



**IKN**  
Institut für  
Kommunikationsnetze

# Performance Comparison of RPR MAC Protocols

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# Proposed fairness mechanisms for RPR

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## **Reactive:**

- Back pressure (Cisco, Nortel)
- Signalling to upstream node (Dynarch, Fujitsu)

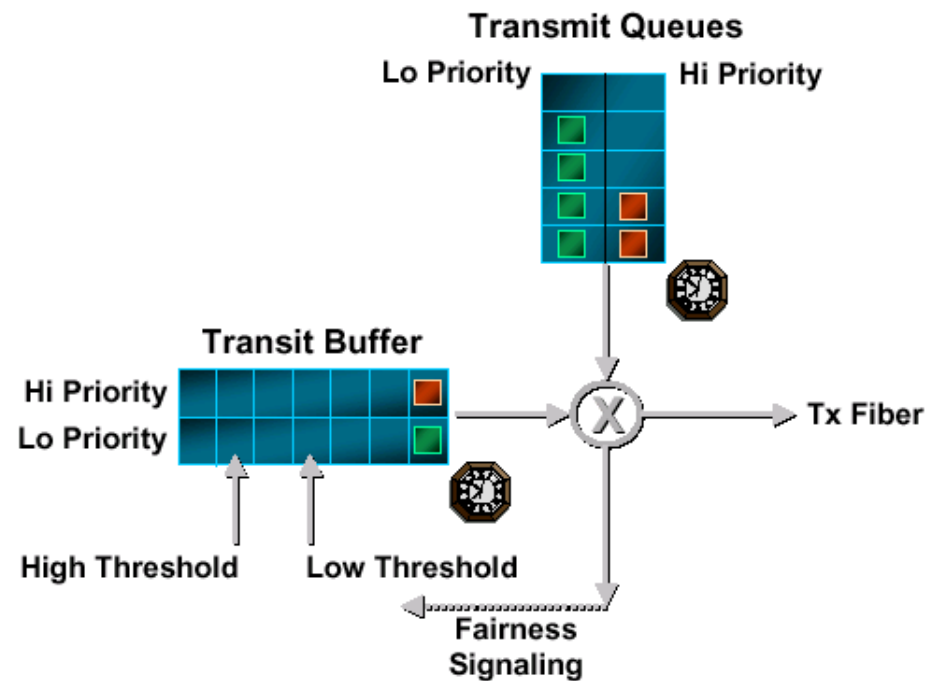
## **Pro-reactive:**

- Control cycle with credit-based rates (Lantern, Vienna Univ: of Technology)
- Control cycle with reservation-based rates (Vienna Univ: of Technology)

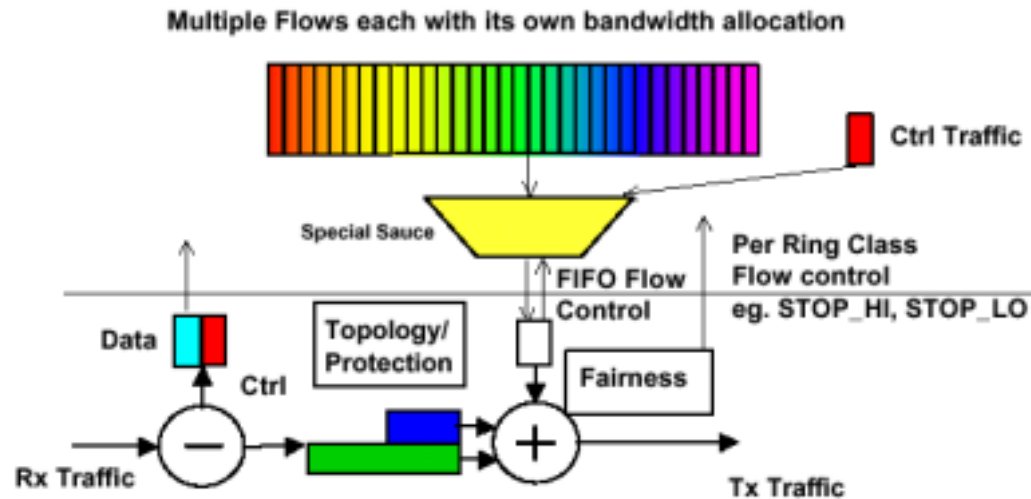
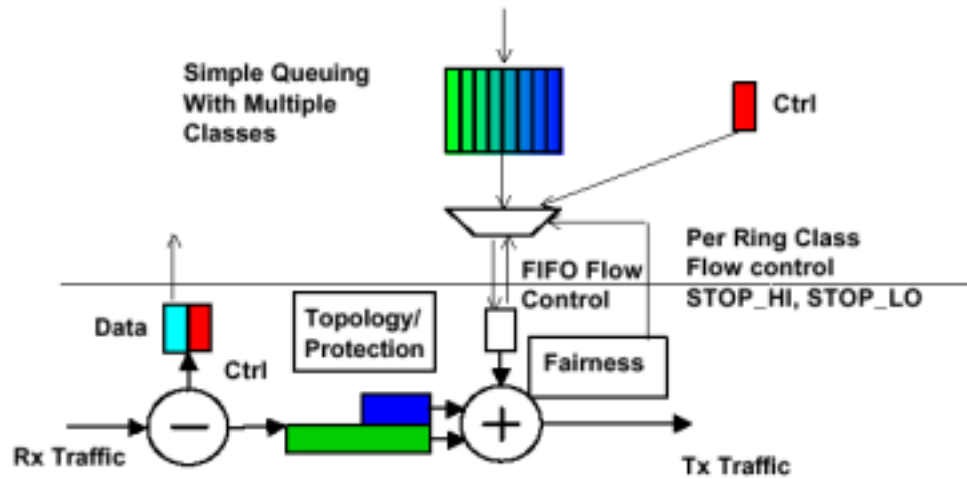
## **Hop-by-hop scheduling (Luminous)**

# Transmit Packet Handling

Priority handling  
Transmit decision algorithm  
Fairness



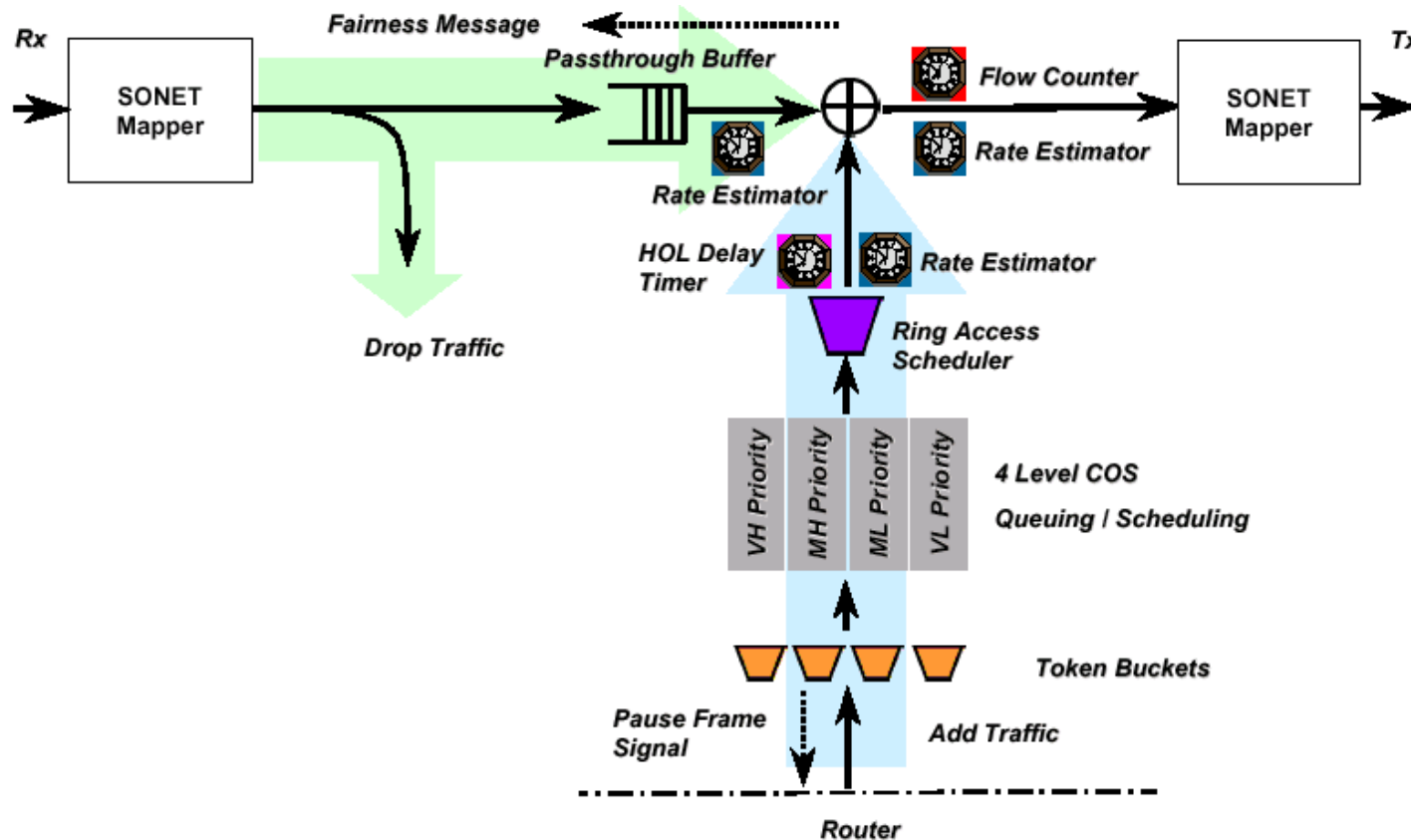
# RPR MAC Model



IEEE March 2000



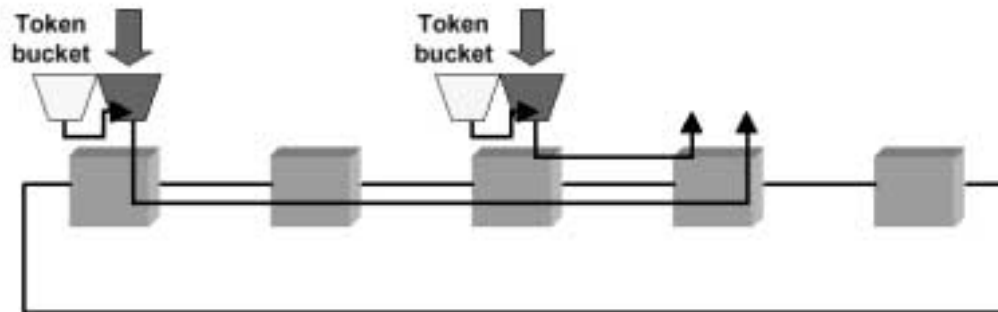
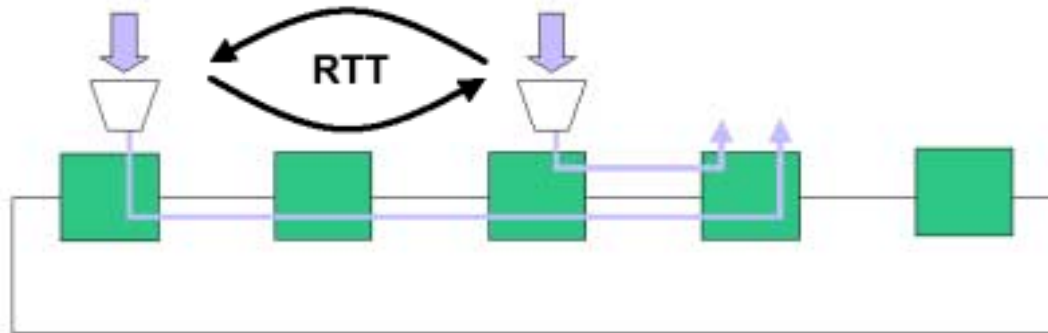
# iPT Fairness Controlled Access Protocol



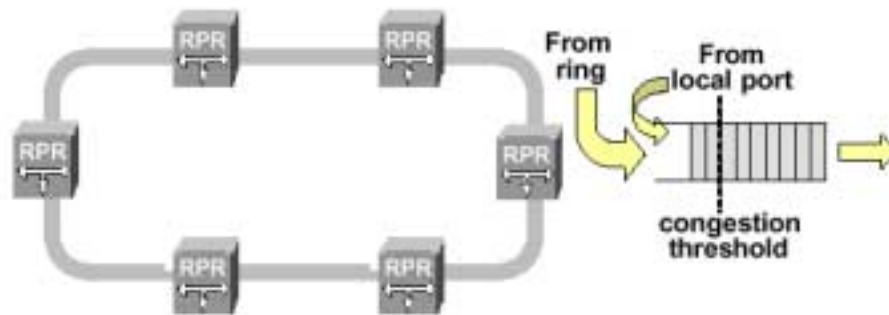
IEEE May 2000



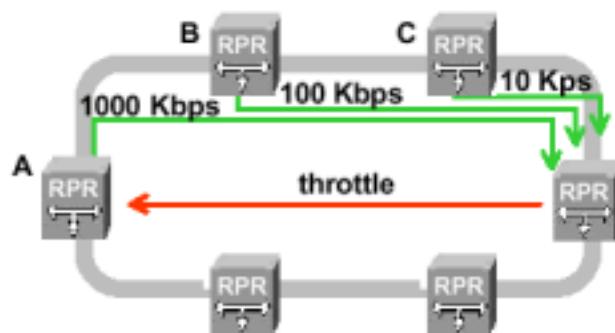
# Distributed resource management



# Congestion management and fairness parameters



Best Effort Model  
Throttle individual sender

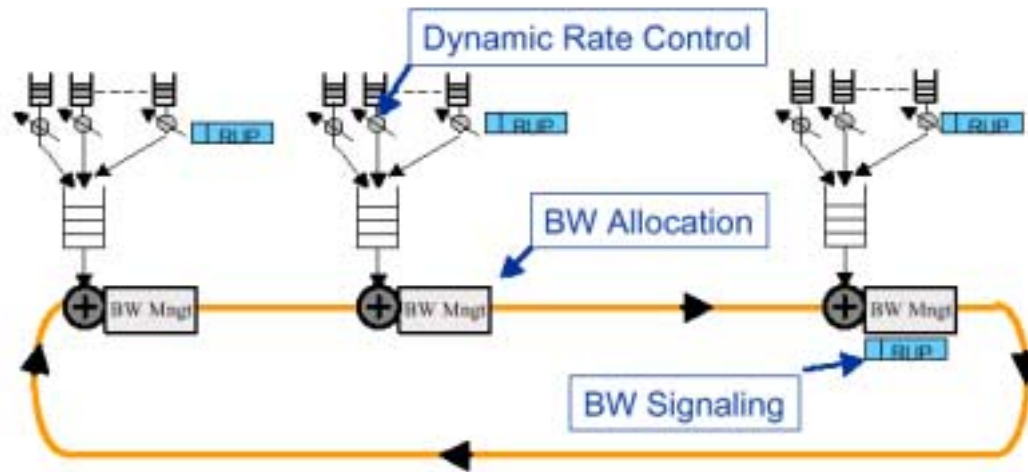


rate cache (per CUG)

sender	rate
A	1000
...	...
...	...
B	100
...	...
...	...

C

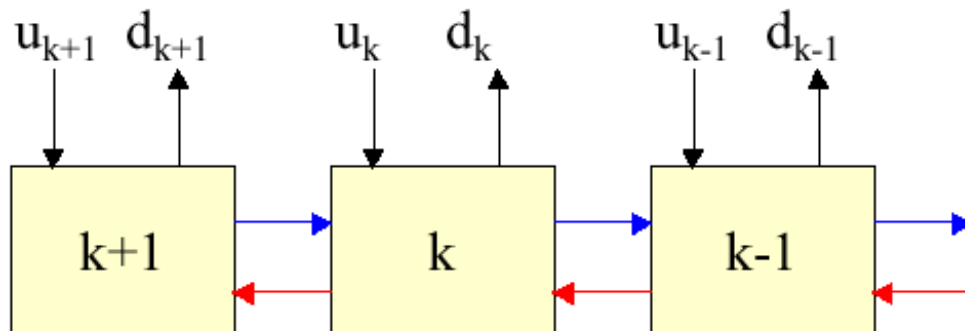
# Ring Access Control



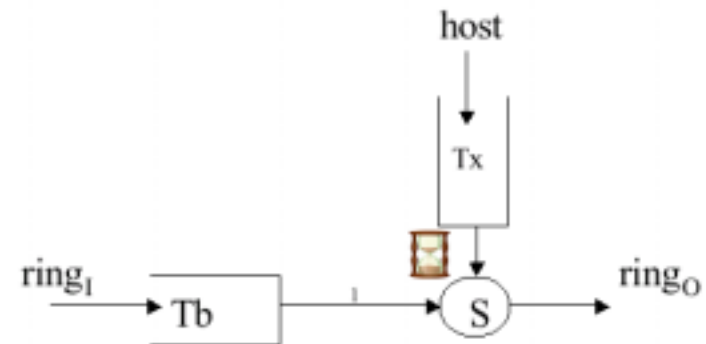
- WFQ Model of bandwidth allocation
- Support per-SLA QoS guarantee
- No upstream unfairness (BW, delay, Jitter)
- High utilization (> 95% per link)
- 100% spatial reuse
- Scalable (#nodes, #rings, # customers)
- Simple and *robust* (aggregate) congestion control signaling



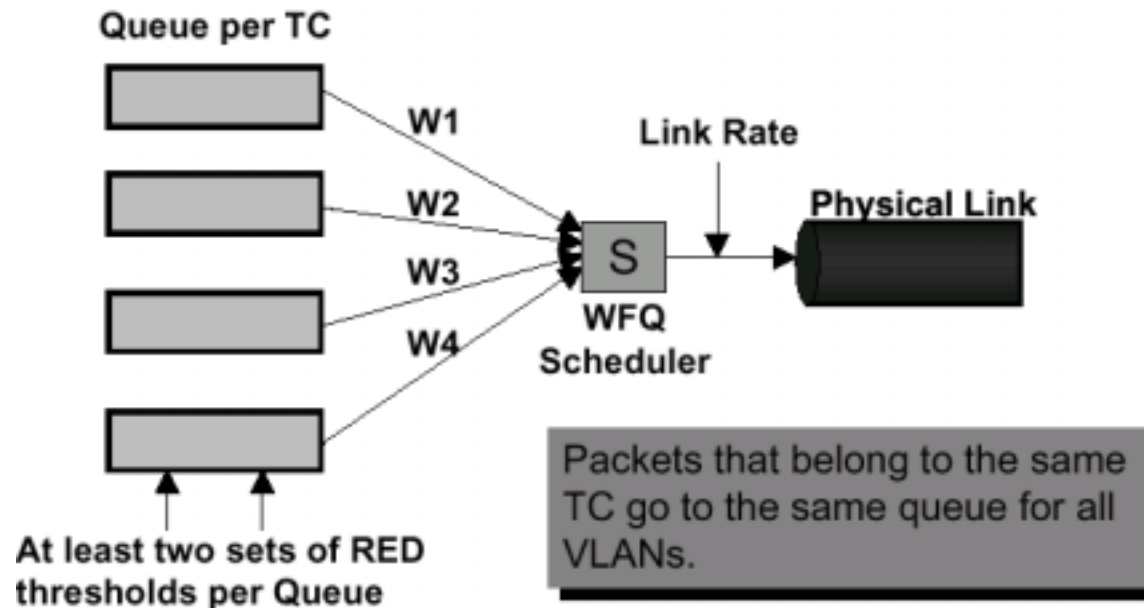
# Congestion management and fairness parameters



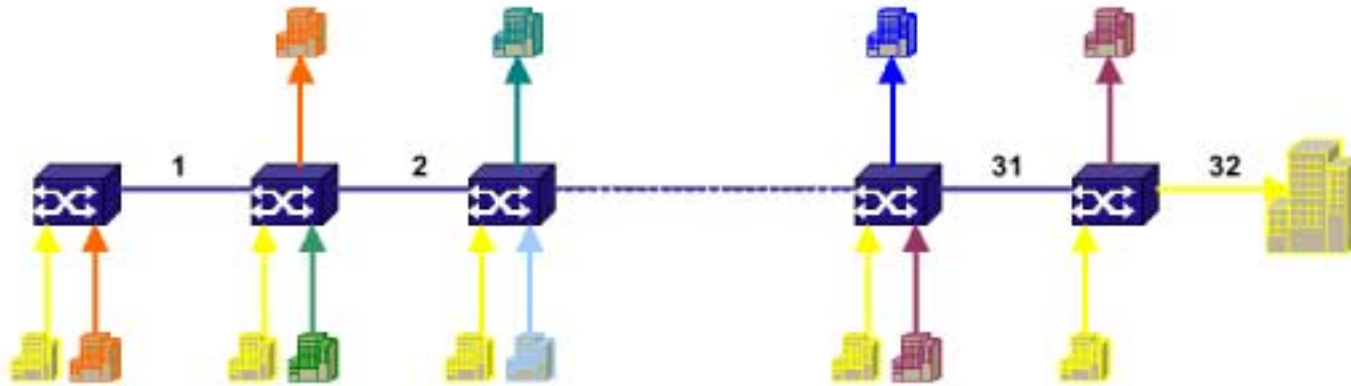
- $u_k$  : actual usage (sourced traffic rate) of node k
- $a_k$  : allowed usage (sourced traffic rate) of node k
- $d_k$  : drop traffic rate at node k
- $f_k$  : actual forward rate from node k+1 to node k-1
- $u_{\max\_k}$  : maximum provisioned usage rate factor for node k
- $u_{\max}$  : maximum provisioned usage rate factor of downstream node
- $u$  : usage value received from downstream node



# DiffServ Switch



# DiffServ Operation



## Per hop Behavior

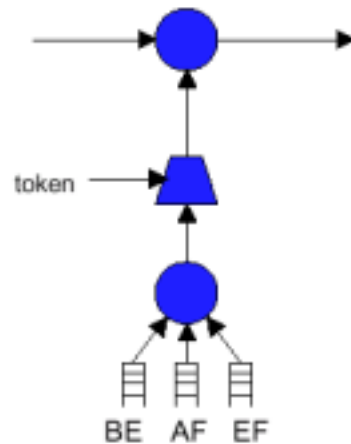
- Expedited Bandwidth
- Assured Bandwidth
- Best Effort

## Waited fair queueing

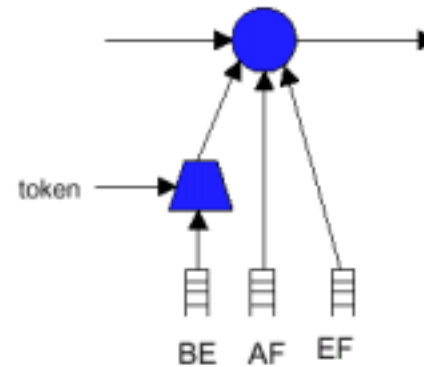
IEEE May 2001



# Nortel OPE-RPR Ring



1-add

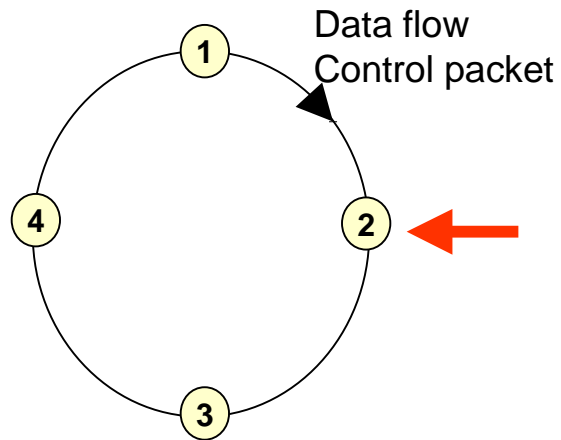


3-add



# Fairness mechanism (2)

Example with single ring



Coordinated table values  
In node 2

Cycle i-1

Flow	High	Low
1 -> 2	H12	B12
1 -> 3	H13	B13
1 -> 4	H14	B14
2 -> 3	H23	B23
2 -> 4	H24	B24
2 -> 1	H21	B21
3 -> 4	H34	B34
3 -> 1	H31	B31
3 -> 2	H32	B32
4 -> 1	H41	B41
4 -> 2	H42	B42
4 -> 3	H43	B43

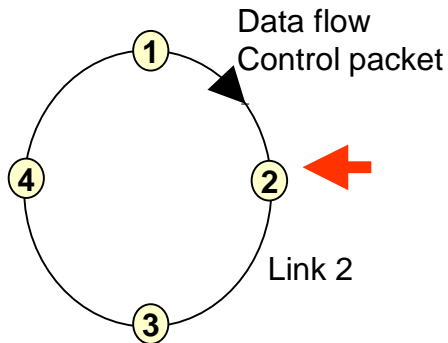
Old table in node 2

Cycle i

Flow	High	Low
1 -> 2	H12	B12
1 -> 3	H13	B13
1 -> 4	H14	B14
2 -> 3	H23	B23
2 -> 4	H24	B24
2 -> 1	H21	B21
3 -> 4	H34	B34
3 -> 1	H31	B31
3 -> 2	H32	B32
4 -> 1	H41	B41
4 -> 2	H42	B42
4 -> 3	H43	B43

New table in node 2

# Fairness mechanism (3)



## Actions in node 2:

- Determine fairness on link 2
  - Correct flows H13, H14  
H23, H24, H21, H43
  - Correct flows L13, L14  
L23, L24, L21, L43
- Determine total amount of coordinated capacity over link 2
- Write new demand of node 2 into control packet
- Send control packet to next node at the scheduled time
- Transmit coordinated flows H23, H24, H21, L23, L24, L21
- Refrain from transmission during rest of the coordinated capacity
- Transmit by immediate access according to the stored rates for each destination

Cycle i-1

Flow	High	Low
1 -> 2		H12
1 -> 3		H13
1 -> 4		H14
2 -> 3		H23
2 -> 4		H24
2 -> 1		H21
3 -> 4		H34
3 -> 1		H31
3 -> 2		H32
4 -> 1		H41
4 -> 2		H42
4 -> 3		H43

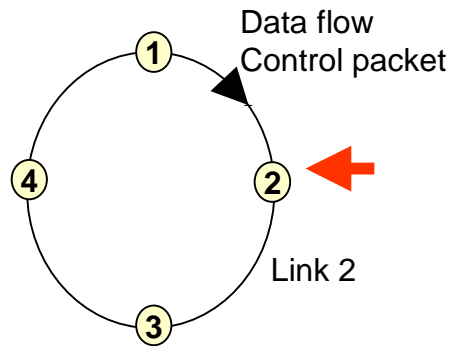
Old table in node 2

Cycle i

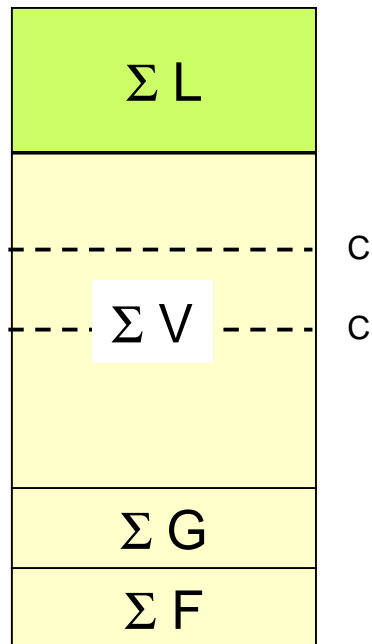
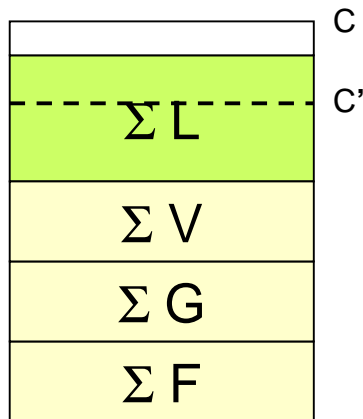
Flow	High	Low
1 -> 2		H12
1 -> 3		H13
1 -> 4		H14
2 -> 3		H23
2 -> 4		H24
2 -> 1		H21
3 -> 4		H34
3 -> 1		H31
3 -> 2		H32
4 -> 1		H41
4 -> 2		H42
4 -> 3		H43

New table in node 2

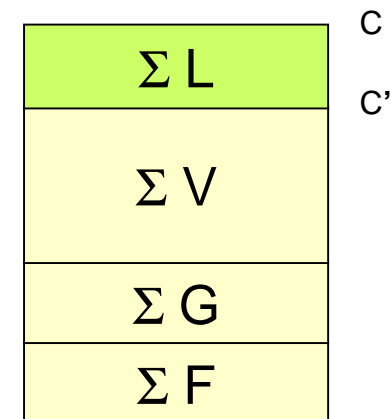
# Fairness mechanism (4)



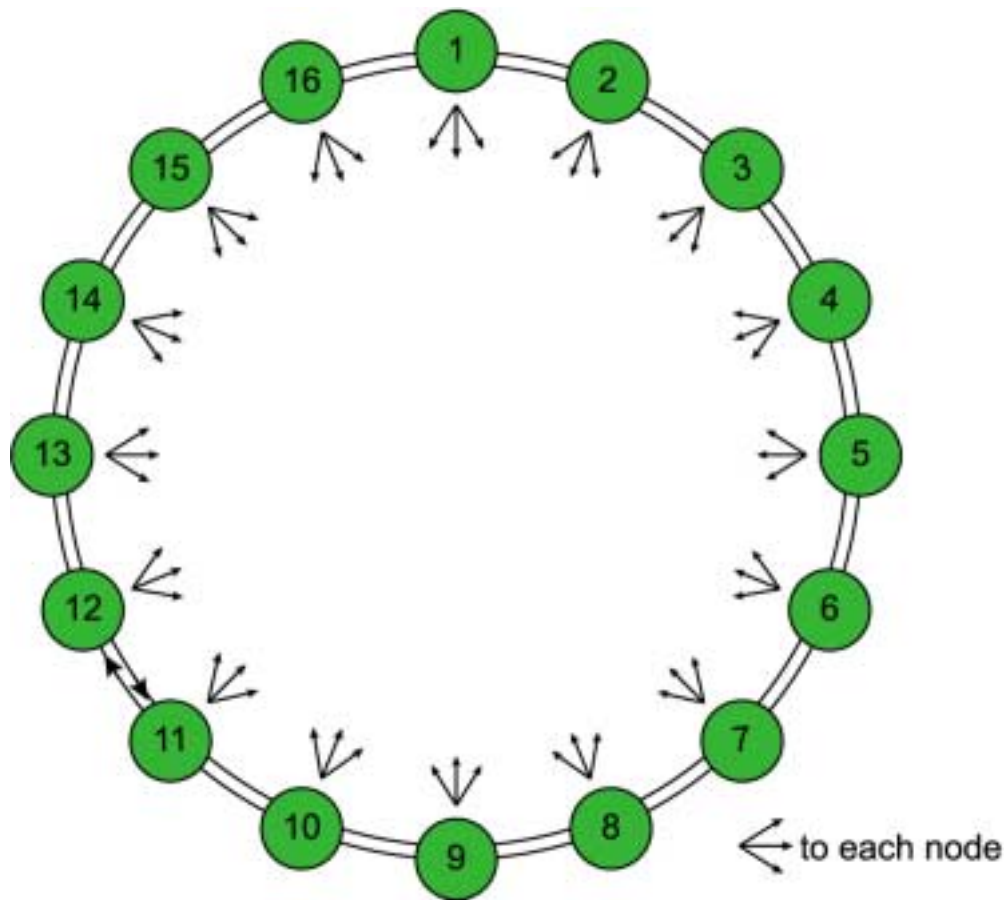
- $\Sigma L$  : all low-traffic flows
- $\Sigma V$  : all non-guaranteed high-traffic flows
- $\Sigma G$  : all guaranteed high-traffic flows
- $\Sigma F$  : all CBR traffic flows
- $V_i = H_i - G_i$  : variable part of high-priority traffic flow



Link capacity  $C$   
 $C - C'$  is minimal capacity  
 for low priority when present



# Dual-Ring - Traffic scenario 1



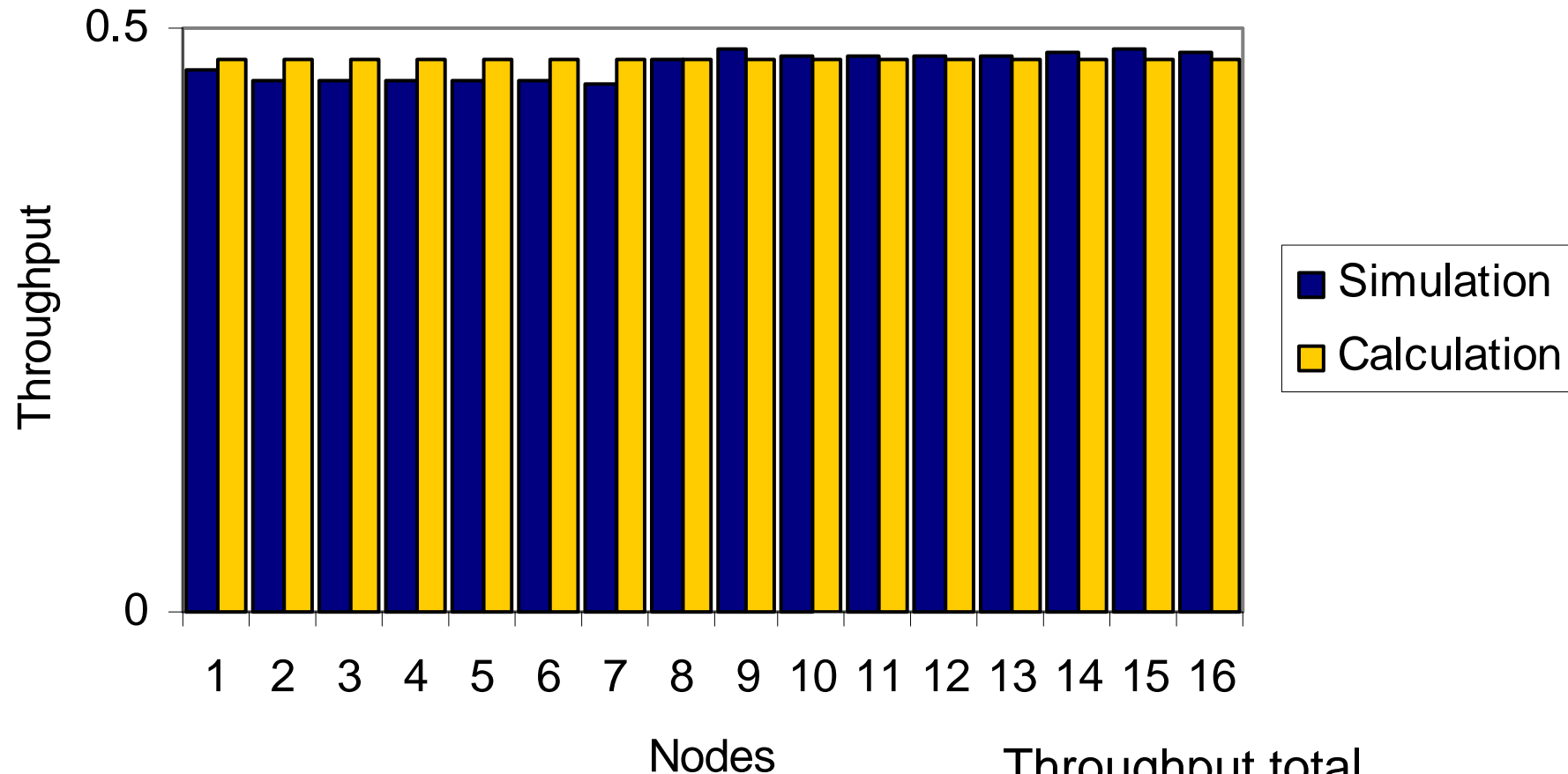
Uniform traffic  
Saturated sources  
16 nodes

Constant packets  
8000 bits

Cyclic reservation protocol

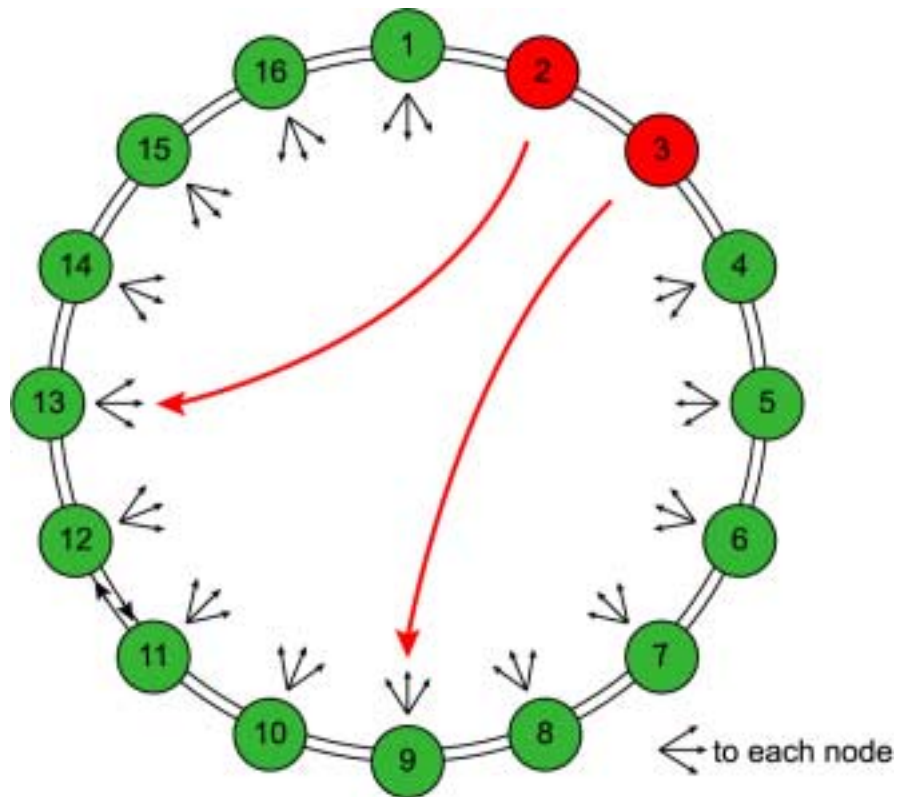


# Dual-Ring - Traffic scenario 1



Throughput total  
Simulation : 7.5  
Calculation : 7.55

# Dual-Ring – Traffic scenario 2

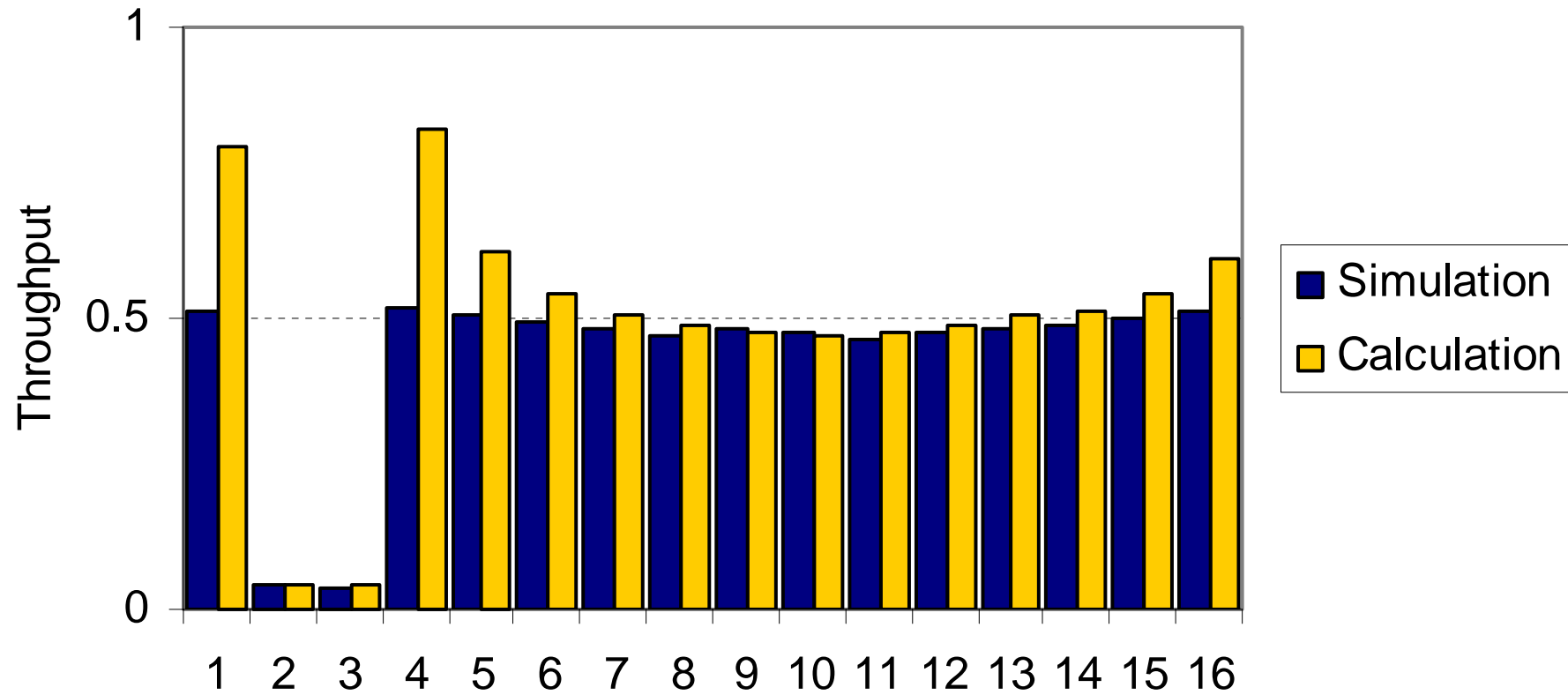


Uniform traffic  
Saturated sources  
16 nodes

Constant packets  
8000 bits

Cyclic reservation protocol

# Dual-Ring – Traffic Scenario 2



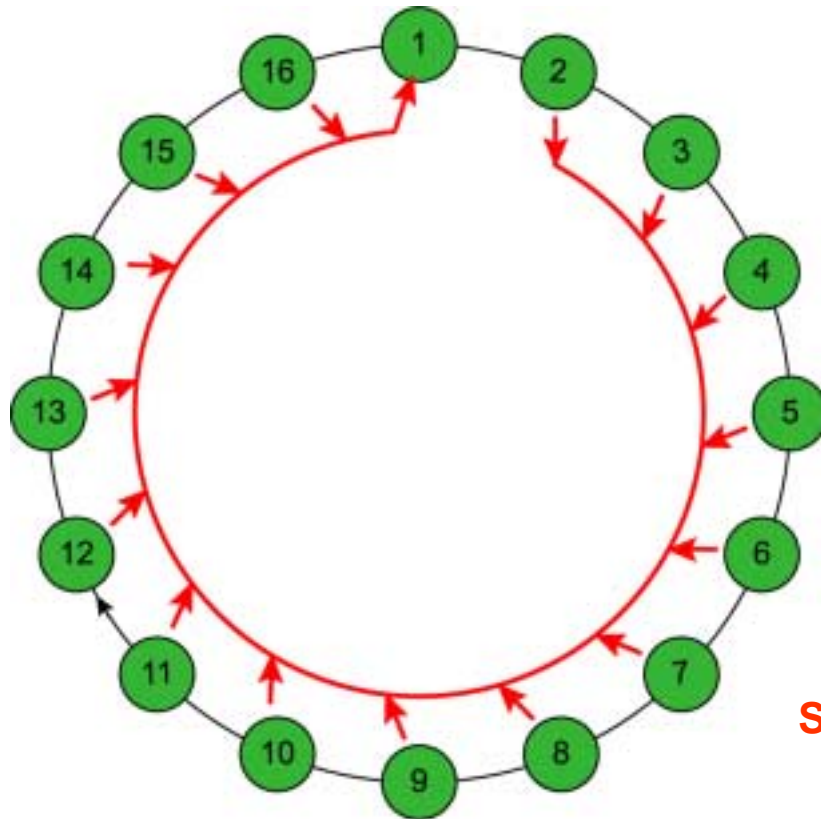
Nodes

Throughput total

Simulation : 6.9

Calculation : 7.9

# Single Ring-Traffic scenario 3



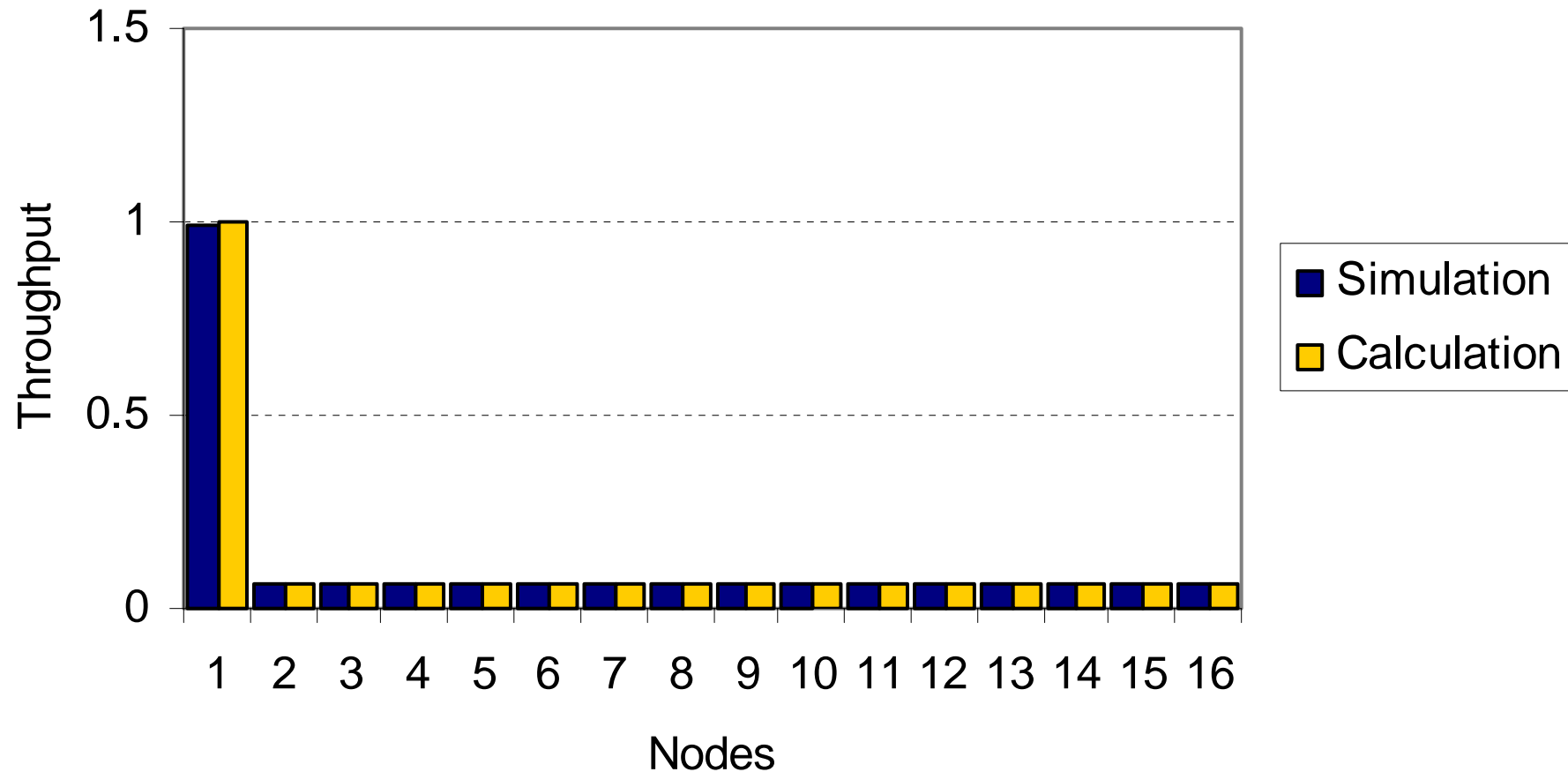
Single ring

Uniform traffic  
Saturated sources  
16 nodes

Constant packets  
8000 bits

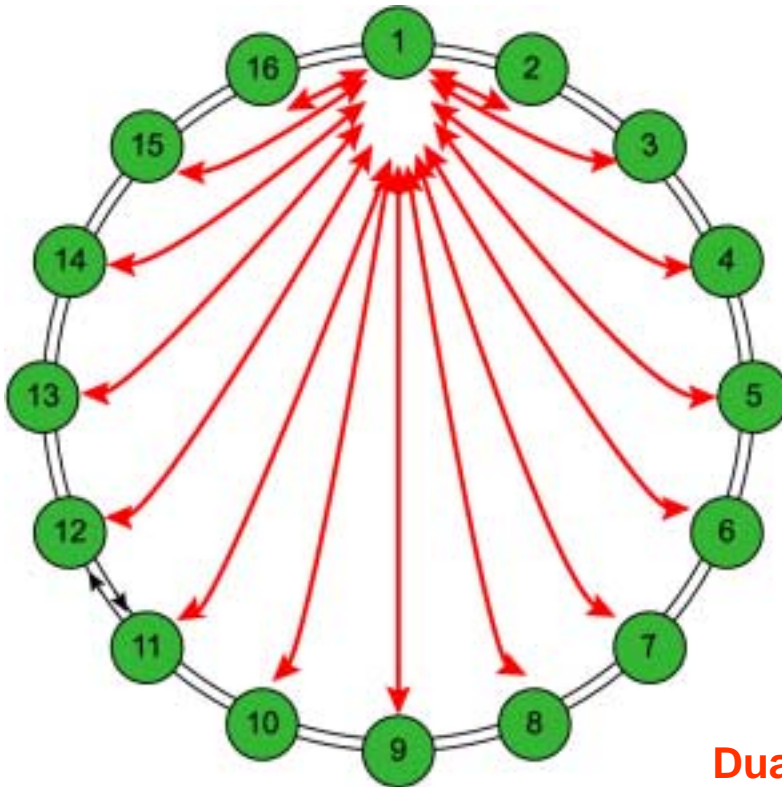
Cyclic reservation protocol

# Single Ring-Traffic scenario 3



Throughput total  
Simulation : 1.98  
Calculation : 1.99

# Dual-Ring – Traffic scenario 4

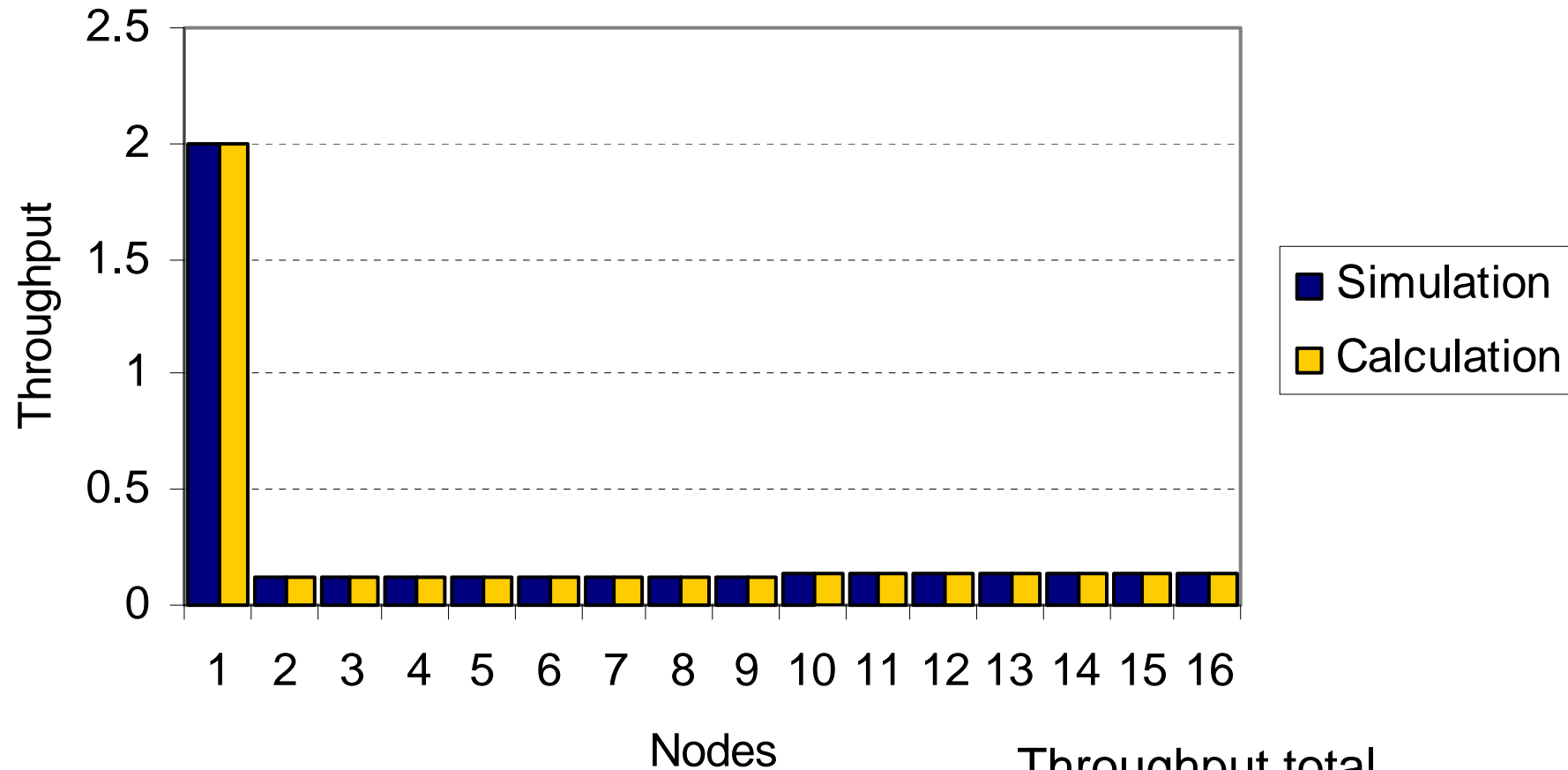


Uniform traffic  
Saturated sources  
16 nodes

Constant packets  
8000 bits

Cyclic reservation protocol

# Dual-Ring – Traffic scenario 4



Throughput total  
Simulation : 3.98  
Calculation : 3.99