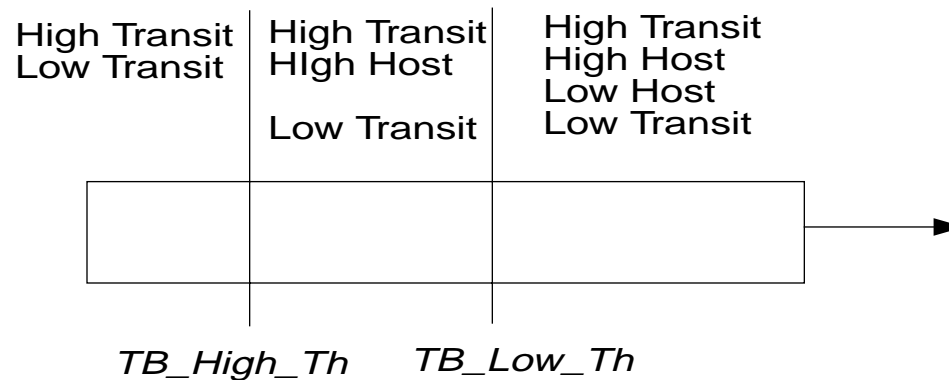
- SRP-fa achieves fair share of ring capacity under various traffic pattern and load scenarios.
- SRP-fa's fair share is based on the maximum throughput achievable with spatial reuse property.
- SRP-fa is effective in both small and large geographical coverage.



## Appendix I: SRP-fa Nutshell (One)

- Layer 2 MAC Transmission Scheduling

❖ Transmission order:



❖ 4 Conditions when Low priority Host traffic is blocked:

- TB occupancy  $\geq$  TB\_Low\_Threshold
- My\_usage > Allow\_usage
- My \_usage > Max\_allow
- TB low priority is non-empty && My\_usage > Fwd\_rate



## Appendix I: SRP-fa Nutshell (Three)

- Updating Allow Usage  
*(once every decay interval)*

Max\_Lrate = AGECOEFF \* Decay\_Interval

if ( Rcvd\_usage != NULL)

    Allow\_usage = Rcvd\_usage;

else

    Allow\_usage + (Max\_Lrate - Allow\_usage)/ LP\_ALLOW

### Reference:

D. Tsiang and G. Suwala, "draft-tsiang-srp-01.txt," IETF Internet Draft, March 1, 2000

## Appendix II: E2E Delay of High Priority Traffic

- **E2E delay** is measured from the time a high priority packet is queued at the source to the time that it is received at the destination. There are three delay componets<sup>1</sup>:

### ❖ Access Delay ( $A_d$ )

$$WC \text{ (Worst Case)} = N \times B_f * HP\_MTU * \text{byte\_delay} + Q_d$$

$$BC \text{ (Best Case)} = HP\_MTU * \text{byte\_delay} + Q_d$$

### ❖ Nodal Delay ( $N_d$ )

$$\begin{aligned} WC &= N * LP\_MTU * \text{byte\_delay} \\ &\quad + N * TB\_High * \text{byte\_delay} \\ &\quad + N * HP\_MTU * \text{byte\_delay} \end{aligned}$$

$$BC = N \times HP\_MTU \times \text{byte\_delay}$$

### ❖ Fiber Delay ( $F_d$ )

Propagation delay through N fiber link segments: 8  $\mu$ S/mile.

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1. N is the number of nodes traversed, while  $B_f$  is the token bucket burst size in # of packets, and  $Q_d$  is the access queueing delay. HP\_MTU and LP\_MTU are high and low priority Max Transmit Unit in bytes.

## Appendix II: E2E Delay of High Priority Traffic *(cont'd)*

### Numerical Examples (SRP delay at OC-12):

$N = 8$  nodes traversed (16 nodes actual)  
 $B_f = 2$  (in HP\_MTU)  
 $Q_d = 50 \mu S$  (estimated)

(1) HP\_MTU = 128, TB\_High = 2700, LP\_MTU=1500 (bytes)

$$\begin{aligned}
 WC &= 77.3 + 460.8 = 538.1 \mu S \\
 BC &= 51.7 + 13.6 = 65.3 \mu S
 \end{aligned}$$

(2) HP\_MTU = 128, TB\_High = 2700, LP\_MTU=9000 (bytes)

$$\begin{aligned}
 WC &= 77.3 + 1261.3 = 1338.6 \mu S \\
 BC &= 51.7 + 13.6 = 65.3 \mu S
 \end{aligned}$$

(2a) Same as (2) but at OC-48

$$\begin{aligned}
 WC &= 56.8 + 315.4 = 372.2 \mu S \\
 BC &= 50.4 + 3.4 = 53.8 \mu S
 \end{aligned}$$

- Simulation results provide similar range of values.