



---

# Multi Choke Point Detection and Virtual Destination Queuing

Necdet UZUN and Mete YILMAZ

9/6/2001



# Agenda

---

- Introduction
- Requirements
- Detailed description
- Example
- Simulation results
- Conclusion



# Introduction

---

- Is it possible to increase link utilization while maintaining fairness?
  - Coupled with cost & complexity increase
- Possible approaches:
  - Static Weighted Fairness
  - Dynamic Weighted Fairness
  - SRP approach for Virtual Destination Queuing (VDQ)



# Requirements

---

- **Virtual destination queues (VDQ).**
- **Multiple node congestion information.**
  - Use of more detailed **choke** (congested) point information provides better **utilization** of network resources.
- **A scheduling policy** that utilizes multi-choke information.



# Usage Packet Handling

---

- Each node generates usage messages to distribute the total usage value of that node
  - when a node is not congested a special message with not congested information will be generated.
  
- A usage packet will be removed from the ring if
  - the node which generated the usage message receives its own usage message back.



## Virtual Destination Queues & Scheduling

---

- This approach may require as many queues as the number of nodes on the ring.
- Upon reception of usage packet node updates allowed usage information for the appropriate choke point.
- Nodes limit the amount of insertion traffic sent through the choke points.



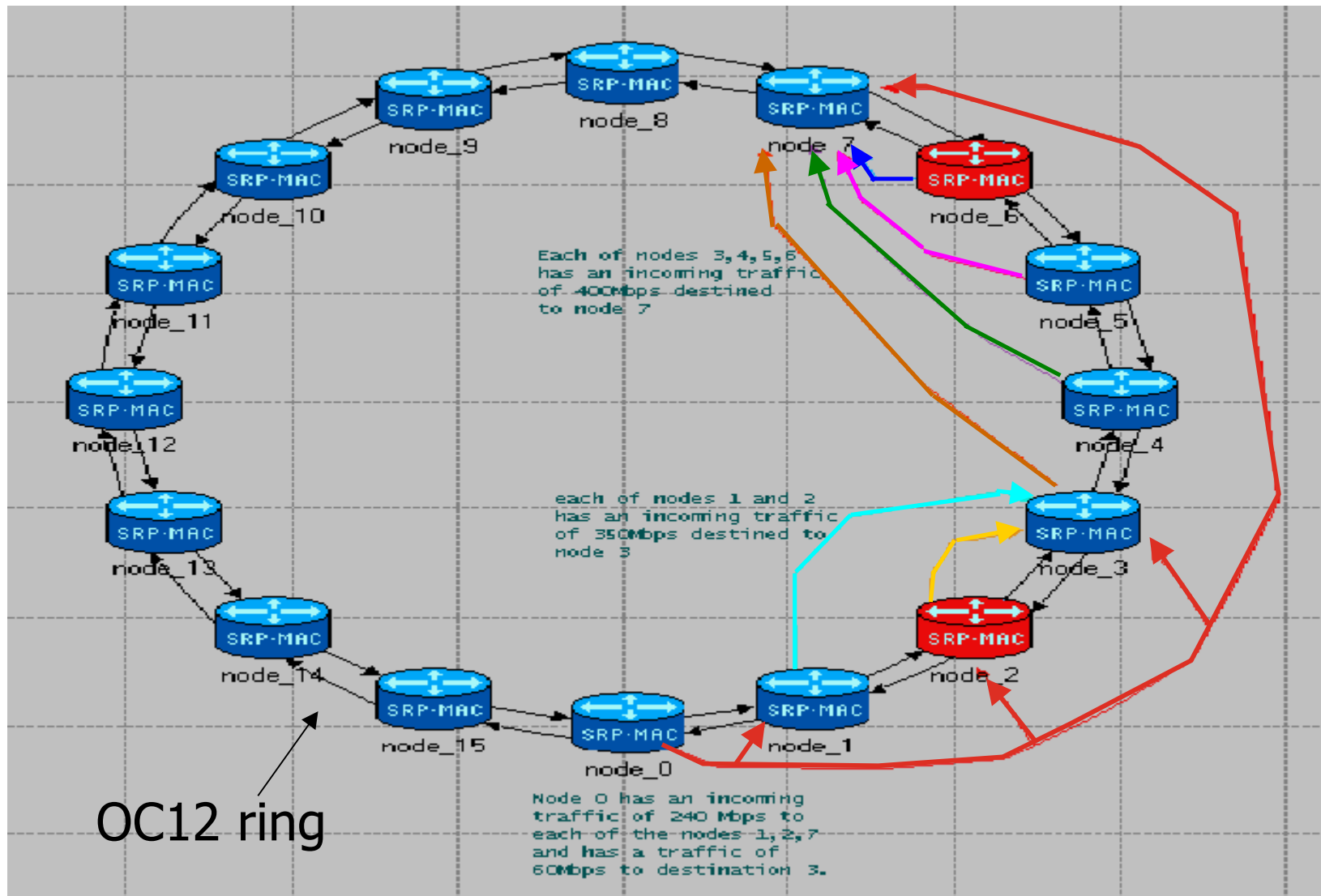
## Details

---

- Choke point information is passed to MAC client and MAC client does the scheduling of VEQ's.
- Usage values are accounted per VEQ basis.
- Usage values and allowed usages are decayed/incremented similar to SRP-fa.
- The number of choke points supported will determine the trade off between implementation complexity and the achievable network utilization.



# A Congested Ring Scenario







## Description

---

- Node 0 receives usage values  $u_2$  and  $u_6$  from node 2 and node 6 respectively.
  
- Node 0 is aware of 3 congestion domains:
  - node 1 and node 2
  - nodes between node 3 and node 6 (inclusive)
  - nodes beyond node 6



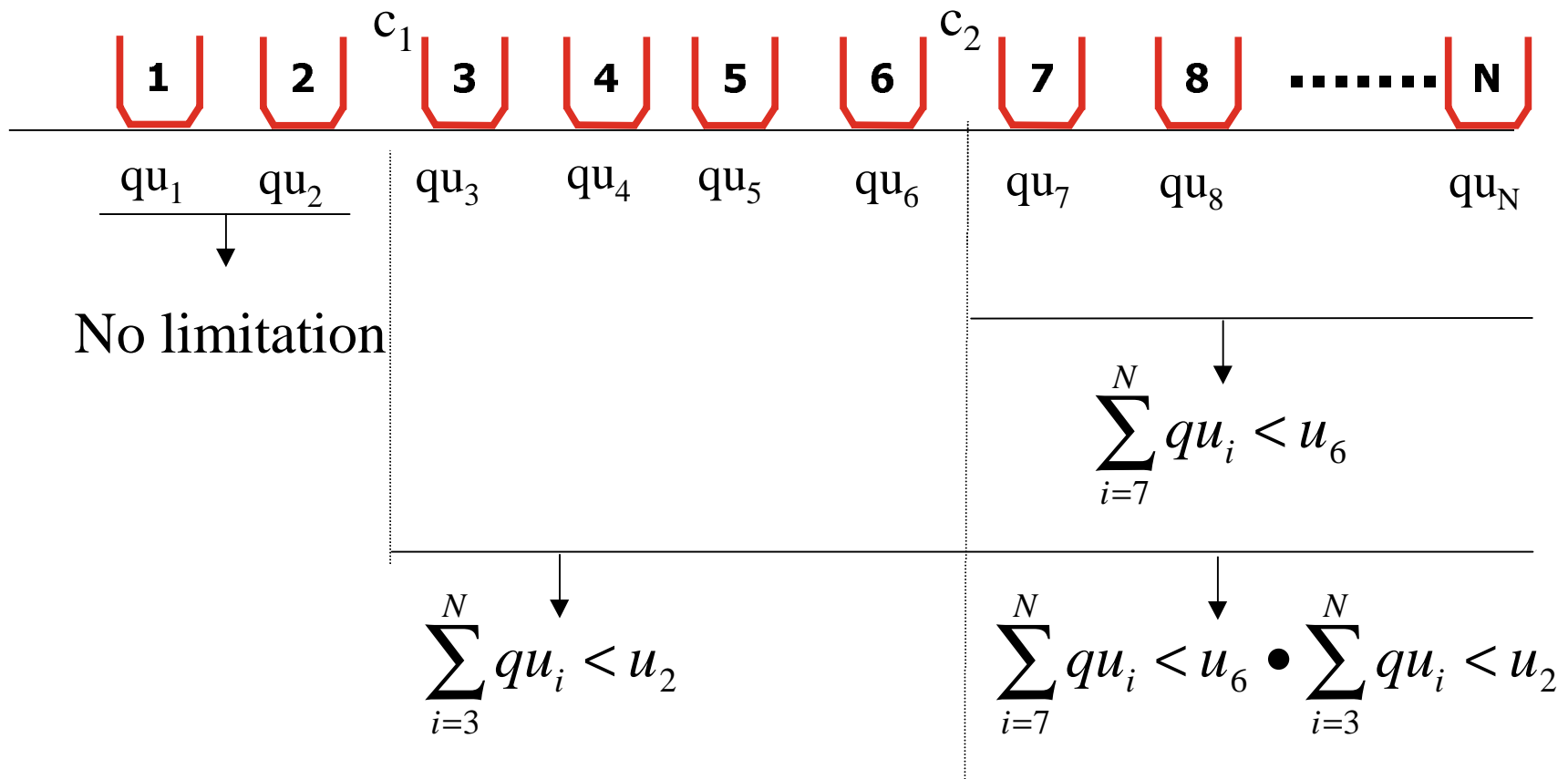
## Description Cont.

---

- Node 0 should obey the following constraints while scheduling its virtual destination queues:
  - Up to line rate for traffic destined to node 1 and node 2.
  - Virtual destination queues for nodes 3,4,5, and 6 can be scheduled as long as the total usage beyond  $VDQ_2$  does not exceed  $u_2$ .
  - Virtual destination queues for nodes beyond 6 can be scheduled as long as the total usage beyond  $VDQ_2$  does not exceed  $u_2$  and the total usage beyond  $VDQ_6$  does not exceed  $u_6$ .

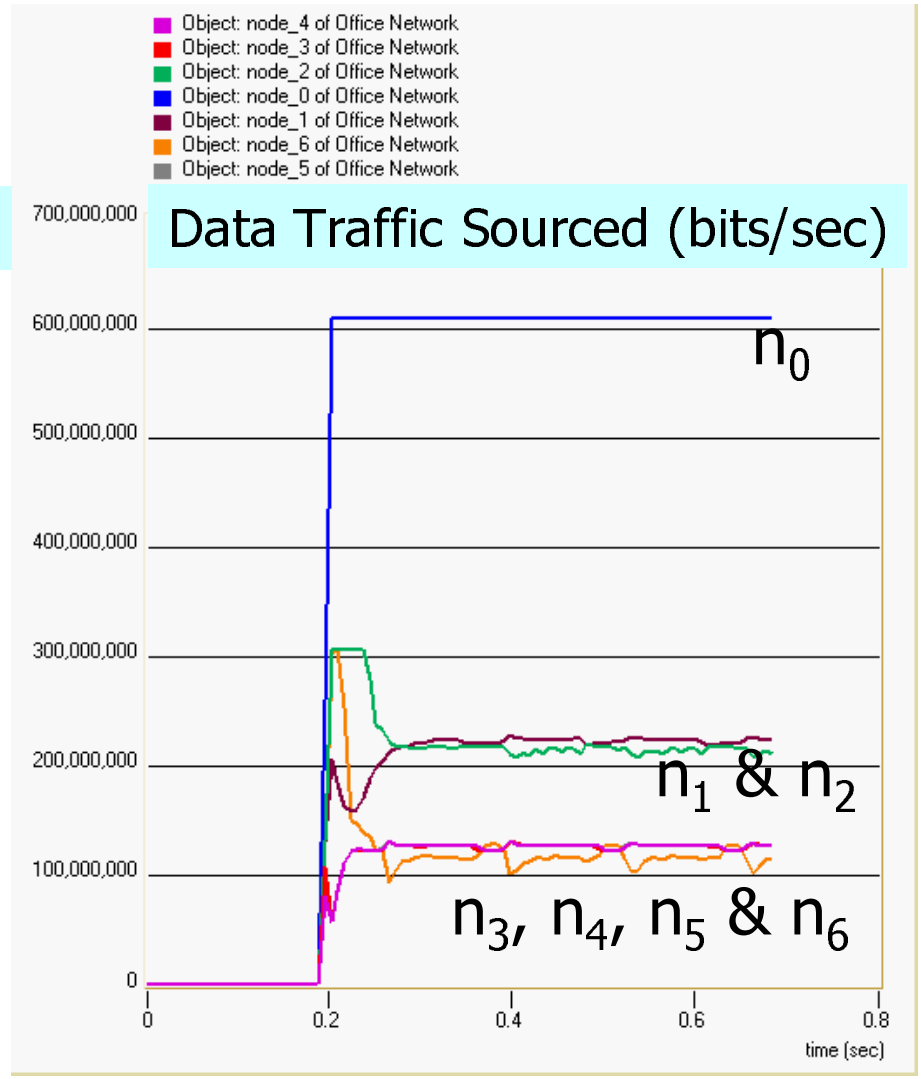
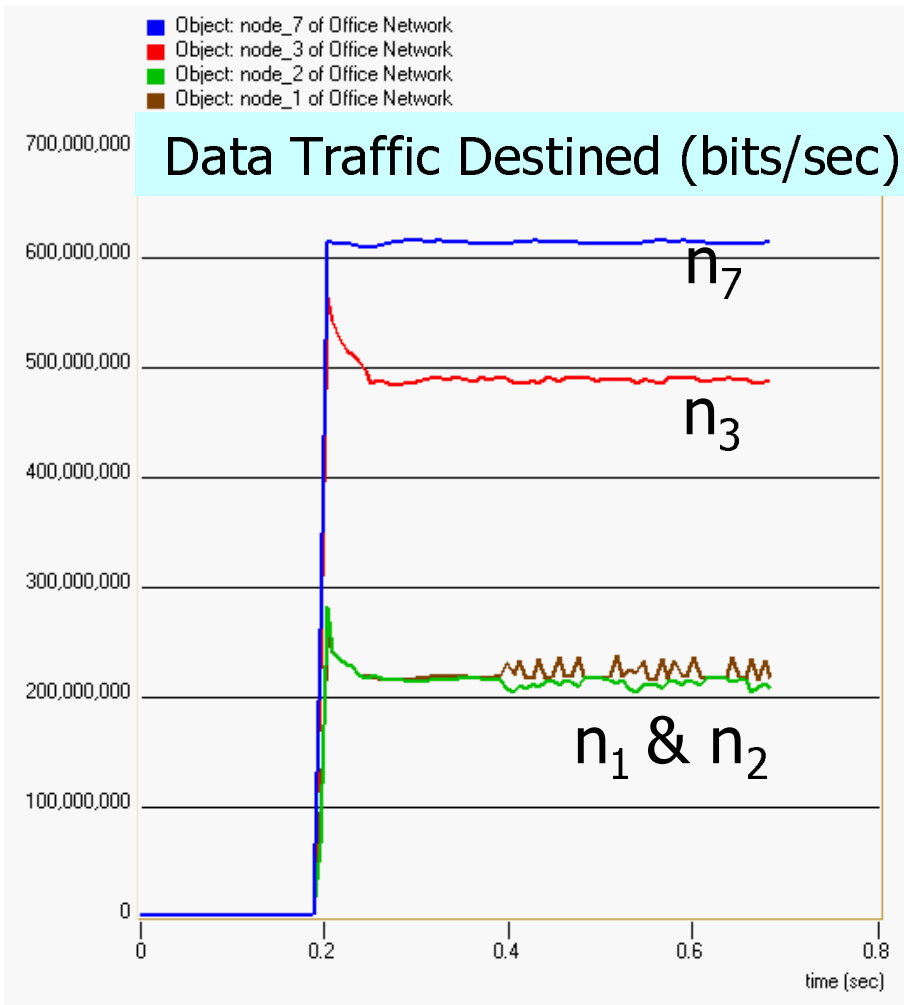
# Description Cont.

- 1 picture worth 1000 words:



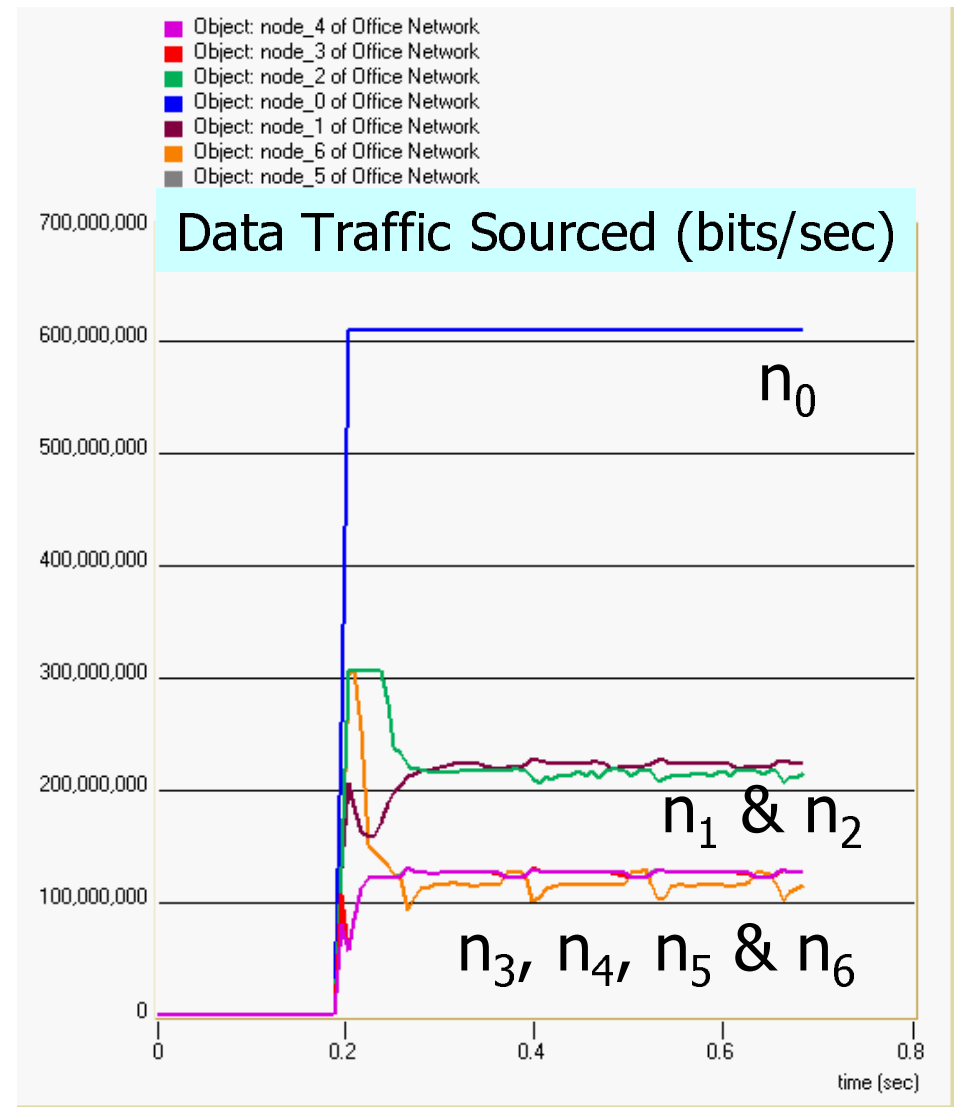
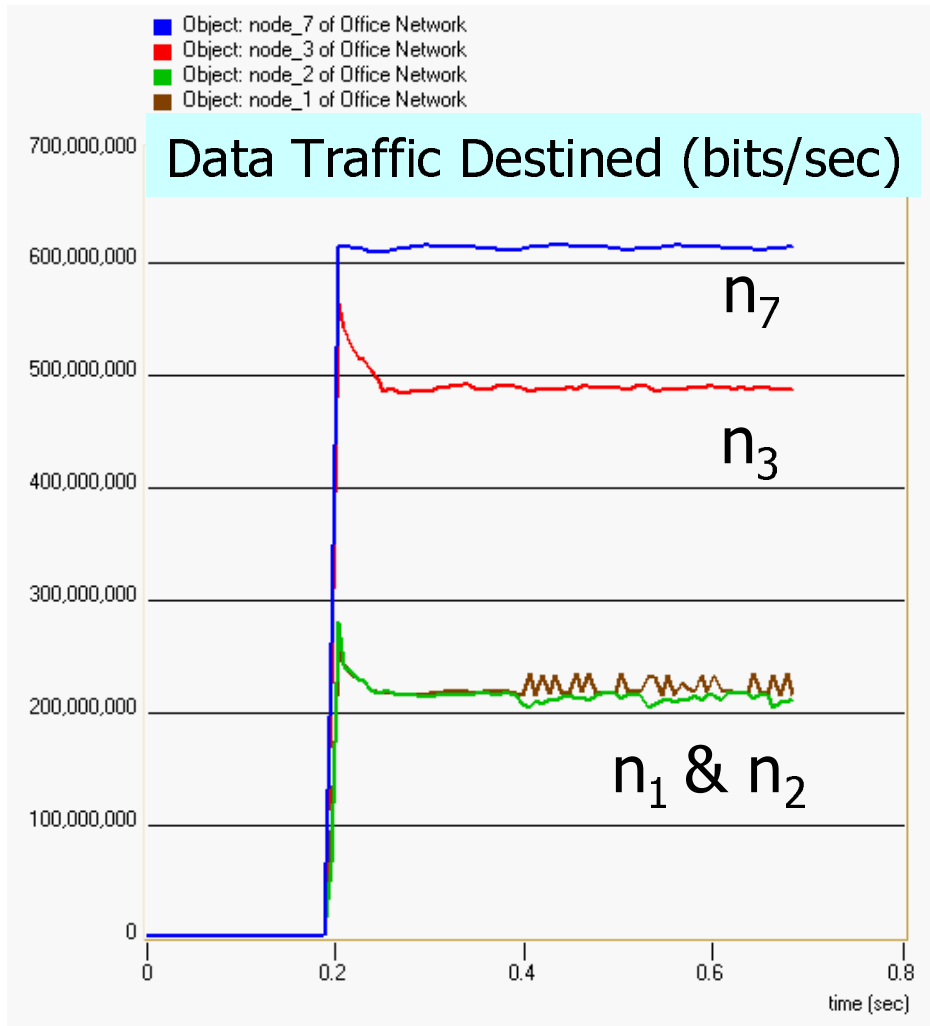


# Simulation I (VDQ, max choke point =4)





# Simulation II (VDQ, max choke point = 1)





## Conclusion

---

- VDAQ increases the total cost of the network.
- In most cases the gain may not be worth the cost increases incurred compared to SRP.
- The number of nodes on the ring will be limited by the number of VDAQ's that can be supported.
- One choke point will be sufficient for many scenarios.