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Re:	IEEE 802.20 Practice Letter Ballot 2	
Abstract	This document proposes text for the LDPC encoding scheme in the IEEE 802.20 standards draft, based on the latest version of the text in the 3GPP2 UMB standard	
Purpose	For adoption into 802.20 standard draft	
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Introduction

Proposed text that describes the LDPC coding design can be found in the subsequent sections of this contribution. This version of the text is based on the most recent working document for the ‘Physical Layer for Ultra Mobile Broadband (UMB) Air Interface Specifications’, 3GPP2 C. P0084 v1.098 [2].

Proposed Text Changes to IEEE 802.20 Draft D0.1m

5.1.7.3 Forward Error Correction

Table 119 Types of Forward Error Correction for the Reverse and Forward Link Channels

Channel	Type of Coding
Reverse Orthogonal Frequency Division Multiple Access Data Channel	Rate-1/5 Turbo, LDPC or Rate-1/3 Convolutional
Forward Data Channel	Rate-1/5 Turbo, LDPC or Rate-1/3 Convolutional

5.1.7.3.4 Low Density Parity Check (LDPC) Encoding

LDPC encoding shall be used to encode the CRC-padded subpackets of the Forward Data Channel if the variable LDPCSupportedFL is set to ‘1’, and if the length of the packet received from the FTC MAC Protocol, as described in Section 4, is greater than or equal to MaxPacketSizeTurboSixInterlace or MaxPacketSizeTurboEightInterlace, for the case when the input FTC MAC packet corresponds to a six or eight interlace transmission respectively, where MaxPacketSizeSixInterlace or MaxPacketSizeEightInterlace are configuration attributes of the FTC MAC Protocol [3]. No LDPC encoding shall be used for interlacing structures involving extended transmissions. The FTC MAC Protocol determines the interlacing structure being used for a given FTC MAC packet.

LDPC encoding shall also be used to encode the CRC-padded subpackets of the Reverse Orthogonal Frequency Division Multiple Access Data Channel if the variable LDPCSupportedRL is set to ‘1’, and if the length of the packet received from the RTC

MAC Protocol, as described in Section 4, is greater than or equal to $\text{MaxRLPacketSizeTurbo}$, except if this packet corresponds to an interlacing structure using extended transmissions. No LDPC encoding shall be used for interlacing structures involving extended transmissions. The RTC MAC Protocol determines whether or not a given packet is transmitted on an interlacing structure using extended transmissions.

5.1.7.3.4.1 Choice of Base Parity Check Matrix

The LDPC code to be used is specified in terms of a base parity check matrix corresponding to different lifting orders. Different base parity check matrices G_i , $0 \leq i < 6$, are specified in 5.1.7.3.4.5.1. Note that these parity check matrices represent a lifted LDPC code, i.e., the entries of the matrices are not binary numbers but rather positive integers representing shift values. In addition, these base matrices contain “NULL locations” which are denoted by solid bullets in 5.1.7.3.4.5.1. The interpretation of these matrices as an LDPC code will be described in 5.1.7.3.4.2.

Each matrix G_i has associated values k_B , n_B , s_B and L_{\max} which are also specified in 5.1.7.3.4.5.1. Here, k_B and n_B determine the size of the matrix G , while L_{\max} denotes the maximum lifting order. The number of columns and rows in G are given by n_B and $n_B - k_B$ respectively. The matrix G_i has associated $k_B = i+6$. s_B denotes the number of “state columns” in the matrix G_i and is equal to 3 for each of the matrices shown. A state column denotes elements of the codeword that are never transmitted. Each of the specified matrices has a maximum lifting order L_{\max} equal to 1024.

Given the CRC-padded input sub-packet of length k , the lifting value L is chosen as $\log_2 L = \lceil \log_2(k/11) \rceil$. Further, k_B is chosen as $\lceil k/L \rceil$. Note that k_B is at least equal to 6 according to this procedure. Based on this, the matrix index i is chosen as $i = k_B - 6 = \lceil k/L \rceil - 6$.

Henceforth, the index i will be dropped and the matrix G_i is referred to only as G .

5.1.7.3.4.2 Generation of the Parity Check Matrix

The base matrix G chosen in the previous section shall be converted to a new base matrix G' , corresponding to the actual lifting order L rather than the maximal lifting order $L_{\max} = 1024$. The matrix G' has the same size as G . An entry g' in G' shall be determined from the entry g at the same location in G according to the formula $g' = \lfloor gL/L_{\max} \rfloor$. NULL locations in G shall remain NULL locations in G' .

The matrix G' shall further be converted to a matrix G'' with twice the number of rows and columns as in G' . This shall be done by replacing each non-NULL entry g' in G' by a 2×2 matrix according to the following procedure:

- If g' is even, replace g' by a 2×2 matrix with first row being given by $[g'/2, \text{"NULL"}]$ and the second row being given by $[\text{"NULL"}, g'/2]$.
- If g' is odd, replace g' by a 2×2 matrix with the first row being given by $[\text{"NULL"}, (g'+1)/2]$ and the second row being given by $[(g'-1)/2, \text{"NULL"}]$.

NULL locations in G' shall be replaced by a 2×2 matrix containing entirely of NULL locations in G'' . The matrix G'' is the base parity check matrix of size $(2(n_B - k_B), 2n_B)$ with a lifting order of $L/2$.

The base matrix G'' shall be converted to a base matrix G''' by applying permutation P_i to the columns of G'' and permutation Q_i to the rows of G'' , where the permutation P_i and Q_i are described in Table 5.1.7.3.4.2-1. The subscript i in P_i and Q_i refers to the subscript $G_i = G$ and thus takes values in $0, \dots, 5$. The numbers in the third column (Q_i) and fourth column (P_i) of Table 5.1.7.3.4.2-1 indicate row index and column index of G'' corresponding to i .

The first (leftmost) $2k_B$ columns of G''' correspond to the information bits V^{in} and $(k_B L - k)$ zero-padded bits. The subsequent K_i columns (K_i depends on G_i) together with the first K_i rows form a lifted parity check matrix that consist of a degree-3 variable node (i.e., a column with three non-NULL elements) followed by $K_i - 1$ degree-2 variable nodes. The degree-2 parity nodes form a dual-diagonal structure and the degree-3 variable node closes the loop of the dual-diagonal structure. Each non-NULL entry of degree-2 variable node in the dual diagonal structure has the lifting parameter zero, corresponding to identity matrix, on both edges. The loop closing edges on the degree-3 node have the same lifting value "a". The non-loop edge of the degree-3 node has lifting parameter zero corresponding to identity matrix so the lifting structure of this degree-3 node is "a-0-a".

The remaining columns in G''' are degree-1 variable nodes. The base matrices G_i , $i=0, \dots, 5$ each contain, up to permutation, a dual diagonal structure, a loop closing degree-3 encoding node and a loop closing degree-2 node (excluding any edges connected to constraints that are in turn connected to degree-1 variable nodes). As the scaled base graph G' is transformed into the expanded graph G'' and permuted into G''' , each of the degree-2 and degree-3 loop closing nodes of G' generates an information node and an encoding node of the same degree, such that the encoding node is part of the core encoding structure of G''' described above. The loop closing degree-2 variable node of

G_i has cyclic shift 0 and $L_{\max}-1$. The corresponding loop-closing edges of the degree-3 variable node have cyclic shifts $3L_{\max}/4 - 1$ and $3L_{\max}/4$. The third non-loop edge of the degree-3 node has cyclic shift 0. As the graph G_i is scaled to generate the graph G' , with values of L from 16 to 1024, the scaled base matrices G' inherit the same structure on the loop-closing degree-2 and degree-3 nodes, with L_{\max} replaced by L . This induces the lifting value of 0 on the loop closing degree-2 information node in G'' , and the “a-0-a” lifting on the loop closing degree-3 encoding node of G''' , where $a = 3L/8$.

An example for the generation of G' , G'' , G''' is shown in 5.1.7.3.4.5.1.

The base matrix G''' shall be converted to a binary parity check matrix H''' by replacing each non-NULL entry in G''' by a $L/2 \times L/2$ square matrix with binary entries. An entry g''' in G''' shall be replaced by a cyclic shift matrix with parameter g''' . The cyclic shift matrix with parameter g''' is defined as the matrix whose value in the location (i,j) is given by ‘1’ if $(i-j) \bmod L/2 = g'''$, and is given by ‘0’ otherwise. Here, the location (i,j) denotes the i 'th row and j 'th column. NULL locations in G''' shall be replaced by an $L/2 \times L/2$ all-zeros matrix.

The CRC-padded input sub-packet of length k shall be extended to length $k_B L$ by inserting in the packet $z_p = k_B L - k$ zeros so that the resulting packet has length $k_B L$. Denote again the original CRC-padded input sub-packet by V^{in} and denote the zero-padded input by a column vector $V^1 = (V^1_0, V^1_1, \dots, V^1_{k'-1})$ where $k' = k_B L$. The locations of the zeros in V^1 are as follows. If V^1 is partitioned into $2k_B$ blocks of size $L/2$, then the zeros are inserted at the ends of blocks $2k_B-4$ and $2k_B-3$. Each block has an equal number of zeros if z_p is even and block $2k_B-3$ has one more zero than block $2k_B-4$ if z_p is odd.

More precisely, define $z_p' = \lfloor z_p/2 \rfloor$ and $z_p'' = \lceil z_p/2 \rceil$. Let the notation V^1_i and V^{in}_i denote the i 'th element of V^1 and V^{in} respectively. The elements of the vector V^1 are given by:

- $V^1_i = V^{\text{in}}_i$ for $i < (2k_B-3)(L/2) - z_p'$.
- $V^1_i = 0$ for $(2k_B-3)(L/2) - z_p' \leq i < (2k_B-3)(L/2)$.
- $V^1_i = V^{\text{in}}_{i-z_p'}$ for $(2k_B-3)(L/2) \leq i < (2k_B-2)(L/2) - z_p''$.
- $V^1_i = 0$ for $(2k_B-2)(L/2) - z_p'' \leq i < (2k_B-2)(L/2)$.
- $V^1_i = V^{\text{in}}_{i-z_p}$ for $i \geq (2k_B-2)(L/2)$.

A vector V''' of length $n_B L$ shall be defined as the vector which satisfies the following conditions:

- $H''V'' = 0$, where the matrix multiplication $H''V''$ is over the binary field.
- The first $k_B L$ entries of V'' are the same as the entries of V^1 .

The vector V'' is of length $n_B L$ and may therefore be viewed as the concatenation of $2n_B$ subsequences each of length $L/2$.

A binary sequence V'' shall be obtained from V'' by permuting the order of sequences of V'' according to the inverse of the permutation P_i .

A binary sequence V^0 of length $n_B L$ shall be obtained from V'' by bit-wise interleaving pairs of subsequences from V'' . More specifically,

$$V^0_{jL+j'} = V''_{jL+(L/2)(j' \bmod 2) + \lfloor j'/2 \rfloor}$$

where $j = 0, 1, \dots, n_B - 1$ and $j' = 0, 1, \dots, L - 1$.

The LDPC output vector V^{out} of length $n = Ln_B - s_B L - (k_B L - k)$ shall be obtained as a sequence of V^0 by deleting the zero padding and state variables from V^0 . The permutation P_i are such that the zero-padded bits appear contiguously in V^0 in positions $k-L$ to $(k_B - 1)L - 1$ and the state variables appear as the first $s_B L$ bits. Note that $s_B = 3$ in all cases.

Thus

$$V^{\text{out}}_i = V^0_{i+3L} \text{ for } 0 \leq i < k-4L$$

$$V^{\text{out}}_i = V^0_{i+3L+z_p} \text{ for } k-4L \leq i < (n_B-3)L - z_p$$

Table 5.1.7.3.4.2-1 Permutation patterns for the construction of dual diagonal structure.

i	K_i	Q_i	P_i
0	12	11 5 3 1 7 9 10 4 2 0 6 8	0 1 2 3 4 5 6 7 8 9 10 12 13 15 21 19 17 41 11 14 20 18 16 40 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65
1	16	7 1 13 15 11 5 9 3 6 0 12 14 10 4 8 2	0 1 2 3 4 5 6 7 8 9 10 11 12 14 15 17 23 51 21 49 19 43 13 16 22 50 20 48 18 42 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 44 45 46 47 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75
2	14	11 3 1 13 9 5 7 10 2 12 8 4 6	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 16 17 19 25 23 59 21 41 15 18 24 22 58 20 40 26 27 28 29 30 31 32 33 34 35 36 37 38 39 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85
3	18	17 15 13 9 11 5 7 1 3	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 18 19 21 53 27 51 25 47 23 43 17 20 52 26 50 24 46 22 42 28 29 30 31 32 33 34 35 36 37 38 39

		16 14 12 8 10 4 6 0 2	40 41 44 45 48 49 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95
4	18	5 17 15 11 13 9 3 7 1 4 16 14 10 12 8 2 6 0	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 20 21 23 53 29 49 27 47 25 45 19 22 52 28 48 26 46 24 44 30 31 32 33 34 35 36 37 38 39 40 41 42 43 50 51 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100 101 102 103
5	18	11 9 17 15 7 13 5 1 3 10 8 16 14 6 12 4 0 2	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 22 23 25 55 31 53 29 51 27 41 21 24 54 30 52 28 50 26 40 32 33 34 35 36 37 38 39 42 43 44 45 46 47 48 49 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100 101 102 103 104 105 106 107 108 109 110 111 112 113 114 115

5.1.7.3.4.3 Encoding

In this section an efficient encoding method is presented according to which the sequence V''' , as defined in Section 5.1.7.3.4.2, is computed from V^l . The method will describe a procedure to generate the sequence V''' . Recall that in Section 5.1.7.3.4.2 it was described how to produce the matrix G''' from the matrix G (which is also G_l). Efficient encoding of V^l to a sequence V''' satisfying $H'''V'''=0$ is described.

The computation of V''' given V^l is particularly simple due to the structure H''' inherits from G''' . The parity check matrix H''' lifted from G''' takes the form

$$H''' = \begin{bmatrix} M_1 & 0 \\ M_2 & I \end{bmatrix}$$

where $M_1 = \begin{bmatrix} A & B & T \\ C & D & E \end{bmatrix}$ is a $(L/2)K_i \times (L/2)(2k_B + K_i)$ matrix where T is lower

triangular, $\begin{bmatrix} B & T \\ D & E \end{bmatrix}$ is invertible and D is $L/2 \times L/2$. The encoding procedure is

composed of two stages. Let $c = (s, p_1, p_2, p_3)$ be a codeword where s denotes systematic part, p_1, p_2 and p_3 are parity parts. In the first stage, a part of codeword p_1, p_2 is obtained using M_1 depending on the systematic information s . In the second

stage, the remaining part of the codeword p_3 is obtained by simple single parity-check coding using $[M_2 \ I]$. The whole procedure for encoding is as follows.

- 1) Obtain $\begin{bmatrix} B & T \\ \phi & 0 \end{bmatrix}$ from Gaussian elimination on $\begin{bmatrix} B & T \\ D & E \end{bmatrix}$, where $\phi = ET^{-1}B + D = I$.
- 2) compute As^T and Cs^T .
- 3) compute $y = T^{-1}As^T$
- 4) compute $p_1^T = Ey + Cs^T$.
- 5) compute p_2^T using $p_2^T = T^{-1}(As^T + Bp_1^T)$
- 6) compute p_3^T by single parity-check coding using $[M_2 \ I]$.

A sequence V''' satisfying $H'''V'''=0$ is obtained from 1 to 6. The sequence V^{out} is then obtained from V''' by permutation as described in Section 5.1.7.3.4.2.

5.1.7.3.4.4 Truncation

For the Forward Data Channel and Forward Superposed Data Channel packets, the truncation operation shall be carried out as described below:

- The truncation operation depends on the packet size $N_{\text{PACKET_BITS}}$ of the packet received from the FTC MAC Protocol, and the variables $\text{MaxRateOneFifthPacketSize}$, $\text{MaxRateOneThirdPacketSize}$ and $\text{MaxRateOneHalfPacketSize}$. $\text{MaxRateOneFifthPacketSize}$ is equal to one of the parameters $\text{MaxRateOneFifthPacketSizeEightInterlaceLDPC}$ or $\text{MaxRateOneFifthPacketSizeSixInterlaceLDPC}$, depending on whether the Forward Data Channel packet or the Forward Superposed Data Channel is transmitted using an eight interlace HARQ structure or a six interlace HARQ structure. $\text{MaxRateOneThirdPacketSize}$ is equal to one of the parameters $\text{MaxRateOneThirdPacketSizeEightInterlaceLDPC}$ or $\text{MaxRateOneThirdPacketSizeSixInterlaceLDPC}$, depending on whether the Forward Data Channel packet or the Forward Superposed Data Channel is transmitted using an eight interlace HARQ structure or a six interlace HARQ structure. $\text{MaxRateOneHalfPacketSize}$ is equal to one of the parameters $\text{MaxRateOneHalfPacketSizeEightInterlaceLDPC}$ or $\text{MaxRateOneHalfPacketSizeSixInterlaceLDPC}$, depending on whether the Forward Data Channel packet or the Forward Superposed Data Channel is

transmitted using an eight interlace HARQ structure or a six interlace HARQ structure. The FTC MAC Protocol determines which HARQ interlacing structure is used for transmitting the Forward Data Channel packet or the Forward Superposed Data Channel.

- MaxRateOneFifthPacketSizeEightInterlaceLDPC,
MaxRateOneFifthPacketSizeSixInterlaceLDPC,
MaxRateOneThirdPacketSizeEightInterlaceLDPC,
MaxRateOneThirdPacketSizeSixInterlaceLDPC,
MaxRateOneHalfPacketSizeEightInterlaceLDPC, and
MaxRateOneHalfPacketSizeSixInterlaceLDPC are configuration attributes of the FTC MAC protocol.
- If $N_{\text{PACKET_BITS}} < \text{MaxRateOneFifthPacketSize}$, the sequence V^{out} is not truncated.
- If $\text{MaxRateOneFifthPacketSize} \leq N_{\text{PACKET_BITS}} < \text{MaxRateOneThirdPacketSize}$, the sequence V^{out} is truncated to length $3k$, i.e., all elements with indices greater than or equal to $3k$ are deleted.
- If $\text{MaxRateOneThirdPacketSize} \leq N_{\text{PACKET_BITS}} < \text{MaxRateOneHalfPacketSize}$, the sequence V^{out} is truncated to length $2L$, i.e., all elements with indices greater than or equal to $2L$ are deleted.
- If $\text{MaxRateOneHalfPacketSize} \leq N_{\text{PACKET_BITS}}$, the sequence V^{out} is truncated to length $3k/2$, i.e., all elements with indices greater than or equal to $2k$ are deleted.
- The output of the truncation operation is denoted by V^{tr} .

For Reverse OFDMA Data Channel packets, the output V^{tr} of the truncation operation shall be equal to V^{out} .

The sequence V^{tr} is a truncated version of the permuted sequence V^{out} and is in the order of bit transmission, i.e., the different HARQ transmissions of this packet are generated in the order specified by V^{tr} .

(2) right-top of G_1

	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37
	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
1	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
2	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•
3	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
4	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
5	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•
6	•	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•
7	•	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•
8	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
9	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
10	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
11	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
12	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
13	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
14	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•

(3) left-bottom of G_1

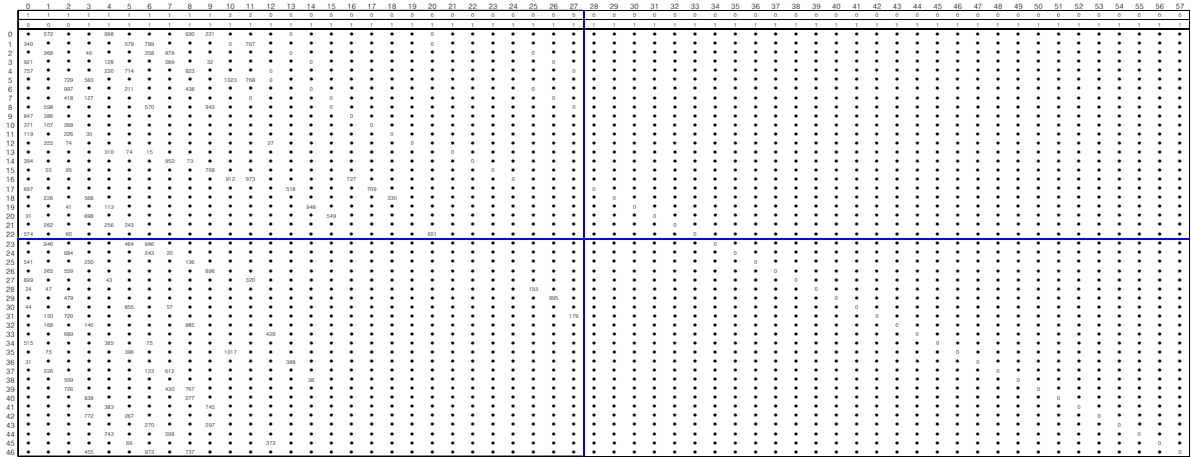
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
	1	1	1	1	1	1	2	2	0	0	0	0	0	0	0	0	0	0	0
	0	0	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
15	•	428	105	•	•	•	929	•	•	•	•	•	•	•	•	•	•	•	•
16	30	264	•	•	•	•	•	•	•	832	•	•	•	•	•	•	•	•	•
17	514	•	410	•	•	•	•	•	•	•	978	•	•	•	•	•	•	•	•
18	•	487	249	•	204	•	•	•	•	•	•	•	•	•	•	•	•	•	•
19	526	126	•	•	•	•	•	•	•	•	•	906	•	•	•	•	•	•	•
20	10	•	90	•	•	889	•	•	•	•	•	•	•	•	•	•	•	•	•
21	•	126	•	•	•	•	•	714	•	•	•	•	•	•	•	•	•	•	•
22	•	•	312	•	•	•	•	•	624	•	•	•	•	•	•	•	•	•	•
23	954	•	•	302	•	•	63	•	•	•	•	•	•	•	•	•	•	•	•
24	33	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
25	524	•	•	752	227	•	•	•	•	•	•	•	•	•	•	•	•	•	•
26	•	647	•	•	•	346	•	•	•	•	•	•	•	•	•	•	•	•	•
27	•	•	918	•	•	•	•	•	•	602	•	•	•	•	•	•	•	•	•
28	14	•	•	131	816	•	•	•	•	•	•	•	•	•	•	•	•	•	•
29	•	216	•	•	•	•	•	103	•	•	•	•	•	•	•	•	•	•	•
30	•	•	893	•	•	•	•	•	•	•	•	771	•	•	•	•	•	•	•

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21
0	0	1	1	1	1	1	1	2	2	0	0	1	1	1	1	1	1	1	1	1	1	1
1	1	1	1	1	1	1	1	2	2	0	0	1	1	1	1	1	1	1	1	1	1	1
2	●	495	●	●	720	413	●	●	●	0	0	●	●	●	●	●	●	●	●	0	0	●
3	63	●	●	●	●	●	163	0	767	●	●	●	●	●	●	●	●	●	●	●	0	●
4	●	629	●	319	818	●	●	●	●	0	●	●	●	●	●	●	●	●	●	●	●	●
5	●	●	247	●	●	●	412	1023	768	0	●	●	●	●	●	●	●	●	●	●	●	●
6	●	●	928	●	●	●	●	0	0	●	0	●	●	●	●	●	●	●	●	●	●	●
7	32	190	●	●	●	●	●	●	●	●	●	●	●	0	●	●	●	●	●	●	●	●
8	243	●	596	●	●	●	●	●	●	●	●	●	●	0	●	●	●	●	●	●	●	●
9	●	880	833	329	●	●	●	●	●	●	●	●	●	●	0	●	●	●	●	●	●	●
10	224	840	●	●	208	●	●	●	●	●	●	●	●	●	●	0	●	●	●	●	●	●
11	●	●	479	●	●	222	●	●	●	17	●	●	●	●	●	●	●	0	●	●	●	●
12	296	856	●	●	●	●	651	●	●	●	●	●	●	●	●	●	●	●	0	●	●	●
13	●	●	926	211	●	●	●	●	167	●	●	●	●	●	●	●	●	●	●	0	●	●
14	764	●	●	●	●	●	●	166	●	●	●	●	387	●	●	●	●	●	●	●	●	0
15	●	238	925	●	●	●	●	●	●	405	●	●	●	●	●	●	●	●	●	●	●	●
16	●	●	●	850	922	●	●	●	●	●	●	●	●	852	●	●	●	●	●	●	●	●

(1) left-top of G_2

G_5 : $k_B = 11$, $n_B = 38$, $s_B = 3$, $L_{\max} = 1024$.

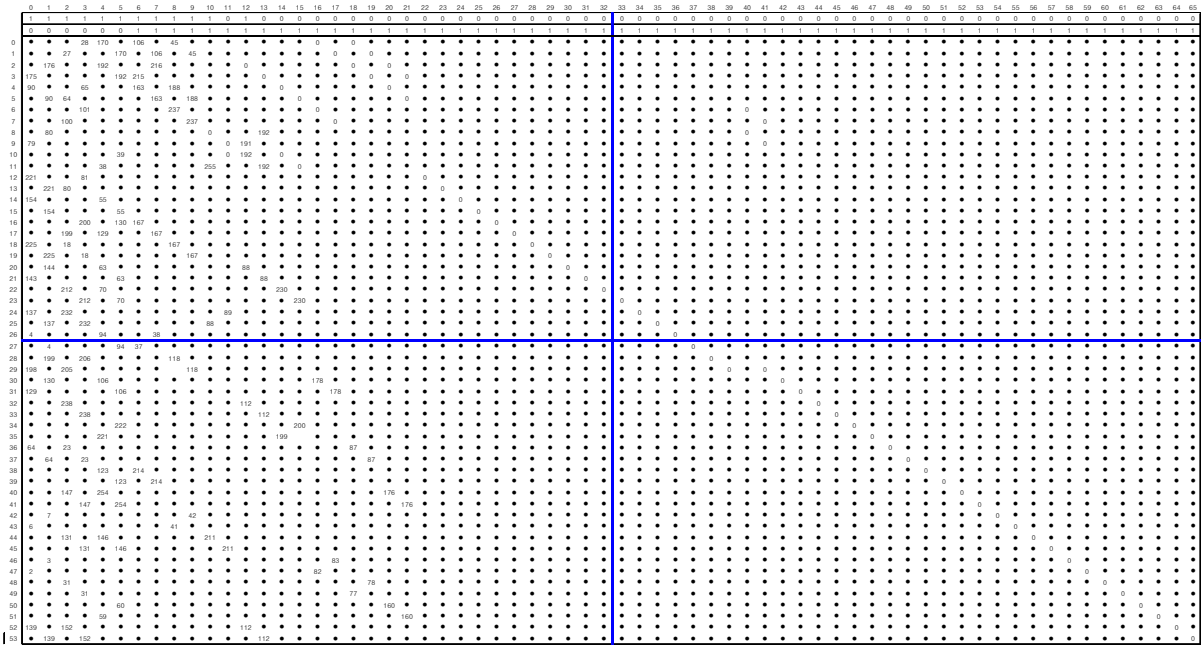
G_5 consists of 4 parts: (1) left-top, (2) right-top, (3) left-bottom, and (4) right-bottom



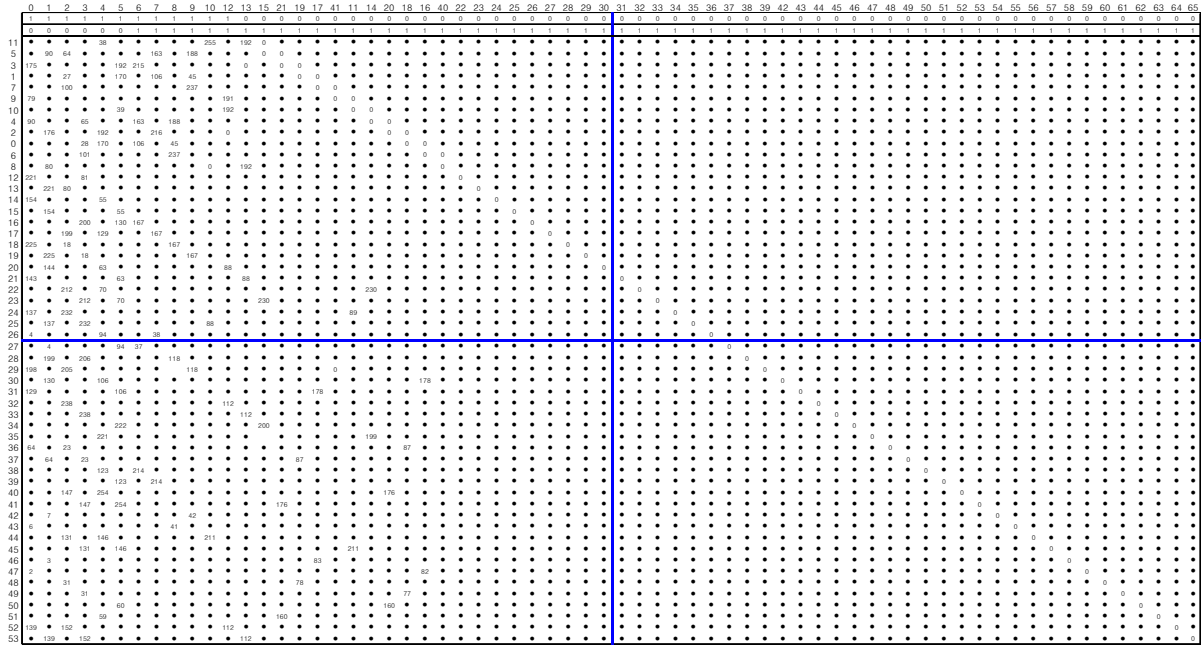
(1) left-top of G_5

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	
0	•	572	•	•	998	•	•	930	221	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
1	340	•	•	•	•	579	789	•	•	0	767	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
2	•	368	•	46	•	•	358	978	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
3	921	•	•	•	128	•	•	589	•	32	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
4	757	•	•	•	•	230	714	•	•	823	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
5	•	•	729	583	•	•	•	•	•	1023	768	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
6	•	•	997	•	•	211	•	•	438	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
7	•	•	418	127	•	•	•	•	•	•	0	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	
8	•	598	•	•	•	•	570	•	943	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•	
9	847	386	•	•	•	•	•	•	•	•	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•	•
10	371	107	268	•	•	•	•	•	•	•	•	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•	•
11	119	•	226	30	•	•	•	•	•	•	•	•	•	•	•	•	•	0	•	•	•	•	•	•	•	•	•	•	•
12	•	355	74	•	•	•	•	•	•	•	•	27	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
13	•	•	•	•	•	310	74	15	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
14	394	•	•	•	•	•	•	952	73	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
15	•	55	95	•	•	•	•	•	768	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
16	•	•	•	•	•	•	•	•	•	912	973	•	•	•	•	•	727	•	•	•	•	•	•	•	•	•	•	•	•
17	697	•	•	•	•	•	•	•	•	•	•	•	518	•	•	•	709	•	•	•	•	•	•	•	•	•	•	•	•
18	•	226	•	568	•	•	•	•	•	•	•	•	•	•	•	•	•	330	•	•	•	•	•	•	•	•	•	•	•
19	•	•	41	•	113	•	•	•	•	•	•	•	•	948	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
20	31	•	•	698	•	•	•	•	•	•	•	•	•	•	549	•	•	•	•	•	•	•	•	•	•	•	•	•	•
21	•	262	•	•	256	243	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•
22	574	•	60	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	651	•	•	•	•	•	•	•	•

G_0 consists of 4 parts: (1) left-top, (2) right-top, (3) left-bottom, and (4) right-bottom.



G_0''' consists of 4 parts: (1) left-top, (2) right-top, (3) left-bottom, and (4) right-bottom



5.1.7.4 Channel Interleaving

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Channel interleaving follows the convolutional or turbo encoding, and consists of a bit-demultiplexing operation followed by a bit permuting operation. Channel interleaving shall not be performed after LDPC encoding, i.e., in this case the encoding operation shall be followed by the sequence repetition operation.

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References

- [1] IEEE P802.20/D0.1m, April 11 2007.
- [2] 'Physical Layer for Ultra Mobile Broadband (UMB) Air Interface Specifications, Version 2.0', C30-20070618-018R2 C. P0084 v1.098, 3GPP2, June 18, 2007.
- [3] C.S0084-0-002, MAC Layer for Ultra Mobile Broadband (UMB) Air Interface Specification, April 2007.