

Interleaver Design for Concatenated Code with the (144,136) Code

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Background

- RS(544,514) has been adopted for 200G/lane AUIs (C2C and C2M).
 - See [dambrosia_3dj_01a_230116.pdf](#) and [motions_3dfdj_230117.pdf](#).
- Concatenated code with 4x interleaved RS(544,514) as the outer code is under discussion.
 - [bliss_3df_01b_2211](#), [farhood_3df_02b_2211](#) both proposed BCH/Hamming inner codes with RS outer code.
- Interleaver between outer and inner code can randomize the errors from inner code decoders, improving overall coding gain, as analyzed in [bliss_3df_01a_220517](#).
 - Convolutional interleaver is usually used for block codes to minimize latency for relatively high interleaving depth.
 - Convolutional interleaver with depth of 12 RS codewords was proposed in [farhood_3df_02b_2211](#) for Hamming(128,120).
- A convolutional interleaver for binary code (144,136) is proposed in this contribution.
 - Questions were raised during Bangkok meeting on how to design the convolutional interleaver on this code.
 - Effective interleaver depth is over 12 RS codewords, with latency of 76.8ns (800 GbE).

Things to be Considered when Designing Interleaver

- **Interleaving depth and performance**

- Hamming(128,120) uses a convolutional interleaver based on number of RS-symbols in an inner code.
- Convolutional interleaver for (144,136) can work on blocks longer than RS-symbols.
- Both codes can have high interleaving depth, enough to randomize error distribution from inner code.

- **Supports 200/400/800/1600 GbE**

- All Ethernet rates that could utilize 200G/lane PMDs should be supported.
- Interleaver design based on the common part across all rates can simplify implementation and specification.

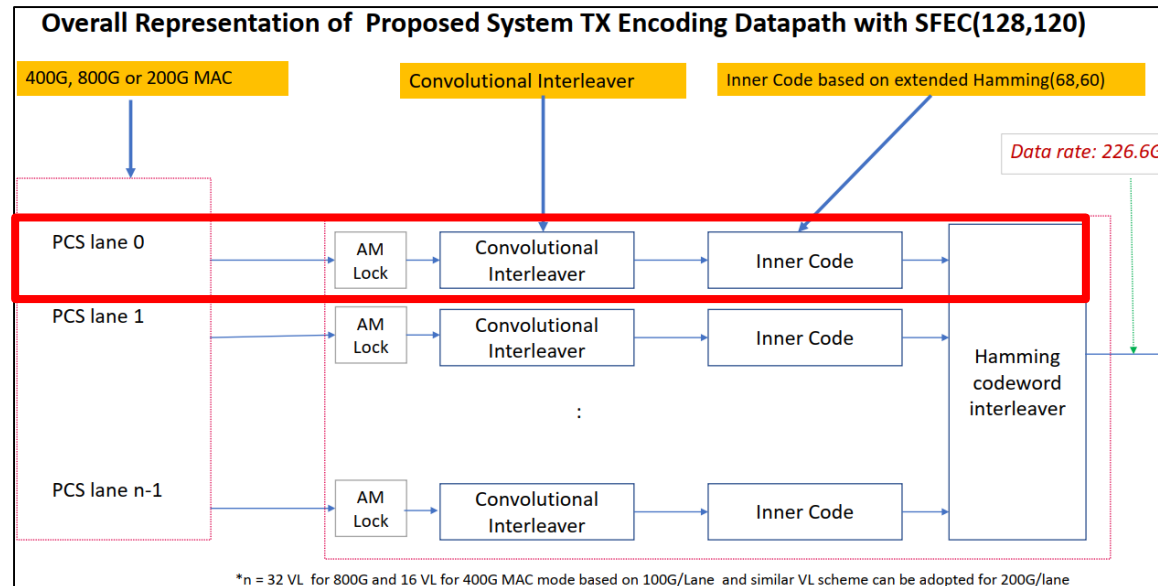
Ethernet Rate (GbE)	PMD Lane Rate (in .3dj)	Number of RS Codewords
200	200G/lane	2 (or 4?)
400		2 (or 4?)
800		4
1600 (TBD)		4

- **Breakout support**

- Minimize the logic required to support breakout.

Issues for (128,120) Interleaver: Designed over 25Gb/s PCS Lanes

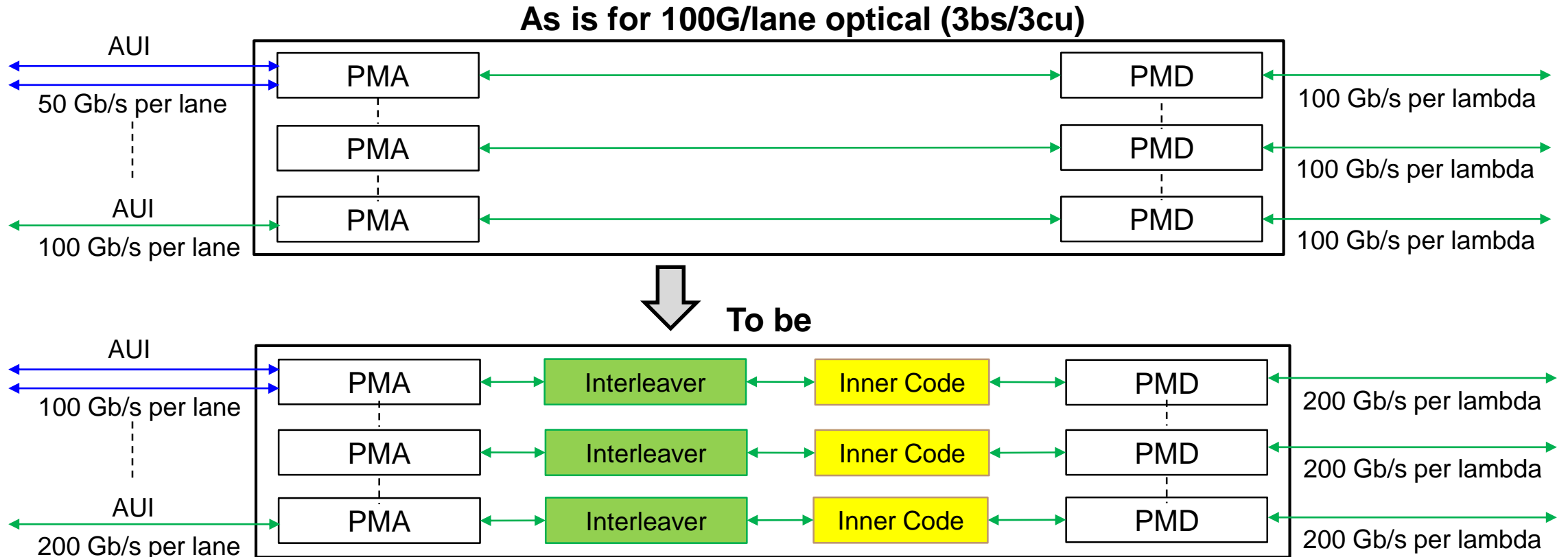
- Interleaver and encoder per PCS lane design
 - Both interleaver and encoder are performed based on 25Gb/s PCS lanes, which is not forward-looking.
 - 1.6 TbE does not have any reason to use 25Gb/s PCS lanes. 100Gb/s PCS lane is more reasonable. ([gustlin 3dj 01b 230206](#))
 - Padding is proposed to have integer PLL design. ([farhood 3dj 01a 230206](#))
- Redesign is required to support potential 100Gb/s PCS lanes.
 - AM locking over 100G/lane PCS is different from 25G/lane.
 - Convolutional interleaver requires redesign for 100G/lane with different delay parameters.



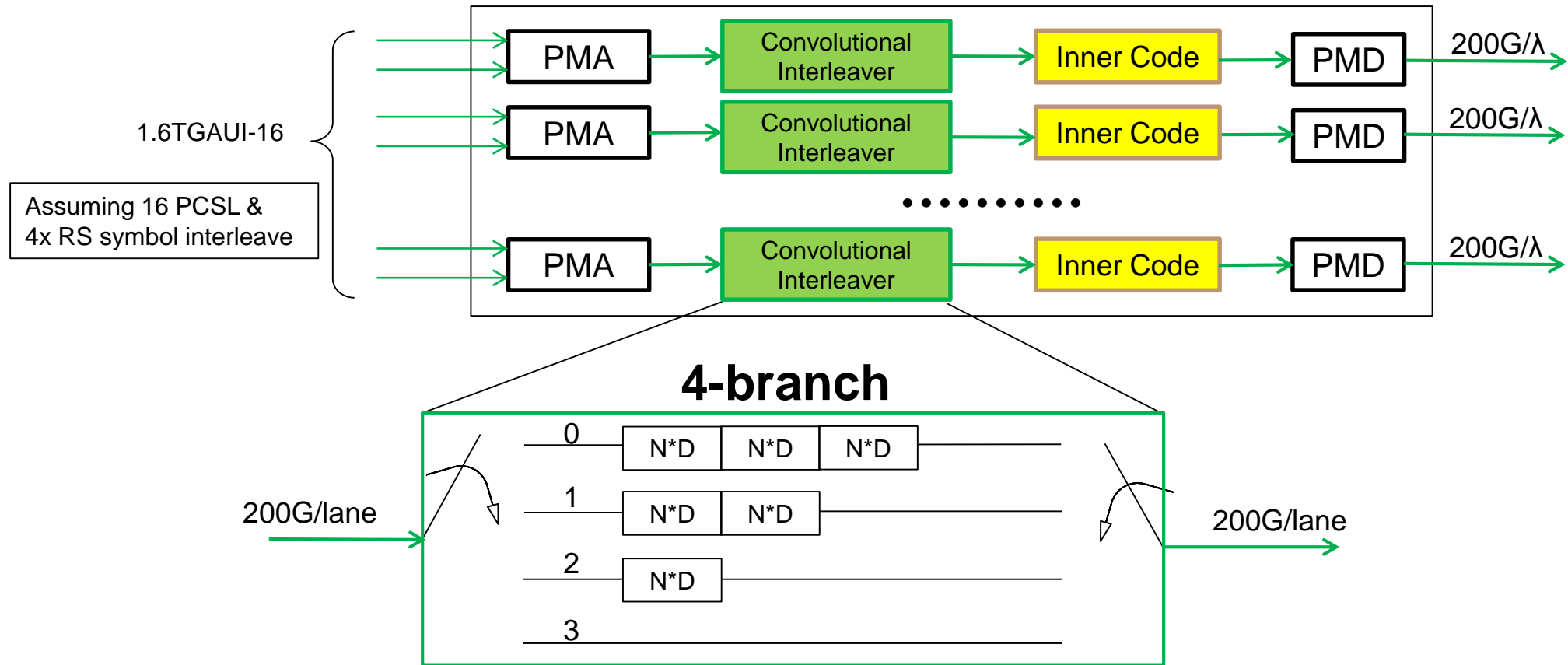
[farhood 3df 02b 2211](#)

Breakout Support of Inner Code Could Work at Per Lambda

- Each 200G/lane PMA/PMD in the module has its own inner code encoder(s)/decoder(s)/interleaver(s).
 - Advantage: Naturally supports breakout as no regrouping/distribution is required over multiple lambdas.
 - Works for both 100G/lane and 200G/lane AUIs, supporting $2 \times 100\text{G} \leftrightarrow 2 \times 100\text{G}$, $2 \times 100\text{G} \leftrightarrow 1 \times 200\text{G}$ and $1 \times 200\text{G} \leftrightarrow 1 \times 200\text{G}$ interop.

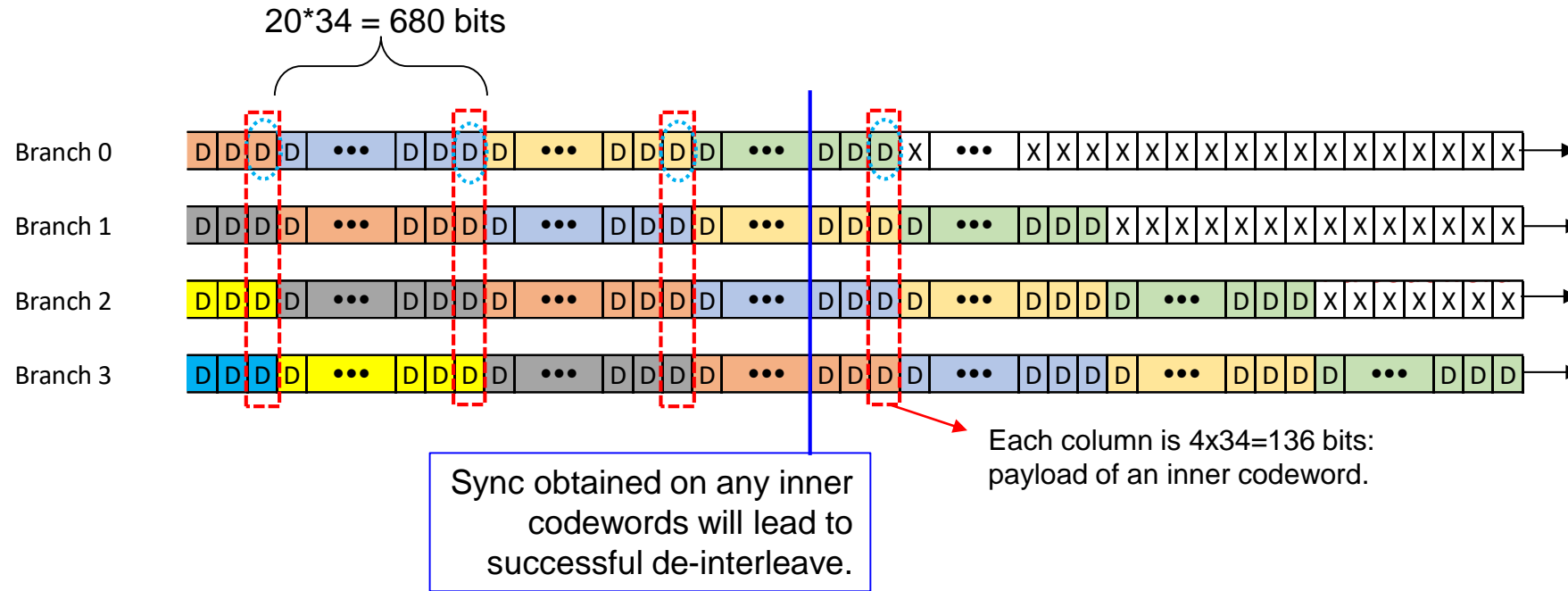


Convolutional Interleaver Design for Binary (144,136)



- Convolutional interleaver is a general interleaving method that could support any block codes.
 - Different code may have different numbers for N, D and branches.
- A 4-branch convolutional interleaver is proposed for (144,136) code.
 - Round-robin distribution based on $D = 34b$ blocks. $N = 2720/34/4 = 20$.
 - For each group of 4 codewords in PCS, each convolutional interleaver gets $4 \cdot 5440/8 = 2720b$.

Convolutional Interleaver Design for Binary (144,136), continued



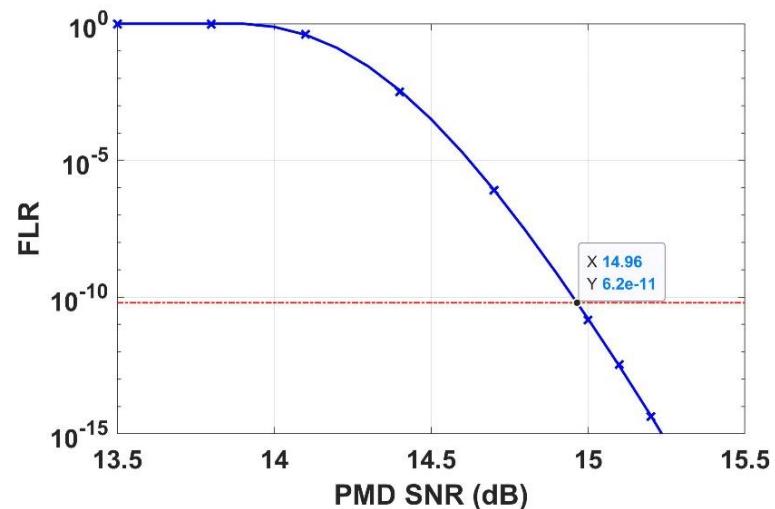
- $D = 34b$, each column of 4 blocks form an inner codeword.
 - If PCS has only 2 codewords, D could be 17 bits and number of branches is increased to 8.
 - Worst case of tailing bits in each D block can still guarantee an equivalent interleaving depth of more than 12 RS codewords.
- Synchronization of inner code can guarantee successful de-interleave.
 - Inner code synchronization can use self-sync methodology similar as in Clause 74.
 - Does not rely on AM from PCS or additional AM inside modules, simplifies logic inside module.

Performance Analysis

- Latency can be evaluated based on number of RS codewords.
 - We recommend to bypass the interleaver for low-latency applications.

Ethernet Rate (GbE)	# of RS CWs in PCS	Interleaver Throughput	# of RS CWs Interleaved	Interleaver Latency, ns	SNR, dB	Pre-FEC BER
200	2	200G	16	358.4	14.96	4.6E-3*
200	4		16	307.2		
400	2		16	179.2		
400	4		16	153.6		
800	4		16	76.8		
1600 (TBD)	4		16	38.4		

* Using **sub-optimal** soft-decoding method for faster simulation.



Performance Analysis

Code	Pre-FEC BER	SNR, dB	Code Rate R (relative to 64B/66B)	NCG Penalty 10log(R), dB
(144,136)	4.6E-3	14.96	103.125/112.5	-0.378
(128,120) + padding	4.8E-3*	14.91	103.125/113.4375	-0.414

*From [farhood_3df_02b_2211](#)

$$\text{NCG difference} = (14.96 - 14.91) - [-0.378 - (-0.414)] = \mathbf{0.014 \text{ dB}}$$

- Due to 1% more overhead, the NCG of “(128,120)+padding” is only 0.014 dB higher than (144,136).
 - 1% higher overhead leads to **performance degradation** of optical transceivers, as raised in [welch_3df_01a_221011](#).
 - Considering the bandwidth limitation, actual performance needs to be analyzed between 225 Gb/s and 226.875 Gb/s for (144,136) and (128,120)+ padding, respectively.
- 1% Higher data rate also leads to **higher power** (optical, AD/DA, etc).
 - Potentially impact future CPO and NPO applications where inner code could be integrated in ASIC.
 - Additional optical transceiver power due to higher overhead could be significantly more than the inner FEC decoder power.
 - It is more economic if we allocate this additional power to boost the soft-decoding gain.
 - Using **more optimized** soft-decoding method will increase the over all coding gain by more than 0.014dB.

Future Integration of Concatenated Code Considerations

- The (128,120) code will result in more complex design.
 - (128,120) has a factor of 3 in the divisor that requires a frac-N PLL.
 - For an oDSP for 4 or 8 lanes, it is not a big deal.
 - For highly integrated ASIC (e.g. 512-lane digital switching chip) it will complicate things.
 - Dividing reference clock (156.25MHz) by 3 will cause worse jitter.
 - Combined with higher power due to 1% higher overhead, it can be problematic for CPO or NPO.
- The (144,136) code enables integer PLL design and has lower power.

Conclusions

- A convolutional interleaver for binary code (144,136) is proposed.
 - It does not rely on 25G/lane PCS lanes.
 - It does not rely on additional alignment method to de-interleave.
 - It supports breakout.
 - The overall performance is on par comparing with Hamming(128,120) + padding.
 - The overall power is lower than (128,120) + padding, due to simpler design and lower data rate.
- We propose to adopt binary code (144,136) as the inner code for concatenated code for 200G/lane optical PMDs.
 - The code supports integer PLL without additional padding.
 - The code is friendly to implementation to both oDSP (for pluggables) and host ASIC (for CPO).

Thank you