Interleaver Design for Concatenated Code with the (144,136) Code

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Background

- RS(544,514) has been adopted for 200G/lane AUIs (C2C and C2M).

  - See dambrosia_3dj_01a_230116.pdf and motions_3dfdj_230117.pdf.

- Concatenated code with 4x interleaved RS(544,514) as the outer code is under discussion.

  - bliss_3df_01b_2211, farhood_3df_02b_2211 both proposed BCH/Hamming inner codes with RS outer code.

- Interleaver between outer and inner code can randomize the errors from inner code decoders, improving overall coding gain, as analyzed in bliss_3df_01a_220517.

  - Convolutional interleaver is usually used for block codes to minimize latency for relatively high interleaving depth.

  - Convolutional interleaver with depth of 12 RS codewords was proposed in farhood_3df_02b_2211 for Hamming(128,120).

- A convolutional interleaver for binary code (144,136) is proposed in this contribution.

  - Questions were raised during Bangkok meeting on how to design the convolutional interleaver on this code.

  - Effective interleaver depth is over 12 RS codewords, with latency of 76.8ns (800 GbE).

Things to be Considered when Designing Interleaver

- **Interleaving depth and performance**

  - Hamming(128,120) uses a convolutional interleaver based on number of RS-symbols in an inner code.

  - Convolutional interleaver for (144,136) can work on blocks longer than RS-symbols.

  - Both codes can have high interleaving depth, enough to randomize error distribution from inner code.

- **Supports 200/400/800/1600 GbE**

  - All Ethernet rates that could utilize 200G/lane PMDs should be supported.

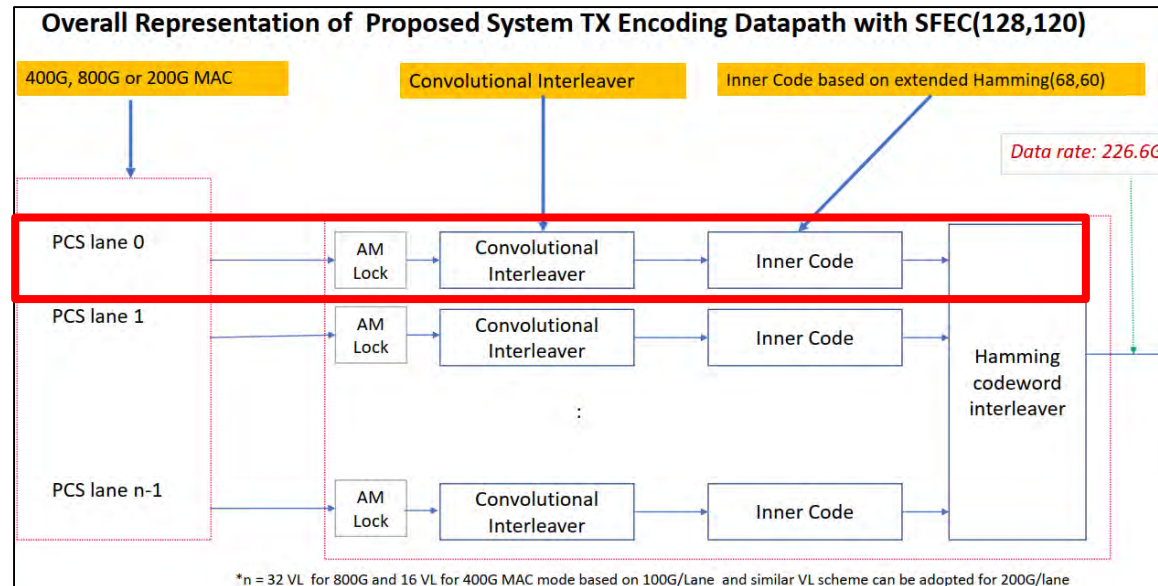 - Interleaver design based on the common part across all rates can simplify implementation and specification.

Ethernet Rate (GbE)	PMD Lane Rate (in .3dj)	Number of RS Codewords
200		2 (or 4?)
400	200G/lane	2 (or 4?)
800		4
1600 *(TBD)*		4

- **Breakout support**

  - Minimize the logic required to support breakout.

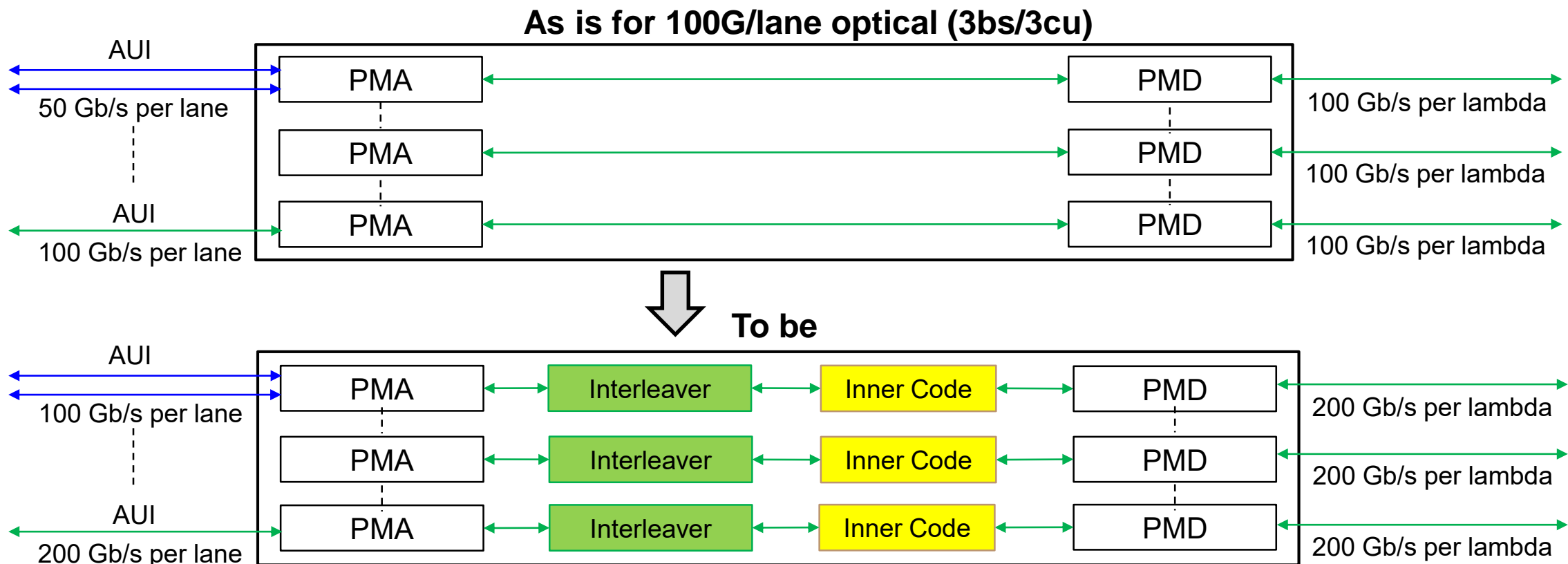# Issues for (128,120) Interleaver: Designed over 25Gb/s PCS Lanes

- Interleaver and encoder per PCS lane design

  - Both interleaver and encoder are performed based on 25Gb/s PCS lanes, which is not forward-looking.

  - 1.6 TbE does not have any reason to use 25Gb/s PCS lanes. 100Gb/s PCS lane is more reasonable. (shrikhande_3dj_01_230206)

  - Padding is proposed to have integer PLL design. (farhood_3dj_01_230206)

- Redesign is required to support potential 100Gb/s PCS lanes.

  - AM locking over 100G/lane PCS is different from 25G/lane.

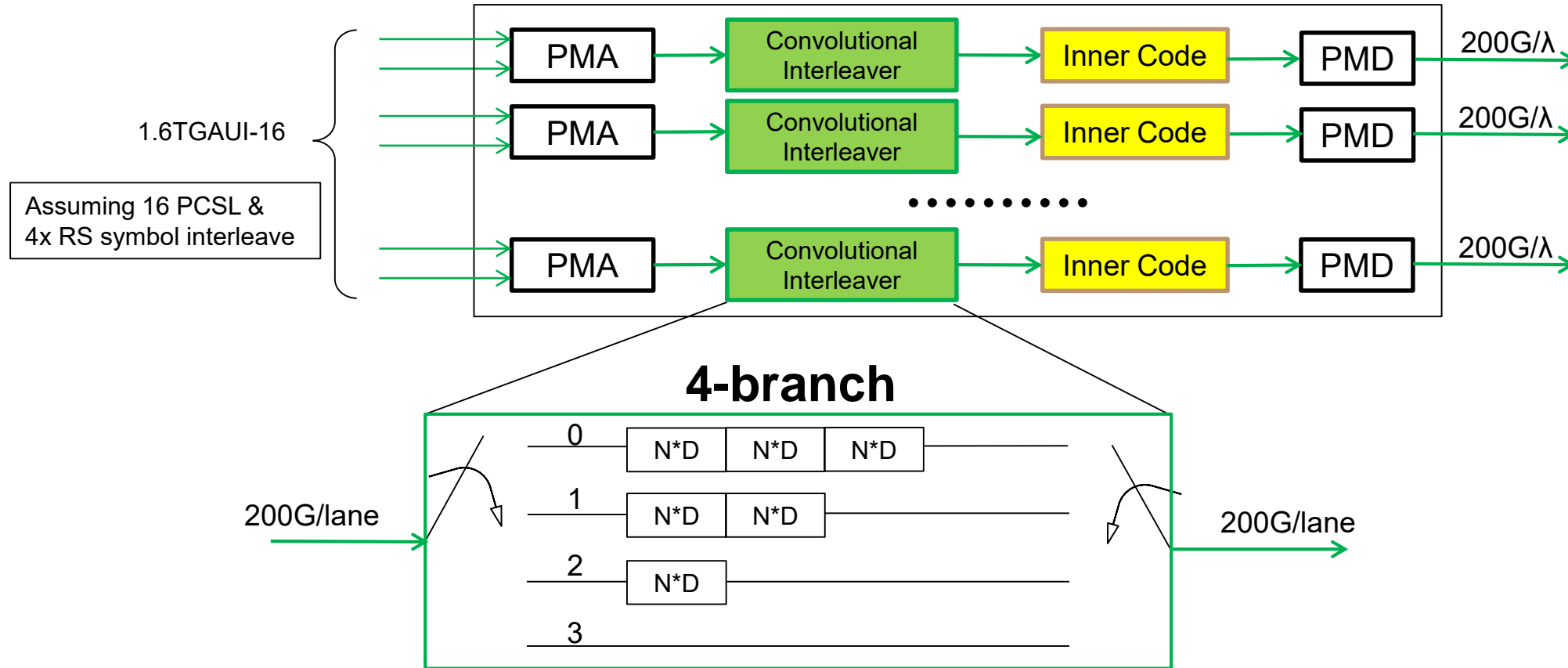 - Convolutional interleaver requires redesign for 100G/lane with different delay parameters.

Overall Representation of Proposed System TX Encoding Datapath with SFEC(128,120)

400G, 800G or 200G MAC   Convolutional Interleaver   Inner Code based on extended Hamming(68,60)

Data rate: 226.6G

PCS lane 0 — AM Lock — Convolutional Interleaver — Inner Code

PCS lane 1 — AM Lock — Convolutional Interleaver — Inner Code

PCS lane n-1 — AM Lock — Convolutional Interleaver — Inner Code

Hamming codeword interleaver

*n = 32 VL for 800G and 16 VL for 400G MAC mode based on 100G/Lane and similar VL scheme can be adopted for 200G/lane

farhood_3df_02b_2211

Breakout Support of Inner Code Could Work at Per Lambda

- Each 200G/lane PMA/PMD in the module has its own inner code encoder(s)/decoder(s)/interleaver(s).

  - Advantage: Naturally supports breakout as no regrouping/distribution is required over multiple lambdas.

  - Works for both 100G/lane and 200G/lane AUIs, supporting 2x100G$\leftrightarrow$2x100G, 2x100G$\leftrightarrow$ 1x200G and 1x200G$\leftrightarrow$1x200G interop.

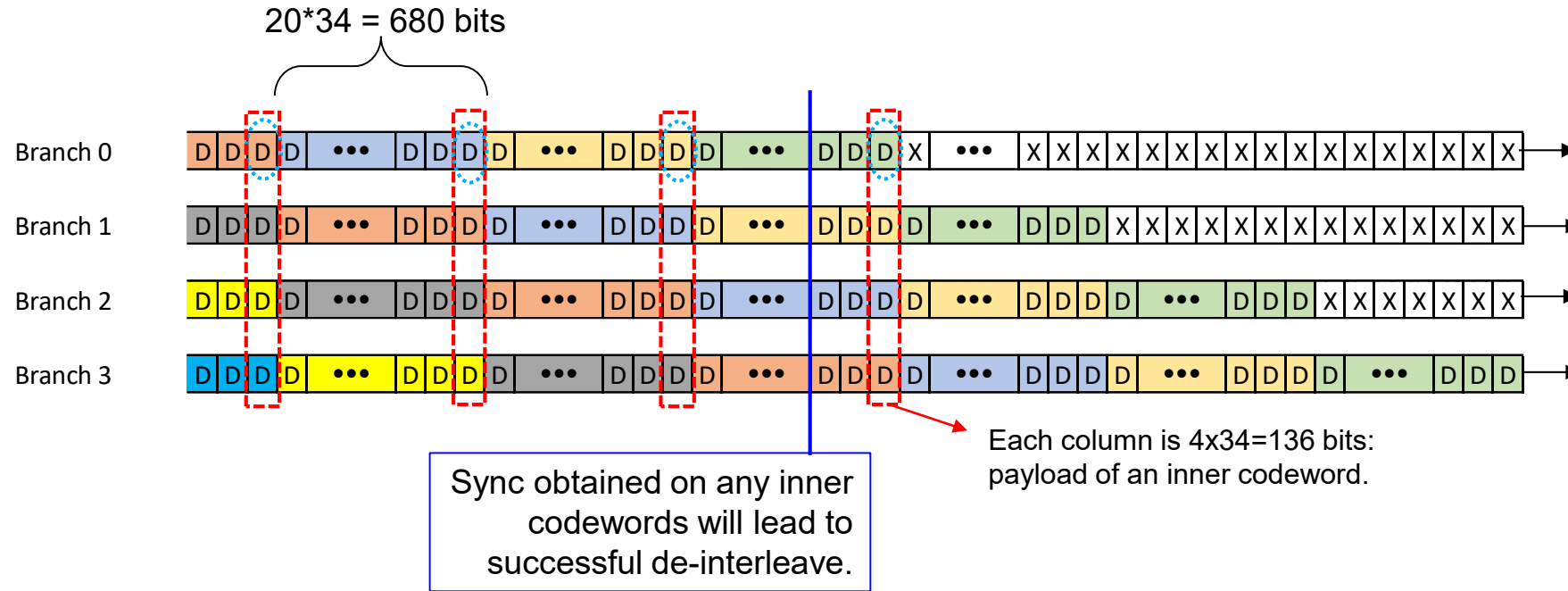**As is for 100G/lane optical (3bs/3cu)**

To be

Convolutional Interleaver Design for Binary (144,136)

- Convolutional interleaver is a general interleaving method that could support any block codes.
  - Different code may have different numbers for N, D and branches.

- A 4-branch convolutional interleaver is proposed for (144,136) code.
  - Round-robin distribution based on D = 34b blocks. N = 2720/34/4 = 20.
  - For each group of 4 codewords in PCS, each convolutional interleaver gets 4*5440/8 = 2720b.

Convolutional Interleaver Design for Binary (144,136), continued

20*34 = 680 bits

Branch 0 | Branch 1 | Branch 2 | Branch 3

Sync obtained on any inner codewords will lead to successful de-interleave.

Each column is 4x34=136 bits: payload of an inner codeword.

- D = 34b, each column of 4 blocks form an inner codeword.
  - If PCS has only 2 codewords, D could be 17 bits and number of branches is increased to 8.
  - Worst case of tailing bits in each D block can still guarantee an equivalent interleaving depth of more than 12 RS codewords.
- Synchronization of inner code can guarantee successful de-interleave.
  - Inner code synchronization can use self-sync methodology similar as in Clause 74.
  - Does not rely on AM from PCS or additional AM inside modules, simplifies logic inside module.

Performance Analysis

- Latency can be evaluated based on number of RS codewords.

  - We recommend to bypass the interleaver for low-latency applications.

Ethernet Rate (GbE)	# of RS CWs in PCS	Interleaver Throughput	# of RS CWs Interleaved	Interleaver Latency, ns	SNR, dB	Pre-FEC BER
200	2		16	358.4		
200	4		16	307.2		
400	2	200G	16	179.2	14.96	4.6E-3*
400	4		16	153.6		
800	4		16	76.8		
1600 *(TBD)*	4		16	38.4		

* Using **sub-optimal** soft-decoding method for faster simulation.
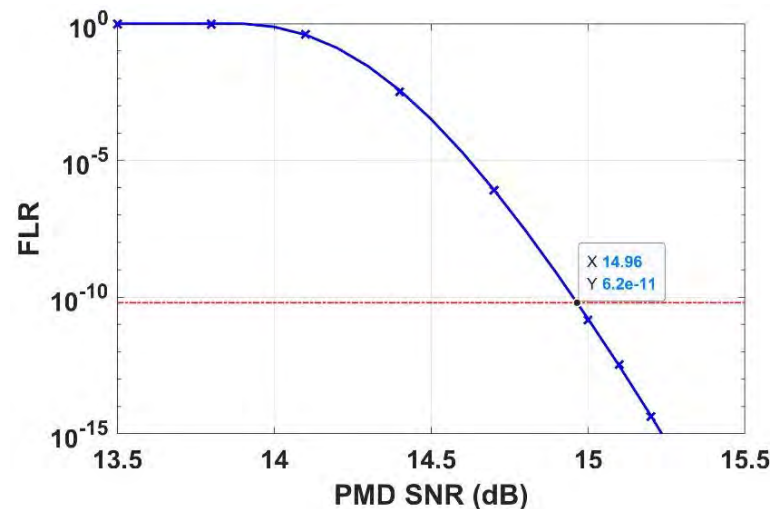
Performance Analysis

Code	Pre-FEC BER	SNR, dB	Code Rate R (relative to 64B/66B)	NCG Penalty 10log(R), dB
(144,136)	4.6E-3	14.96	103.125/112.5	-0.378
(128,120) + padding	4.8E-3*	14.91	103.125/113.4375	-0.414

*From farhood_3df_02b_2211

NCG difference = (14.96 – 14.91) – [-0.378 – (-0.414)] = **0.014 dB**

- Due to 1% more overhead, the NCG of "(128,120)+padding" is only 0.014 dB higher than (144,136).
  - 1% higher overhead leads to **performance degradation** of optical transceivers, as raised in welch_3df_01a_221011.
  - Considering the bandwidth limitation, actual performance needs to be analyzed between 225 Gb/s and 226.875 Gb/s for (144,136) and (128,120)+ padding, respectively.

- 1% Higher data rate also leads to **higher power** (optical, AD/DA, etc).
  - Potentially impact future CPO and NPO applications where inner code could be integrated in ASIC.
  - Additional optical transceiver power due to higher overhead could be significantly more than the inner FEC decoder power.
  - It is more economic if we allocate this additional power to boost the soft-decoding gain.
    - Using **more optimized** soft-decoding method will increase the over all coding gain by more than 0.014dB.

Future Integration of Concatenated Code Considerations

- The (128,120) code will result in more complex design.

  - (128,120) has a factor of 3 in the divisor that requires a frac-N PLL.

    - For an oDSP for 4 or 8 lanes, it is not a big deal.

    - For highly integrated ASIC (e.g. 512-lane digital switching chip) it will complicate things.

      - Dividing reference clock (156.25MHz) by 3 will cause worse jitter.

    - Combined with higher power due to 1% higher overhead, it can be problematic for CPO or NPO.

- The (144,136) code enables integer PLL design and has lower power.

Conclusions

- A convolutional interleaver for binary code (144,136) is proposed.

  - It does not rely on 25G/lane PCS lanes.

  - It does not rely on additional alignment method to de-interleave.

  - It supports breakout.

  - The overall performance is on par comparing with Hamming(128,120) + padding.

  - The overall power is lower than (128,120) + padding, due to simpler design and lower data rate.

- We propose to adopt binary code (144,136) as the inner code for concatenated code for 200G/lane optical PMDs.

  - The code supports integer PLL without additional padding.

  - The code is friendly to implementation to both oDSP (for pluggables) and host ASIC (for CPO).

Thank you