RS(544,514) FEC performance

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Introduction

The IEEE 802.3bs Task Force has adopted RS(544,514) FEC with interleaving of FEC symbols from two FEC codewords to give good burst error tolerance.

Concerns over the latency of this scheme has led to proposals for either non-multiplexed or symbol multiplexed FEC schemes for 50 Gb/s and next generation 100 Gb/s Ethernet.

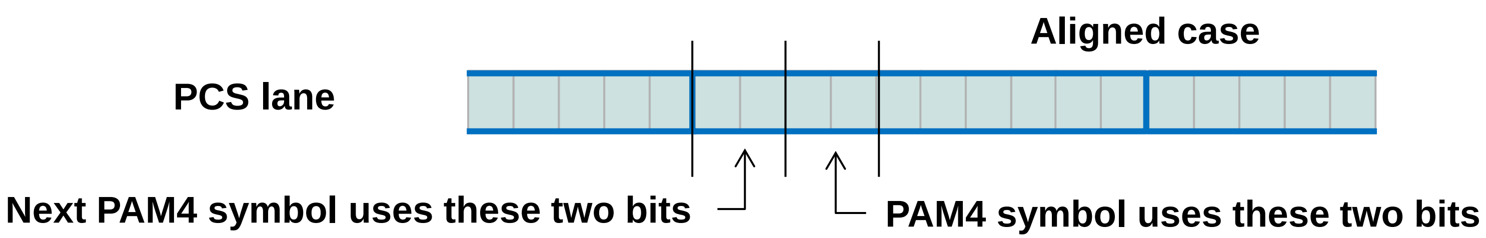
This presentation analyses the performance of such schemes using a development of the principles explained for the NRZ case in Annex 1 of anslow_3bs_02_1114.

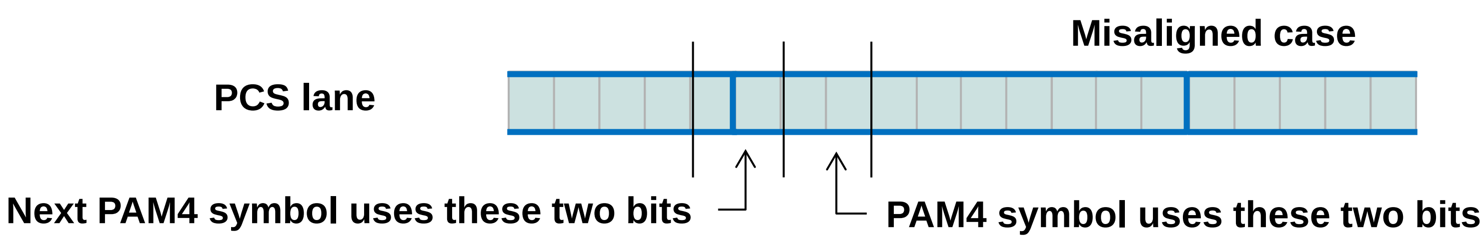
Signal structure

Assuming:

- A single PCS lane or multiple PCS lanes formed by round robin distribution of FEC symbols to the PCS lanes
- RS(544,514) FEC (which has 10-bit symbols)
- Gray coding (see P802.3bs D1.3 120.5.7)

There are two ways that the PAM4 coding can occur:

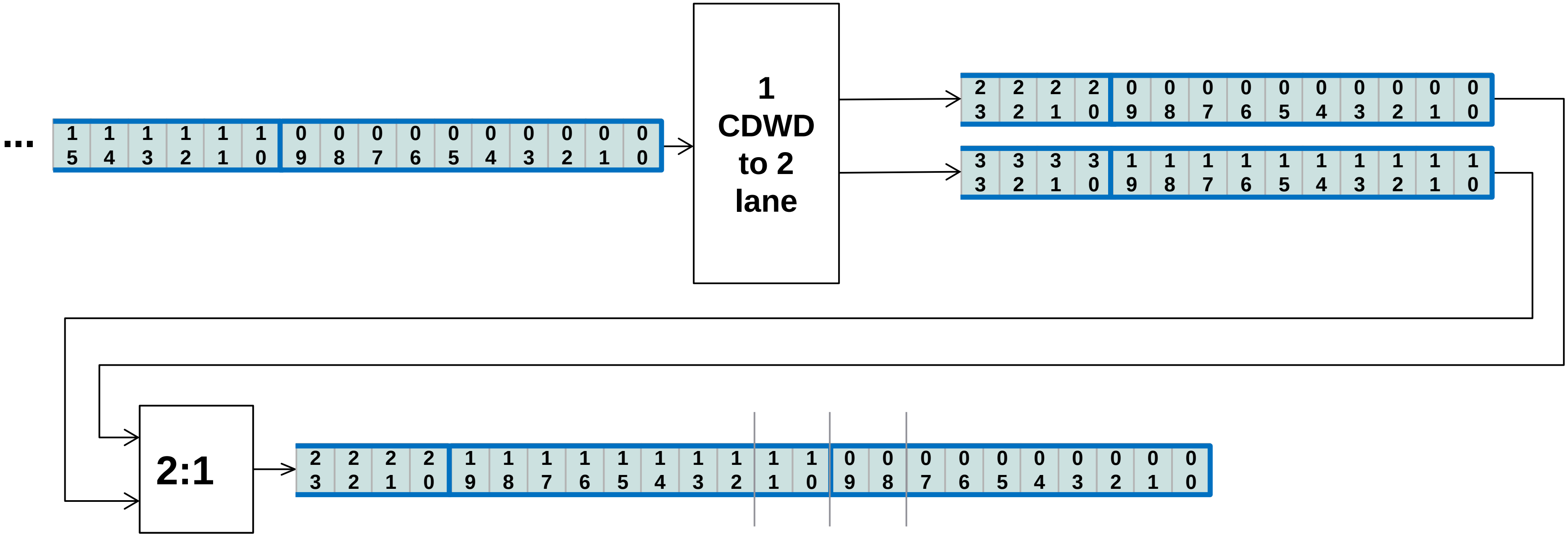




The analysis in anslow_3bs_03_0515 showed no difference in performance between the two cases.

Symbol multiplexing

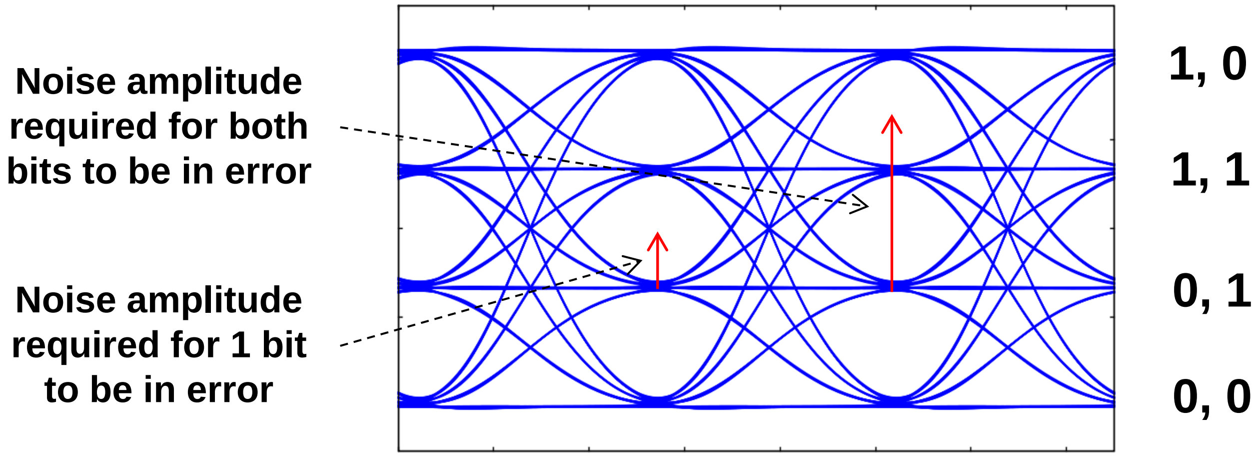
Round robin distribution of FEC symbols to the PCS lanes. Symbol multiplexing in the PMA.



The 2:1 PMA must find FEC symbol boundaries. If not totally de-skewed, the symbol order may be changed, but performance is the same as without multiplexing.

Gray coding

Assume the use of Gray coding (see P802.3bs D1.3 120.5.7) as illustrated below:



If noise causes any of the 4 levels to be mistaken for an adjacent level, this causes one of the two bits to be in error.

If there is just enough Gaussian noise to cause a BER of 3.8E-4* due to single level errors, then the probability of that noise causing both bits to be in error is 2E-24.

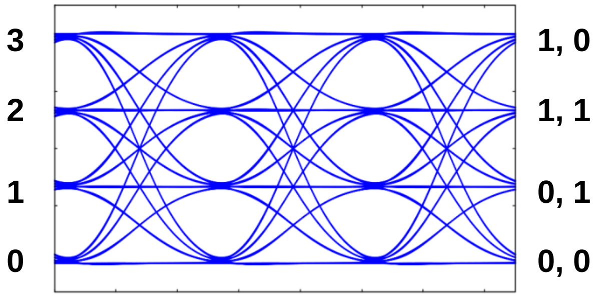
This analysis therefore assumes that only one of the two bits is in error.

* FLR = 6.2E-10 (equivalent to BER = 1E-12 with random errors) after RS(544,514) FEC

Burst error model 1

The NRZ burst analysis in anslow_3bs_02_1114 page 12 assumed that if a **bit** is in error, the worst case probability that the next **bit** is also in error is 0.5. If we assume for Gray coded PAM4 that an error in a particular symbol only causes the decision on the next symbol to move up or down one level, then the possibilities are:

Correct level	Received level		Error pattern	
	One up	One down	One up	One down
3	3	2	✓, ✓	✓, ✗
2	3	1	✓, ✗	✗, ✓
1	2	0	✗, ✓	✓, ✗
0	1	0	✓, ✗	✓, ✓



Since two of the eight possibilities result in both bits being correct, these states terminate the burst. Therefore for Gray coded PAM4, if a **symbol** is in error, the worst case probability that the next **symbol** is also in error is 0.75.

Burst error model 2

The second aspect of this table is that of the six possibilities giving bits in error, two have errors in the first bit while four have errors in the second bit.

Correct level	Received level		Error pattern	
	One up	One down	One up	One down
3	3	2	✓, ✓	✓, ✗
2	3	1	✓, ✗	✗, ✓
1	2	0	✗, ✓	✓, ✗
0	1	0	✓, ✗	✓, ✓

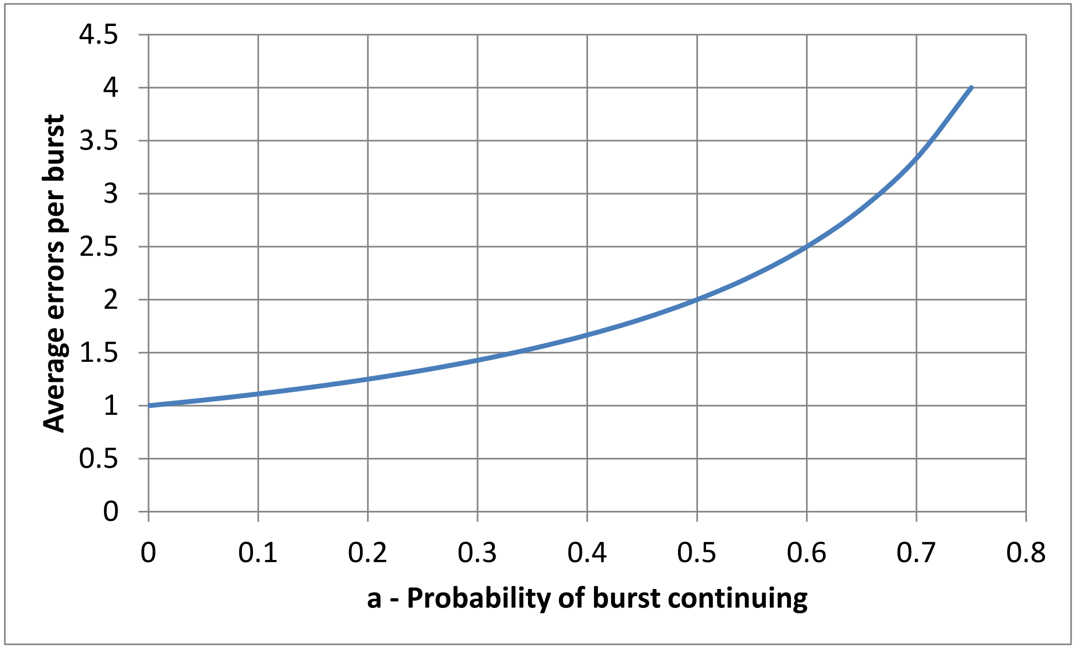
The analysis in the remainder of this contribution therefore assumes that if a given symbol is in error, the probability of a bit error in the first bit is 1/3 and in the second bit is 2/3.

Burst error model 3

The "SNR" shown on the X axis of the following results slides is related to the noise induced input SER via the following equation:

$$SER_{in} = \frac{3}{4} erfc\left(\sqrt{\frac{SNR}{2}}\right) \tag{1}$$

Which does not include the additional errors due to the bursts. The average number of errors in a burst is related to the probability of the burst continuing "a" as shown below:



For a = 0.75, the BER_{in} including bursts is 4 x the BER_{in} due to noise.

Single burst bound

*

As pointed out in anslow_01_0815_logic, for a non-interleaved scheme, a single burst that lasts for ~74 PAM4 symbols has a high probability of causing errors in 16 FEC symbols (which is uncorrectable). With a = 0.75, the probability of a burst this long is 0.75^74 = 5.7E-10. When this is combined with the probability that the codeword has at least one error in it, a simple lower bound for the FLR can be calculated.

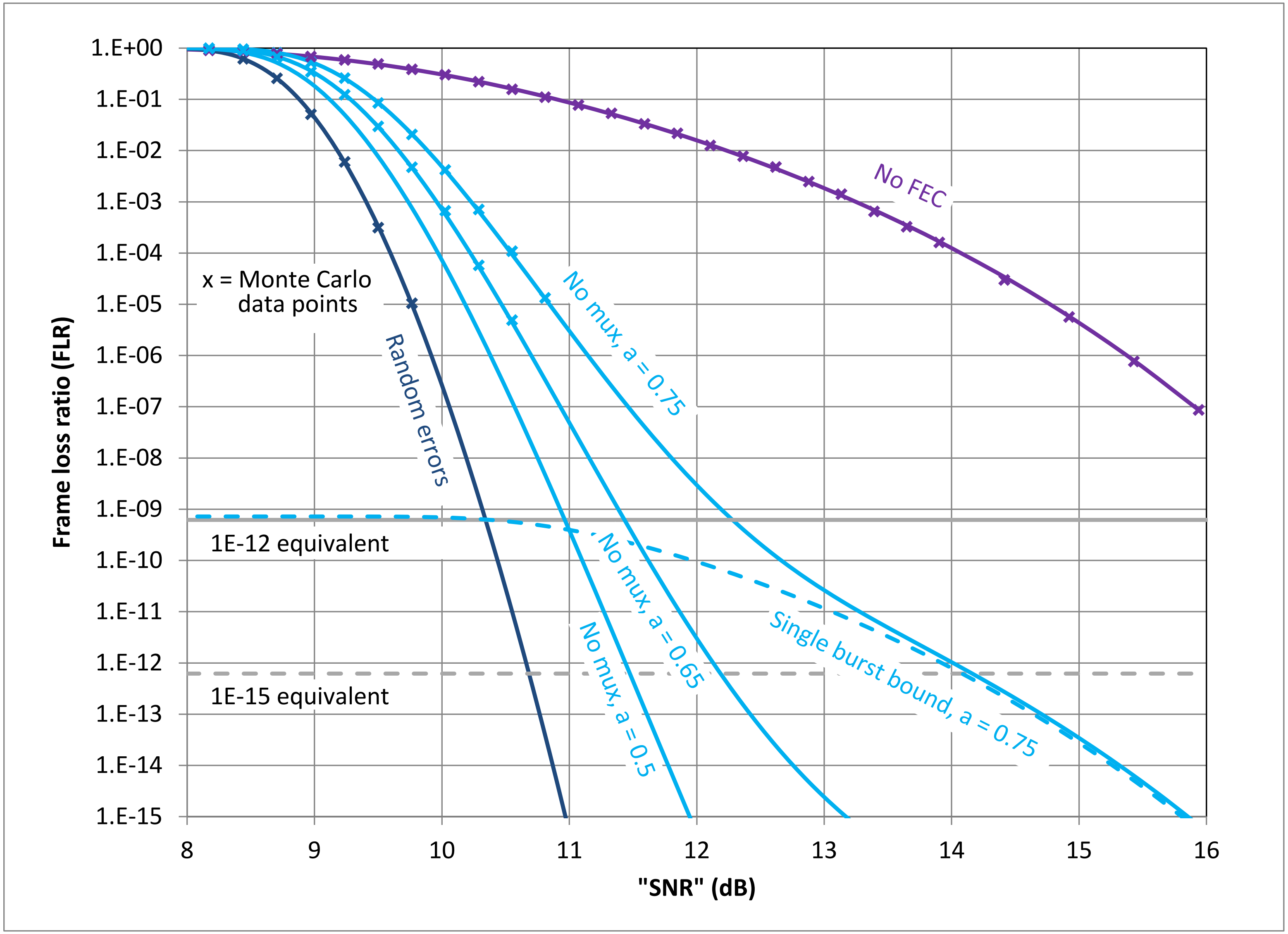
If a is the probability of the burst continuing, a more accurate calculation for the probability that a single burst is uncorrectable is:

$$P_{uncorr} = 1/5*a^{72}*(1-a) + 2/5*a^{73}*(1-a) + 3/5*a^{74}*(1-a) + 4/5*a^{75}*(1-a) + a^{76}*(1-a) + a^{77}*(1-a) + a^{78}*(1-a) + \ldots$$

For a = 0.75, this evaluates to 6.4E-10.

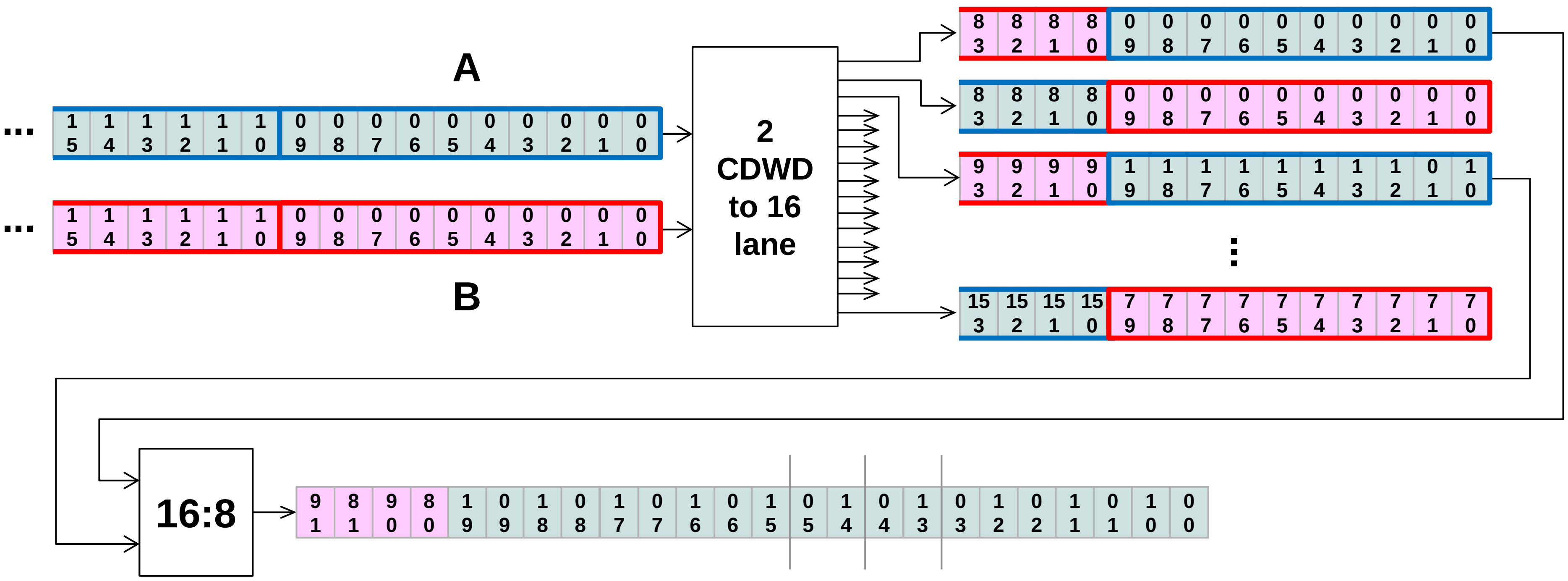
This bound is plotted as a dashed line on the next page.

RS(544,514) no mux or symbol mux



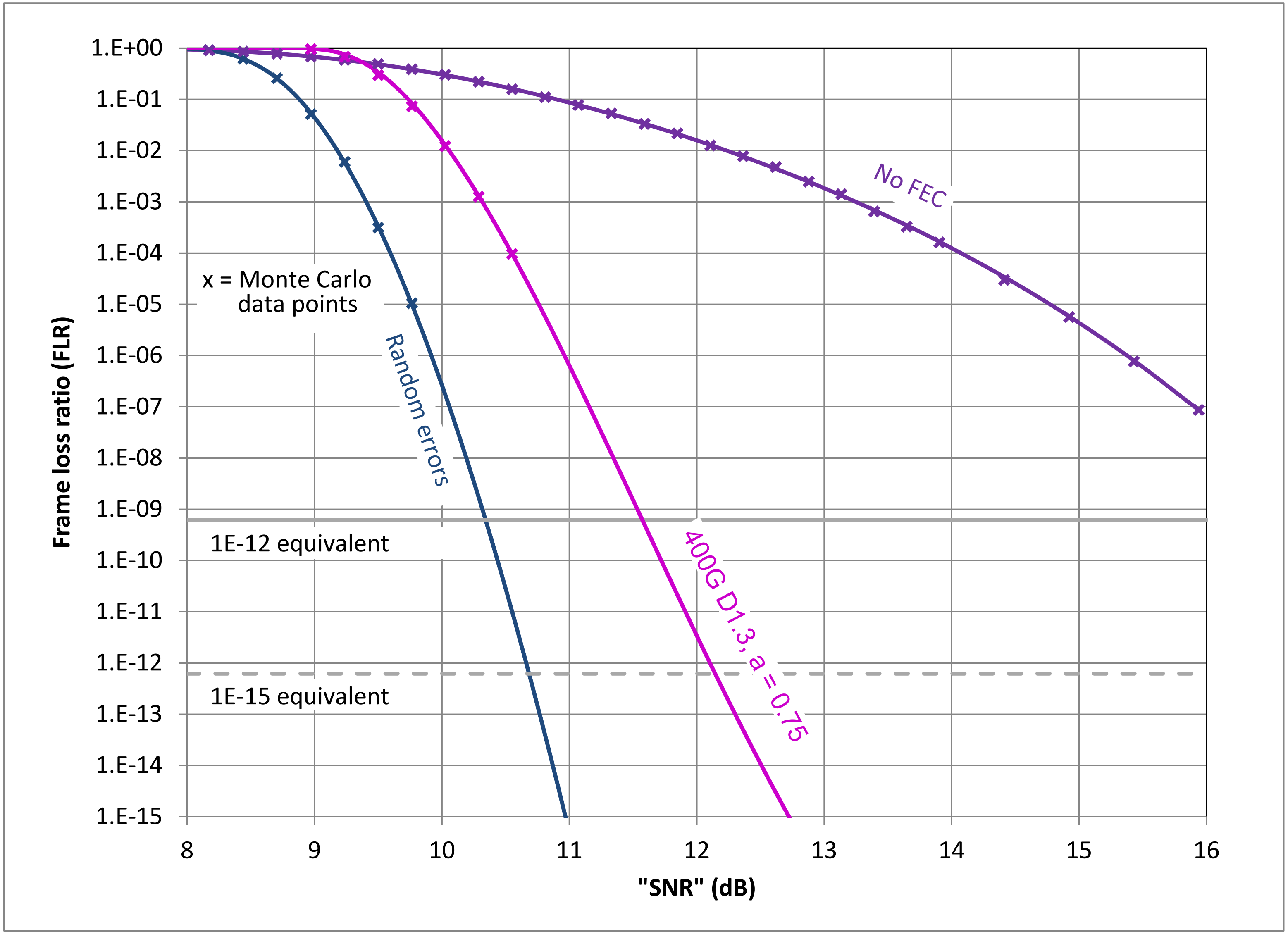
P802.3bs D1.3 scheme

Symbol interleave from 2 FEC codewords. Bit multiplex in the PMA.

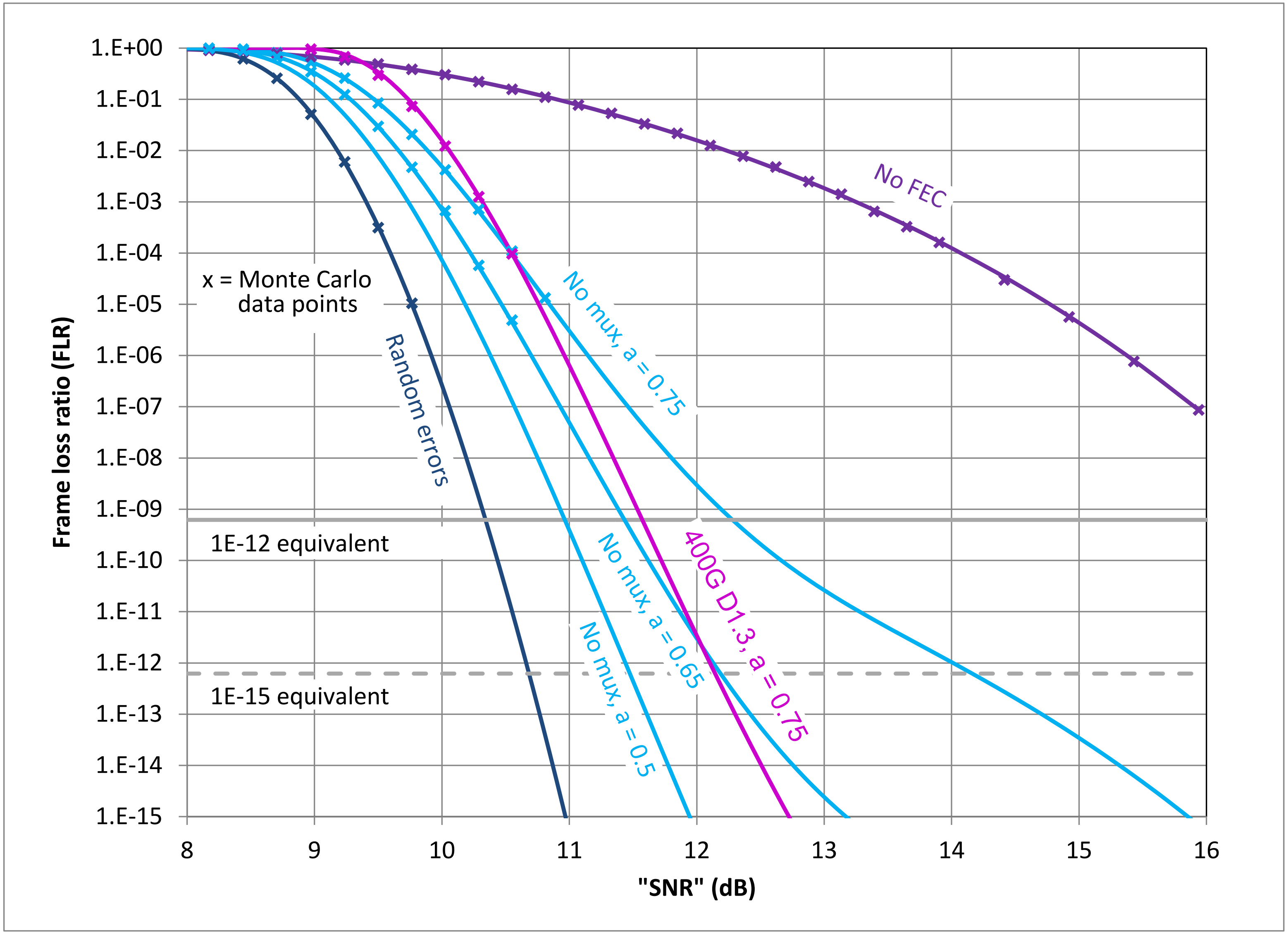


If one codeword is uncorrectable, the other is marked bad also.

P802.3bs D1.3 performance



All curves



1.E+00
1.E-01
1.E-02
1.E-03
1.E-04
1.E-05
1.E-06
1.E-07
1.E-08
1.E-09
1.E-10
1.E-11
1.E-12
1.E-13
1.E-14
1.E-15
Frame loss ratio (FLR)
8
9
10
11
12
13
14
15
16
"SNR" (dB)
x = Monte Carlo
data points
Random errors
No FEC
No mux, a = 0.75
No mux, a = 0.65
No mux, a = 0.5
400G D1.3, a = 0.75
1E-12 equivalent
1E-15 equivalent

Results for RS(544,514) all gain used for PAM4

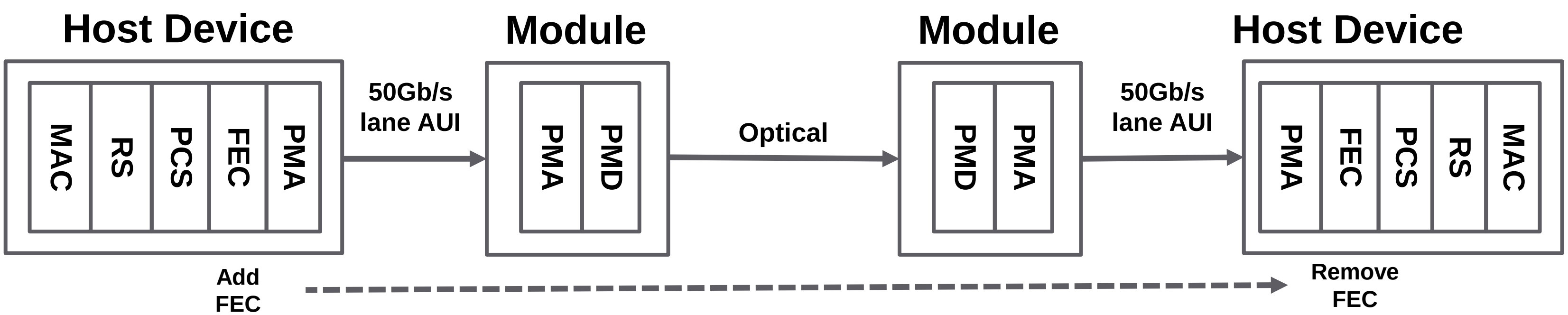
From the curves shown on the previous slide, if all of the coding gain were to be used for the PAM4 link, the BERs at the FEC input required to give FLRs equivalent to that of a BER of 1E-12 and 1E-15 are:

	RS(544,514)	
	FLR = 6.2E-10	FLR = 6.2E-13
No FEC	1E-12	1E-15
No mux, a = 0.75	5.9E-5*	4.9E-7*
No mux, a = 0.65	2.1E-4*	5.1E-5*
P802.3bs D1.3, a = 0.75	2.3E-4*	7.8E-5*
No mux, a = 0.5	3.1E-4*	1.3E-4*
Random errors	3.8E-4	2.3E-4

Note – these values are the BER **including** the additional errors due to the bursts. To account for burst errors, the values marked with “*” have been multiplied by 4 when a = 0.75, 2.9 when a = 0.65, 2 when a = 0.5.

Multi-part links with FEC

If the FEC bytes are added at the source FEC sublayer and then the correction is applied only at the destination FEC sublayer as in:



Then the worst case input BER for the FEC decoder must be met by the concatenation of all of the sub-links.

In the case of CDAUI-8 -> FR8 -> CDAUI-8, the worst case BER for each lane of the electrical sub-links is 1E-5. Even though there may be two additional CDAUI-8 C2C sub-links, this is tolerated on the basis that it is extremely unlikely that all four sub-links will be at the worst case BER at the same time given that each sub-link BER is averaged over 8 lanes.

The results for multiple sub-links sharing the same RS(544,514) protection is shown on the next slide.

Multi-part link results

The BER of the electrical sub-links for a BER of 2.4E-4 in the optical sub-link are shown in the table below (0.16 dB optical penalty).

An additional row is included for “no mux, a = 0.75” where the electrical BER is 2E-5 to allow 1E-5 for each of two AUI sub-links (0.34 dB optical penalty).

	RS(544,514) FLR = 6.2E-10			
	Electrical		Optical	
No mux, a = 0.75	Burst	6.3E-6*	Random	2.4E-4
No mux, a = 0.75	Burst	2E-5*	Random	1.4E-4
No mux, a = 0.65	Burst	5.7E-5*	Random	2.4E-4
P802.3bs D1.3, a = 0.75	Burst	6.3E-5*	Random	2.4E-4
No mux, a = 0.5	Burst	1E-4*	Random	2.4E-4
Random errors	Random	1.4E-4	Random	2.4E-4

Note – these values are the BER **including** the additional errors due to the bursts. To account for burst errors, the values marked with “*” have been multiplied by 4 when a = 0.75, 2.9 when a = 0.65, 2 when a = 0.5.

Conclusion

Assuming either non-multiplexed or symbol multiplexed RS(544,514) FEC schemes for 50 Gb/s and next generation 100 Gb/s Ethernet:

- If the probability of a burst continuing (a) is allowed to be as high as 0.75, either:
  - the electrical sub-link BER has to be ~ 3E-6 with consequent reduction in capability for the electrical links
  - or the optical BER has to be ~ 1.4E-4 with 0.34 dB optical penalty
- but the margin required to achieve an FLR equivalent to 1E-15 is much larger than expected.

- If the probability of a burst continuing (a) is restricted to 0.65, then an electrical sub-link BER of 1E-5 and an optical sub-link BER of 2.4E-4 as per 400GBASR-R PHYs seems viable.

Thanks!