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# Transit Path Design and Inter-Operability

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# Agenda

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- Introduction
- Requirements
- Simulation Results
- Conclusion



# Introduction

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- Some architectures have 1 transit buffer
- Some architectures have 2 transit buffers
- Buffers operating in store-and-forward or cut-through mode
- Each architecture has a unique congestion control and fairness algorithm with:
  - Unique messaging infrastructure
  - Unique control parameters
- Transit path design is critical for performance



# Introduction

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- Questions:
  - Can different architectures live in the same ring?
  - If yes, can we get the same or similar performance in a heterogeneous ring compared to a homogenous ring?
- Answers:
  - Yes!
  - And yes!



# Requirements

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- The mass of the problem revolves around congestion control
- Need to make sure no one is favored based on a particular architecture or location on the ring
- How?
  - Need to investigate on a case-by-case basis
  - Who wants to talk to whom?
- A single control message format needs to be defined



# Toy Scenario

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- Simple Single Transit Buffer (SSTB)
  - One transit buffer (cut-through)
  - XON/XOFF to stop low priority traffic
- SRP
  - Two transit buffers (store-and-forward)
  - Usage messages
- Traffic pattern:
  - All traffic destined from the nodes to the hub
  - ~%30 high priority, ~%70 low priority
- Homogeneous ring performance comparison presented in May 2001

[http://www.ieee802.org/17/documents/presentations/may2001/nu\\_ctvst\\_02.pdf](http://www.ieee802.org/17/documents/presentations/may2001/nu_ctvst_02.pdf)



# Interaction

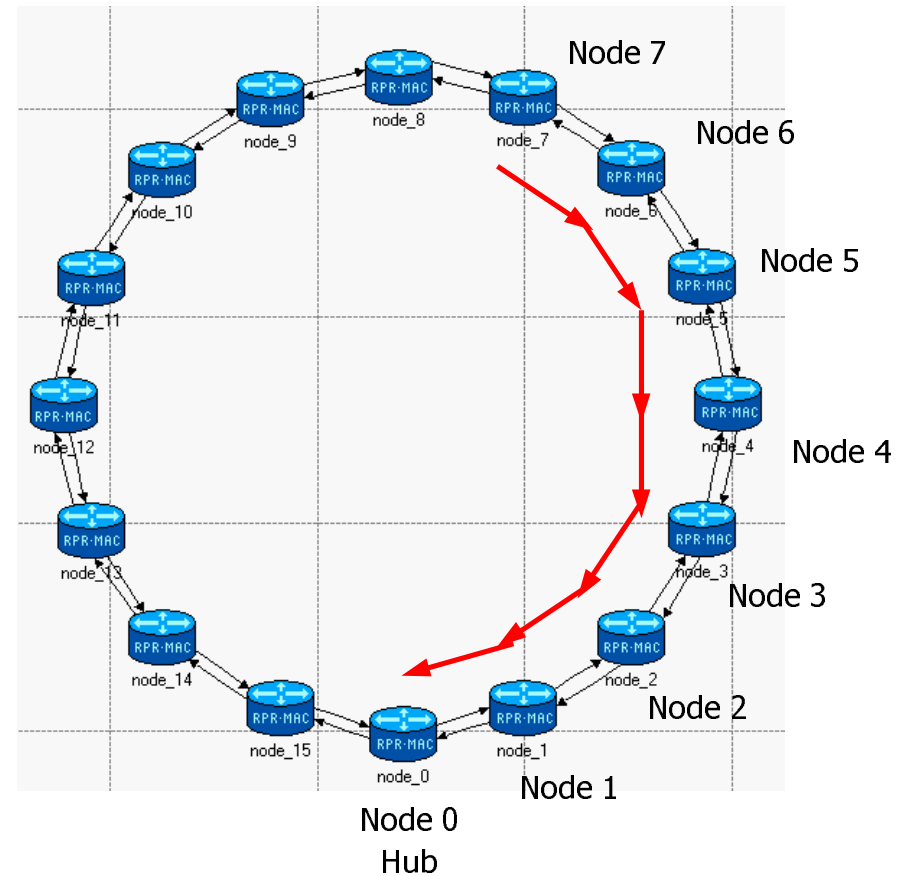
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- Relay control messages originated from other type of nodes upstream
- Simple Single Transit Buffer (SSTB)
  - If (value==NULL\_USAGE) then XON; else XOFF
  - Token bucket for rate shaping of low priority TX
- SRP
  - Cut allowed usage by a % when XOFF (used 50%)
  - Limiting data rate on the ring to a percentage of the link rate (used 95%)
  - Token bucket for rate shaping of low priority TX



# Different combinations

	Single Tb	Dual Tb
Mix 1	7,5,3,1	6,4,2,0
Mix 2	6,4,2,0	7,5,3,1
Mix 3	7,6,5,4	3,2,1,0
Mix 4	3,2,1,0	7,6,5,4
Mix 5	7,4,3,0	6,5,2,1
Mix 6	6,5,2,1	7,4,3,0







# Scenarios

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- CT: all single TB nodes
  - CT\_rt: single TB with rate shaper for Low Priority
- SF: all dual TB nodes
  - SF\_rt: dual TB with rate shaper for Low Priority
- Mix#: single and dual TB nodes, each with rate shaper for Low Priority



## Details

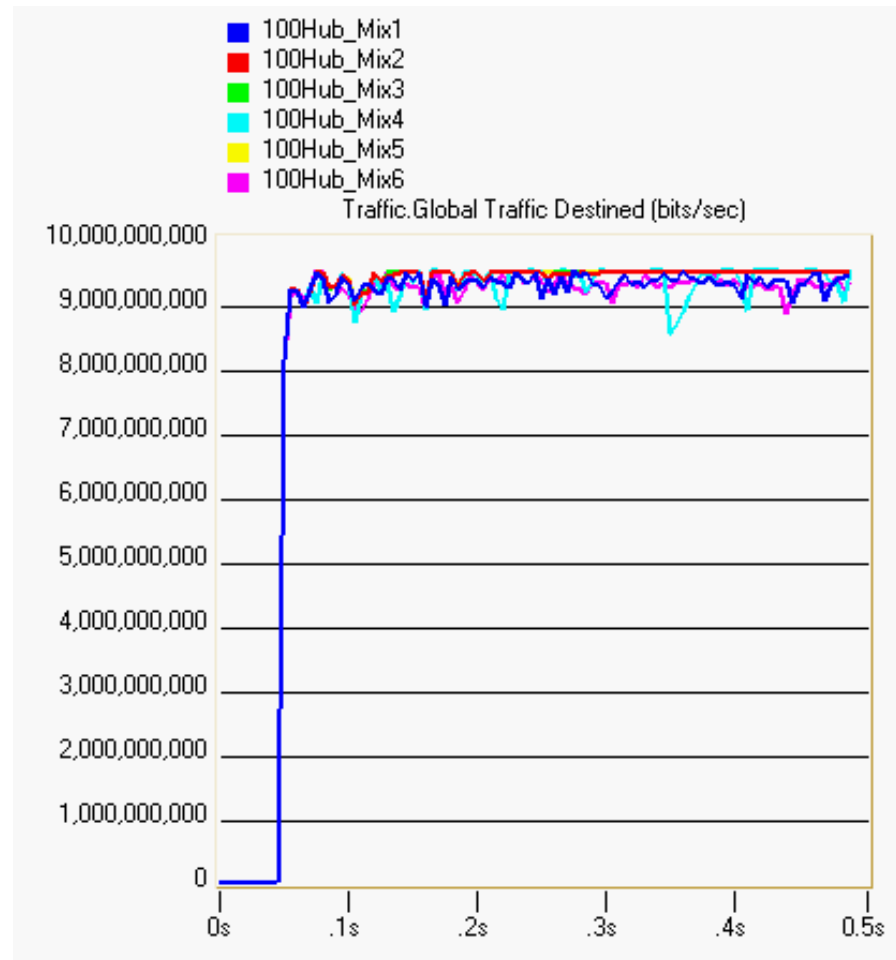
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- High Priority TX CAR = 420 Mbps
  - (CT nodes only)
- Low Priority TX shaping rate = 930 Mbps
  - (CT and SF nodes)
- SF nodes limited to %95 of link rate
  - (no limit for CT nodes)



# Throughput (1), Heterogeneous Rings

## Traffic Destined (bps)

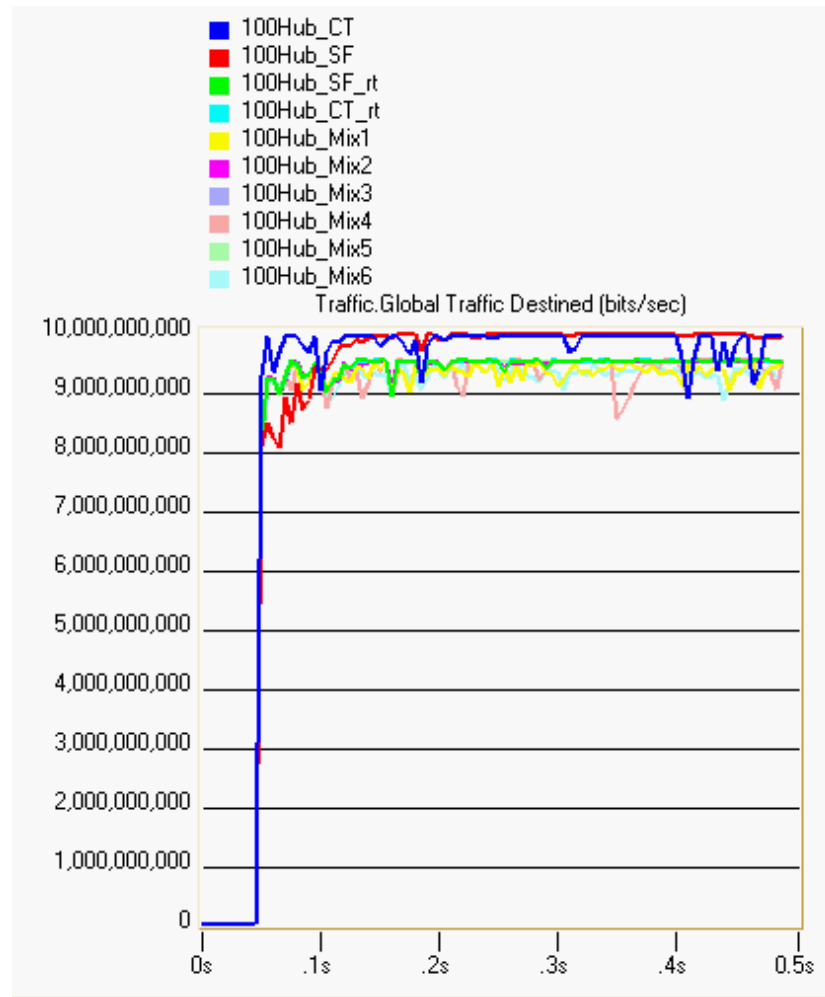




# Throughput (2)

## Heterogeneous Rings compare with Homogeneous Rings

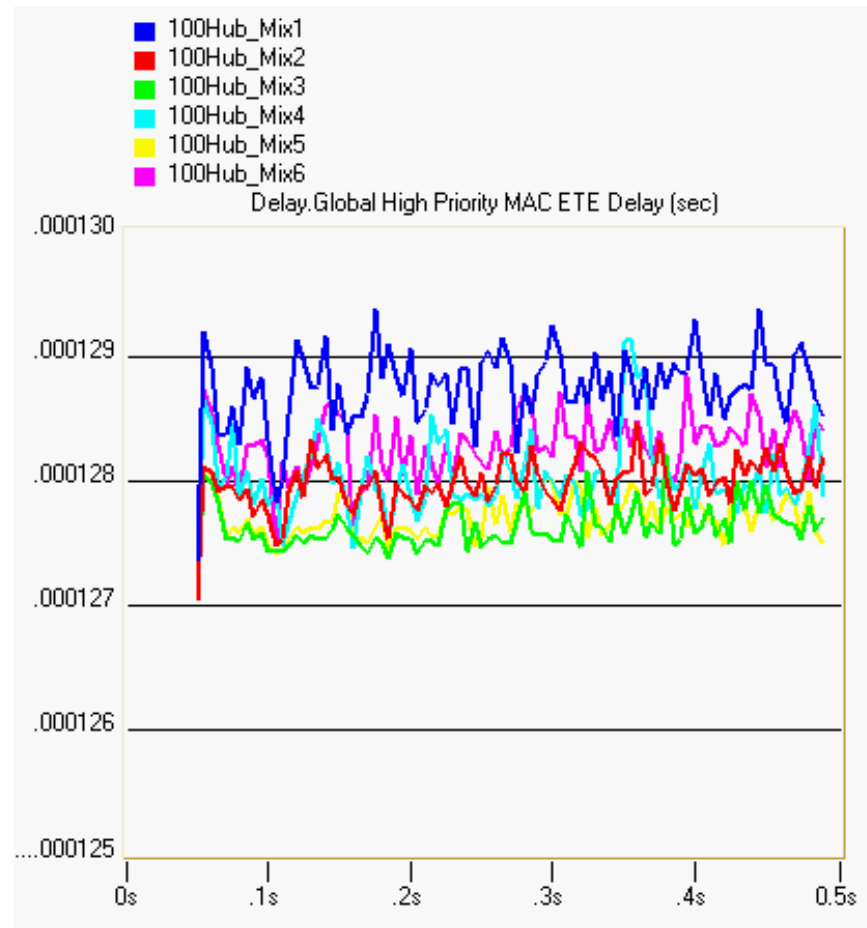
### Traffic Destined (bps)





# Delay (1), Heterogeneous Rings

## High Priority MAC ETE Delay (sec)





# Delay (2), Heterogeneous Rings

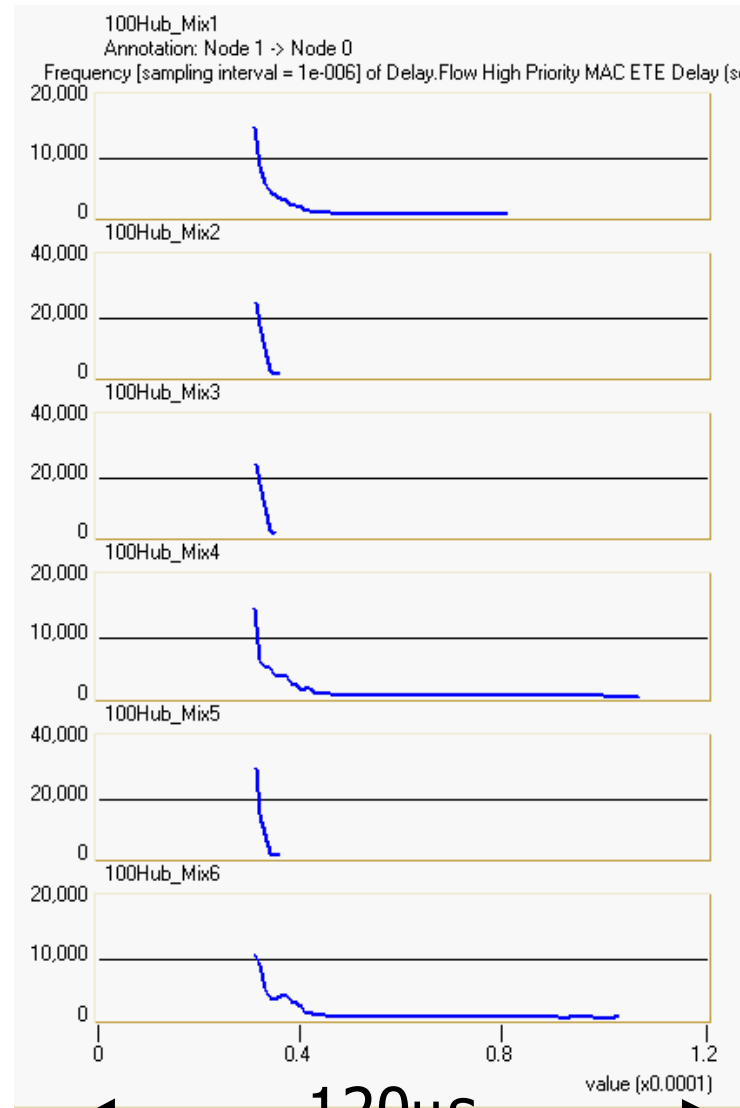
High Priority

MAC ETE Delay (sec)

Histogram

Max Jitter observed:

72  $\mu$ sec





## Conclusion

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- Inter-operability is possible!
- Need cooperation of all nodes involved in a heterogeneous ring
- Need to agree on common control messaging format and parameters