ENCODING/DECODING SCHEME TO ACHIEVE STANDARD COMPLIANT MEAN TIME TO FALSE PACKET ACCEPTANCE



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PROBABILITY OF ERROR FOR HARD DECISION DECODING



- Most hard-decision (HD) decoding of block is bounded distance decoding
 - BCH and RS codes
 - Berlekamp-Massey decoding algorithm
- Two kinds outputs for HD decoding
 - 1) Decoding failure
 - 2) A valid codeword.
- Output 2) can be a wrong codeword, i.e. not the transmitted codeword
- Two kinds of error rates for HD decoding
 - a) P_f : probability of decoding failure (or uncorrectable)
 - b) P_m : probability of miscorrection (or undetectable)
- The final probability of error is

 $P_e = P_f + P_m$ (usually $P_f >> P_m$)

PROBABILITY OF ERROR FOR SOFT DECISION DECODING



Optimal decoding: Maximum Likelihood Decoding (MLD)

- Always output a codeword
 - The codeword that maximizes the conditional probability of a codeword given a received word
- No decoding failure for MLD
- Probability of error is the probability of miscorrection $P_e = P_m$

Iterative decoding of LDPC codes

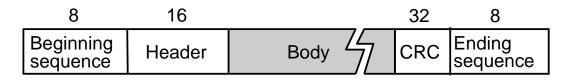
- Add a syndrome check for power saving in every iteration
- Always output a word
 - The decoder will stop iterating if the estimated codeword passes the syndrome check
 - Otherwise it will keep iterating until the maximum number of iterations is reached and output a word
 - The information bits of the output word can be correct even though they may not pass the syndrome check due to errors in the (less protected) parity bits only
 - Therefore the LDPC decoder does not declare a decoding failure in this case

Probability of bit error is the probability of miscorrection on information bits

$$P_e = P_{m,info}$$

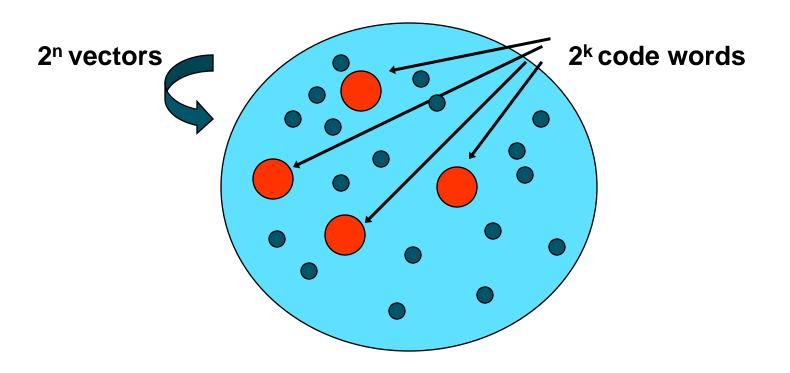
PACKET ERROR DETECTION WITH CYCLIC REDUNDANCY CHECK (CRC)

 Add n-k bits of extra data (the CRC field) to an k-bit message to provide error detection function (i.e. an (n,k) binary cyclic code)

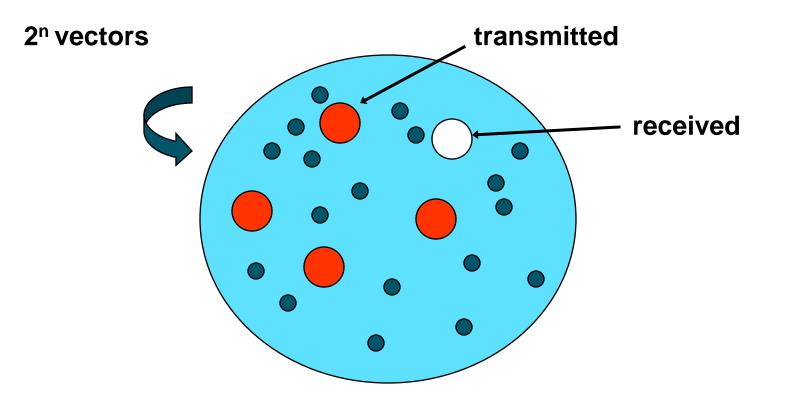


- For efficiency, n-k << n</p>
 - e.g., n-k = 32 for Ethernet and k = 12,000 (1500 bytes)

Replace k information bits by a unique n bit code word

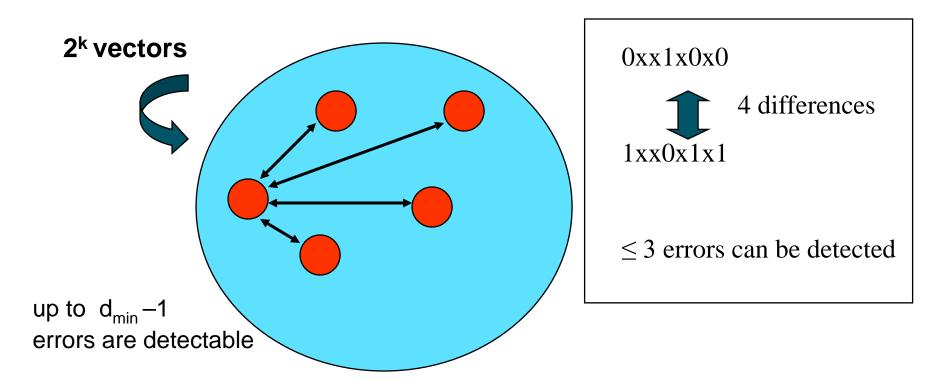


Received vector not equal to one of the 2^k code words



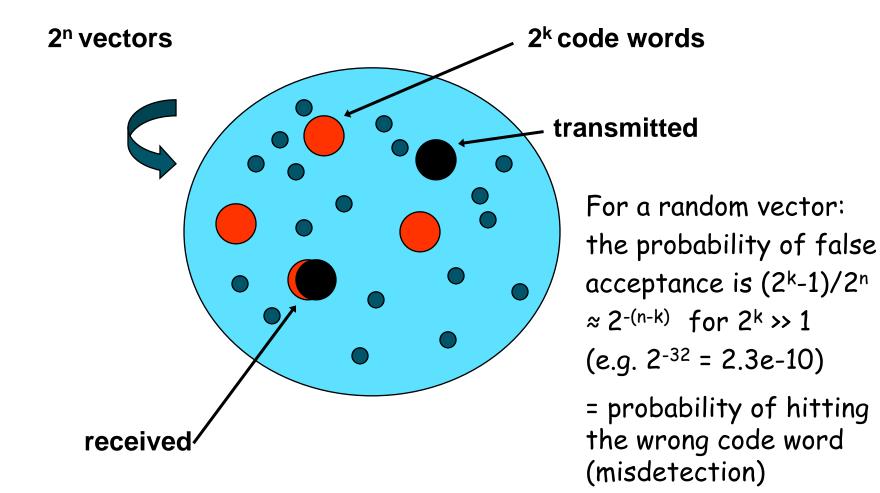


code words differ in at least d_{min} positions



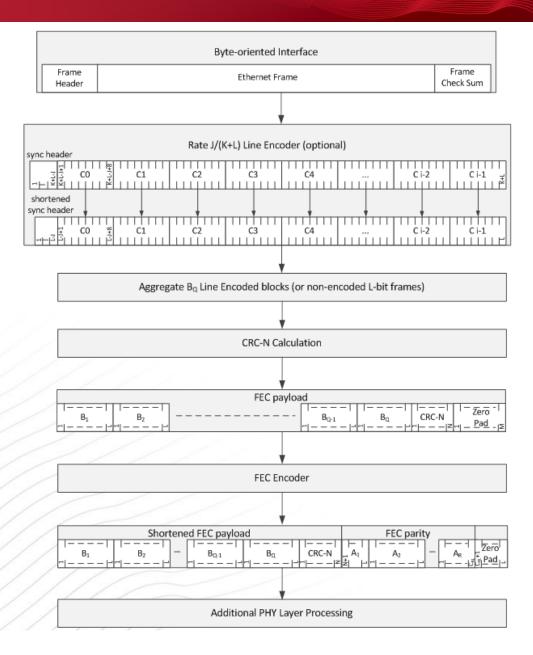
MISDETECTION





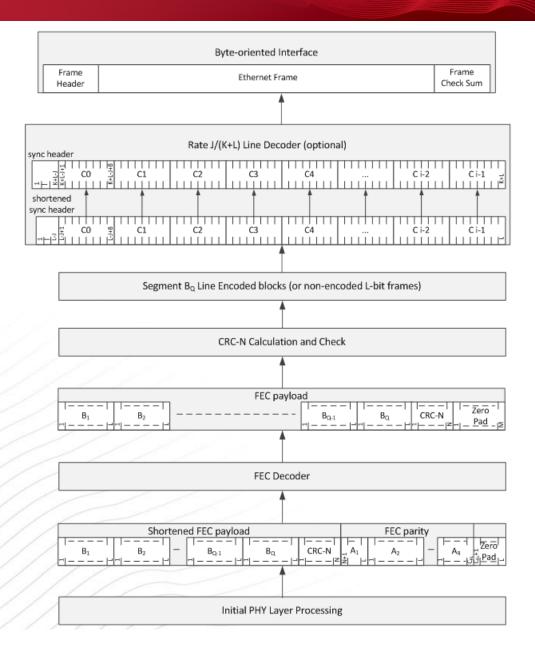
GENERAL PHY ENCODING SCHEME





GENERAL PHY DECODING SCHEME







 For the number of bits N in the CRC plus the number of bits in the FCS (which is 32 for the Ethernet CRC-32 FCS):

$$FPAR = FLR * F / 2^{(N+32)}$$

where F is the number of minimum size Ethernet frames including header and FCS per FEC codeword information payload given by:

$$F = \frac{FEC \ Payload}{[8*(64+H)]} = \left[\frac{\lfloor k/L \rfloor}{(64+H)}*\frac{L}{8}\right] + \begin{cases} 0, & k \ mod \ L = 0\\ 1, & k \ mod \ L = 1\\ 2, & k \ mod \ L \ge 2 \end{cases}$$

and FLR is the Frame Loss Ratio (= Frame Error Ratio).

 This assumes the worst case where codeword information bit errors corrupt every Ethernet frame in the (n, k) FEC codeword payload.



- The length N of the CRC is chosen to be long enough to insure the MTTFPA is achieved.
- The MTTFPA is the inverse of the False Packet Acceptance Ratio (FPAR) times the Ethernet frame bit rate R:

MTTFPA = 1/(FPAR * R)

- The minimum MTTFPA occurs at the maximum packet rate since they are inversely related.
- For a fixed bit rate B over the PHY layer, the packet rate is maximized for the smallest size Ethernet frames (64 bytes plus any additional header):

$$a = \frac{B}{8 * (64 + H + IFG)}$$

 where the number of header bytes H plus the number of bytes for the inter-frame gap IFG are added to the minimum Ethernet payload size of 64 bytes.



 Therefore the MTTFPA which is the inverse of the FPAR times the Ethernet frame rate R is given by:

$$MTTFPA = 1/(FPAR * R) = 1/(FLR * F * 2^{-(N+32)} * \frac{B}{8 * (64 + H + IFG)})$$

 So the minimum value N for a CRC to achieve a desired MTTFPA can be calculated solving for N above:

$$N \ge \frac{\log[MTTFPA * FLR * F * \frac{B}{8 * (64 + H + IFG)} * 2^{-32}]}{\log[2]}$$

The preferred value of the MTTFPA in IEEE 802.3 standards is the age of the universe which is estimated to be 14 billion years (4.4 x 10¹⁷ seconds).



The number of minimum size Ethernet frames including header and FCS per long FEC codeword information payload given by:

$$F = \left[\frac{\lfloor 14400/65 \rfloor}{(64+8)} * \frac{65}{8}\right] + 2 = 26$$

• For an FLR of 1e-6 in the downstream at a 10 Gbps bit rate,

$$N \ge \frac{\log \left[4.4 \times 10^{17} * 10^{-6} * 26 * \frac{10 \times 10^9}{8 * (64 + 8 + 12)} * 2^{-32} \right]}{\log[2]} = 35.2$$

- A minimum 36 bit CRC is needed to achieve the required MTTFPA
- Choose a standard 40 bit CRC used in GSM control channels (with a FIRE code for FEC)

• CRC generator polynomial $x^{40} + x^{26} + x^{23} + x^{17} + x^3 + 1$

EPOC MEAN TIME TO FALSE PACKET ACCEPTANCE (DOWNSTREAM)

For a 40 bit CRC plus the 32 bits in the FCS, the Mean Time To False Packet Acceptance is:

$$MTTFPA = 1/\left(10^{-6} * 26 * 2^{-(40+32)} * \frac{10 \times 10^9}{8 * (64+8+12)}\right) = 1.22 \times 10^{19}$$

 This is greater than the required age of the universe MTTFPA of 4.4 x 10¹⁷ seconds.

EPOC MEAN TIME TO FALSE PACKET ACCEPTANCE (UPSTREAM)

 For the same 40 bit CRC plus the 32 bits in the FCS, the Mean Time To False Packet Acceptance is:

$$MTTFPA = 1/\left(5 \ge 10^{-5} \le 26 \le 2^{-(40+32)} \le \frac{10 \ge 10^9}{8 \le (64+8+12)}\right) = 2.44 \ge 10^{17}$$

- This is less than the required age of the universe MTTFPA of 4.4 x 10¹⁷ seconds.
- Approaches to increase the upstream MTTFPA
 - Use a longer CRC
 - Reduce the Frame Loss Ratio
 - Reduce the maximum PHY layer bit rate
- Best approach is to reduce the upstream bit rate

EPOC MEAN TIME TO FALSE PACKET ACCEPTANCE (UPSTREAM) - CONTINUED

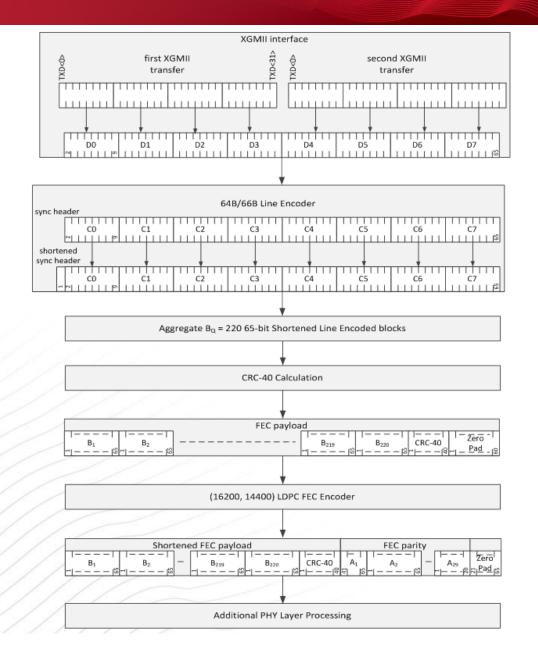
- Reduce the upstream maximum PHY layer bit rate to 5 Gbps
- For the same 40 bit CRC plus the 32 bits in the FCS, the Mean Time To False Packet Acceptance is:

$$MTTFPA = 1/\left(5 \ge 10^{-5} \le 26 \le 2^{-(40+32)} \le \frac{5 \ge 10^9}{8 \le (64+8+12)}\right) = 4.88 \ge 10^{17}$$

 This is now greater than the required age of the universe MTTFPA of 4.4 x 10¹⁷ seconds.

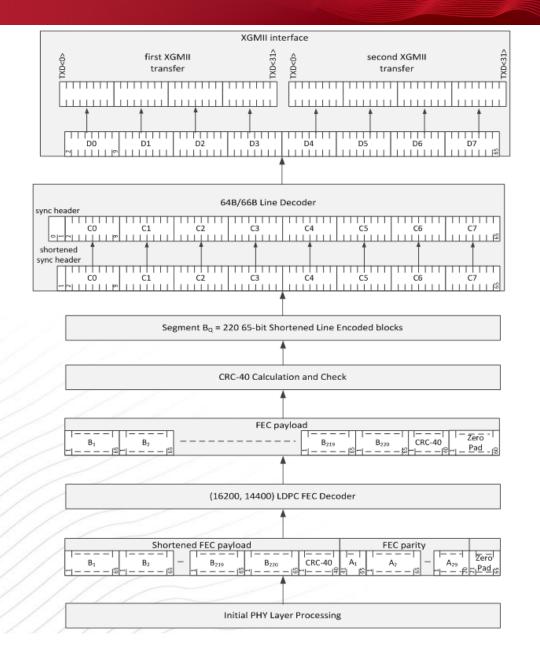
EPOC PHY ENCODING SCHEME





EPOC PHY DECODING SCHEME







• For each FEC Payload received (or equivalent) (Similar to 74.7.4.5 FEC decoder)

- If Information CRC check is good
 - For <u>each</u> block contained in B_Q:
 - Set Bit 0 of expanded 64b/66b Sync Header to inverse of 65b Bit 0 (XOR)
 - Append remaining 65 bits from block
- If Information CRC check is bad
 - For <u>every</u> block contained in B_Q:
 - Set Bit 0 and Bit 1 of expanded 64b/66b Sync Header to '11' to indicate invalid
 - Discard Bit 0 of 65b block
 - Append remaining 64 bits of 65b block

Higher layer 64b/66b decoding will process as normal

 A "invalid" Sync Header gets flagged as "sh_invalid" and processed/counted accordingly to ensure that detected errors are signaled to the MAC for every frame containing an error.



- An encoding and decoding scheme to achieve standard compliant Mean Time to False Packet Acceptance has been presented
- Analysis of this approach utilizing a 40 bit CRC exceeds the required MTTFPA equal to the age of the universe
- The use of a standardized 40 bit CRC used in the GSM standard is proposed for this scheme
- The maximum bit rates for the PHY layer depend on the Ethernet Frame Loss Ratios
 - 10 Gbps downstream with a 1e-6 Frame Loss Ratio
 - 5 Gbps upstream with a 5e-5 Frame Loss Ratio



Move to:

Incorporate the MTTFPA encoding/decoding scheme as presented in prodan_3bn_02_0913.pdf with the adopted FEC codes in the downstream and upstream EPoC PHY layer.

Moved:

Second:

Technical decision, 75% or greater Yes: No: Abstain:



Thank You